

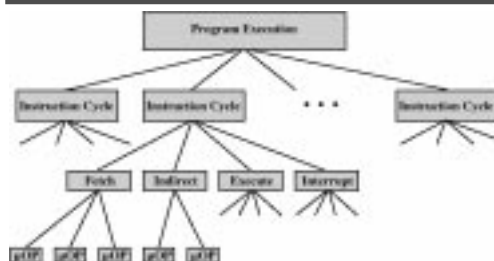
William Stallings Computer Organization and Architecture

Chapter 14 Control Unit Operation

Micro-Operations

- A computer executes a program
- Fetch/execute cycle
- Each cycle has a number of steps
 - see pipelining
- Called micro-operations
- Each step does very little
- Atomic operation of CPU

Constituent Elements of Program Execution



Fetch - 4 Registers

- Memory Address Register (MAR)
 - Connected to address bus
 - Specifies address for read or write op
- Memory Buffer Register (MBR)
 - Connected to data bus
 - Holds data to write or last data read
- Program Counter (PC)
 - Holds address of next instruction to be fetched
- Instruction Register (IR)
 - Holds last instruction fetched

Fetch Sequence

- Address of next instruction is in PC
- Address (MAR) is placed on address bus
- Control unit issues READ command
- Result (data from memory) appears on data bus
- Data from data bus copied into MBR
- PC incremented by 1 (in parallel with data fetch from memory)
- Data (instruction) moved from MBR to IR
- MBR is now free for further data fetches

Fetch Sequence (symbolic)

- t1: MAR \leftarrow (PC)
- t2: MBR \leftarrow (memory)
- PC \leftarrow (PC) + 1
- t3: IR \leftarrow (MBR)
- (tx = time unit/clock cycle)
- or
- t1: MAR \leftarrow (PC)
- t2: MBR \leftarrow (memory)
- t3: PC \leftarrow (PC) + 1
- IR \leftarrow (MBR)

Rules for Clock Cycle Grouping

- Proper sequence must be followed
 - $MAR \leftarrow (PC)$ must precede $MBR \leftarrow (memory)$
- Conflicts must be avoided
 - Must not read & write same register at same time
 - $MBR \leftarrow (memory)$ & $IR \leftarrow (MBR)$ must not be in same cycle
- Also: $PC \leftarrow (PC) + 1$ involves addition
 - Use ALU
 - May need additional micro-operations

Indirect Cycle

- $MAR \leftarrow (IR_{address})$ - address field of IR
- $MBR \leftarrow (memory)$
- $IR_{address} \leftarrow (MBR_{address})$
- MBR contains an address
- IR is now in same state as if direct addressing had been used
- (What does this say about IR size?)

Interrupt Cycle

- t1: $MBR \leftarrow (PC)$
- t2: $MAR \leftarrow \text{save-address}$
- $PC \leftarrow \text{routine-address}$
- t3: $memory \leftarrow (MBR)$
- This is a minimum
 - May be additional micro-ops to get addresses
 - N.B. saving context is done by interrupt handler routine, not micro-ops

Execute Cycle (ADD)

- Different for each instruction
- e.g. ADD R1,X - add the contents of location X to Register 1 , result in R1
- t1: $MAR \leftarrow (IR_{address})$
- t2: $MBR \leftarrow (memory)$
- t3: $R1 \leftarrow R1 + (MBR)$
- Note no overlap of micro-operations

Execute Cycle (ISZ)

- ISZ X - increment and skip if zero
 - t1: $MAR \leftarrow (IR_{address})$
 - t2: $MBR \leftarrow (memory)$
 - t3: $MBR \leftarrow (MBR) + 1$
 - t4: $memory \leftarrow (MBR)$
 - if $(MBR) == 0$ then $PC \leftarrow (PC) + 1$
- Notes:
 - if is a single micro-operation
 - Micro-operations done during t4

Execute Cycle (BSA)

- BSA X - Branch and save address
 - Address of instruction following BSA is saved in X
 - Execution continues from X+1
 - t1: $MAR \leftarrow (IR_{address})$
 - $MBR \leftarrow (PC)$
 - t2: $PC \leftarrow (IR_{address})$
 - $memory \leftarrow (MBR)$
 - t3: $PC \leftarrow (PC) + 1$

Functional Requirements

- Define basic elements of processor
- Describe micro-operations processor performs
- Determine functions control unit must perform

Basic Elements of Processor

- ALU
- Registers
- Internal data paths
- External data paths
- Control Unit

Types of Micro-operation

- Transfer data between registers
- Transfer data from register to external
- Transfer data from external to register
- Perform arithmetic or logical ops

Functions of Control Unit

- Sequencing
 - Causing the CPU to step through a series of micro-operations
- Execution
 - Causing the performance of each micro-op
- This is done using Control Signals

Control Signals (1)

- Clock
 - One micro-instruction (or set of parallel micro-instructions) per clock cycle
- Instruction register
 - Op-code for current instruction
 - Determines which micro-instructions are performed

Control Signals (2)

- Flags
 - State of CPU
 - Results of previous operations
- From control bus
 - Interrupts
 - Acknowledgements

Control Signals - output

- Within CPU
 - Cause data movement
 - Activate specific functions
- Via control bus
 - To memory
 - To I/O modules

Example Control Signal Sequence - Fetch

- $MAR \leftarrow (PC)$
 - Control unit activates signal to open gates between PC and MAR
- $MBR \leftarrow (\text{memory})$
 - Open gates between MAR and address bus
 - Memory read control signal
 - Open gates between data bus and MBR

Internal Organization

- Usually a single internal bus
- Gates control movement of data onto and off the bus
- Control signals control data transfer to and from external systems bus
- Temporary registers needed for proper operation of ALU

Hardwired Implementation (1)

- Control unit inputs
- Flags and control bus
 - Each bit means something
- Instruction register
 - Op-code causes different control signals for each different instruction
 - Unique logic for each op-code
 - Decoder takes encoded input and produces single output
 - n binary inputs and 2^n outputs

Hardwired Implementation (2)

- Clock
 - Repetitive sequence of pulses
 - Useful for measuring duration of micro-ops
 - Must be long enough to allow signal propagation
 - Different control signals at different times within instruction cycle
 - Need a counter with different control signals for t_1 , t_2 etc.

Problems With Hard Wired Designs

- Complex sequencing & micro-operation logic
- Difficult to design and test
- Inflexible design
- Difficult to add new instructions

Required Reading

■ Stallings chapter 14
