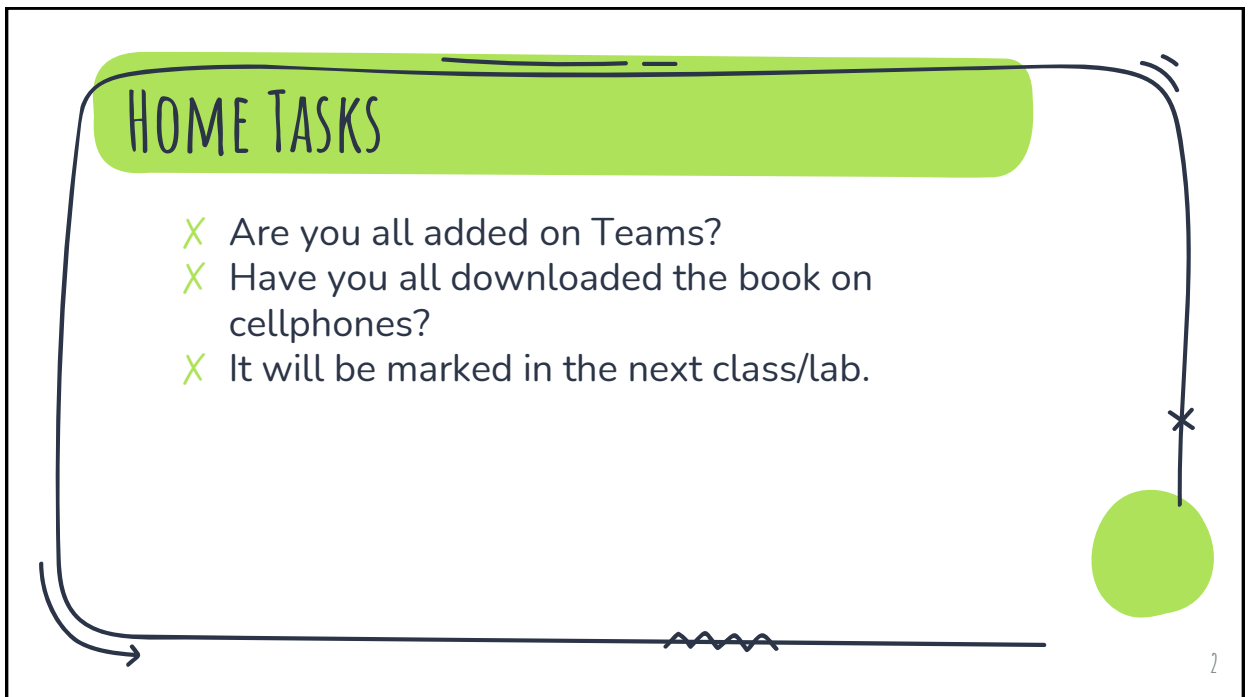
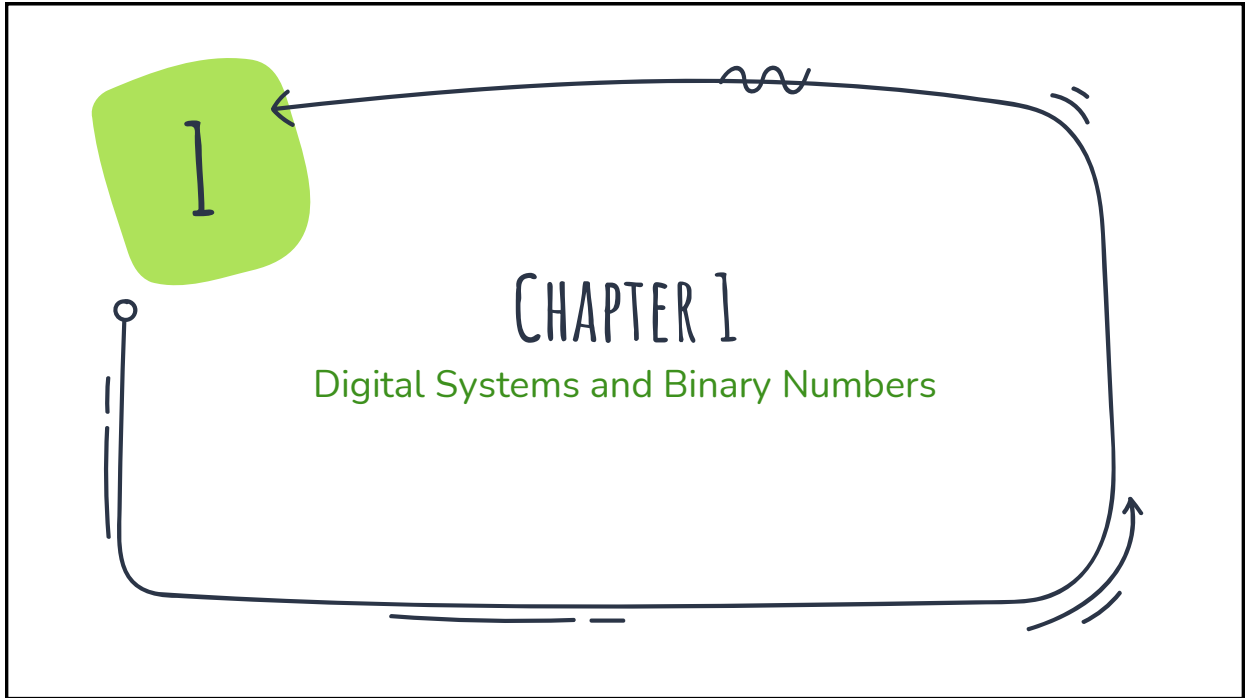


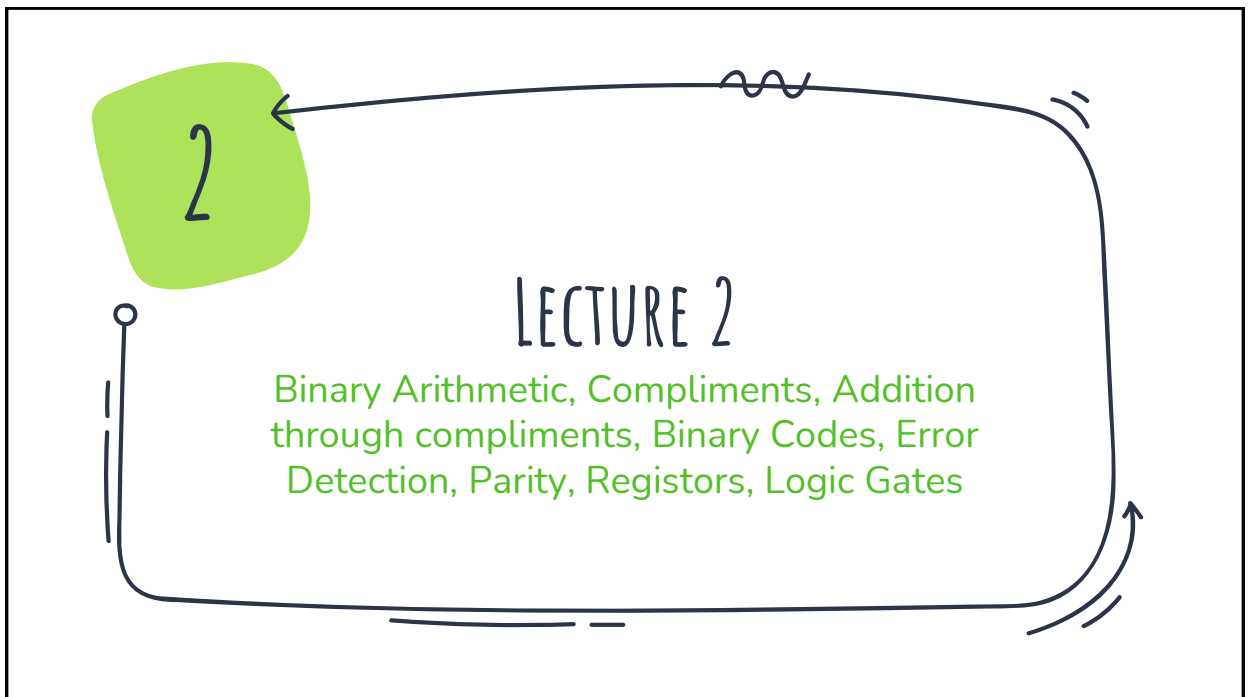
1



2



3



4

ARITHMETIC - ADDITION

X Binary addition is similar to decimal arithmetic

| | | |
|-------------|-------------|---------|
| No carries | 1 0 1 1 0 0 | Carries |
| 0 1 1 0 0 | 1 0 1 1 0 | |
| + 1 0 0 0 1 | + 1 0 1 1 1 | |
| ----- | ----- | |
| 1 1 1 0 1 | 1 0 1 1 0 1 | |

1+1 is 2 (or 10_2), which results in a carry

5

ARITHMETIC - SUBTRACTION

| | | |
|-------------|-------------|---------|
| No borrows | 0 0 1 1 0 | Borrows |
| 1 0 1 1 0 | 1 1 1 1 0 | |
| - 1 0 0 1 0 | - 1 0 0 1 1 | |
| ----- | ----- | |
| 0 0 1 0 0 | 0 1 0 1 1 | |

0 - 1 results in a borrow
Borrow makes it $(10)_2 = (2)_{10}$

6

ARITHMETIC -- MULTIPLICATION

$$\begin{array}{r}
 1011 \\
 \times 101 \\
 \hline
 1011 \\
 0000 \\
 1011 \\
 \hline
 110111
 \end{array}$$

7

COMPLEMENTS

- X Conventional addition (using carry) is easily implemented in digital computers.
- X However; subtraction by borrowing is difficult and inefficient for digital computers
- X Much more efficient to implement subtraction using **ADDITION** of the **COMPLEMENTS** of numbers

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TWO TYPES OF COMPLEMENTS OF RADIX R

X R's Complement

- In Binary 2's complement
- In Decimal 10's complement

X (R-1) Complement aka Diminished Radix

- In Binary 1's complement
- In Decimal 9's Complement

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SUMMARY: HOW TO TAKE COMPLIMENTS

- X **9's Complement:** Subtract each digit from 9
- X **1's Complement:** Subtract each digit from 1
- X OR Invert each bit.
- X **10's Complement:** Take 9's Complement and add 1
- X **2's Complement:** Take 1's Complement and add 1
- X OR Keep the last 1 and invert the rest of bits

10

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SUMMARY: PROOF WITH THE FORMULAE

X Given a number N in base r having n digits

X r 's Complement =

$$(r^n - N)$$

X $R-1$'s Complement =

$$(r^n - 1) - N$$

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DIMINISHED RADIX (R-1) COMPLEMENT

X Example: Take the number $N = 546700$

X Here $n=6$ so put value in formula

$$X = (10^6 - 1) - 546700$$

$$X = (1000000 - 1) - 546700$$

$$X = 999999 - 546700$$

$$X = 453299 \quad \text{So, 9's complement is 453299}$$

X In simple words subtract each digit from 9

$$(r^n - 1) - N$$

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DIMINISHED RADIX (R-1) COMPLEMENT - ACTIVITY

- X 9's Complement
- X Find the 9's complement of
- X 546700 and 12389

The 9's complement of 546700 is
 $999999 - 546700 = 453299$

And the 9's complement of 12389 is
 $99999 - 12389 = 87610$.

| | | | | | | |
|---|---|---|---|---|---|---|
| | 9 | 9 | 9 | 9 | 9 | 9 |
| - | 5 | 4 | 6 | 7 | 0 | 0 |
| | 4 | 5 | 3 | 2 | 9 | 9 |
| | | 9 | 9 | 9 | 9 | 9 |
| - | | 1 | 2 | 3 | 8 | 9 |
| | | 8 | 7 | 6 | 1 | 0 |

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DIMINISHED RADIX (R-1) COMPLEMENT - BASE 2

- X Example: Take the number $N = 1010$
- X Here $n=4$, then put values in formula
- X $(2^4 - 1) - 1010$ Note: $(15)_{10} = (1111)_2$
- X $= (16 - 1)_{10} - (1010)_2$
- X $= (1111)_2 - (1010)_2$
- X Note: $1-0=1$ and $1-1=0$ (Bit Changes)
- X In simple words **just change the bits**

$$(r^n - 1) - N$$

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DIMINISHED RADIX (R-1) COMPLEMENT - ACTIVITY

✕ Find out the 1's Complement of 1011001 & 0001111

The complement 1's of 1011001 is 0100110

The 1's complement of 0001111 is 1110000

| | | | | | | |
|---|---|---|---|---|---|---|
| 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| - | 1 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 |

| | | | | | | |
|---|---|---|---|---|---|---|
| 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| - | 0 | 0 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 |

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RADIX (R) COMPLEMENT - 10'S COMPLEMENT

$$(r^n - N)$$

✕ Find the 10's complement of 546700 and 12389

✕ The 10's complement of 546700 is

✕ $1000000 - 546700 = 453300$

✕ The 10's complement of 12389 is

✕ $100000 - 12389 = 87611$.

| | | | | | | |
|---|---|---|---|---|---|---|
| 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| - | 5 | 4 | 6 | 7 | 0 | 0 |

| | | | | | |
|---|---|---|---|---|---|
| 4 | 5 | 3 | 3 | 0 | 0 |
|---|---|---|---|---|---|

| | | | | | |
|---|---|---|---|---|---|
| 1 | 0 | 0 | 0 | 0 | 0 |
|---|---|---|---|---|---|

| | | | | | |
|---|---|---|---|---|---|
| - | 1 | 2 | 3 | 8 | 9 |
|---|---|---|---|---|---|

| | | | | |
|---|---|---|---|---|
| 8 | 7 | 6 | 1 | 1 |
|---|---|---|---|---|

Notice that it is the same as 9's complement + 1.

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RADIX (R) COMPLEMENT - 2'S COMPLEMENT

- ✗ Find the 2's complement of 1011001 & 0001111

The 2's complement of 1011001 is 0100111

The 2's complement of 0001111 is 1110001

| | | | | | | |
|-------|---|---|---|---|---|---|
| 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| - | 1 | 0 | 1 | 1 | 0 | 0 |
| <hr/> | | | | | | |
| 0 | 1 | 0 | 0 | 1 | 1 | 1 |

| | | | | | | |
|-------|---|---|---|---|---|---|
| 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| - | 0 | 0 | 0 | 1 | 1 | 1 |
| <hr/> | | | | | | |
| 1 | 1 | 1 | 0 | 0 | 0 | 1 |

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FAST COMPUTATION OF 2'S COMPLEMENT

- ✗ **Method 1:**
- ✗ The 2's complement of binary number is obtained by adding 1 to the 1's complement value.

Example:

1's complement of 101100 is 010011 (invert the 0's and 1's)

2's complement of 101100 is 010011 + 1 = 010100

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FAST COMPUTATION OF 2'S COMPLEMENT

X Method 2

- X The 2's complement can be formed by leaving all least significant 0's and the first 1 unchanged, and then replacing 1's by 0's and 0's by 1's in all other higher significant bits.

X Example:

- X The 2's complement of 1101100 is

X 0010100

- X Leave the two low-order 0's and the first 1 unchanged, and then replacing 1's by 0's and 0's by 1's in the four most significant bits.

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SUBTRACTION WITH R-COMPLEMENT

Subtract N from M : $M - N$

r's complement of $N = r^n - N$
Add M to r's complement
Result = $M + (r^n - N)$

- (1) if $M \geq N$, simply ignore the carry
- (2) if $M < N$, The answer is negative. Take the r's complement of sum and place negative sign in front of sum.

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EXAMPLE 1 $M > N$

Perform the subtraction $72532 - 13250 = 59282$.

$M > N$: Case 1 Discard carry

The 10's complement of 13250 is 86750. Therefore:

$$\begin{array}{rcl} M & = & 72532 \\ \text{10's complement of } N & = & +86750 \\ \text{Sum} & = & 159282 \end{array}$$

Discard end carry

$$\text{Answer} = 59282$$

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EXAMPLE 2: $M < N$ --- CASE 2

Perform the subtraction $13250 - 72532 = -59282$.

$$\begin{array}{rcl} M & = & 13250 \\ \text{10's complement of } N & = & +27468 \\ \text{Sum} & = & 40718 \end{array}$$

Now Take 10's complement of Sum. That means: $100000 - 40718$
 $= 59282$

Place negative sign in front of the number: -59282

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Exercise

- Subtract $(3250 - 72532)$ using 10's Complement.

M = 03250

N = 72532

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Exercise

- Using 2's complement subtract 1000100 from 1010100

M = 1010100

N = 1000100

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Exercise

- Subtract (1000100 – 1010100) using 2's Complement

M = 1000100

N = 1010100

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SIGNED BINARY NUMBERS

- X It is usual to represent the sign with a bit placed in the leftmost position of the binary number



The Most common notations are:

- Signed-magnitude system.
- Signed-complement system.

Sign bit 0 positive
Sign bit 1 negative

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SIGNED MAGNITUDE SYSTEM

Example 1

01001 → + 9

11001 → - 9

Example 2

3-bit binary pattern

| Bit Pattern | 000 | 001 | 010 | 011 | 100 | 101 | 110 | 111 |
|-----------------------------------|-----|-----|-----|-----|-----|-----|-----|-----|
| Signed-Magnitude Decimal Value | +0 | +1 | +2 | +3 | -0 | -1 | -2 | -3 |

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SIGNED COMPLEMENT SYSTEM

- ✗ In this system, a negative number is indicated by its complement.
- ✗ The signed-complement system can use either the 1's complement or the 2's complement notations

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EXAMPLE

- X Assuming the representation of the number 9 in binary with 8-bits, we have the following cases:
- X Unsigned 9 or +9 has the same representation in both signed-magnitude and signed-complement systems which is: **00001001**
- X -9 has the signed-magnitude representation: **10001001**
- X -9 has the signed-1's complement representation: **11110110**
- X -9 has the signed-2's complement representation: **11110111**

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SIGNED COMPLEMENT SYSTEM

- X The signed-complement conversion table of a 3-bit binary pattern is as follows:

| Bit Pattern | 000 | 001 | 010 | 011 | 100 | 101 | 110 | 111 |
|-------------------------------------|-----|-----|-----|-----|-----|-----|-----|-----|
| Signed 1's complement decimal value | +0 | +1 | +2 | +3 | -3 | -2 | -1 | -0 |
| Signed 2's complement decimal value | +0 | +1 | +2 | +3 | -4 | -3 | -2 | -1 |

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SIGNED COMPLEMENT SYSTEM

4-bit system

X Positive Values

| Value | Sign-and-Magnitude | 1s Comp. | 2s Comp. |
|-------|--------------------|----------|----------|
| +7 | 0111 | 0111 | 0111 |
| +6 | 0110 | 0110 | 0110 |
| +5 | 0101 | 0101 | 0101 |
| +4 | 0100 | 0100 | 0100 |
| +3 | 0011 | 0011 | 0011 |
| +2 | 0010 | 0010 | 0010 |
| +1 | 0001 | 0001 | 0001 |
| +0 | 0000 | 0000 | 0000 |

X Negative values

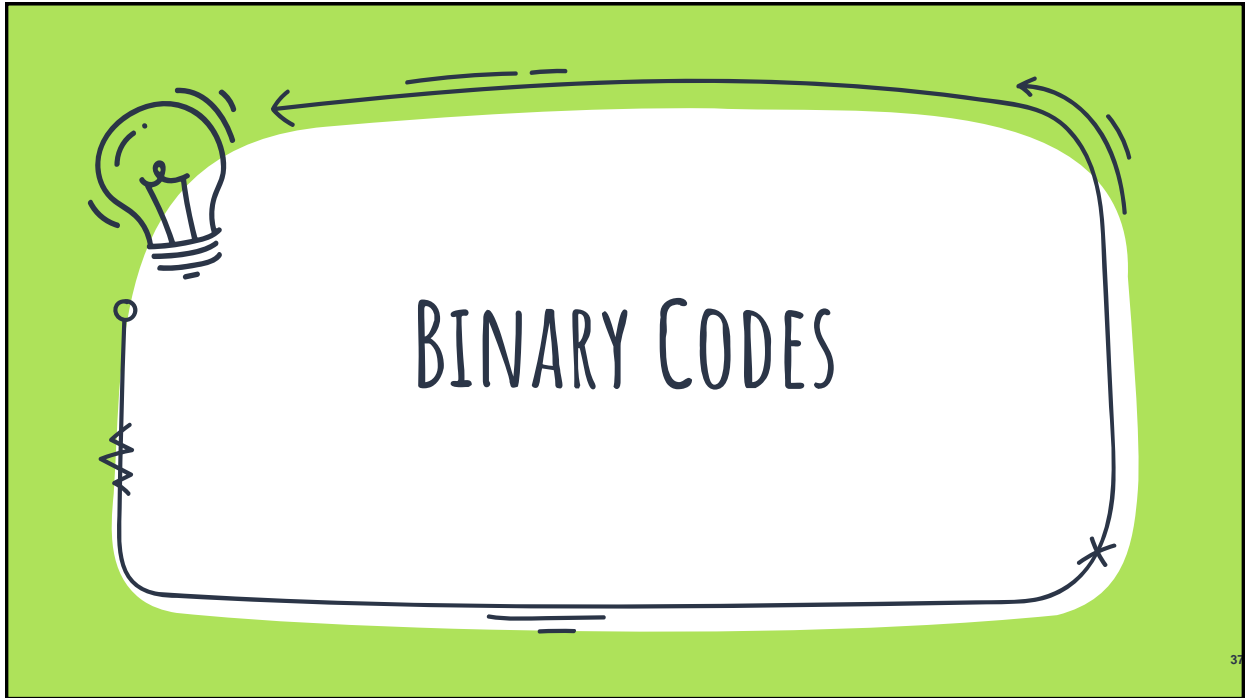
| Value | Sign-and-Magnitude | 1s Comp. | 2s Comp. |
|-------|--------------------|----------|----------|
| -0 | 1000 | 1111 | - |
| -1 | 1001 | 1110 | 1111 |
| -2 | 1010 | 1101 | 1110 |
| -3 | 1011 | 1100 | 1101 |
| -4 | 1100 | 1011 | 1100 |
| -5 | 1101 | 1010 | 1011 |
| -6 | 1110 | 1001 | 1010 |
| -7 | 1111 | 1000 | 1001 |
| -8 | - | - | 1000 |

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COMPLEMENT OF FRACTIONS

- We can extend the idea of complement on fractions.
- Examples:
 - Negate 0101.01 in 1s-complement
Answer: 1010.10
 - Negate 111000.101 in 1s-complement
Answer: 000111.010
 - Negate 0101.01 in 2s-complement
Answer: 1010.11

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THE 8421 BCD

- X BCD stands for Binary-Coded Decimal.
- X A BCD number is a four-bit binary group that represents one of the ten decimal digits 0 through 9

| Decimal Digit | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|---------------|------|------|------|------|------|------|------|------|------|------|
| BCD | 0000 | 0001 | 0010 | 0011 | 0100 | 0101 | 0110 | 0111 | 1000 | 1001 |

Note: 1010, 1011, 1100, 1101, 1110, and 1111 are INVALID CODE!

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BCD: *Example:*

X Convert Decimal to BCD

Decimal number 4926 4 9 2 6

8421 BCD coded number 0100 1001 0010 0110

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ACTIVITY

X Convert the BCD coded number

X 1000 0111 0001 into decimal.

BCD Coded Number 1000 0111 0001

Decimal Number 8 7 1

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BCD

- X People understand decimal system better
- X BCD makes it easy to replace a decimal number with an individual binary code
- X Decimal 15 is BCD 0001 0101 in Binary it was 1111
- X Since most computers store data in eight-bit bytes
 - X Ignore 4 extra bits
 - X One can store two digits per byte, called "packed" BCD

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BCD ADDITION

- X BCD is a numerical code and can be used in arithmetic operations. Here is how to add two BCD numbers:

1. Add the two BCD numbers, using the rules for basic binary addition.
2. If a 4-bit sum is equal to or less than 9, it is a valid BCD number.
3. If a 4-bit sum > 9, or if a carry out of the 4-bit group is generated it is an invalid result. Add 6 (0110) to a 4-bit sum in order to skip the six the invalid states and return the code to 8421. If a carry results when 6 is added, simply add the carry to the next 4-bit group.

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BCD ADDITION - EXAMPLES

X 4+5

$$\begin{array}{r} 0100 \\ + 0101 \\ \hline 1001 \end{array}$$

X 4+8

$$\begin{array}{r} 0100 \\ + 1000 \\ \hline 1100 \\ + 0110 \\ \hline 10010 \end{array}$$

Sum is greater than 1001

Add $(0110)_2$ or $(6)_{10}$

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BCD ADDITION - EXAMPLES - ACTIVITY

X 184+576

$$\begin{array}{r} 1 1 \\ 0001 1000 0100 184 \\ + 0101 0111 0110 576 \\ \hline 0111 0000 1010 760 \\ 0110 0110 0000 \end{array}$$

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THE EXCESS-3 CODE

- X Add 3 to each digit of decimal and convert to 4-bit binary form

| Decimal | Binary | +3 | Excess-3 |
|---------|--------|------|----------|
| 0 | 0000 | 0011 | 0011 |
| 1 | 0001 | 0011 | 0100 |
| 2 | 0010 | 0011 | 0101 |
| 3 | 0011 | 0011 | 0110 |
| 4 | 0100 | 0011 | 0111 |
| 5 | 0101 | 0011 | 1000 |
| 6 | 0110 | 0011 | 1001 |
| 7 | 0111 | 0011 | 1010 |
| 8 | 1000 | 0011 | 1011 |
| 9 | 1001 | 0011 | 1100 |

Sample Problem:

| Decimal | 3 | 5 | 9 |
|----------|----------------|---|---|
| | ↓ | ↓ | ↓ |
| Excess-3 | 0110 1000 1100 | | |

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SOME WEIGHTED CODES

- X Weight codes – Different position has different weight
 X 8421 code $0101 = 0 \times 8 + 1 \times 4 + 0 \times 2 + 1 \times 1 = 5$
 X 84-2-1 $0101 = 0 \times 8 + 1 \times 4 + 0 \times (-2) + 1 \times (-1) = 3$

Four Different Binary Codes for the Decimal Digits

| Decimal Digit | BCD 8421 | 2421 | Excess-3 | 8, 4, -2, -1 |
|---------------|----------|------|----------|--------------|
| 0 | 0000 | 0000 | 0011 | 0000 |
| 1 | 0001 | 0001 | 0100 | 0111 |
| 2 | 0010 | 0010 | 0101 | 0110 |
| 3 | 0011 | 0011 | 0110 | 0101 |
| 4 | 0100 | 0100 | 0111 | 0100 |
| 5 | 0101 | 1011 | 1000 | 1011 |
| 6 | 0110 | 1100 | 1001 | 1010 |
| 7 | 0111 | 1101 | 1010 | 1001 |
| 8 | 1000 | 1110 | 1011 | 1000 |
| 9 | 1001 | 1111 | 1100 | 1111 |

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WARNING: CONVERSION OR CODING



- X Do NOT mix up conversion of a decimal number to a binary number with coding a decimal number with a BINARY CODE.
- X $13_{10} = 1101_2$ (This is conversion)
- X $13 \Leftrightarrow 0001|0011$ (This is coding)

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CHARACTER CODES - ASCII

- X Many applications require handling of not only numbers but letters and special characters
- X ASCII – American Standard Code for Information Interchange
- X 7 Bits to store 128 characters
- X In ASCII, every letter, number, and punctuation symbol has a corresponding number, or ASCII code

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ASCII TABLE

| Dec | Hex | Char | Dec | Hex | Char | Dec | Hex | Char | Dec | Hex | Char |
|-----|-----|------|-----|-----|------|-----|-----|------|-----|-----|------|
| 64 | 40 | @ | 78 | 4E | N | 96 | 60 | ` | 109 | 6D | m |
| 65 | 41 | A | 79 | 4F | O | 97 | 61 | a | 110 | 6E | n |
| 66 | 42 | B | 80 | 50 | P | 98 | 62 | b | 111 | 6F | o |
| 67 | 43 | C | 81 | 51 | Q | 99 | 63 | c | 112 | 70 | p |
| 68 | 44 | D | 82 | 52 | R | 100 | 64 | d | 113 | 71 | q |
| 69 | 45 | E | 83 | 53 | S | 101 | 65 | e | 114 | 72 | r |
| 70 | 46 | F | 84 | 54 | T | 102 | 66 | f | 115 | 73 | s |
| 71 | 47 | G | 85 | 55 | U | 103 | 67 | g | 116 | 74 | t |
| 72 | 48 | H | 86 | 56 | V | 104 | 68 | h | 117 | 75 | u |
| 73 | 49 | I | 87 | 57 | W | 105 | 69 | i | 118 | 76 | v |
| 74 | 4A | J | 88 | 58 | X | 106 | 6A | j | 119 | 77 | w |
| 75 | 4B | K | 89 | 59 | Y | 107 | 6B | k | 120 | 78 | x |
| 76 | 4C | L | 90 | 5A | Z | 108 | 6C | l | 121 | 79 | y |
| 77 | 4D | M | | | | | | | 122 | 7A | z |

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ERROR DETECTION

- ✗ Errors can occur during data transmission. They should be detected, so that re-transmission can be requested.
- ✗ With binary numbers, usually single-bit errors occur.
- ✗ Example: 0010 erroneously transmitted as 0011 or 0000 or 0110 or 1010.

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ERROR DETECTION

- ✗ To detect errors, an eighth bit is sometimes added to the ASCII character to indicate its parity.
- ✗ We insert an extra bit in the leftmost position of the code
- ✗ Two types of parity

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EVEN PARITY

X Even parity – set bit to make number of 1's even

X Examples

1000001 with even parity is 01000001

1000011 with even parity is 11000011

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ODD PARITY

X Similar except make the number of 1's odd

X Examples

1000001 with odd parity is 11000001

1000011 with odd parity is 01000011

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ERROR DETECTION

- Parity bit
 - Even parity: additional bit added to make total number of 1's even.
 - Odd parity: additional bit added to make total number of 1's odd.
- Example of odd parity on ASCII values.

Parity bits

| Character | ASCII Code |
|-----------|------------|
| 0 | 1 0110000 |
| 1 | 0 0110001 |
| ... | ... |
| 9 | 1 0111001 |
| : | 1 0111010 |
| A | 1 1000001 |
| B | 1 1000010 |
| ... | ... |
| Z | 1 1011010 |
| [| 0 1011011 |
| \ | 1 1011100 |

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ERROR DETECTION PROCEDURE

- ✗ The parity bit is helpful in detecting errors during the transmission of information from one location to another.
- ✗ This function is handled by generating an even parity bit at the sending end for each character.
- ✗ The eight-bit characters that include parity bits are transmitted to their destination.
- ✗ The parity of each character is then checked at the receiving end.
- ✗ If the parity of the received character is not even, then at least one bit has changed value during the transmission.

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ERROR DETECTION

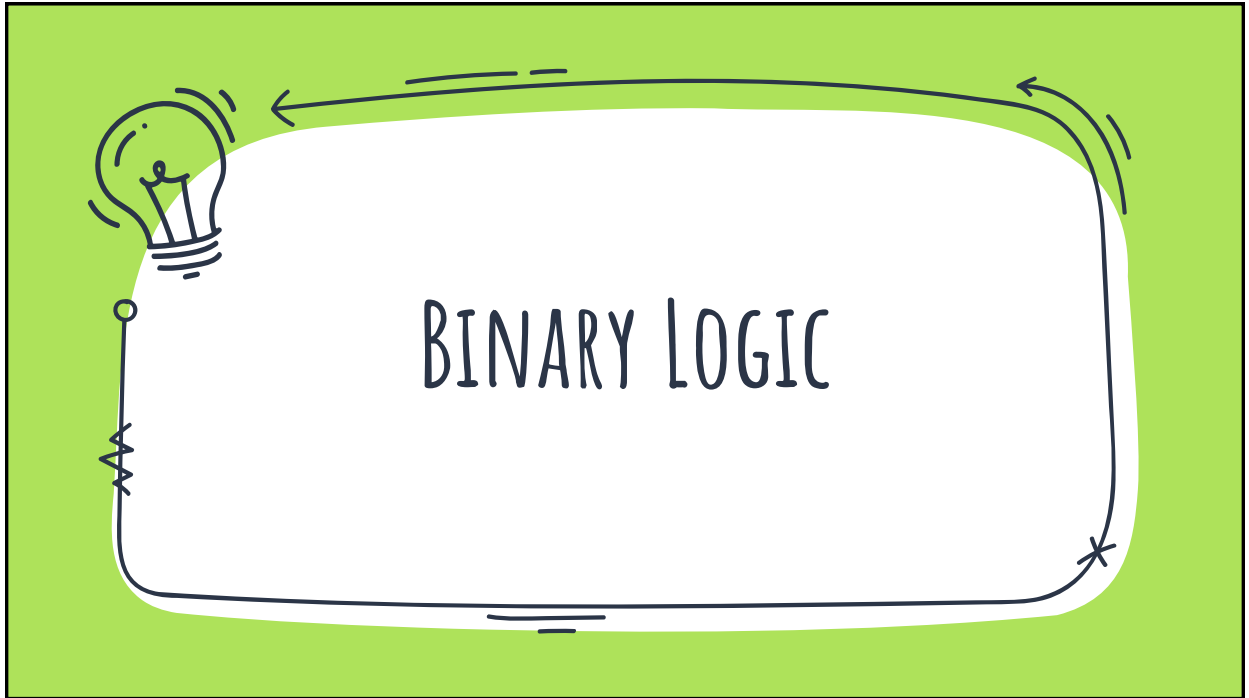
- ✗ Keeping even parity is more common.
- ✗ Parity bit can detect odd number of errors but not even number of errors.
- ✗ Example: Assume odd parity,
- ✗ $10011 \rightarrow 10001$ (detected)
- ✗ $10011 \rightarrow 10101$ (not detected)

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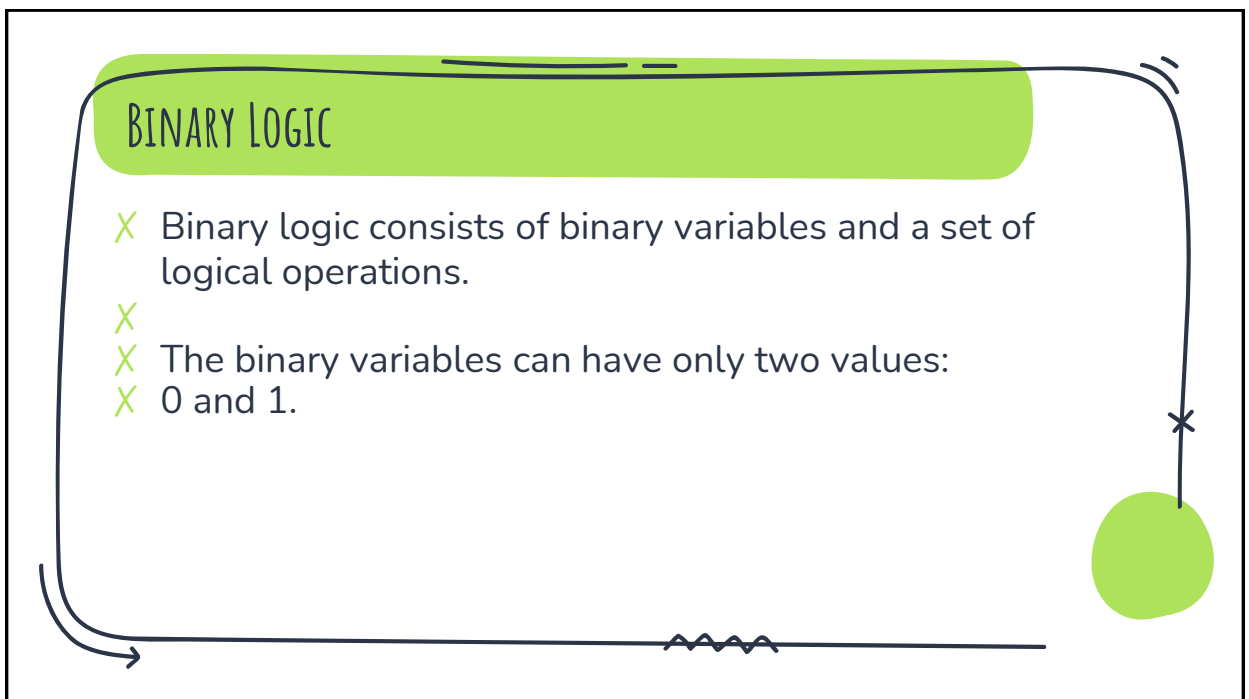
READING ASSIGNMENT BINARY STORAGE AND REGISTERS



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BINARY LOGIC

- X The basic logical operations are AND, OR, and NOT.
- X For x and y binary variables the three logical operations are shown as:

X AND operation \cdot $x \cdot y$

X OR operation $+$ $x + y$

X NOT operation $'$ x'

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THE AND OPERATION

- X The circuit represents **AND** operation and is called **AND** gate



- X Two or more input bits produce one output bit.
- X Both inputs must be true (1) for the output to be true.
- X Otherwise the output is false (0).

| x | y | z |
|---|---|---|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

Truth Table

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THE OR OPERATION

- The circuit represents **OR** operation and is called **OR** gate



| x | y | z |
|---|---|---|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

Truth Table

- Two or more input bits produce one output bit.
- Either inputs must be true (1) for the output to be true.

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THE NOT OPERATION - INVERTER

- The circuit represents **NOT** operation and is called **NOT** gate or **INVERTER**



Symbol

| x | x' |
|---|----|
| 0 | 1 |
| 1 | 0 |

Truth Table

- One bit as input produces its opposite as output.

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TRUTH TABLES OF LOGIC OPERATIONS



| A | B | A & B |
|---|---|-------|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |



| A | B | A B |
|---|---|-------|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

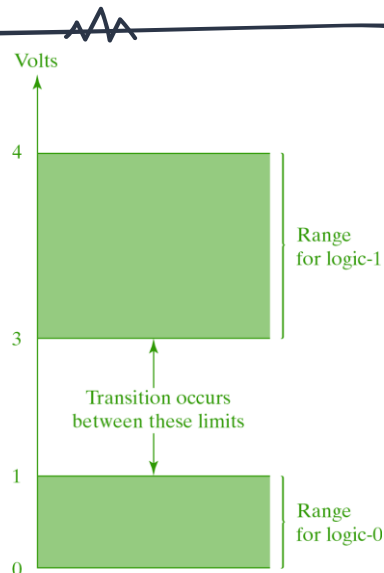


| A | \bar{A} |
|---|-----------|
| 0 | 1 |
| 1 | 0 |

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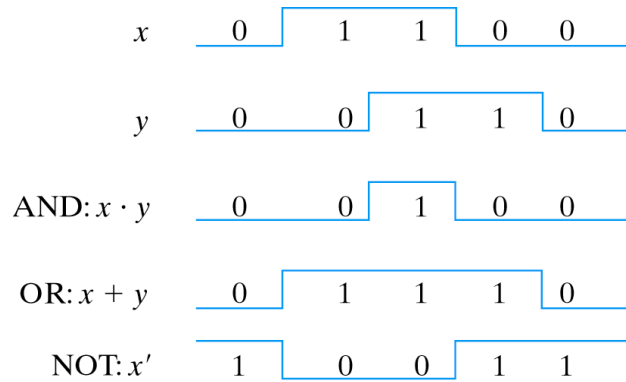
REPRESENTATION OF BINARY VARIABLES

- X Different Digital Systems represent 0 and 1 differently
- X Logical 0 as 0 volts.
Logical 1 as 4 volts
- X Range



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BINARY LOGIC FUNCTIONS



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REFERENCES

- X Chapter 1 – Digital Design Morris Mano
- X Digital Logic and Computer Design – M. Singh, University of North Carolina
- X Digital Design – O. Ozturk, Bilknet University
- X Template is taken from slides carnival.

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CREDITS

These slides are adapted
from Morris Mano Book
5th/6th Edition

