

A Few Open Problems in Neural Theorem Proving

(in Lean)

Sean Welleck

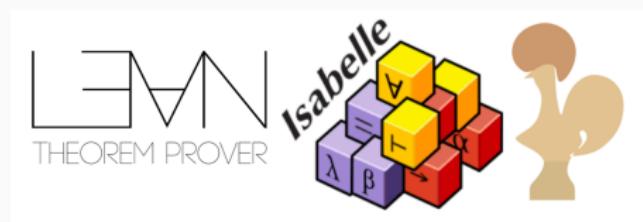
September 5, 2024

Carnegie Mellon University

Neural theorem proving

Use neural networks to:

- Generate proofs in an interactive proof assistant



Neural theorem proving | Rapid progress

Rapid progress in methods based on language models:

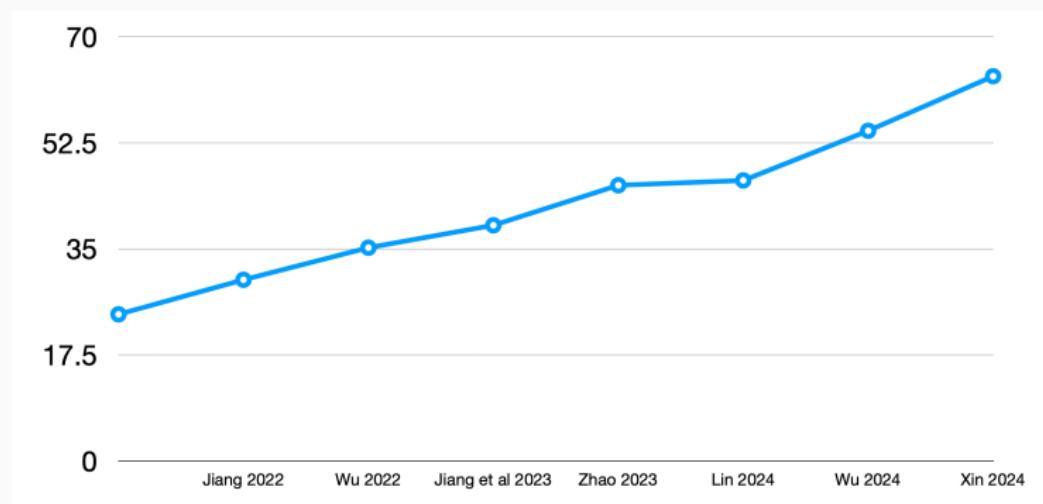


Figure 1: miniF2F benchmark performance, 2022-2024

Neural theorem proving | Rapid progress

```
theorem imo_1960_p2 (x : ℝ) (h₀ : 0 ≤ 1 + 2 * x) (h₁ : (1 - Real.sqrt (1 + 2 * x)) ^ 2 ≠ 0)
  (h₂ : 4 * x ^ 2 / (1 - Real.sqrt (1 + 2 * x)) ^ 2 < 2 * x + 9) : -(1 / 2) ≤ x ∧ x < 45 / 8 := by
  norm_num at h₀ h₁ h₂
  have h₃ : 0 ≤ 1 + 2 * x := by linarith
  have h₄ : 0 < 1 + Real.sqrt (1 + 2 * x) := by
    nlinarith [Real.sqrt_nonneg (1 + 2 * x)]
  have h₅ : 4 * x ^ 2 / (1 - Real.sqrt (1 + 2 * x)) ^ 2 < 2 * x + 9 := by
    linarith
  have h₆ : 1 - Real.sqrt (1 + 2 * x) ≠ 0 := by
    intro h
    apply h₁
    nlinarith
  have h₇ : 4 * x ^ 2 / (1 - Real.sqrt (1 + 2 * x)) ^ 2 = (1 + Real.sqrt (1 + 2 * x)) ^ 2 := by
    field_simp [h₆]
    nlinarith [sq_sqrt (show 0 ≤ 1 + 2 * x by linarith)]
  rw [h₇] at h₅
  constructor < ;> nlinarith [sq_sqrt (show 0 ≤ 1 + 2 * x by linarith)]
```

Figure 2: Generated International Math Olympiad solution in Lean
(DeepSeek Prover-1.5B, Xin et al 2024)

Why talk about Lean?

- Increasing interest from the mathematical community
- Increasing interest from the AI community
- For AI research, the choice of proof assistant matters (**not ideal!**)

This talk

3 open problems in neural theorem proving in Lean:

- Going beyond human data
- Going beyond competition problems
- Going beyond mathematics

1. Going beyond human data

Language model-based proving:

- **Train** a model $p_\theta(y|x)$ on a dataset $\mathcal{D} = \{(x, y)\}$, e.g.,
 - x : proof state
 - y : next tactic (next “step”)
 - \mathcal{D} : extracted from human-written theorems and proofs

1. Going beyond human data

Language model-based proving:

- Train a model $p_\theta(y|x)$ on a dataset $\mathcal{D} = \{(x, y)\}$, e.g.,
 - x : proof state
 - y : next tactic (next “step”)
 - \mathcal{D} : extracted from human-written theorems and proofs
- Generate proofs:

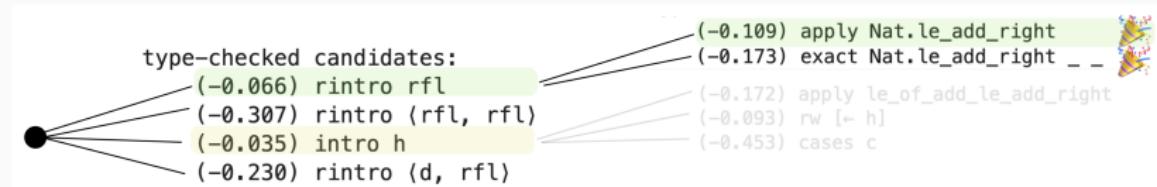


Figure 3: Best-first search

1. Going beyond human data

- Some models are already trained on \approx all Lean projects!
 - E.g., Lean-GitHub [5]: data from 237 Lean 4 repos
- More human-written data will help, but difficult to scale¹

¹Please don't stop making more publicly available formal mathematics data!

1. Going beyond human data

Open problem I: how do we *synthesize* useful data?

- Proofs
- Theorems
- Augmentations (formal, informal, ...)
- ...

1. Going beyond human data

Not a new problem; common methods:

- Statement autoformalization [Wu et al 2022 [4]]
 - Informal theorem → formal theorem
- Expert iteration [Polu et al 2022 [3]]
 - Generate proofs with a model, train on successful ones, iterate

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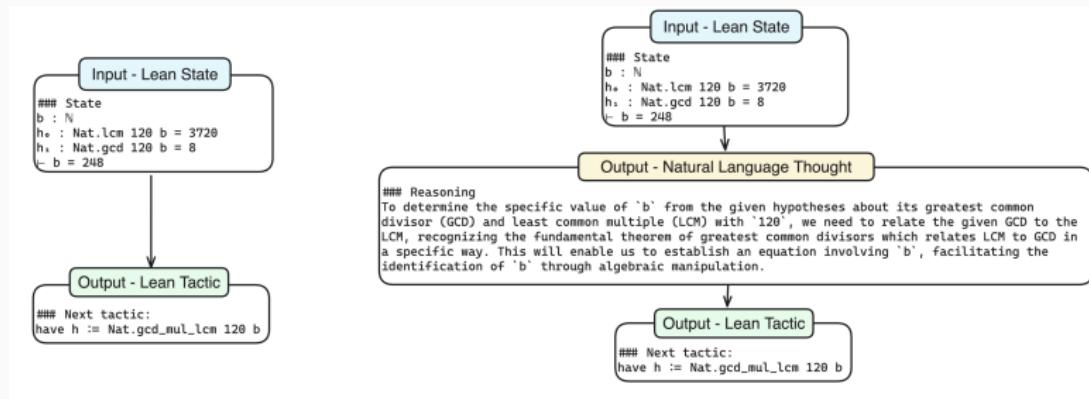
Used in several state-of-the-art methods, e.g. DeepSeek-Prover 1.5, AlphaProof

1. Going beyond human data | Lean-STaR

Lean-STaR: Learning to Interleave Thinking and Proving
Haohan Lin, Zhiqing Sun, Yiming Yang, Sean Welleck

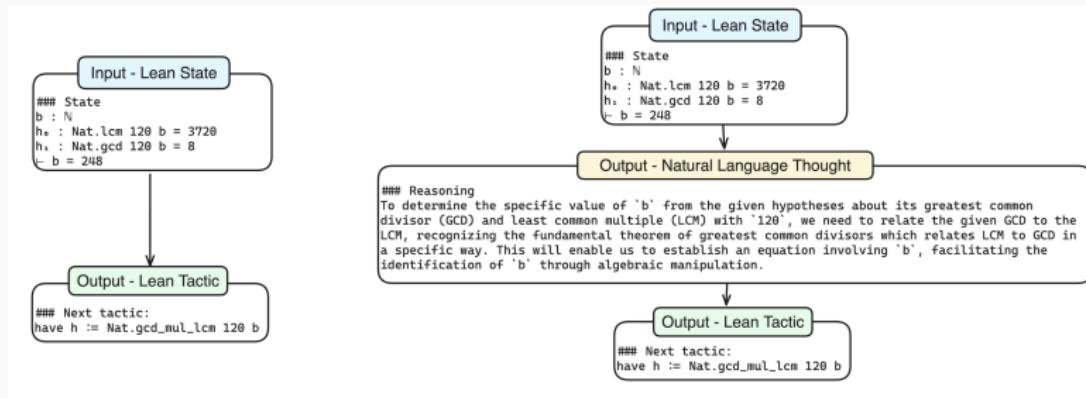
<https://arxiv.org/abs/2407.10040>

1. Going beyond human data | Lean-STaR



Can we do better by interleaving *informal* steps of reasoning? (right)

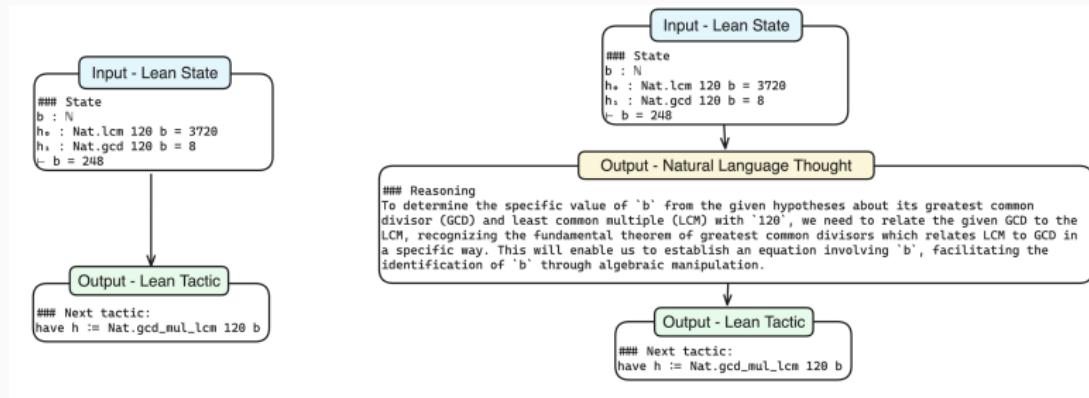
1. Going beyond human data | Lean-STaR



Why?

- Plan proof steps
- Diversify search space
- More tokens can give more computational capacity

1. Going beyond human data | Lean-STaR



Why?

- Plan proof steps
- Diversify search space
- More tokens can give more computational capacity

Data doesn't exist! We need to synthesize it.

1. Going beyond human data | Lean-STaR

Lean-STaR (Self-taught reasoner²)

Step 1: generate an informal “thought” with an off-the-shelf language model retrospectively

- (state, tactic) → thought

Train an initial model on a dataset of such examples:

- $p_\theta^0(\text{thought}, \text{tactic}|\text{state})$

²Inspired by STaR: Bootstrapping Reasoning with Reasoning, Zelikman et al 2022

1. Going beyond human data | Lean-STaR

Step 2: generate proofs with the model

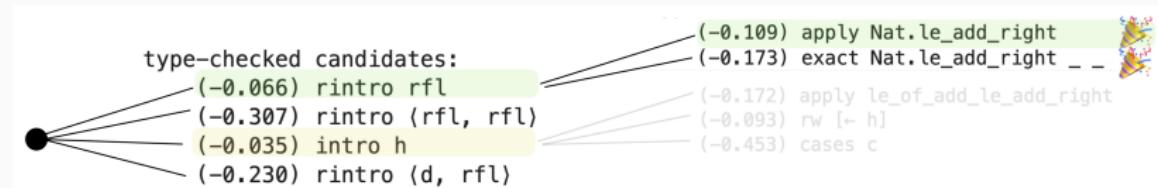


Figure 4: Best-first search: difficult to score (thought, tactic) candidates

1. Going beyond human data | Lean-STaR

Step 2: generate proofs with the model

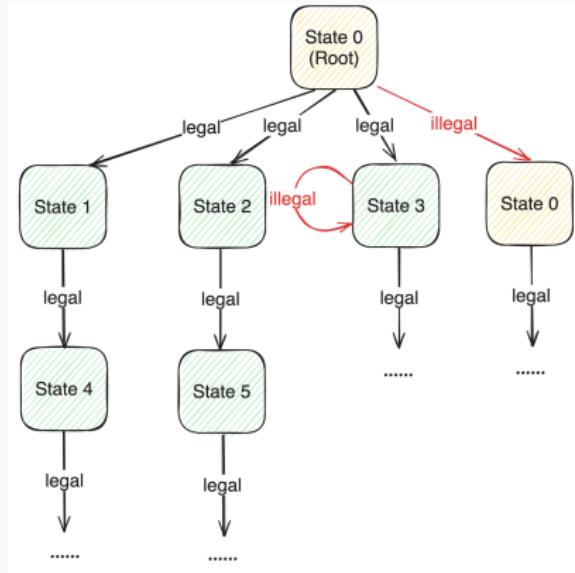


Figure 5: New sampling method

1. Going beyond human data | Lean-STaR

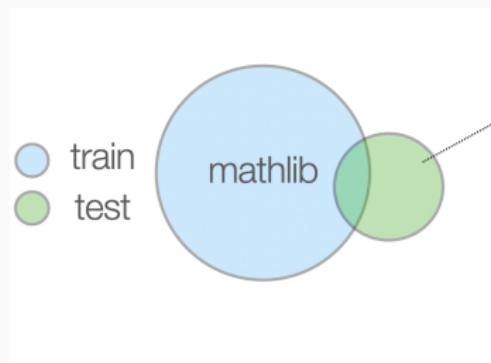
Step 3: train on the successful proofs, and repeat:³

- Collect (state, thought, tactic) from successful proofs
- Train a new model $p_\theta^1(\text{thought}, \text{tactic}|\text{state})$
- Generate proofs
- ...

³I.e. Expert Iteration [Polu et al 2022 [3]]

1. Going beyond human data | Lean-STaR

- miniF2F [7]: competition problems (AMC, AIME, IMO)



A Venn diagram consisting of two overlapping circles. The left circle is light blue and labeled "train". The right circle is light green and labeled "test". They overlap in the center.

Problem 1959 IMO Problems/Problem 1

Prove that the fraction $\frac{21n + 4}{14n + 3}$ is irreducible for every natural number n .

theorem imo_1959_p1
(n : ℕ)
(h₀ : 0 < n) :
Nat.gcd (21*n + 4) (14*n + 3) = 1 := by sorry

1. Going beyond human data | Lean-STaR

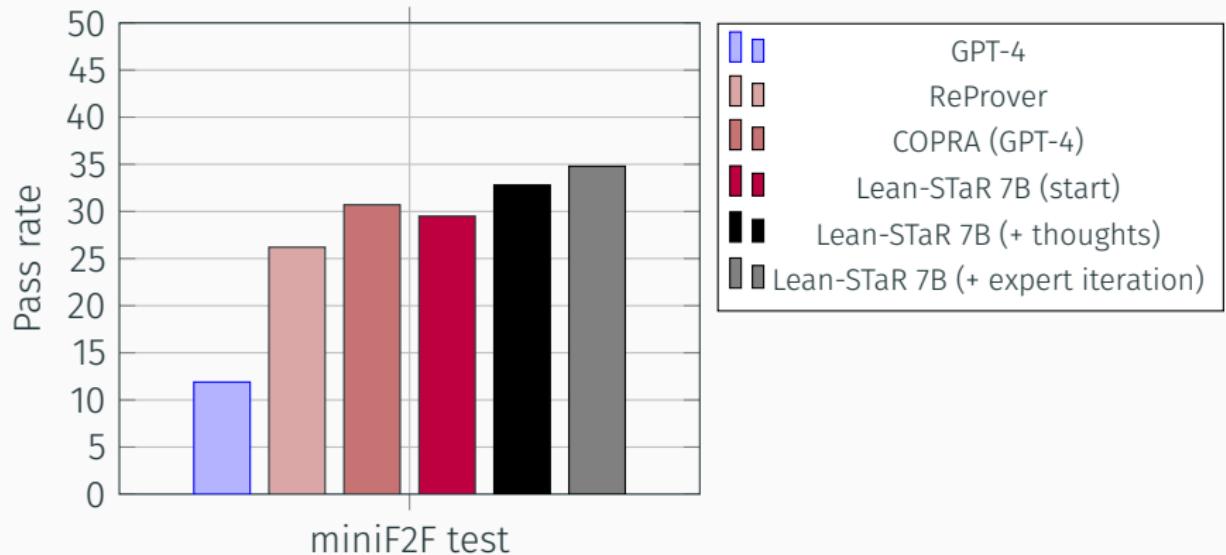


Figure 6: MiniF2F test

1. Going beyond human data | Lean-STaR

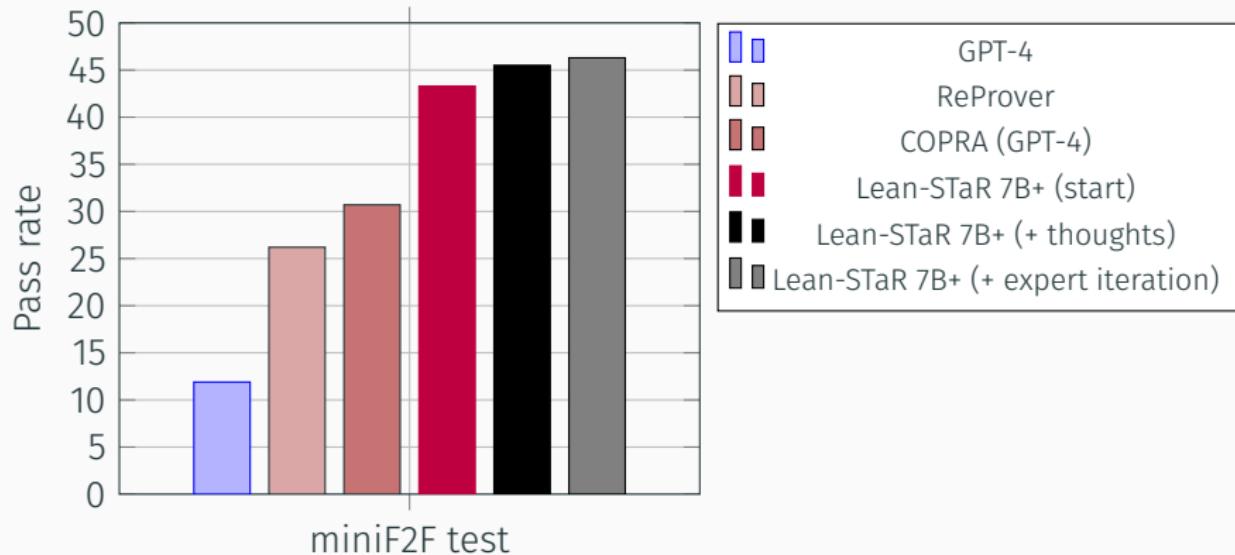


Figure 7: MiniF2F test

1. Going beyond human data | Lean-STaR

```
theorem amc12a_2002_p6 (n : ℕ) (h₀ : 0 < n)
  : ∃ m, m > n ∧ ∃ p, m * p ≤ m + p :=
use n + 1, by linarith only [h₀] -- Reasoning: To demonstrate the existence of
such an `m` and `p` that satisfy the given inequalities, we leverage the
property of natural numbers that ensures an increase in `n` from `0` leads to
the possibility of finding such `m` and `p`. This approach is supported by the
fundamental property of natural numbers and their arithmetic operations,
suggesting that for a given `n` greater than `0`, there exists a successor `m`
and a number `p` that satisfy the conditions outlined.
exact ⟨0, by simp ⟩ -- Reasoning: To demonstrate that there exists an  $\{n \in \mathbb{N}\}$ 
such that  $\{(n + 1) * p \leq n + 1 + p\}$  holds for some  $\{p \in \mathbb{N}\}$ , we
recognize that setting  $\{p = 0\}$  satisfies the inequality by the properties of
zero in arithmetic. This is because  $\{n + 1 > 0\}$  by assumption, ensuring the
left-hand side is non-negative and the right-hand side simplifies to  $\{n + 1\}$ , making the inequality trivially true.
```

Figure 8: Example generated thoughts and proof from Lean-STaR

1. Going beyond human data | Lean-STaR

```
theorem aimes_1983_p2 (x p : ℝ) (f : ℝ → ℝ) (h₀ : 0 < p ∧ p < 15) (h₁ : p ≤ x ∧ x
≤ 15)
  (h₂ : f x = abs (x - p) + abs (x - 15) + abs (x - p - 15))
  (h₃ : 15 ≤ f x :=)
cases1 h₁ with h₂ h₃ -- Reasoning: Given a function '|f|' defined over the real
numbers, where '|f(x)|' equals the sum of the absolute values of the
differences '(x - p)', '|(x - 15)|', and '|(x - p - 15)|', along with conditions
on '|f(p)|' and '|f(15)|' that ensure '|f(x)|' lies between '|f(p)|' and '|f(15)|' and '|f(p)|'
is positive, we aim to show that '|f(x)| ≥ 15|. To demonstrate this, we
consider all possibilities for the contributions of '|f(p)|' relative to '|f(15)| and
'|f(15)|, specifically at '|f(p)| itself, between '|f(p)| and '|f(15)|, and at '|f(15)|.
By evaluating '|f(x)|' for each case, we can establish a lower bound of
'|f(x)| that is greater than or equal to '|f(15)|.
simp [h₀, abs_of_nonneg 2 h₁], abs_of_nonneg (sub_nonneg 2 h₂, *)
-- Reasoning: The objective is to show that '|f(x)| is at least 15 for '|(x)| within the bounds '|p| (ie  $x \leq 15$ )', where '|(p)| is positive but less than
15. The provided expression '|f(x)| suggests the properties of absolute
values and inequalities. By simplifying '|f(x)| using the properties of absolute
value functions, we can demonstrate the desired inequality '|f(x)| ≥ 15|.
simp only [abs_of_nonneg] at * -- Reasoning: The objective is to establish an
inequality involving an expression with absolute values and a sum, where all
terms are functions of '|(x)|. By recognizing that the absolute value functions
are always non-negative, we can simplify the given expression by removing the
absolute value symbols, thus focusing on the arithmetic manipulation of the
terms to align both sides of the inequality.
tag only [abs_of_nonneg, sub.add_eq.add_sub] at * -- Reasoning: The objective
is to demonstrate that the expression ' $15 \leq x - p + |x - 15| + |x - p - 15|$ ' holds given the conditions on 'p', 'x', and the function 'f'. Since the
expression in question involves absolute values which inherently involve
non-negativity, manipulating these expressions to highlight their
non-negativity or equality is a logical initial step. This involves
recognizing that the absolute value of a difference can be rewritten as the
difference of absolute values, which simplifies the manipulation of the
expression to prove the inequality.
congr 1 -- Reasoning: To establish the inequality ' $15 \leq x - |x - 15| + |x - p - 15| - p$ ', we recognize that the left-hand side of the inequality involves a
mix of additions and subtractions of real numbers. By introducing a congruence
transformation that simplifies operations involving absolute values, we aim to
align both sides of the inequality in a form that directly reflects the
properties of real numbers under these operations.
simp only [abs_of_nonpos (mult_pos_of_le h₂), abs_of_nonneg (sub_nonneg_of_le
h₃)], abs_of_nonpos -- Reasoning: The objective is to prove an inequality involving absolute
values and subtractions that aligns the given conditions and constraints.
Simplification using the properties of absolute values and the provided bounds
can directly lead to the desired conclusion by expressing the absolute values
in terms of their underlying expressions and applying the given constraints.
simp [abs_of_pos, *] -- Reasoning: To prove the given inequality involving
absolute values, simplification using algebraic properties and given
hypotheses will directly lead us to the desired inequality by recognizing the
algebraic manipulation and application of the known bounds ' $h₁$ ' and ' $h₂$ ', as
well as the non-negativity of ' $p$ '.
simp [abs_of_nonpos] -- Reasoning: The objective is to demonstrate that ' $|x - 15| \leq 15$ 
+  $|x - p - 15| - p$ ' under the given constraints on 'p' and 'x'. Given ' $p < 15$ ' and ' $p \leq x \leq 15$ ', it's apparent that ' $p - 15$ ' is non-positive. This
observation allows us to apply the property that the absolute value of a
non-positive number is equal to its negative value, thereby simplifying the
expression on the right-hand side of the inequality.
abs_of_nonpos -- Reasoning: The objective is to demonstrate a relationship
between expressions involving inequalities and subtractions. The simplification
of these expressions into a form that directly compares their numerical values
can lead to a straightforward application of known inequalities and algebraic
properties, demonstrating the inequality's validity under the given conditions.
```

Figure 9: Example generated thoughts and proof from Lean-STaR

1. Going beyond human data | Lean-STaR

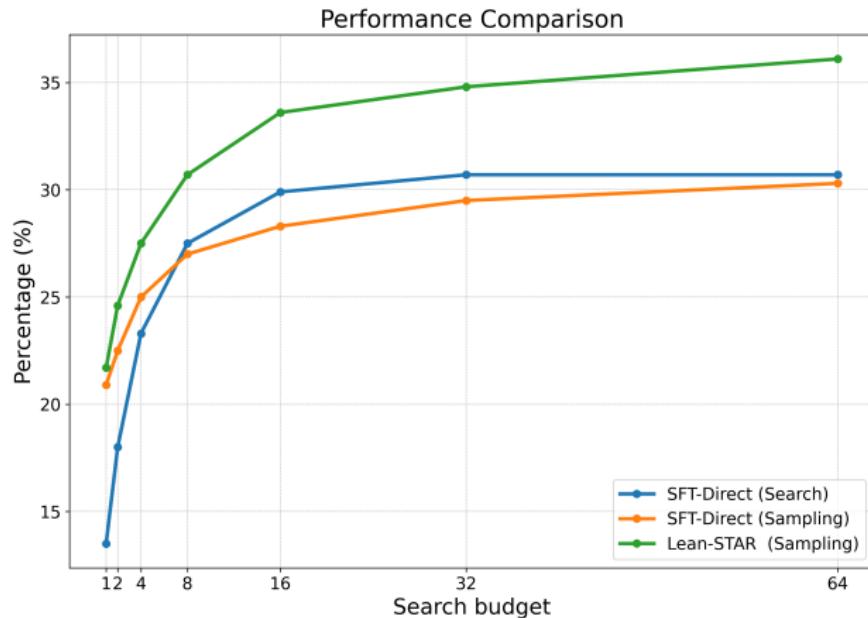


Figure 10: Increasing the search budget is more effective with thoughts

3 open problems in neural theorem proving in Lean:

- Going beyond human data
 - Synthesizing data: problems, proofs, plans, ...
- **Going beyond competition problems**
- Going beyond mathematics

2. Going beyond competition problems

Lots of exciting progress! Some methods can solve IMO problems!
However, not much impact on proving in practice.

2. Going beyond competition problems

Accessibility gap:

- Some methods are hard to integrate into tools
 - Not open-source (AlphaProof, ...)
 - Expensive to run (MCTS, ...)

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Accessibility gap:

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 - Not open-source (AlphaProof, ...)
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However, there are model-agnostic tools available to plug into!

2. Going beyond competition problems

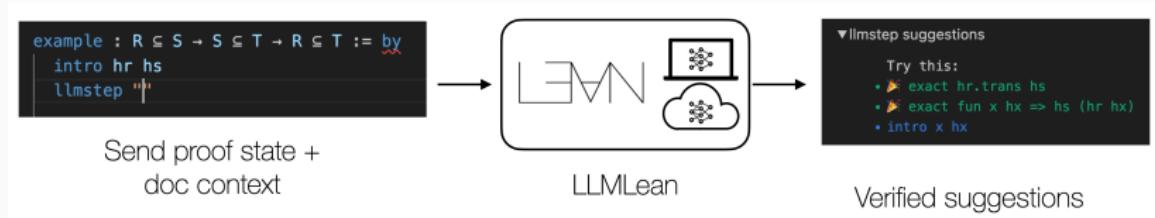


Figure 11: <https://github.com/cmu-l3/llmlean>

2. Going beyond competition problems

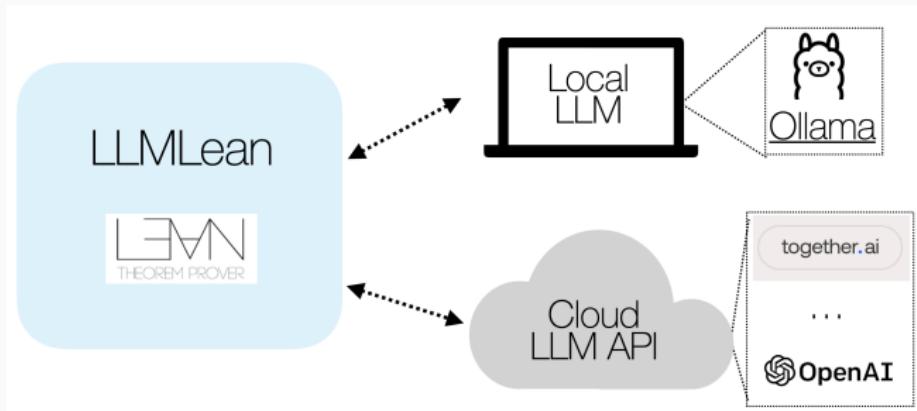


Figure 12: <https://github.com/cmu-l3/llmlean>

2. Going beyond competition problems



ForMathlib > Entropy > E Basic.lean > {} ProbabilityTheory > {} entropy

```
namespace ProbabilityTheory
section entropy
```

— `H[X | Y=y] = $\sum_s P[X=s | Y=y] \log 1/(P[X=s | Y=y])$. -/
lemma entropy_cond_eq_sum (hx : Measurable X) (μ : Measure α) [IsProbabilityMeasure hx | Y = y ; μ] = $\sum_s x, \text{negMulLog } (\mu[Y = y]).map X (x)).toReal := by$

```
by_cases hy : μ(Y⁻¹{y}) = 0
  rw [entropy_def, cond_eq_zero_of_meas_eq_zero _ hy]
  simp
  have : IsProbabilityMeasure (μ[Y = y]) := cond_isProbabilityMeasure _ hy
  rw [entropy_eq_sum hx]
```

lemma entropy_cond_eq_sum_finiteRange (hx : Measurable X) (μ : Measure α) [IsProbabilityMeasure hx | Y = y ; μ] = $\sum_{x \in \text{FiniteRange.toFInset } X}, \text{negMulLog } (\mu[Y = y]).map X (x)).toReal$

```
lmqed
```

— If 'X', 'Y' are 'S'-valued and 'T'-valued random variables, and 'Y = f(X)' some injection 'f : S ↪ T', then 'H[Y] = H[X]'.

One can also use "entropy_of_comp_eq_of_comp" as an alternative if verifying in For the upper bound only, see "entropy_comp_le". -/
lemma entropy_comp_of_injective (μ : Measure α) (hx : Measurable X) (f : S → T) (hf : Function.Injective f)
 (hf : X ; μ) = H[X ; μ] := by
 have hf_m : Measurable f := measurable_of_countable f
 rw [entropy_def, ← Measure.map_map hf m hx, measureEntropy_map_of_injective]

y : T
isIt : FiniteRange X
H[X | Y = y ; μ] = $\sum_{x \in \text{FiniteRange.toFInset } X}, (\mu[Y = y]).map X (x))$

▼ LLMLean suggestions

- Try this:
 - rw [entropy_cond_eq_sum hx μ y]
apply tsum_eq_sum
intro s hs
convert negMulLog_zero
convert BNNReal.zero_toReal
convert measure_mono_null _ (full_measure_of_finiteRange hx)
simp [hs]
 - rw [entropy_cond_eq_sum hx μ y]
convert tsum_eq_sum _ using 1
intro x hx
convert negMulLog_zero
convert BNNReal.zero_toReal
convert measure_mono_null _ (full_measure_of_finiteRange hx)
simp [hx]
 - by_cases hy : μ(Y⁻¹{y}) = 0
rw [entropy_def, cond_eq_zero_of_meas_eq_zero _ hy]
simp
have : IsProbabilityMeasure (μ[Y = y]) := cond_isProbabilityMeasure _ hy
rw [entropy_eq_sum_finiteRange hx μ := μ[Y = y]]

Figure 13: Example on Polynomial Freiman Rusza Conjecture project
<https://github.com/cmu-l3/llmlean>

2. Going beyond competition problems

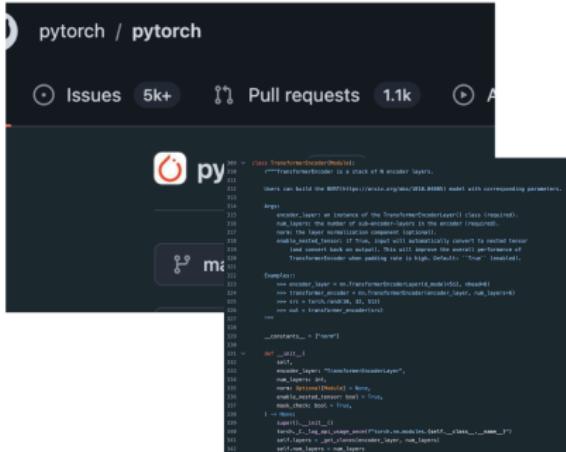
Benchmarking gap:

- Benchmark improvements (e.g., on competition problems) do not measure improvement in real-world proving conditions

2. Going beyond competition problems

```
Example Problem (Finding the Maximum Sum Subarray):  
Python  
  
def max_subarray_sum(arr):  
    max_so_far = 0  
    max_ending_here = 0  
  
    for i in range(len(arr)):  
        max_ending_here = max_ending_here + arr[i]  
        if max_so_far < max_ending_here:  
            max_so_far = max_ending_here  
  
        if max_ending_here < 0:  
            max_ending_here = 0  
  
    return max_so_far  
  
if __name__ == "__main__":  
    arr = list(map(int, input().split()))  
    print(max_subarray_sum(arr))
```

Python
Interview Question



Code in a real repository

Figure 14: Interview questions \neq real code development

2. Going beyond competition problems

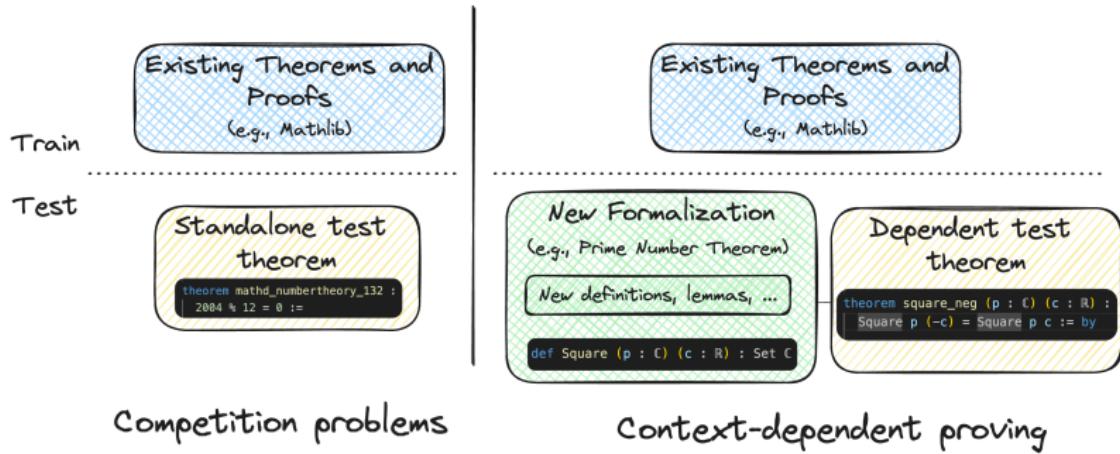


Figure 15: Competition problems \neq real proof development

2. Going beyond competition problems

Real-world proving is **context-dependent**:

- (context, theorem) → proof
 - Context: repository of code, new definitions, auxiliary lemmas

2. Going beyond competition problems

Generalization to new contexts is studied in other proof assistants,
e.g., online setting⁴, testing on held-out repositories⁵

Not a focus for state-of-the-art models/benchmarks in Lean!

⁴Tactician [2], Graph2Tac [1]

⁵CoqGym [6]

2. Going beyond competition problems | miniCTX

miniCTX: Neural Theorem Proving with (Long-)Contexts
Jiewen Hu, Thomas Zhu, Sean Welleck
<https://www.arxiv.org/abs/2408.03350>

2. Going beyond competition problems | miniCTX

miniCTX:

Collect (context, theorem) examples from real Lean projects:⁶

- “Future mathlib”: theorems added after a time cutoff
- Recent projects: PFR, PrimeNumberTheorem
- Textbook exercises: How To Prove It, Math 2001

⁶+ tools for easily adding new projects: <https://github.com/cmu-l3/ntp-toolkit>

2. Going beyond competition problems | miniCTX

miniCTX:

Collect (context, theorem) examples from real Lean projects:⁶

- “Future mathlib”: theorems added after a time cutoff
- Recent projects: PFR, PrimeNumberTheorem
- Textbook exercises: How To Prove It, Math 2001

Goal: generalize to new theorems/contexts/repositories

⁶+ tools for easily adding new projects: <https://github.com/cmu-l3/ntp-toolkit>

2. Going beyond competition problems | miniCTX

Context:

- Preceding code in the file
- All accessible premises
- Repository metadata (to recover any other code)

2. Going beyond competition problems | miniCTX

Does context actually matter? A simple experiment.

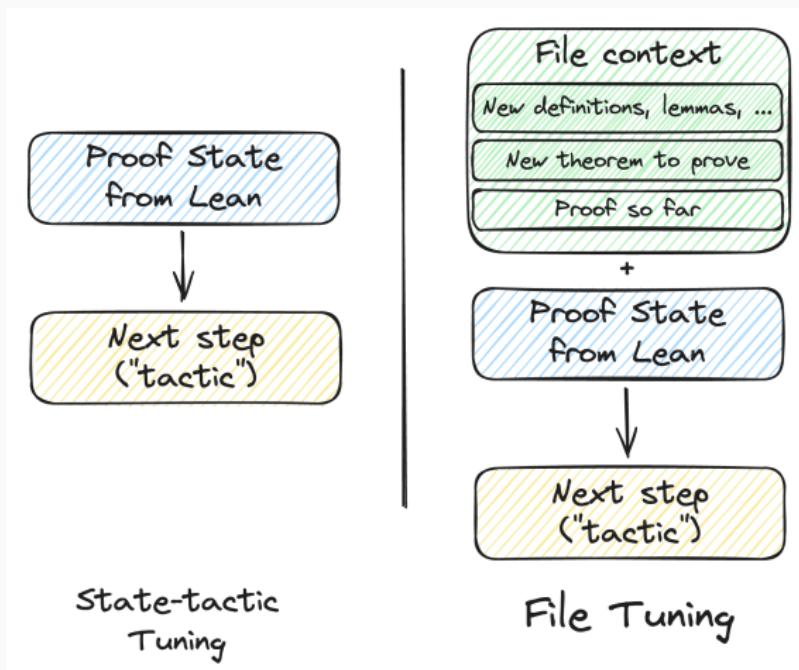


Figure 16: “File tuning”: train on (preceding code, state, next-tactic) examples

2. Going beyond competition problems | miniCTX

Two methods can have similar performance on competition problems, but vastly difference performance on actual projects:

Models	MiniF2F		MiniCTX				
	Test	Prime	PFR	Mathlib	HTPI	Avg.	
GPT-4o (full proof)	-	1.15%	5.56%	2.00%	9.73%	5.59%	
GPT-4o (+ context)	-	13.79%	1.85%	18.00%	31.89%	22.07%	
State-tactic prompting	28.28%	19.54%	5.56%	16.00%	19.15%	20.61%	
State-tactic tuning	32.79%	11.49%	5.56%	22.00%	5.95%	9.31%	
File tuning	33.61%	32.18%	5.56%	34.00%	38.38%	31.65%	

2. Going beyond competition problems | deployment

File-tuned model is deployed in LLMLean:

LLM on your laptop:

1. Install [ollama](#).
2. Pull a language model:

```
ollama pull wellecks/ntpctx-llama3-8b
```



Figure 17: <https://github.com/cmu-l3/llmlean>

2. Going beyond competition problems | deployment

Several open-source artifacts:

- Data/models: <https://huggingface.co/l3lab>
- Data extraction: <https://github.com/cmu-l3/ntp-toolkit>
- Evaluation: <https://github.com/cmu-l3/minictx-eval>

2. Going beyond competition problems | miniCTX

Many approaches to explore in the future:

- “File tuning”: context is preceding code
- Premise selection: context is a set of definitions and theorems
- Full repo: context is all other code in the repository
- ...

2. Going beyond competition problems

Many other potential tools beyond proof completion!

3 open problems in neural theorem proving in Lean:

- Going beyond human data
 - Synthesizing data
- Going beyond competition problems
 - Have actual tools as a goal
- Going beyond mathematics

3. Going beyond mathematics | miniCodeProps

miniCodeProps: a Minimal Benchmark for Proving Code Properties
Evan Lohn, Sean Welleck

<https://arxiv.org/abs/2406.11915>

3. Going beyond mathematics

Interactive theorem provers

- Mathematics:
 - Math as code
 - Guarantees on proof correctness
- Code:
 - Prove properties of code

3. Going beyond mathematics

Formally verified code

Code

```
inductive MyTree {α: Type} where
| leaf : MyTree α
| node : MyTree α → α → MyTree α

def tree_size : MyTree α → N
| .leaf => 1
| .node l _x r => 1 + (tree_size l) + (tree_size r)

def balanced : MyTree α → Prop
| .leaf => true
| .node l _x r => ((tree_size l) =
    (tree_size r)) ∧ (balanced l) ∧ (balanced r)
```

Tree implementation

+

Property

```
-- The size of a balanced tree is odd
theorem balanced_tree_size_odd
| (t: MyTree α) (hb: balanced t): Odd (tree_size t) :=
```

"The size of a balanced tree is odd"

Proof

```
by
cases t with
| leaf => simp [tree_size]
| node p x q =>
  unfold tree_size
  unfold balanced at hb
  simp [hb.1]
```



3. Going beyond mathematics

[AWS Open Source Blog](#)

Lean Into Verified Software Development

by Kesha Hietala and Emrina Torlak | on 08 APR 2024 | in [Amazon Verified Permissions](#), [Open Source](#), [Security](#), [Identity](#), & [Compliance](#), [Technical How-to](#) | [Permalink](#) | [Comments](#) | [Share](#)

Some software components are really important to get right, like your application's access control policies and core

business logic. There are a growing number of tools that can help you verify these components using [automated reasoning](#). In development, the Lean proof assistant is a great tool for formalizing and verifying mathematical proofs.

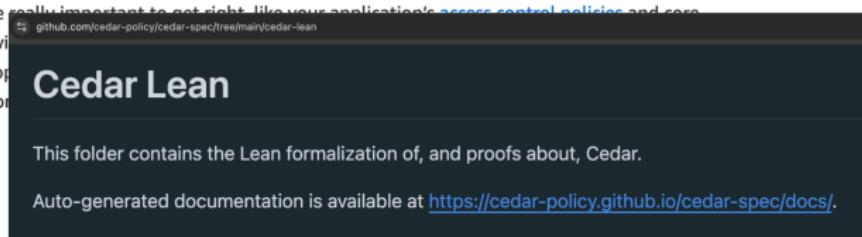


Figure 18: <https://aws.amazon.com/blogsopensource/lean-into-verified-software-development/>

3. Going beyond mathematics

AI/neural theorem proving for program verification is actively studied in other proof assistants, such as Coq and Isabelle.

Not in Lean!

3. Going beyond mathematics

Our question:

- What is the simplest program verification scenario that:
 - Is a subproblem of the full ‘verification problem’
 - Breaks current neural theorem proving methods

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“Simple”:

- Self-contained, no complex dependencies
- Relatively small (fast, cheap evaluation)

3. Going beyond mathematics

Formally verified code

Code

```
inductive MyTree (α: Type) where
| leaf : MyTree α
| node : MyTree α → α → MyTree α

def tree_size : MyTree α → N
| .leaf => 1
| .node l _x r => 1 + (tree_size l) + (tree_size r)

def balanced : MyTree α → Prop
| .leaf => true
| .node l _x r => ((tree_size l) =
  (tree_size r)) ∧ (balanced l) ∧ (balanced r)
```

Tree implementation

Property

```
-- The size of a balanced tree is odd
theorem balanced_tree_size_odd
| (t: MyTree α) (hb: balanced t): Odd (tree_size t) :=
```

"The size of a balanced tree is odd"

Proof

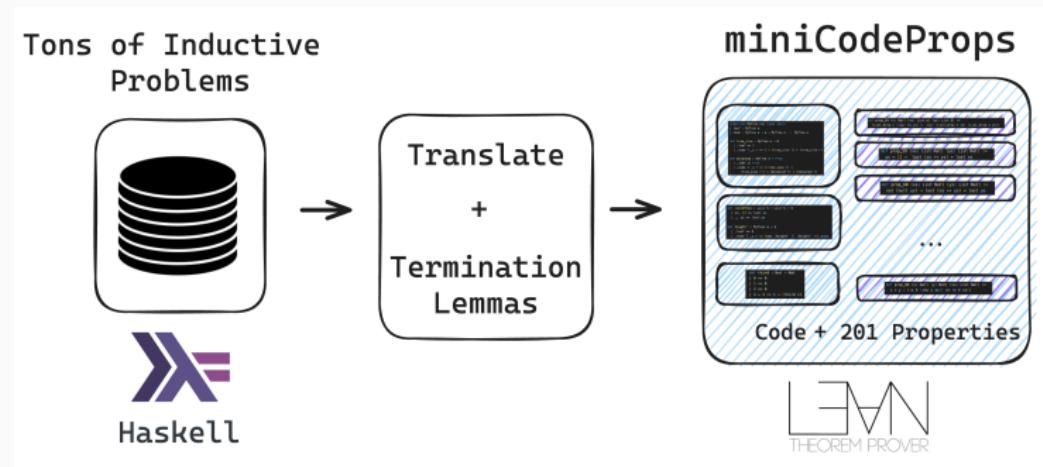
```
by
cases t with
| leaf => simp [tree_size]
| node p x q =>
  unfold tree_size
  unfold balanced at hb
  simp [hb.1]
```



Subproblem: theorem proving! Given (code, property), generate proof

3. Going beyond mathematics | miniCodeProps

Code blocks and 201 properties from *Tons of Inductive Problems*⁷, translated from Haskell to Lean.



⁷<https://tip-org.github.io/>, Claessen et al 2015

3. Going beyond mathematics | miniCodeProps

MiniCodeProps

- Implementation + properties about lists, trees, and heaps
- Classified into difficulties:
 - Easy: Data structure properties
 - Medium: Termination properties
 - Hard: Sorting algorithm properties

3. Going beyond mathematics | miniCodeProps

Evaluation:

- Given property and all dependent code, generate a proof

Models:

- **GPT-4o**: generate full proof, 32 attempts + 1 round of refinement
- **ntp-ctx**: generate a proof via best-first search

<https://github.com/cmu-l3/minicodeprops-eval>

3. Going beyond mathematics | miniCodeProps

Model	Easy	Medium & Hard	Overall
GPT-4o (32 samples)	75.6% (65/86)	4.34% (5/115)	34.8% (70/201)
+ refinement	77.9% (67/86)	6.96% (8/115)	37.3% (75/201)
ntp-context-1.3B	72.1% (62/86)	8.69% (10/115)	35.8% (72/201)

Figure 19: Baselines perform well on easy properties

3. Going beyond mathematics | miniCodeProps

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Figure 20: Poor performance on medium/hard properties

3. Going beyond mathematics | miniCodeProps

```
import Mathlib

def butlast : List α → List α
| [] => []
| [_x] => []
| x::xs => x::(butlast xs)

def butlastConcat : List α → List α → List α
| xs, [] => butlast xs
| xs, ys => xs ++ butlast ys

theorem prop_49 (xs: List Nat) (ys: List Nat) :
  (butlast (xs ++ ys) = butlastConcat xs ys) := by
  induction ys generalizing xs
  · case nil => [redacted]
  · simp [butlast, butlastConcat]
  · case cons y ys ih => [redacted]
  · simp [butlast, butlastConcat]
  · induction xs
  · case nil => [redacted]
  · · simp [butlast, butlastConcat]
  · case cons x xs ih' => [redacted]
    · simp [butlast, butlastConcat, List.cons_append, ih']
```

Figure 21: Successful proof (GPT-4o)

3. Going beyond mathematics | miniCodeProps

```
import Mathlib

def filter : List Nat → (Nat → Bool) → List Nat
| [], _f => []
| x::xs, f => if f x then x::(filter xs f) else (filter xs f)

theorem qsort_term2 (x:Nat) (xs: List Nat) :
List.length (filter xs fun y => decide (y > x)) < Nat.succ (List.length xs) := by
-- induction xs with
-- | nil => []
-- | ·simp [filter, Nat.zero_lt_succ]
-- | cons y ys ih =>
-- | ·simp only [filter]
-- | ·split_ifs with h
-- | ··simp only [List.length]
-- | ··exact Nat.succ_lt_succ ih
-- | ··simp only [List.length]
-- | ··exact Nat.lt_succ_of_lt ih
```

Figure 22: Successful proof (GPT-4o)

3. Going beyond mathematics | miniCodeProps

```
def pairwise : List (List Nat) → List (List Nat)
| xs::ys::xs' => lmerge xs ys :: pairwise xs'
| xs' => xs'

def pairwise_len : List (List Nat) → Nat
| xs => if (Odd xs.length) then xs.length/2 + 1 else xs.length/2

lemma len_pairwise (xs : List (List Nat)) : 2 * (pairwise xs).length = (if (Odd xs.length) then xs.length + 1 else xs.length) := by
  generalize hxl : xs.length = xl
  split_ifs with hxl
  case pos;
    induction xl using Nat.strong_induction_on generalizing xs with
    | h : xs = h & rw [← h] at h
      cases xs with
      | nil => simp at h
      | cons head2 tail2 =>
        cases tail2 with
        | nil => simp at h
        | cons head3 tail3 =>
          cases head3 with
          | nil => unfold pairwise; simp at h; rw [← h]; simp
          | cons head4 tail4 =>
            unfold pairwise;
            rw [← h]; List.length_cons; List.length_cons; Nat.succ_eq_add_one ..; Nat.succ_eq_add_one ..
            ring_nf
            simp
            have hodd : Odd tails.length := by
              rw [List.length_cons, List.length_cons, Nat.succ_eq_add_one .., Nat.add_assoc, Nat.odd_eqd] at h
              apply h.2; simp
            have tmap := h tails.length (by rw [← h]; simp; linearIh); tails rfl heven
            ring_nf at tmap
            rw [tmap];
            ring_nf
            ring_nf
    case neg;
    induction xl using Nat.strong_induction_on generalizing xs with
    | h : xs = h & rw [← h] at h
      cases xs with
      | nil => simp at h
      | cons head2 tail2 =>
        cases tail2 with
        | nil => simp at h
        | cons head3 tail3 =>
          cases head3 with
          | nil => unfold pairwise; simp [← h]
          | cons head4 tail4 =>
            unfold pairwise;
            rw [← h]; List.length_cons; List.length_cons; Nat.succ_eq_add_one ..; Nat.succ_eq_add_one ..
            ring_nf
            simp
            simp at h
            have heven : Even tails.length := by
              rw [Nat.add_assoc, Nat.even_eqd] at h
              apply h.2; simp
            simp at h
            have tmap := h tails.length (by rw [← h]; simp; linearIh); tails rfl heven
            ring_nf at tmap
            exact tmap
```

Figure 23: Human-written proof showing potential length of proofs

This talk

3 open problems in neural theorem proving in Lean:

- Going beyond human data
 - Synthesizing data: problems, proofs, plans, ...
- Going beyond competition problems
 - Have actual tools as a goal
- Going beyond mathematics
 - Program verification

Thank you!

Haohan Lin (Tsinghua)

Evan Lohn (CMU)

Jiewen Hu (CMU)

Zhiqing Sun (CMU)

Yiming Yang (CMU)

Thomas Zhu (CMU)

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Sean Welleck (CMU)

Learning, Language, and Logic (L3) Lab

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