

1 Existential quantification

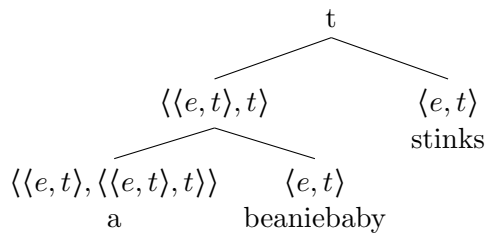
$\llbracket \text{a beaniebaby stinks} \rrbracket$

Lexical entries:

1. $\llbracket \text{beaniebaby} \rrbracket = \lambda y[BEANIEBABY(y)]$
2. $\llbracket \text{stink} \rrbracket = \lambda z[STINK(z)]$
3. $\llbracket \text{a} \rrbracket = \lambda f_{\langle e, t \rangle}[\lambda g_{\langle e, t \rangle}[\exists x[f(x) \wedge g(x)]]]$

★ If you used other variables for any lexical entry, that's also right! e.g., $\lambda z[BEANIEBABY(z)] = \lambda y[BEANIEBABY(y)]$, etc.

Tree annotated with types:



★ Common mistake with types: be careful with how many brackets you have, and where!

- ➡ $\langle \langle e, t \rangle, \langle \langle e, t \rangle, t \rangle \rangle$: ill-formed! A closing bracket is missing.
- ➡ $\langle e, t, t \rangle$: ill-formed! There should be one input, one output for each pair of brackets (e.g., $\langle \langle e, t \rangle, t \rangle$).
- ➡ $\langle e, \langle t, t \rangle \rangle$: not the same thing as $\langle \langle e, t \rangle, t \rangle$!
- ➡ $\langle \langle \langle e, t \rangle, \langle e, t \rangle \rangle, t \rangle$: not the same thing as $\langle \langle e, t \rangle, \langle \langle e, t \rangle, t \rangle \rangle$!

Step-by-step computation:

1. $\llbracket \text{a beaniebaby} \rrbracket$

- (a) $= \llbracket \text{a} \rrbracket (\llbracket \text{beaniebaby} \rrbracket)$
- (b) $= \lambda f_{\langle e, t \rangle} [\lambda g_{\langle e, t \rangle} [\exists x [f(x) \wedge g(x)]]] (\llbracket \text{beaniebaby} \rrbracket)$
- (c) $= \lambda g_{\langle e, t \rangle} [\exists x [\llbracket \text{beaniebaby} \rrbracket (x) \wedge g(x)]]$
- (d) $= \lambda g_{\langle e, t \rangle} [\exists x [\lambda y [BEANIEBABY(y)](x) \wedge g(x)]]$
- (e) $= \lambda g_{\langle e, t \rangle} [\exists x [BEANIEBABY(x) \wedge g(x)]]$

2. $\llbracket \text{a beaniebaby stinks} \rrbracket$

- (a) $= \llbracket \text{a beaniebaby stink} \rrbracket$
- (b) $= \llbracket \text{a beaniebaby} \rrbracket (\llbracket \text{stink} \rrbracket)$
- (c) $= \lambda g_{\langle e, t \rangle} [\exists x [BEANIEBABY(x) \wedge g(x)]] (\llbracket \text{stink} \rrbracket)$
- (d) $= \exists x [BEANIEBABY(x) \wedge \llbracket \text{stink} \rrbracket (x)]$
- (e) $= \exists x [BEANIEBABY(x) \wedge \lambda z [STINK(z)](x)]$
- (f) $= \text{T iff } \exists x [BEANIEBABY(x) \wedge STINK(x)]$

★ Common mistake: Don't forget to state the very final step as a truth condition! Remember, what you're calculating is the denotation of a sentence, which is its truth value/condition.

2 Types summary

1. $\langle e, t \rangle$: type for one-place predicates
2. $\langle \langle e, t \rangle, t \rangle$: type for quantifiers (also sometimes called generalized quantifiers¹); e.g., *every beaniebaby*, *everyone*, *a/some beaniebaby*, *someone*, *all beaniebabies*, *few/most beaniebabies*, etc.)
3. $\langle \langle e, t \rangle, \langle \langle e, t \rangle, t \rangle \rangle$: type for quantificational determiners; e.g., *every*, *all*, *a*, *some*, *few*, *most*, etc.

 What other types are possible? What might that kind of function do?

¹It's called "generalized" because at a certain point in history, logicians/semanticists discovered that all quantificational phrases (e.g., many beaniebabies, two beaniebabies, few beaniebabies) are all the same thing: a relation between sets. The treatment of those quantifiers was "generalized" from the analysis of \forall and \exists , hence, "generalized".