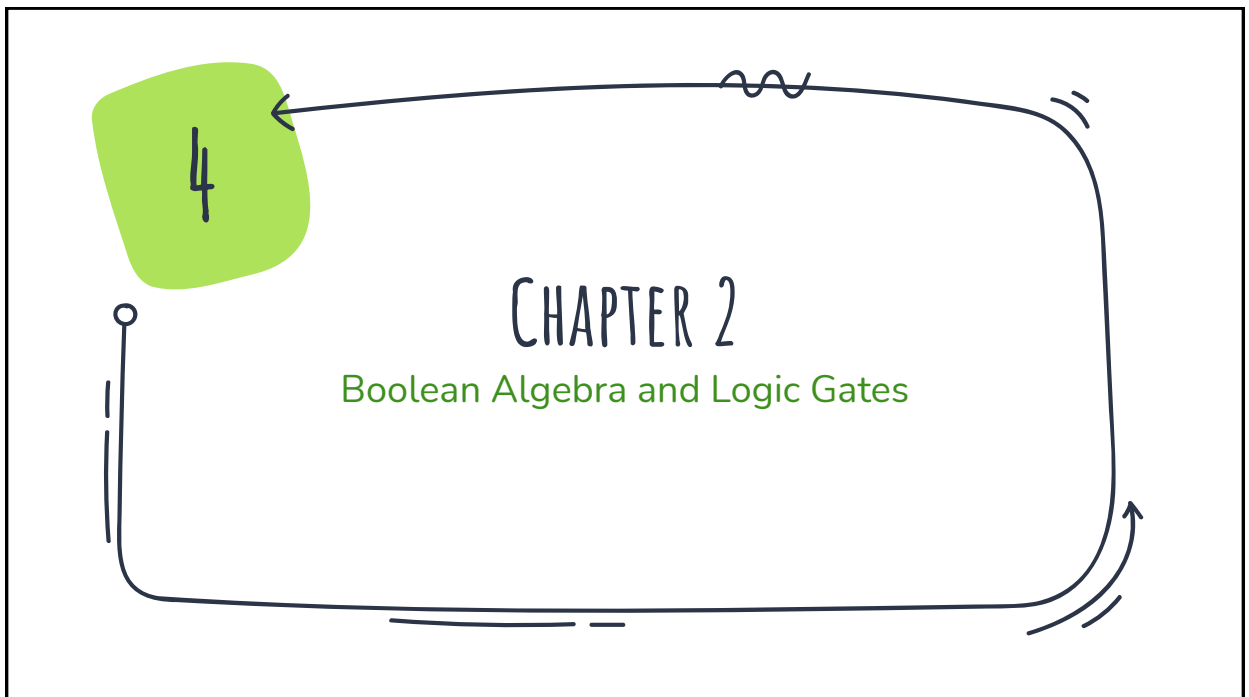


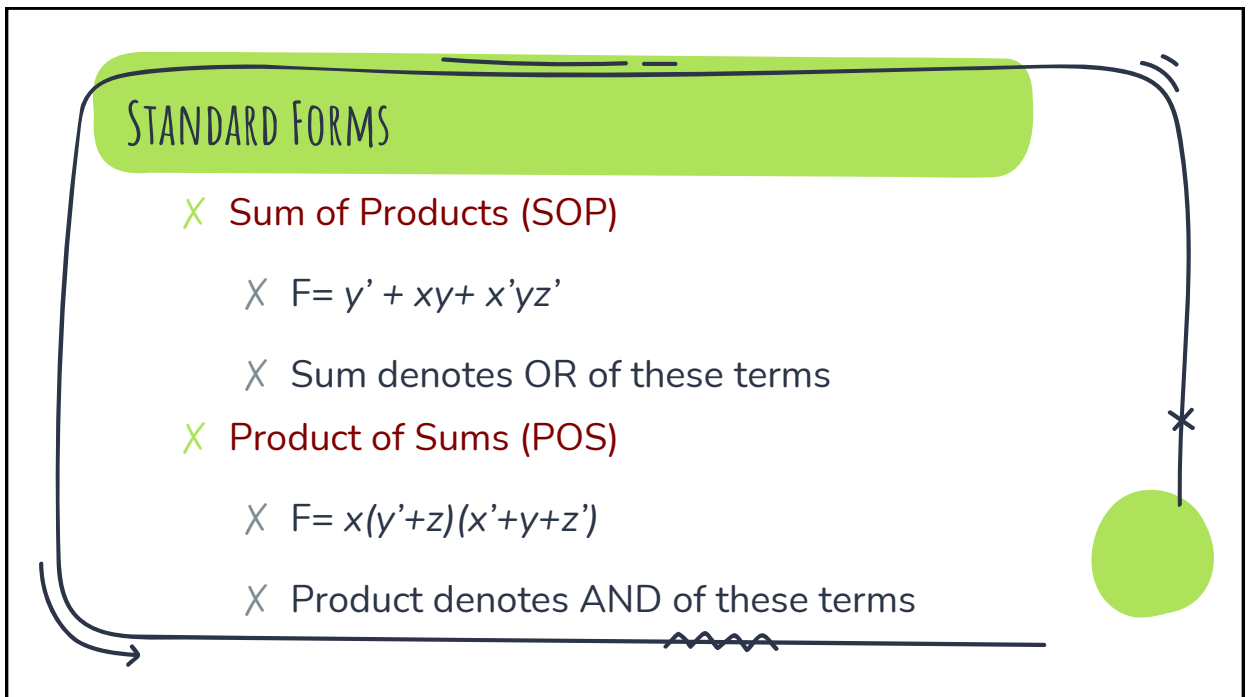
1



3



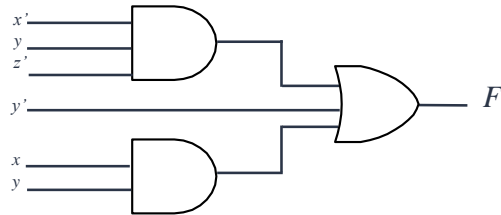
9



10

SUM OF PRODUCTS

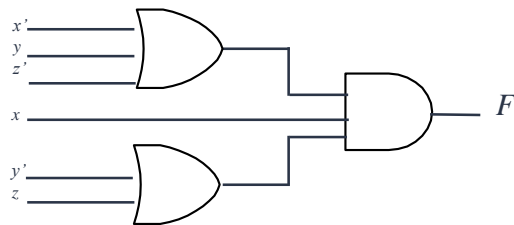
$$F = y' + xy + x'yz'$$



11

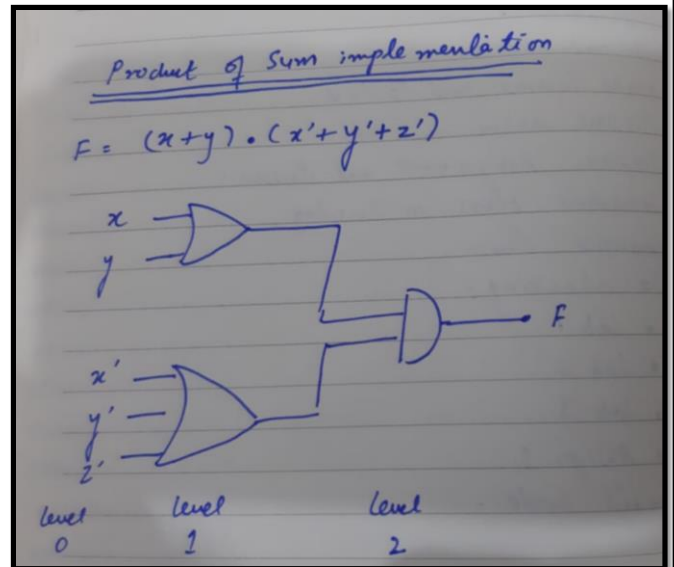
PRODUCT OF SUMS

$$F = x(y' + z)(x' + y + z')$$



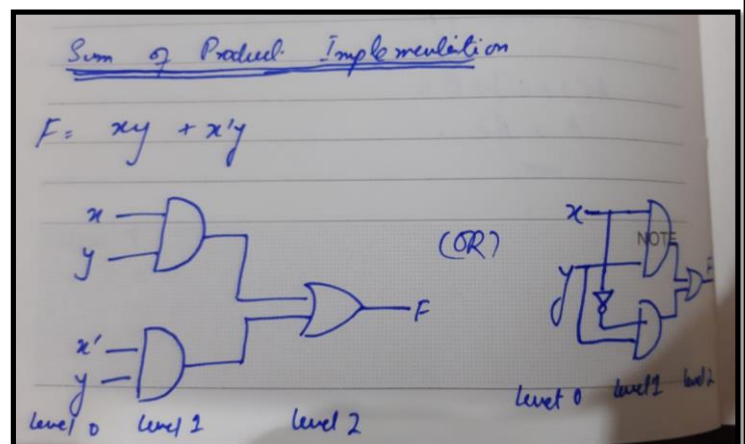
12

2-LEVEL IMPLEMENTATION POS



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2-LEVEL IMPLEMENTATION SOP



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TWO-LEVEL IMPLEMENTATIONS

- X Product of sums expressions can be implemented with two-level circuits as following
 - X literals and their complements at the “0th” level
 - X **OR gates at the first level**
 - X a single **AND gate at the second level**

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TWO-LEVEL IMPLEMENTATIONS

- X Sum of product expressions can be implemented with two-level circuits as following
 - X literals and their complements at the “0th” level
 - X **AND gates at the first level**
 - X a single **OR gate at the second level**

16

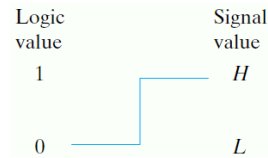
POSITIVE AND NEGATIVE LOGIC

X Signal levels as high (H) or low (L)

X Positive logic

X $H \rightarrow 1$

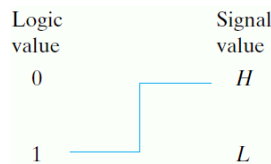
X $L \rightarrow 0$



X Negative logic

X $H \rightarrow 0$

X $L \rightarrow 1$



17

POSITIVE AND NEGATIVE LOGIC

X Truth Table for H and L

<i>x</i>	<i>y</i>	<i>F</i>
<i>L</i>	<i>L</i>	<i>L</i>
<i>L</i>	<i>H</i>	<i>L</i>
<i>H</i>	<i>L</i>	<i>L</i>
<i>H</i>	<i>H</i>	<i>H</i>

X In positive logic: AND

<i>x</i>	<i>y</i>	<i>z</i>
0	0	0
0	1	0
1	0	0
1	1	1



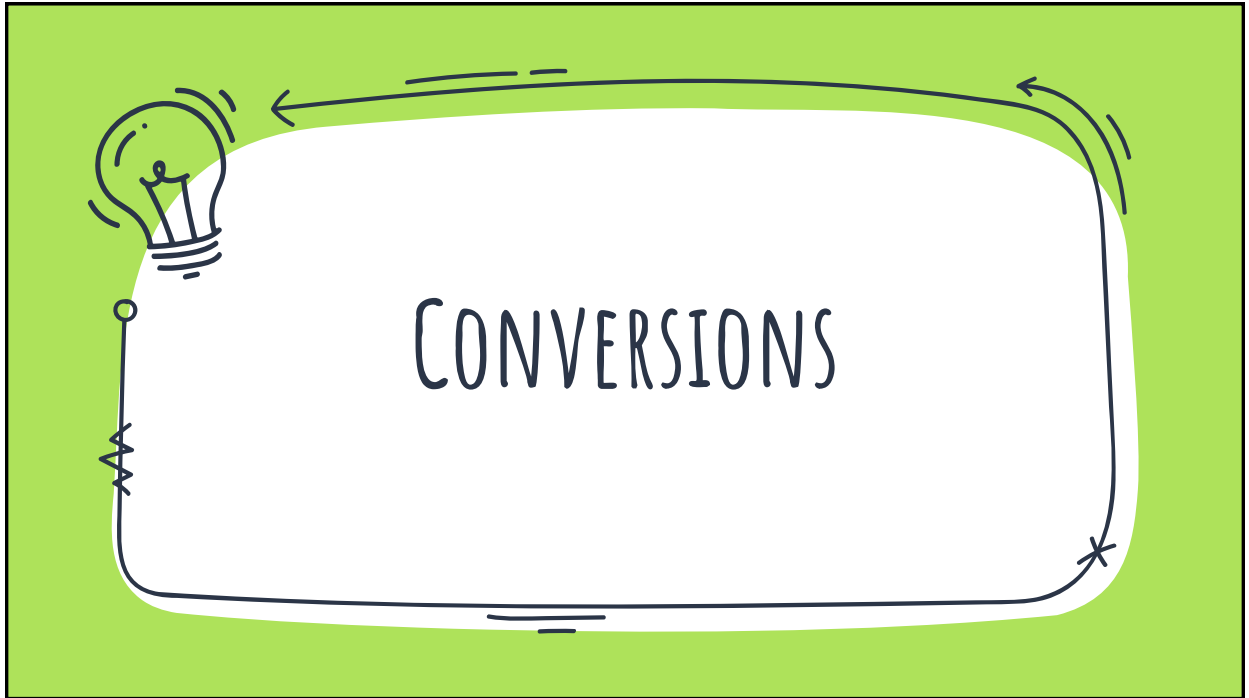
X In negative logic: OR

<i>x</i>	<i>y</i>	<i>z</i>
1	1	1
1	0	1
0	1	1
0	0	0

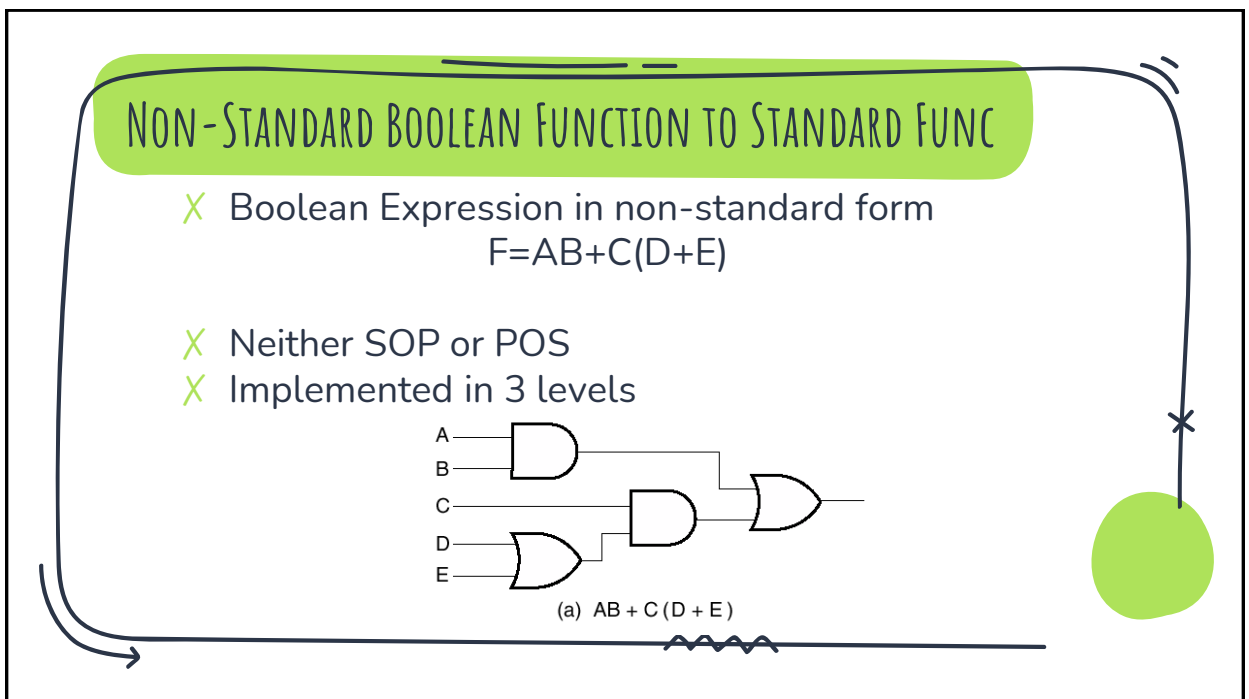


Arrows indicate negative logic

18



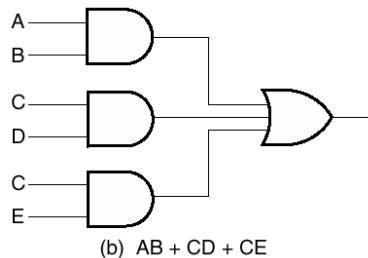
19



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NON-STANDARD BOOLEAN FUNCTION TO STANDARD FUNC

- X $F = AB + C(D + E)$
- X Convert to Standard Form by distributive law
- X $F = AB + C(D + E) = AB + CD + CE$
- X Now it becomes a 2 Level Standard SOP Circuit



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CONVERSION B/W CANONICAL FORMS

- X Complement of a function expressed as sum of minterms equals sum of minterms **missing from the original function**

$$F(A, B, C) = \sum(1, 4, 5, 6, 7) = m_1 + m_4 + m_5 + m_6 + m_7$$

- X Complement is

$$F'(A, B, C) = \sum(0, 2, 3) = m_0 + m_2 + m_3$$

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Example: Conversion from standard to Canonical Form: Sum of Minterms

Express the Boolean function $F = x + y'z$ in a sum of minterms.

AIM: Convert it to 3 variable AND Terms and take their SUM

$$x = x(y + y') = xy + xy'$$

$$xy = xy(z + z') = xyz + xyz'$$

$$xy' = xy'(z + z') = xy'z + xy'z'$$

$$y'z = y'z(x + x') = xy'z + x'y'z$$

Adding all terms and excluding recurring terms:

$$F(x, y, z) = x'y'z + xy'z' + xy'z + xyz' + xyz$$

$$F(x, y, z) = m_1 + m_4 + m_5 + m_6 + m_7 = \sum(1, 4, 5, 6, 7)$$

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Example: Conversion from standard to Canonical Form: Product of Maxterms

Convert the standard equation to canonical form.

$$F = (x' + y)(x + z)(y + z)$$

$$x' + y + zz' = (x' + y + z)(x' + y + z')$$

$$x + z + yy' = (x + z + y)(x + z + y')$$

$$y + z + xx' = (y + z + x)(y + z + x') \quad (\text{Remove any recurring terms})$$

$$F = (x + y + z)(x + y' + z)(x' + y + z)(x' + y + z')$$

$$M_0 M_2 M_4 M_5 = F(x, y, z) = \prod(0, 2, 4, 5)$$

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Example: Conversion from standard to Canonical Form: Product of Maxterms

Express the Boolean function $F = xy + x'z$ in a Product of Maxterms.

AIM: Convert it to 3 variable OR Terms and take their Product

$$F = xy + x'z = (xy + x')(xy + z) \quad (\text{Distributive law})$$

$$(xy + x') = (x + x')(y + x')$$

$$(xy + z) = (x + z)(y + z)$$

$$x' + y + zz' = (x' + y + z)(x' + y + z')$$

$$x + z + yy' = (x + z + y)(x + z + y')$$

$$y + z + xx' = (y + z + x)(y + z + x') \quad (\text{Remove any recurring terms})$$

$$F = (x + y + z)(x + y' + z)(x' + y + z)(x' + y + z')$$

$$M_0 M_2 M_4 M_5 = F(x, y, z) = \Pi(0, 2, 4, 5)$$

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SUMMARY OF CONVERSION

- X To convert into canonical form:
- X SOP: multiply each term with $(x+x')$ (x is missing literal)
- X POS: Add with each term (zz') (z is missing literal)
- X To convert from SOP to POS:
- X Use distributive law
 - a) $x + (y.z) = (x+y).(x+z)$
 - b) $x.(y + z) = (x.y) + (x.z)$

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SOP \leftrightarrow POS USING TRUTH TABLE

- X Find Sum of Minterms and Product of MaxTerms of $F = xy + x'z$
- X Find Sum of Minterms & Product of Maxterms from truth table

X $F(x,y,z) = \sum(1,3,6,7)$

X $F(x,y,z) = \prod(0,2,4,5)$

X	Y	Z	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

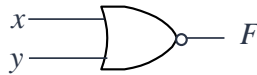
28

SOME MORE GATES

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NOR GATE

✕ 2-input NOR logic gate:



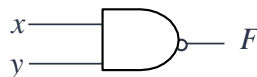
$$F = (X + Y)'$$

X	Y	F
0	0	1
0	1	0
1	0	0
1	1	0

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NAND GATE

✕ 2-input NAND logic gate:



$$F = (X.Y)'$$

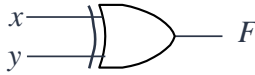
X	Y	F
0	0	1
0	1	1
1	0	1
1	1	0

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EXCLUSIVE OR GATE

X 2-input exclusive-OR (XOR) logic gate:

$$F = X \oplus Y$$



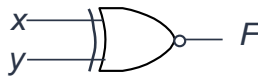
X	Y	F
0	0	0
0	1	1
1	0	1
1	1	0

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EXCLUSIVE NOR GATE

X 2-input exclusive-NOR logic gate:

$$F = (X \oplus Y)'$$



X	Y	F
0	0	1
0	1	0
1	0	0
1	1	1

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Summary

Inputs		Outputs
x	y	xy
0	0	0
0	1	0
1	0	0
1	1	1

$xy = x \text{ AND } y = x * y$
AND is true only if
both inputs are true

NOR is NOT of OR

x	y	x NOR y
0	0	1
0	1	0
1	0	0
1	1	0

Inputs		Outputs
x	y	$x + y$
0	0	0
0	1	1
1	0	1
1	1	1

$x + y = x \text{ OR } y$
OR is true if **either**
inputs are true

NAND is NOT of AND

x	y	x NAND y
0	0	1
0	1	1
1	0	1
1	1	0

Inputs	Outputs
x	\bar{x}
0	1
1	0

$x \text{ bar} = \text{NOT } x$
NOT inverts the bit
We will denote x bar as $\sim x$

XOR is true if both inputs **differ**

x	y	x XOR y
0	0	0
0	1	1
1	0	1
1	1	0

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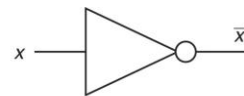
Summary



AND Gate



OR Gate

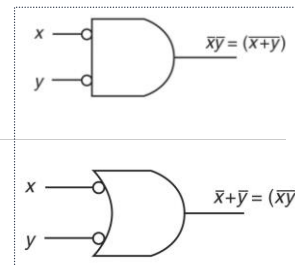
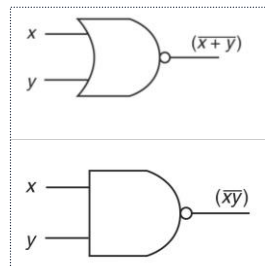


NOT Gate

Here we see the logic gates
that represent the boolean
operations previously
discussed



XOR looks like OR
but with the added
curved line



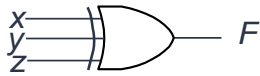
We typically represent NOR and NAND by the two
on the left, but the two on the right are also correct

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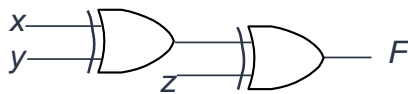
3 INPUT XOR

3-input exclusive-OR (XOR) logic gate:

$$F = X \oplus Y \oplus Z$$



Implemented by 3 Input Gates



Implemented by 2 Input Gates

X	Y	Z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

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NAND AND NOR CONVERSION

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We can make every gate using
NAND or NOR gate

Universal Gates!

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USING ONLY NAND

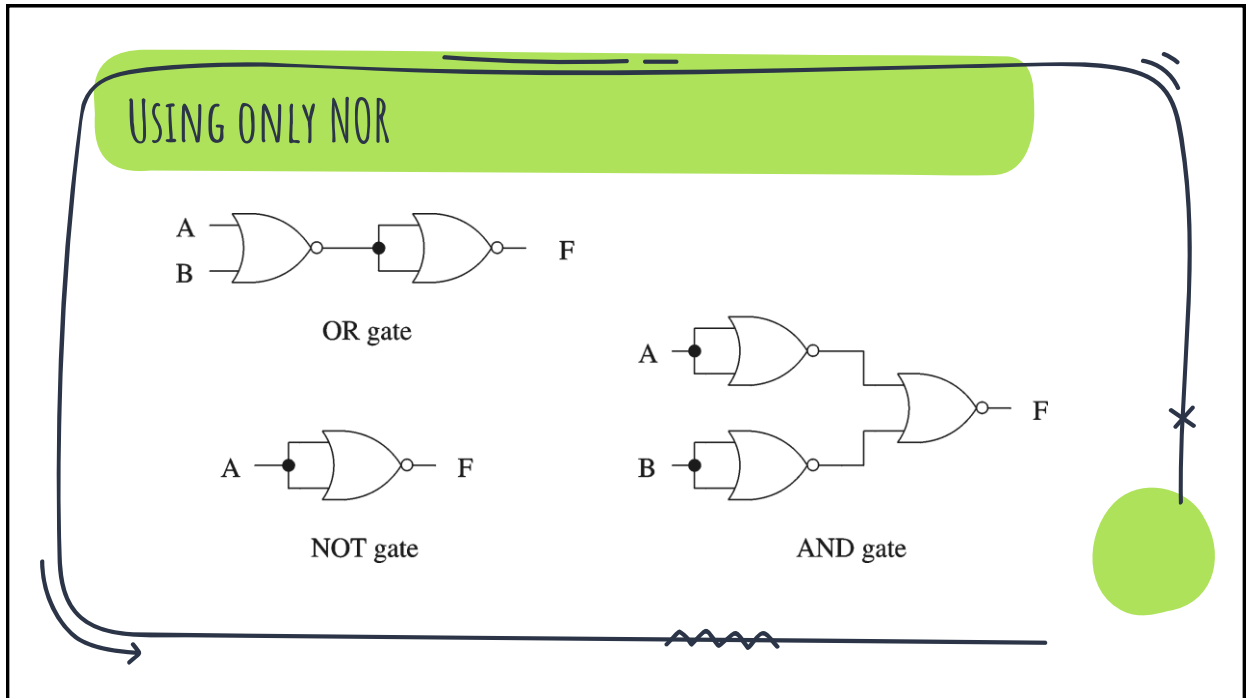
NAND and NOR gates are called **Universal Gates**
Using them we can create circuits that compute AND, OR and NOT

OR Gate

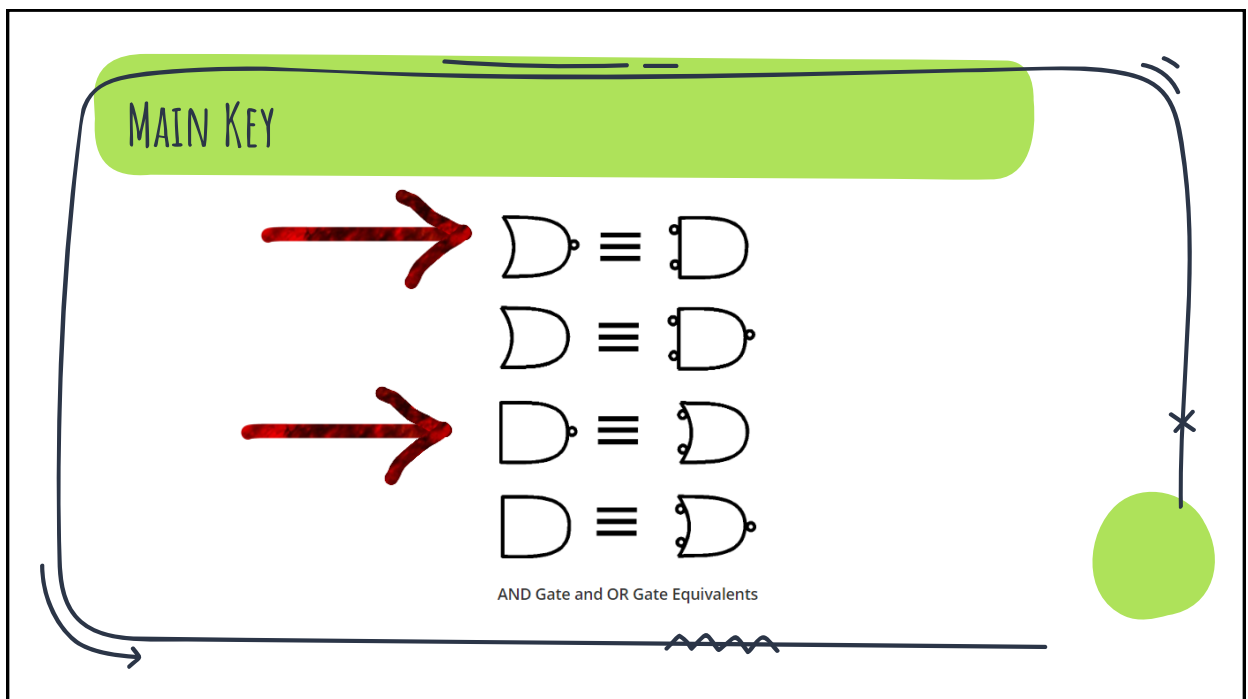
AND Gate

NOT Gate

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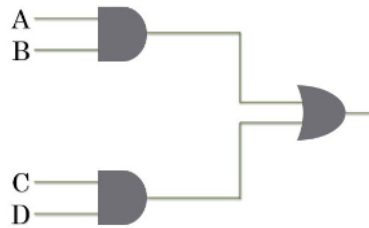


46



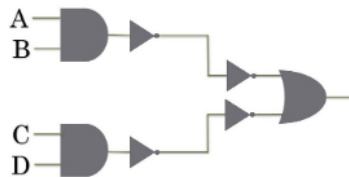
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IMPLEMENT USING ONLY NAND: EXAMPLE 1

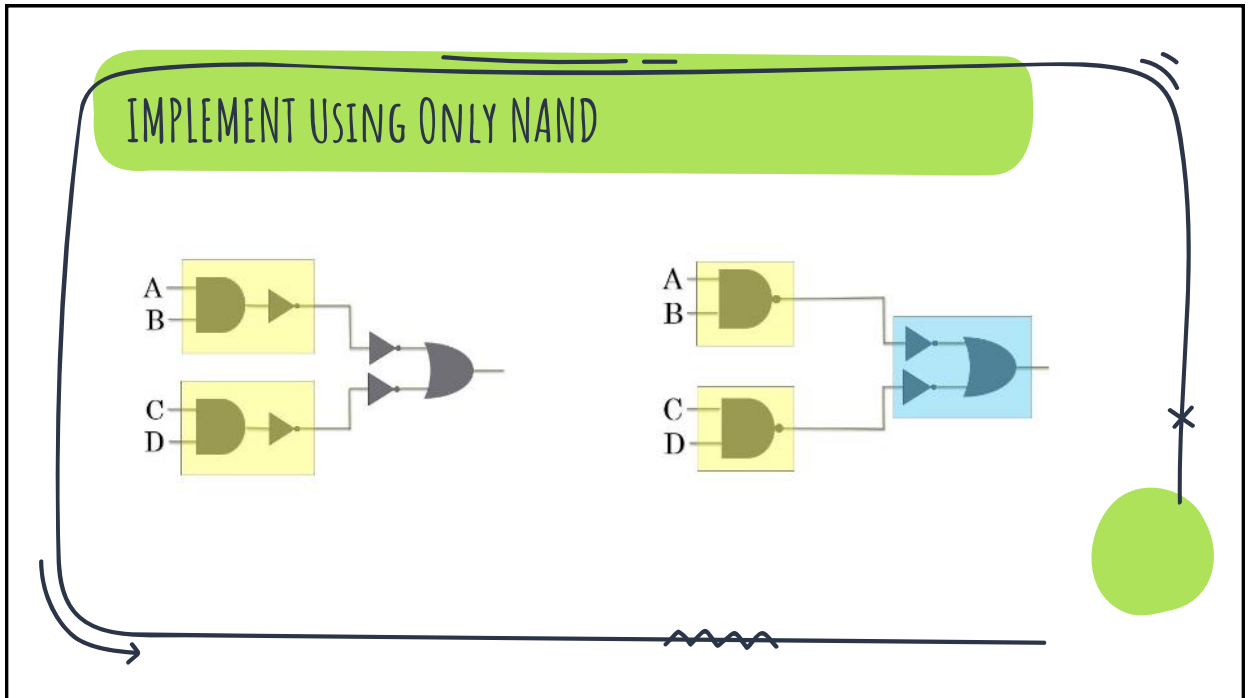


50

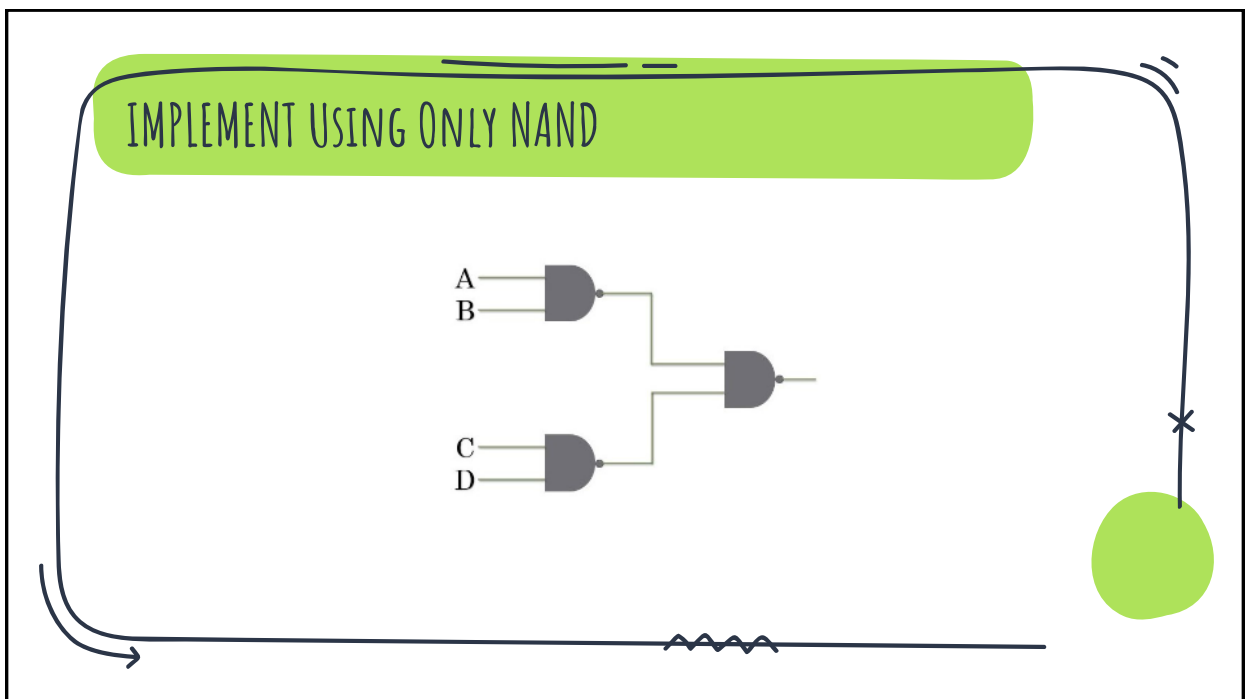
IMPLEMENT USING ONLY NAND



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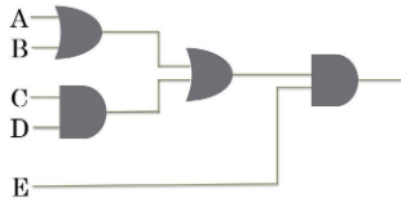


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EXAMPLE 2: IMPLEMENT USING ONLY NAND



$$(A+B+CD)E$$

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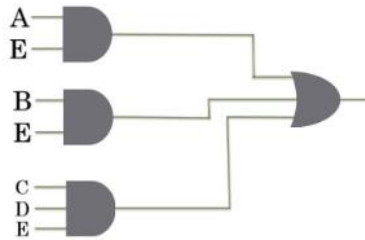
EXAMPLE 2:

- ✗ The original circuit implements the boolean function :
 $(A+B+CD)E$
- ✗ Convert to Sum of Products(SOP).
- ✗ In this case, simply open parenthesis.
- ✗ We get: $AE + BE + CDE$.
- ✗ Since this Boolean equation is in SOP form, the circuit for this equation will be in a standard two-level implementation

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EXAMPLE 2: IMPLEMENT USING ONLY NAND – METHOD 2

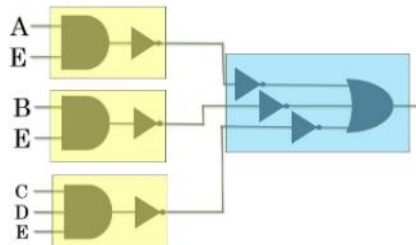
$$X = AE + BE + CDE.$$



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EXAMPLE 2: IMPLEMENT USING ONLY NAND – METHOD 2

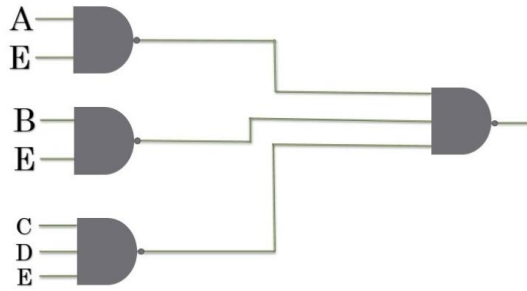
$$X = AE + BE + CDE.$$



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EXAMPLE 2: IMPLEMENT USING ONLY NAND – METHOD 2

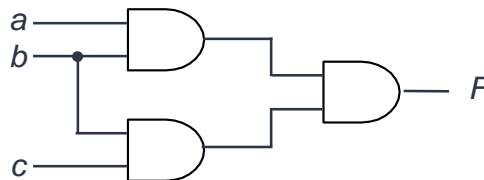
$$X \text{ } AE + BE + CDE.$$



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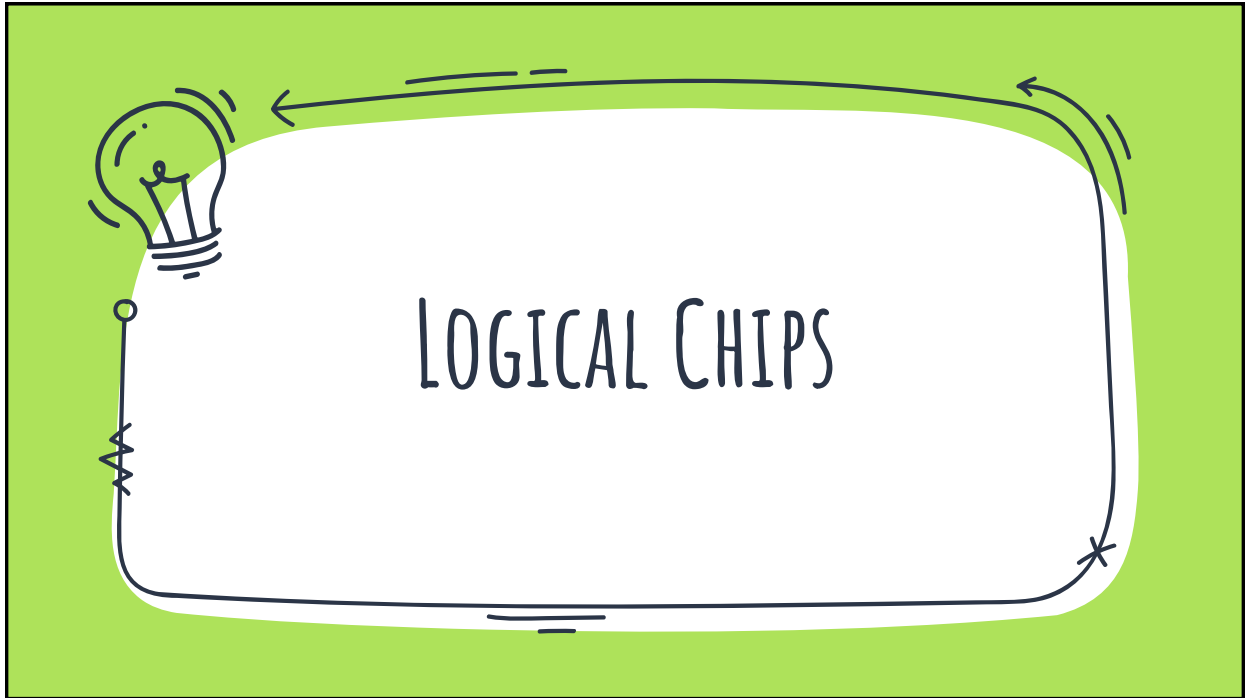
Home Task

Using ONLY NAND gates, draw a schematic for the following function: $F = (a.b) + (b.c)$

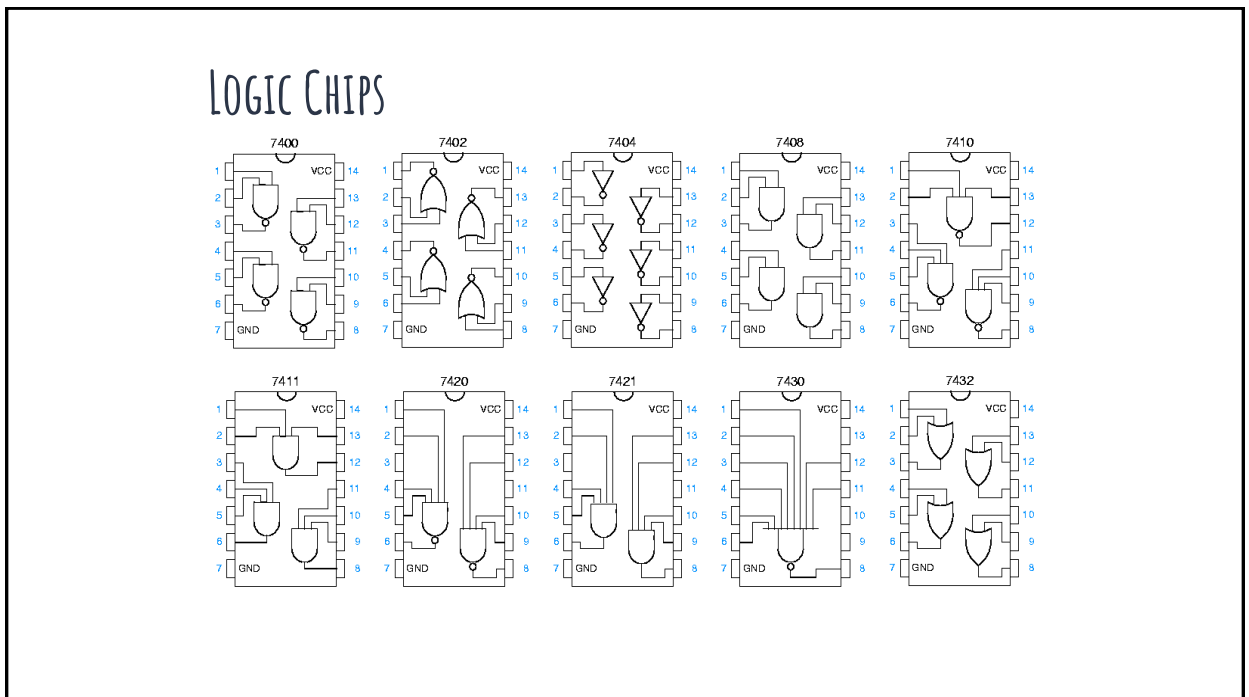


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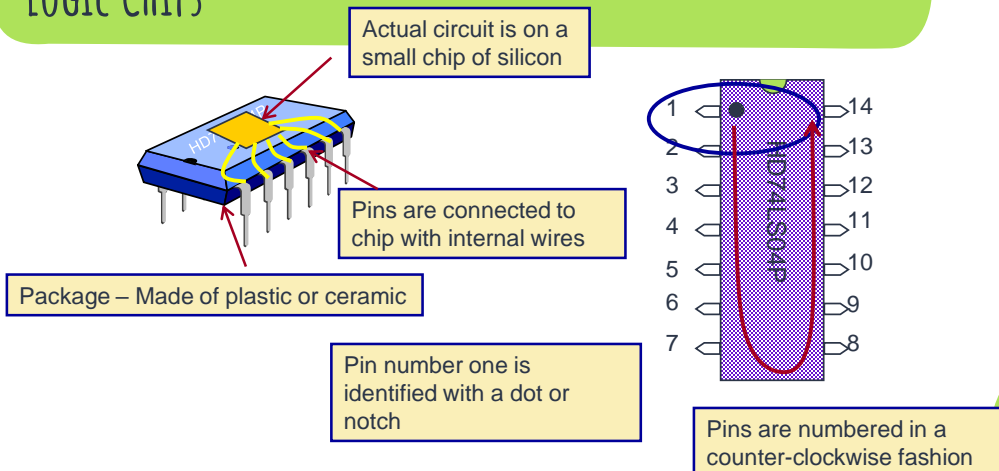
70

LOGIC CHIPS

- X Digital Electronics devices are usually in a chip format.
- X The chip is identified with a part number or a model number.
- X A standard series starts with numbers 74, 4, or 14.
 - X 7404 is an inverter
 - X 7408 is an AND
 - X 7432 is an OR
 - X 4011B is a NAND

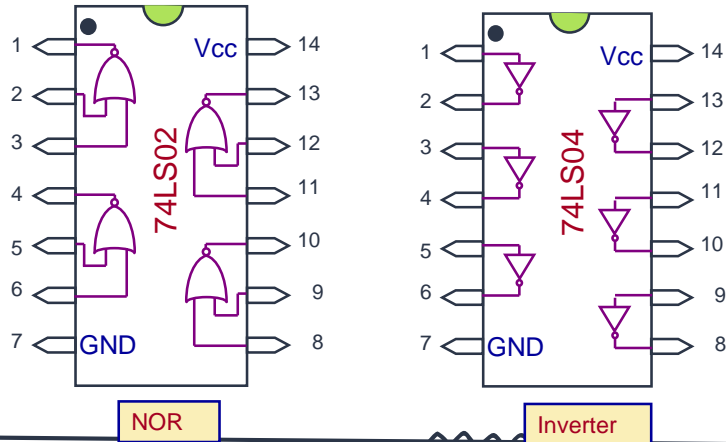
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LOGIC CHIPS



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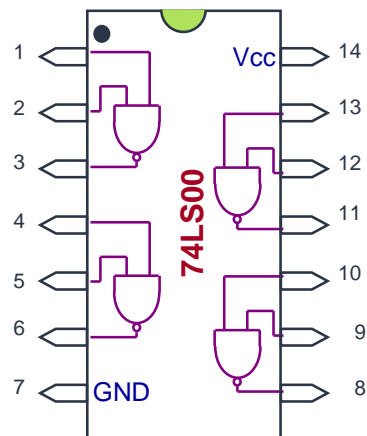
LOGIC CHIPS: EXAMPLES



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LOGIC CHIPS: EXAMPLES

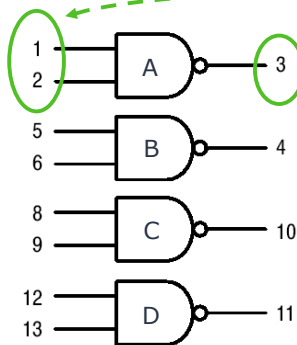
NAND Gates



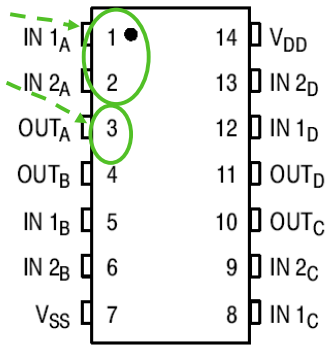
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LOGIC CHIPS

Quad 2-Input NAND Gate

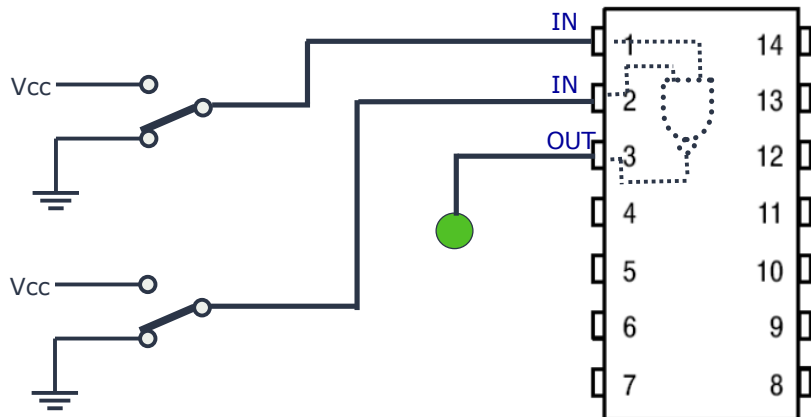


Quad 2-Input NAND Gate



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LOGIC CHIPS



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INTEGRATED CIRCUITS (IC)

- X Small Scale Integration (SSI)
 - X 10s of logic gates (previous examples) (late 1960s)
- X Medium Scale Integration (MSI)
 - X 10s to 1000 of logic gates (Introduced in late 1960s)
- X Large Scale Integration (LSI)
 - X 10000s of logic gates (Introduced in early 1970s)
- X Very Large Scale Integration (VLSI)
 - X Millions of logic gates (Introduced in early 1970s)
- X Ultra Large Scale Integration (ULSI)

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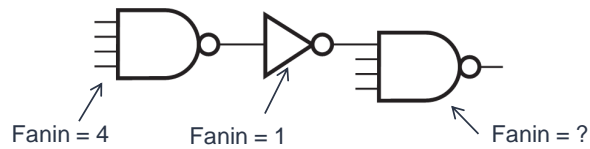
DIGITAL LOGIC FAMILIES

- X TTL
 - X Transistor-Transistor Logic
- X ECL
 - X Emitter Coupled Logic
- X MOS
 - X Metal Oxide Semi Conductor
- X CMOS
 - X Complementary Metal Oxide Semi Conductor

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FAN-IN

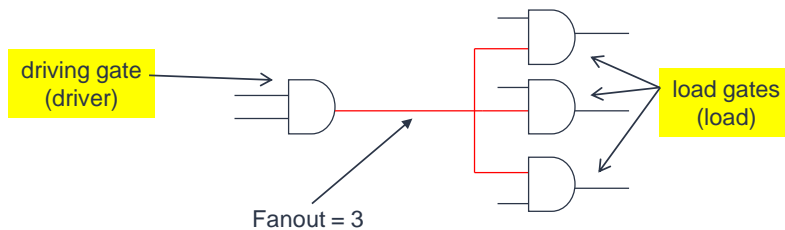
- ✗ **Fan in** of a gate is the number of inputs to the gate
- ✗ A 3-input OR gate has a fan-in = 3



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FAN-OUT

- ✗ **Fan out** of a gate is the number of gates that it can drive
- ✗ The driven gate is called a load



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REFERENCES

- X Chapter 2 – Digital Design Morris Mano
- X Digital Logic Circuits – V.D. Agrawal, Auburn University, Auburn
- X Logic System Design – Seattle Pacific University
- X Template is taken from slides carnival.

Slides Carnival