Estimating the Probability of a Conjecture to be a Theorem in PLN for Inference Control

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SingularityNET

Artificial Intelligence and Theorem Proving 2025 (AITP-25)

Outline

Probability of Conjecture to be Theorem (in PLN)

Use such Estimates to Guide Reasoning

State of the Art (notable papers):

- Logical Prior Probability, Abram Demski (2016)
- Uniform Coherence, Scott Garrabrant et al (2016)
- Logical Induction, Scott Garrabrant et al (2016)

Ternary predicate relating theories, proofs and propositions

 Θ : Theory \times Proof \times Proposition \rightarrow Bool



Θ – − ► Predictive Patterns – − ► Estimate Conjectures



Non-Axiomatic Logic (NAL), Pei Wang, 2013



Probabilistic Logic Networks (PLN), Ben Goertzel et al, 2008



Subjective Logic, Audun Jøsang, 2016

PLN Call

Traditional Logic:

<u>PLN:</u>

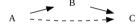
$$\Gamma \vdash T$$

 \rightarrow

 $\Gamma \vdash T$

 \rightleftarrows

Deduction:



$$\begin{array}{c}
B \Rightarrow C & A \Rightarrow B \\
A \Rightarrow C
\end{array}$$

Deduction:

Induction:



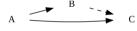
$$A \xrightarrow{B} C$$

$$\begin{array}{c}
B \Rightarrow C & A \Rightarrow B \\
A \Rightarrow C
\end{array}$$

$$A \Rightarrow C \qquad A \Rightarrow B$$
$$B \Rightarrow C$$

Deduction:

Induction:



$$\begin{array}{c} B \Rightarrow C & A \Rightarrow B \\ \hline A \Rightarrow C & \end{array}$$

$$\frac{A \Rightarrow C \qquad A \Rightarrow B}{B \Rightarrow C}$$

$$\frac{A \Rightarrow C \qquad B \Rightarrow C}{A \Rightarrow B}$$

Truth Value:

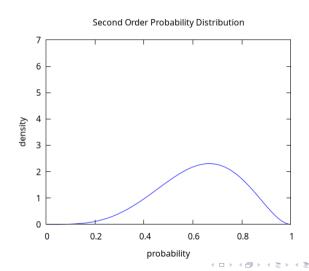
$$A \Rightarrow B \stackrel{\text{m}}{=} \text{TV}$$

TV

Second Order Probability
Distribution

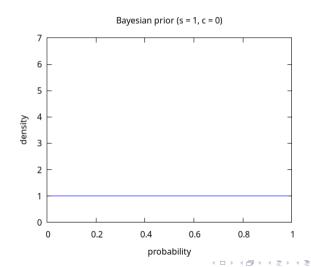
 \approx

P(B|A)



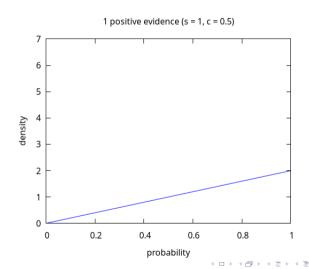
$$A \Rightarrow B \stackrel{\text{m}}{=} \langle s, c \rangle$$

- s = strength
- c = confidence
- Beta Distribution



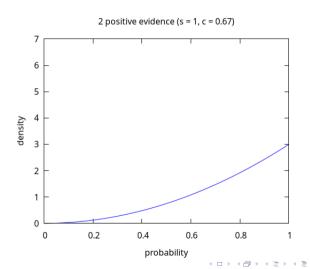
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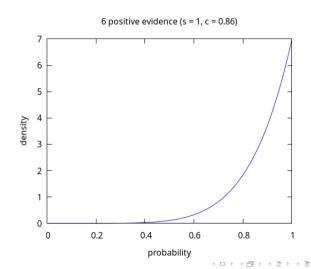
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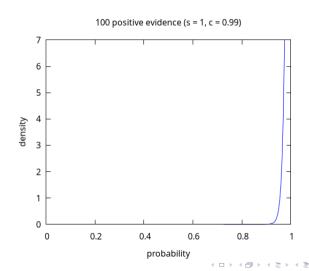
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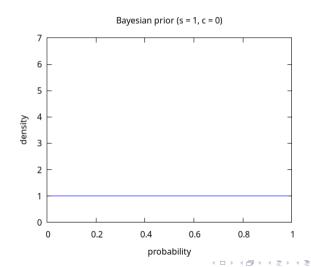
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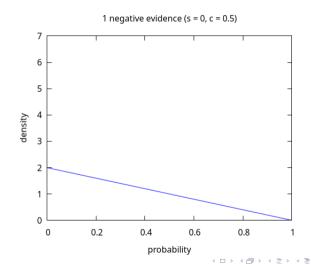
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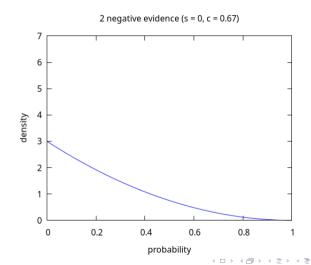
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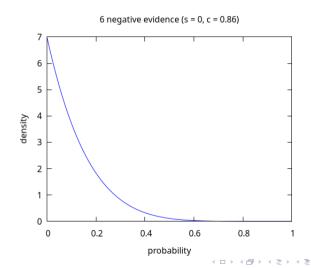
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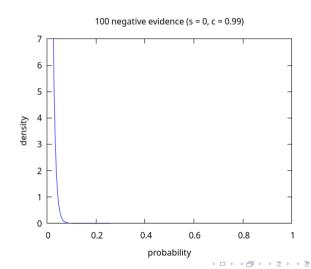
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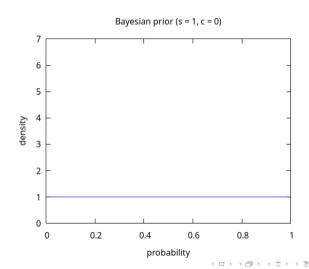
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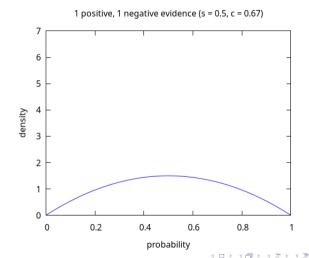
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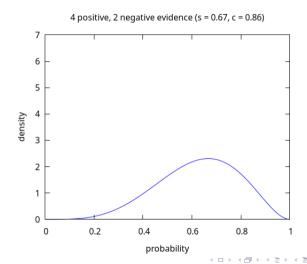
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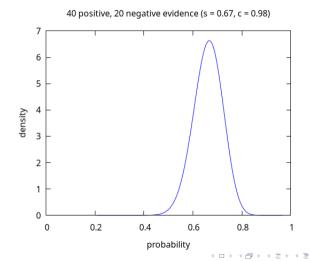
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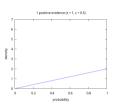
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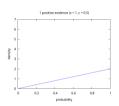
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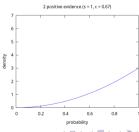


Revision:

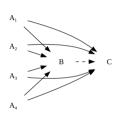
$$\frac{A \Rightarrow B \quad (e) \qquad A \Rightarrow B \quad (f) \qquad e \perp f}{A \Rightarrow B}$$

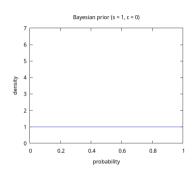






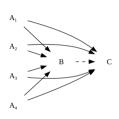
Induction + Revision:

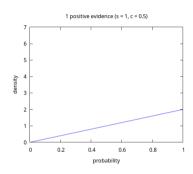




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Induction + Revision:

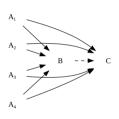


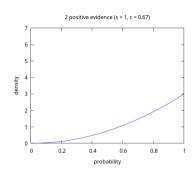


$$\cfrac{A_1 \Rightarrow C \qquad A_1 \Rightarrow B}{B \Rightarrow C \qquad (\text{Ind})} \qquad \cfrac{A_2 \Rightarrow C \qquad A_2 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_3 \Rightarrow C \qquad A_3 \Rightarrow B}{B \Rightarrow C \qquad (\text{Ind})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Rev})} \qquad \cfrac{A_4 \Rightarrow C \qquad A_4 \Rightarrow B}{B \Rightarrow C \qquad (\text{Re$$

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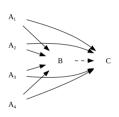
Induction + Revision:

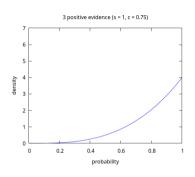




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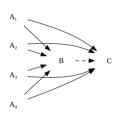
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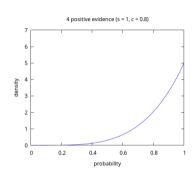


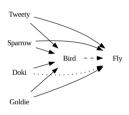


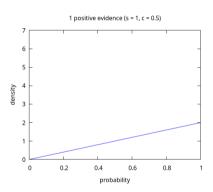
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Induction + Revision:

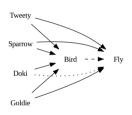


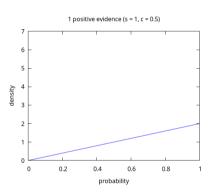




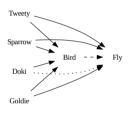


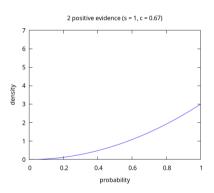
$$\frac{\text{Tweety} \Rightarrow \text{Fly } \triangleq <1,1>}{\text{Bird} \Rightarrow \text{Fly } \triangleq <1,0.5>} \text{Tweety} \Rightarrow \text{Bird } \triangleq <1,1>} \text{(Ind, t)}$$



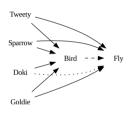


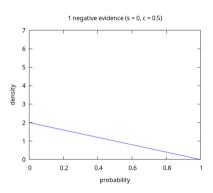
$$\frac{\text{Sparrow} \Rightarrow \text{Fly } \triangleq <1,1>}{\text{Bird} \Rightarrow \text{Fly } \triangleq <1,0.5>} \text{Sparrow} \Rightarrow \text{Bird } \triangleq <1,1>} \text{(Ind, s)}$$





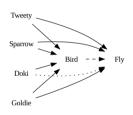
$$\frac{\overline{\mathsf{Bird}}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<1,0.5>}{\mathsf{Bird}}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<1,0.5>}\ (\mathsf{s}) \qquad \qquad t\perp s \\ \overline{\mathsf{Bird}}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<1,0.67>}$$

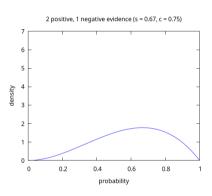




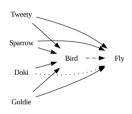
$$\frac{\text{Doki} \Rightarrow \text{Fly} \triangleq <0,1>}{\text{Bird} \Rightarrow \text{Fly} \triangleq <0,0.5>} \frac{\text{Doki} \Rightarrow \text{Bird} \triangleq <1,1>}{\text{Clnd},d)}$$

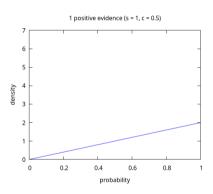






$$\frac{\overline{\mathsf{Bird}}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<1,0.67>\ \ (\mathsf{t,s})\ \ \overline{\mathsf{Bird}}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<0,0.5>\ \ (\mathsf{d})}{\mathsf{Bird}\Rightarrow\mathsf{Fly}\ \stackrel{\underline{=}}{=}<0,67,0.75>}\ \ (\mathsf{Rev},t,s,d)$$

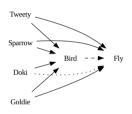


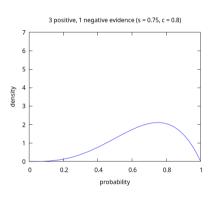


$$\frac{\text{Goldie} \Rightarrow \text{Fly} \stackrel{\underline{m}}{=} <1,1> \qquad \text{Goldie} \Rightarrow \text{Bird} \stackrel{\underline{m}}{=} <1,1>}{\text{Bird} \Rightarrow \text{Fly} \stackrel{\underline{m}}{=} <1,0.5>} \text{(Ind, g)}$$

PLN Recall

Induction + Revision example:





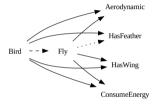
$$\frac{\text{Bird} \Rightarrow \text{Fly} \triangleq \langle 0.67, 0.75 \rangle}{\text{Bird} \Rightarrow \text{Fly} \triangleq \langle 1, 0.5 \rangle} \frac{\text{(g)}}{\text{Bird} \Rightarrow \text{(Rev}, t, s, d, g)}$$

$$\frac{\text{Bird} \Rightarrow \text{Fly} \triangleq \langle 0.75, 0.8 \rangle}{\text{Bird} \Rightarrow \text{Fly} \triangleq \langle 0.75, 0.8 \rangle} \frac{\text{(hev}, t, s, d, g)}{\text{(hev}, t, s, d, g)}$$

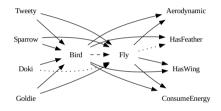
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PLN Recall

Abduction + Revision:



Induction + Abduction + Revision:



PLN also has:

- Quantifiers ∃, ∀
- Traditional Connectors ∧, ∨, ¬
- Composite Predicates: Bird ∧ ¬Penguin ⇒ Fly
- Probabilistic Computational Model

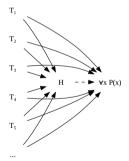
Uncertain Reasoning:

$$H \Rightarrow \forall x \ P(x)$$
?

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$$H \Rightarrow \forall x \ P(x)$$
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Induction + Revision:

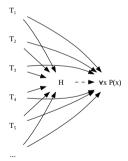


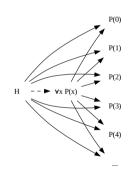
Uncertain Reasoning:

$$H \Rightarrow \forall x P(x)$$
?

Induction + Revision:

Abduction + Revision:

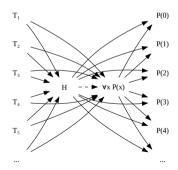




Uncertain Reasoning:

$$H \Rightarrow \forall x \ P(x)$$
?

Induction + Abduction + Revision:



Ternary predicate relating theories, proofs and propositions:

$$\Theta$$
: Theory \times Proof \times Proposition \rightarrow Bool

• Theory: Typing relationships encoding axioms and inference rules.

• Proof: *Inhabitant* of a type.

• Proposition: Type.

Nat

$$\Theta(\{Z : \text{Nat}, S : \text{Nat}->\text{Nat}\}, (S (S (S Z))), \frac{m}{\text{Nat}}) \stackrel{\text{\tiny m}}{=} <1,1>$$

Example (Propositional Calculus):

Instances:

Example (Propositional Calculus):

Instances:

Crisp Patterns:

• Modus ponens:

$$\Theta(\Gamma, f, a \to b) \land \Theta(\Gamma, x, a) \Rightarrow \Theta(\Gamma, f(x), b)$$

$$\stackrel{\text{m}}{=}$$
<1,1>

Existential Quantification Introduction:

$$\begin{split} \Theta(\Gamma, \Pi, \mathbf{T}) \wedge (\operatorname{cl} \Gamma) \wedge (\operatorname{cl} \Pi) \wedge (\operatorname{cl} \mathbf{T}) \Rightarrow \exists \pi \; \Theta(\Gamma, \pi, \mathbf{T}) \\ &\stackrel{\boxplus}{=} \\ <1.1> \end{split}$$

Uncertain Patterns:

• Marginal estimate:

$$\exists \pi \ \Theta(\mathsf{PC}, \pi, \tau)$$
 $\stackrel{=}{=}$
 $< 0.001, 0.8 >$

Conditional estimate:

$$P(\tau) \Rightarrow \exists \pi \ \Theta(PC, \pi, \tau)$$

$$\stackrel{\text{m}}{=}$$
 $< 0.2, 0.7 >$

• How likely is there a *proof* of T in Γ :

$$\exists \pi \ \Theta(\Gamma, \pi, T) \triangleq \$TV$$

• How likely is there a *proof* of T in Γ :

$$\exists \pi \ \Theta(\Gamma, \pi, T) \stackrel{\text{m}}{=} \$TV$$

• How likely is Π proving a *theorem* in Γ :

$$\exists \tau \ \Theta(\Gamma, \Pi, \tau) \triangleq \$TV$$

• How likely is there a *proof* of T in Γ :

$$\exists \pi \ \Theta(\Gamma, \pi, T) \stackrel{\text{m}}{=} \$TV$$

• How likely is Π proving a *theorem* in Γ :

$$\exists \tau \ \Theta(\Gamma, \Pi, \tau) \stackrel{\text{m}}{=} \$ \text{TV}$$

• How likely is there a *theory* in which Π proves T:

$$\exists \gamma \ \Theta(\gamma, \Pi, T) \triangleq \$TV$$

• How likely is there a *proof* of T in Γ :

$$\exists \pi \ \Theta(\Gamma, \pi, T) \stackrel{\text{m}}{=} \$TV$$

• How likely is Π proving a *theorem* in Γ :

$$\exists \tau \ \Theta(\Gamma, \Pi, \tau) \triangleq \$TV$$

• How likely is there a *theory* in which Π proves T:

$$\exists \gamma \ \Theta(\gamma, \Pi, T) \triangleq \$TV$$

• How likely is there a *proof* of a *theorem* in a *theory* with certain *properties*:

$$\exists \gamma, \pi, \tau \ \Theta(\gamma, \pi, \tau) \land P(\gamma) \land Q(\pi) \land R(\tau) \land S(\gamma, \pi, \tau) \stackrel{\text{m}}{=} \$TV$$



(bc PARAMS THEORY (: \$proof PROP))

```
(bc PARAMS THEORY (: $proof PROP)) \downarrow \quad [\$proof] = z (bc PLN_PARAMS PLN_THEORY (: \$pln_proof \ (\stackrel{m}{=} (\exists \ z \ (\Theta \ [THEORY] \ z \ [PROP])) \ \$TV))
```

```
(bc PARAMS THEORY (: $proof PROP))

↓ [$proof] = z

(bc PLN_PARAMS PLN_THEORY (: $pln_proof (= (∃ z (Θ [THEORY] z [PROP])) $TV))

↓

$TV = ?
```

```
(bc PARAMS THEORY (: $proof PROP))

↓ [$proof] = z

(bc PLN_PARAMS PLN_THEORY (: $pln_proof (= (∃ z (Θ [THEORY] z [PROP])) $TV))

↓

$TV = ?
```

● Complete ignorance: \$TV = <1,0>

```
(bc PARAMS THEORY (: $proof PROP))

↓ [$proof] = z

(bc PLN_PARAMS PLN_THEORY (: $pln_proof (= (∃ z (⊕ [THEORY] z [PROP])) $TV))

↓

$TV = ?
```

- Complete ignorance: \$TV = <1,0>
- Complete certainty: \$TV = <1,1>



```
(bc PARAMS THEORY (: $proof PROP))

↓ [$proof] = z

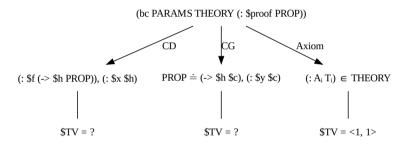
(bc PLN_PARAMS PLN_THEORY (: $pln_proof (= (∃ z (Θ [THEORY] z [PROP])) $TV))

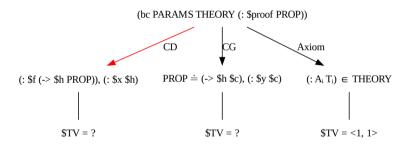
↓

$TV = ?
```

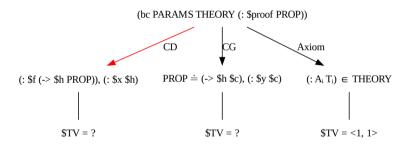
- Complete ignorance: \$TV = <1,0>
- Complete certainty: \$TV = <1,1>
- Partial certainty: \$TV = < 0.7, 0.8 >





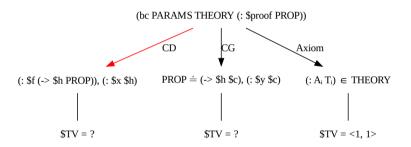


```
;; Backward Chainer
(: bc (-> $a
                                          : Knowledge base space
          Nat
                                          : Maximum depth
          $b
                                 (bc PARAMS ; HOuery Sproof PROP))
          $b))
                                          : Result
:: Base case
  (bc $kb $_ (: $prf $ccln)) (match $kb (: $prf $ccln) (: $prf $ccln)))
  Recusive step (Condensed Detachment)
(= (bc $kb (S $k) (: ($prfabs $prfarg) $ccln))
   (let* (((: $prfabs (-> $prms $ccln))) (bc $kb $k (: $prfabs (-> $prms $ccln))))
          ((: $prfarg $prms) (bc $kb $k (: $prfarg $prms))))
     (: ($prfabs $prfarg) $ccln)))
```



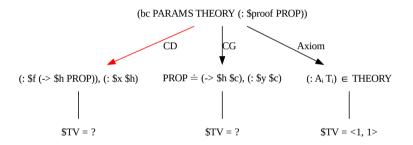
PLN_PARAMS PLN_THEORY

Estimate Probability of Conjecture to be Theorem as Guide



```
(: $pln_proof (≞ (∃ f h x (Λ (Θ [THEORY] f (-> h [PROP]))
(Θ [THEORY] x h)) $TV)))
```

(bc



```
(bc PLN_PARAMS PLN_THEORY (: pln_proof (= (\exists x (\Theta [THEORY] x H))))
```

Conclusion

- Early experiment (https://github.com/trueagi-io/chaining/tree/main/experimental/pln-inf-ctl)
 - Small MetaMath corpus
 - 1th run: exhaustive search, populate Θ

```
 \begin{array}{l} (\triangleq \ (\Theta \ (\text{Cons} \ ([:] \ [\text{ali}.1] \ [\varphi]) \ [\text{PC}]) \ [\text{ali}.1] \ [\varphi]) \ (\text{STV 1 1})) \\ (\triangleq \ (\Theta \ [\text{PC}] \ [\text{ax-3}] \ ([-] \ ([-] \ [\varphi]) \ ([-] \ [\psi])) \ ([-] \ [\psi] \ [\varphi]))) \ (\text{TV 1 1})) \\ (\triangleq \ (\Theta \ (\text{Cons} \ ([:] \ [\text{ali}.1] \ [\varphi]) \ [\text{PC}]) \ [\text{ax-1}] \ ([-] \ [\varphi] \ ([-] \ [\psi] \ [\varphi]))) \ (\text{STV 1 1})) \\ (\triangleq \ (\Theta \ (\text{Cons} \ ([:] \ [\text{ali}.1] \ [\varphi]) \ [\text{PC}]) \ ([\text{ax-mp}] \ [\text{ali}.1] \ [\text{ax-1}]) \ ([-] \ [\psi] \ [\varphi])) \ (\text{STV 1 1})) \\ \dots \end{array}
```

- 2nd run: speed-up via PLN existential reasoning
- Glorified Memoizer
- Future work
 - Large MetaMath corpus
 - Induction, abduction and more