ACM/CS 114 Parallel algorithms for scientific applications

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Unstructured grids

- solving differential equations using structured grids
 - ▶ finite differences
 - finite volumes
- ▶ finite elements take a different approach
 - subdivide the domain into elements
 - construct basis functions over each element
 - ⊳
- a simplicial mesh
- what follows is an overview of very common but surprisingly hard problems that occur in simulations of physical systems

Finite elements

- ▶
- 1

Meshing

- creating simplicial meshes is a surprisingly hard problem
- ▶ in two dimensions
 - there is a well posed mathematical problem with some guarantees that a triangulation can be generated for sufficiently well defined domains
 - given a set of vertices, one can always find a *Delaunay* triangulation where no vertex falls in the interior of the circumcircle of any triangle
 - the quality of the triangulation can be improved by adding extra vertices
 - for practical purposes, this is sufficient to handle most cases
 - the problem is rather challenging numerically, with high incidence of round-off errors, overflows and underflows
 - certain geometrical tests are so sensitive that robust packages employ exact arithmetic
- similar guarantees exist when triangulating surfaces embedded in three dimensional space
- ▶ in three dimensions the problem is theoretically intractable at this point
 - meshing packages are full of heuristics to avoid the known pathological cases
- no known parallel implementations



Parallelization

- typically, the finest grain task is the element update that models the response of the material to the deformation induced by the motion of the vertices
 - ▶ that's where all the physics is
 - computationally intensive for non-trivial responses
 - usually involves solving a complicated variational problem, such as energy minimization
 - little or no interaction with neighboring elements, although non-local updates are becoming more common
 - e.g. thin shells, fracture based on deformation energy in a neighborhood
- coarse grain tasks are defined through graph partitioning that decomposes the mesh into smaller subgraphs
 - with uniform element and node distributions
 - well characterized boundaries
- performance characteristics
 - computational cost grows with the number of elements
 - communication cost grows with the number of nodes on partition boundaries



Graph partitioning

- mesh domain decomposition is mathematically identical to graph partitioning
 - a hard computational problem
 - with wide applicability to many practical problems
 - hence, there are many excellent algorithms and software packages that solve most problems of practical interest
- algorithms fall into two broad categories
 - ▶ those that use geometrical information, such as nodal coördinates
 - those that rely only on graph connectivity
 - breadth first search, spectral bisection, Kernighan/Lin
- comparison metrics
 - speed of partitioning: important since meshes tend to be large
 - number of cut edges
 - for finite elements: important to minimize the number of nodes on partition boundaries, since it is directly linked to the communication cost of the calculation
- ▶ METIS from George Karypis' lab at the University of Minnesota
 - multi-level Kernighan/Lin algorithms, other heuristics
 - ParMETIS is the parallel implementation



Uniform subdivision

- high fidelity simulations require high resolution
 - ▶ i.e. small elements, which implies large element counts
 - many interesting problems require tens or hundreds of millions of elements before they converge
 - creating input meshes of this size sequentially is impractical
- one possibility is parallel uniform subdivision
 - ▶ the input mesh is partitioned and distributed among processes
 - each process subdivides its simplices
 - edges get split in two, triangles in four, tetrahedra in eight
 - mesh quality is not affected significantly: only two new classes of simplices are introduced, regardless of the number of subdivision levels
 - the formation of new entities on partition boundaries implies the recalculation of the communications maps among processes







- ▶ can be done in linear time
 - scales well and can create enormous meshes (billions of elements)

Non-uniform subdivision and adaptive remeshing

- ▶ there are many good non-uniform subdivision algorithms
 - e.g. longest edge subdivision: each tetrahedron that has an edge longer than some threshold gets subdivided by splitting that edge in two and connecting the new node to the two non-adjacent vertices
- they tend to be difficult to parallelize because they induce one-sided topological changes on the partition boundaries that must be communicated to the neighboring process
 - hence they are rarely used in large scale calculations
 - this is an architectural limitation, not an algorithmic one!
 - it requires restructuring the solver to support non-trivial communication with its neighbors
- similar considerations apply to adaptive remeshing
 - where a process discovers that it requires a higher element density to satisfy some correctness criterion
 - but difficult to implement when the refinement reaches the partition boundary

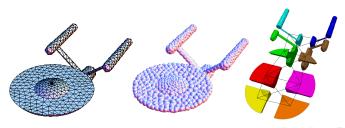


Contact detection

- modeling impact requires detecting when bodies come into contact and computing appropriate forces
 - both are difficult problems, especially for non-smooth, distributed meshes
 - ▶ the computations involve only elements on the free boundary
- the element updates and contact search have very different communication patterns
 - contact is global, since any two surface triangles can collide
 - ▶ and it is expensive, since it involves geometrical queries
 - ▶ it is best solved as an auxiliary parallel problem with its own partitioning
 - boundary partitioned using geometric criteria, such as axis-aligned cutting planes
 - requires frequent re-balancing as mesh geometry and topology change
- ▶ in *surface driven* contact detection
 - the boundary is partitioned and the centers of all the faces are inserted in some kind of spatial data structure
 - so that coarse calculations about contact candidates can be made efficiently
 - typically bucket or sparse-bucket arrays that support orthogonal range queries (ORQ)
 - candidates can be checked pairwise for actual contact using ray-triangle and edge-triangle intersection algorithms

Volume based contact detection

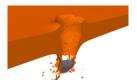
- volume based approaches replace the elements on the boundary with equivalent spheres
- a variety of geometrical criteria can be used:
 - ► equivalent volume
 - diameter determined by the longest edge
- contact detection can take place very efficiently using ORQ algorithms
- simple spring models resolve the contact forces
 - based on simple potentials that take into account the sphere inter-penetration



Fracture and fragmentation

- when the material is under sufficient tension, modeling the physics correctly requires topological changes to the mesh
 - the material fractures when it is energetically favorable for an opening to form between two elements that are being pulled apart
- modeling this typically involves cohesive elements
 - prismatic elements that are inserted in the place of what used to be the common face of two tetrahedra and absorb excess energy
- when some threshold is exceeded, the cohesive element fractures and becomes free surface
- cohesive element insertion involves node, edge and face splitting
- difficulties arise as cracks run up to process boundaries and must be propagated across
- eventually, fragments form as cracks disconnect portions of the mesh





Element erosion

- mesh distortions induced by large deformations
 - degrade the numerical accuracy of the calculation
 - force the timestep sizes to impractical values
 - eventually cause the simulation to fail





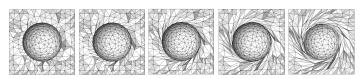
- ▶ a strategy in wide use is to *erode* badly distorted elements
 - using physically motivated erosion criteria that model fracture mechanics
 - trivial to implement sequentially
 - less so in parallel, as it involves geometrical and topological changes to the mesh
- candidates are selected based on a geometric criterion of the element quality
- erosion takes place when the deformation energy in a suitable neighborhood exceeds the fracture energy

Mesh optimization

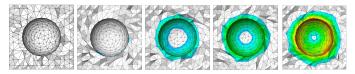
- ▶ the numerical accuracy of finite element calculations is determined in part by the quality of the mesh
 - the initial triangulation generated by the mesher
 - distortions induced by large deformations
- element quality is intuitively related to its deviation from an equilateral tetrahedron
 - there are a dozen or so pathological cases that involve either extremely small or extremely large internal angles
- optimizing mesh quality requires
 - an element quality metric that places a high penalty on the pathological cases
 - geometrical optimization: redistributing the mesh nodes so that it is possible to construct a defect-free triangulation
 - topological optimization: adjusting the way nodes are connected to form simplices
 - control over the element size, which is linked directly to both the numerical accuracy and the computational cost of the simulation
 - a sufficiently well characterized description of the boundary so that when new nodes are inserted near the boundary they do not induce distortions

Mesh optimization examples

- consider a mesh with an embedded sphere that is undergoing rigid rotation
- a quarter of a turn is sufficient to distort the mesh



 continuous remeshing maintains good mesh quality; after three whole revolutions



- ► challenges:
 - keep numerical diffusion under control
 - implement in parallel



Implicit representations of surfaces

- simulating phenomena that involve the interaction of solids and fluids has its own challenges
 - fluid solvers are typically Eulerian, with structured grid that is fixed in space while the fluid flows through it
 - whereas solid solvers are typically Lagrangian, with an unstructured grid that moves along with the body it describes
- coupling requires some form of information exchange that enables
 - the solid solver to inform the fluid about the position and velocity of its evolving boundary
 - the fluid solver to exert a force on the solid through its pressure field on the boundary
- the solution involves constructing level sets, and computing the distance to the solid, the closest node and the closest face for each point on the grid







Summary of open problems

- many interesting physics problems require making topological changes to the simplicial meshes, and propagating them across process boundaries
 - good sequential algorithms exist, but they are difficult to parallelize
 - naïve parallelization is almost always prohibitively expensive
 - because maintaining coherence and consistency of the global information requires a lot of communication
- changing the solver architecture may be key to efficient parallel implementations
 - evolving mesh topology requires complicated interactions among processes
 - event based solutions enable such interactions
 - and improve simulation control by external agents
 - better monitoring through probes and sensors
 - instrumentation for virtual experiments