

Module 7 - Recursion

Example 1

$$T(n) = \begin{cases} 1 & \text{if } n == 0 \\ T(n-1) + 2 & \text{if } n > 0 \end{cases}$$

$\left\{ \begin{array}{l} T(n) = T(n-1) + 2 \\ T(n-1) = T(n-2) + 2 \\ T(n-2) = T(n-3) + 2 \\ \vdots \end{array} \right.$

$$\begin{aligned} T(n) &= T(n-1) + 2 = T(n-1) + 1 \times 2 && \text{step 1} \\ &= T(n-2) + 2 + 2 = T(n-2) + 2 \times 2 && \text{step 2} \\ &= T(n-3) + 2 + 2 + 2 = T(n-3) + 3 \times 2 && \text{step 3} \\ &= \dots && \vdots \\ &= T(n-k) + k \cdot 2 && \text{step K} \\ &= \dots && \vdots \\ &= T(n-n) + n \cdot 2 && \text{step n} \\ &\quad \parallel \\ &\quad T(0) = 1 \end{aligned}$$

$$T(n) = T(0) + n \cdot 2 = 1 + 2n$$

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$$T(n) = 2n + 1 \Rightarrow T(n) = O(n)$$

Example 2:

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 2T(\frac{n}{2}) + 2n & \text{if } n > 1 \end{cases}$$

$$\begin{aligned} T(2^k) &= 2T(2^{k-1}) + 2 \cdot 2^k \\ T(2^{k-1}) &= 2T(2^{k-2}) + 2 \cdot 2^{k-1} \\ T(2^{k-2}) &= 2T(2^{k-3}) + 2 \cdot 2^{k-2} \\ &\vdots \end{aligned}$$

Let $n = 2^k$, then $\frac{n}{2} = 2^{k-1}$

$$\begin{aligned} T(n) &= 2T(2^{k-1}) + 2 \cdot 2^k = 2^1 T(2^{k-1}) + 1 \cdot 2 \cdot 2^k && \text{step 1} \\ &= 2(2T(2^{k-2}) + 2 \cdot 2^{k-1}) + 2 \cdot 2^k = 2^2 T(2^{k-2}) + 2 \cdot 2 \cdot 2^k && \text{step 2} \\ &= 2^3(2T(2^{k-3}) + 2 \cdot 2^{k-2}) + 2 \cdot 2 \cdot 2^k = 2^3 T(2^{k-3}) + 3 \cdot 2 \cdot 2^k && \text{step 3} \\ &= \dots \\ &= \underbrace{2^k}_{\downarrow n} T(\underbrace{2^{k-k}}_{\downarrow 1}) + \underbrace{k}_{\downarrow \log n} \cdot \underbrace{2 \cdot 2^k}_{\downarrow n} && \text{step k} \end{aligned}$$

if $n = 2^k$
then $k = \log n$

$$T(n) = 2n \log n + n \Rightarrow T(n) = O(n \log n)$$