ACM/CS 114 Parallel algorithms for scientific applications

Michael A. G. Aïvázis

California Institute of Technology

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Point to point communication

to send a message

```
int MPI_Send(
void* buffer, int count, MPI_Datatype datatype,
int destination, int tag, MPI_Comm communicator
);
```

▶ to receive a message

```
int MPI_Recv(
void* buffer, int count, MPI_Datatype datatype,
int source, int tag, MPI_Comm communicator
);
```

- the tag enables choosing the order you may receive pending messages
- but for a given (source,tag,communicator) messages are received in the order they were sent
- ► receiving via wildcards: MPI_ANY_SOURCE and MPI_ANY_TAG
- in standard communication mode, sending and receiving messages are blocking, so the function does not return until you can safely access the buffer
 - ▶ to read, free, etc.



Communication modes

- in standard mode, the specification does not explicitly mention buffering strategy
 - buffering messages would remove some of the access constraints but it requires time and storage for the multiple copies
 - portability across implementations implies conservative assumptions about the order of initiation of sends and receives to avoid deadlock
- in ready mode, you must post a receive before the matching send can be initiated
 - ▶ MPI_Rsend, MPI_Rrecv
- in buffered mode, sends can be initiated, and may complete, regardless of when the matching receive is initiate
 - ▶ MPI_Bsend, MPI_Brecv
- in synchronous mode, sends can be initiated regardless of whether the matching receive has been initiated, but the send will not return until the message has been received
 - ▶ MPI_Ssend.MPI_Srecv



Asynchronous communication

▶ there are non-blocking versions of all these

```
int MPI_Isend(
void* buffer, int count, MPI_Datatype datatype,
int destination, int tag,
MPI_Comm communicator, MPI_Request* request
);
```

- faster, but you must take care to not access the message buffers until the messages have been delivered
- ▶ more details later in the course, as needed
- for sends
 - standard mode: MPT_Isend
 - ▶ ready mode: MPI_Irsend
 - buffered mode: MPI_Ibsend
 - ▶ synchronous mode: MPI_Issend
- ▶ only one call for receives: MPI_Irecv
- extra request argument to check for completion of the request
 - ► MPI_Test, MPI_Wait and their relatives



Creating communicators and groups

- communicators and groups are intertwined
 - you cannot create a group without a communicator
 - you cannot create a communicator without a group
- ▶ the cycle is broken by MPI_COMM_WORLD

```
#include <mpi.h>
  int main(int argc, char* argv[]) {
      /* declare a communicator and a couple of groups */
      int loner = 0;
      MPI Comm workers;
      MPI_Group world_grp, workers_grp;
9
      /* initialize MPI; for brevity all status checks are omitted */
10
      MPI Init(&argc, &argv);
      /* get the world communicator to build its group */
      MPI Comm group (MPI COMM WORLD, &world grp);
14
      /* build another group by excluding a process */
16
      MPI Group excl (world grp, 1, &loner, &workers grp);
18
      /* now build a communicator out of the processes in workers grp */
      MPI Comm create (MPI COMM WORLD, worker grp, &workers);
20
      /* etc.... */
      /* shut down MPI */
24
      MPI Finalize():
26
      return 0:
```

Manipulating communicators and groups

releasing resources

```
int MPI_Group_free(MPI_Group* group);
int MPI_Comm_free(MPI_Comm* communicator);
int MPI_Comm_disconnect(MPI_Comm* communicator);
```

you can make a new group by adding or removing processes from an existing one

```
int MPI_Group_incl(
    MPI_Group grp, int n, int* ranks, MPI_Group* new_group);
int MPI_Group_excl(
    MPI_Group grp, int n, int* ranks, MPI_Group* new_group);
```

or by using set operations

```
int MPI_Group_union(
    MPI_Group grp1, MPI_Group grp2, MPI_Group* new_group);
int MPI_Group_intersection(
    MPI_Group grp1, MPI_Group grp2, MPI_Group* new_group);
int MPI_Group_difference(
    MPI_Group grp1, MPI_Group grp2, MPI_Group* new_group);
```