



CPT-281 - Introduction to Data Structures with C++

Module 5

Stacks

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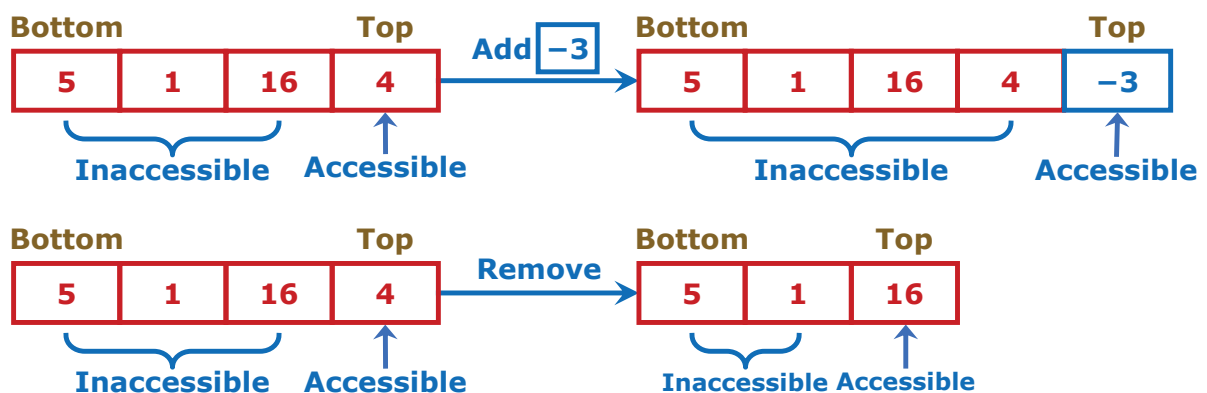
• Stacks

→ **Stacks** are abstract data structures (ADT) with the property that only the top element is accessible.

- Each time, you add a new element onto the top of the stack.
- Only the top element is accessible.
- Each time, you can only remove an element from the top of the stack.



→ Stacks are based on the **LIFO (Last-In-First-Out)** principle.



- **Class-Member Functions in Stack Class**

→ Theoretically, stacks only support 7 functions.

Functions	Behavior
<code>size_t size() const;</code>	Returns the number of elements stored in the stack.
<code>bool empty() const;</code>	Tests whether the stack is empty.
<code>T& top();</code>	Returns the top element in the stack (l-value).
<code>const T& top() const;</code>	Returns the top element in the stack (r-value).
<code>void push(const T&);</code>	Adds a new element onto the top of the stack.
<code>void pop();</code>	Removes the top element from the stack.
<code>void clear();</code>	Removes all elements from the stack.

- **Using vector to implement stack**

→ The top of the stack is the rear end of the vector.

- `.size()`, `.empty()`, and `clear()` functions are the same as vector.
- `.top()` ➔ `.back()`
- `.pop()` ➔ `.pop_back()`
- `.push(item)` ➔ `.push_back(item)`

[See sample code]

- **Using linked list to implement stack**

→ Singly-linked list or doubly-linked list?

Singly-linked list

→ In class data fields, we need to keep reference of **top node only** or **top and bottom nodes**?

Top only

[See sample code]

- **Discussion**

→ What are the applications of stacks?

Undo/redo system

Back/forward system

Code compilation

→ Graph traversal

Depth-First Search (DFS)

Stack is an **intermediate data structure** in DFS.

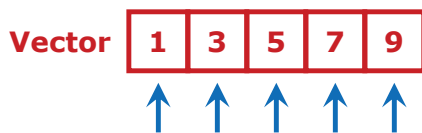
Breadth-First Search (BFS)

Queue is an **intermediate data structure** in BFS.

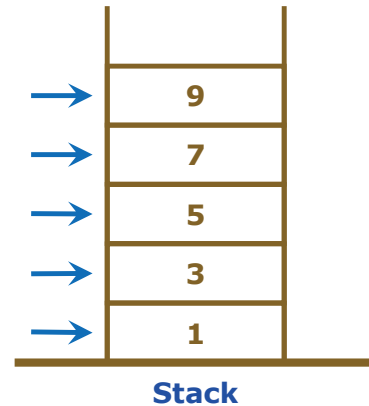
→ Stacks are useful to solve algorithm problems.

- **[Example 1] Reversing elements in a linear container**

→ Using a stack to reverse a vector of integers



- 1) Create an empty stack.
- 2) Push each element in the vector onto the stack.
- 3) Pop each element from the stack to form the reversed vector.



```

1  /** Reverses a vector of integers.
2      @param vec: vector of integers to reverse
3  */
4  void reverse(vector<int>& vec) {
5      List_Stack<int> stk;
6      for (size_t i = 0; i < vec.size(); i++) { stk.push(vec.at(i)); }
7      for (size_t j = 0; j < vec.size(); j++) {
8          vec[j] = stk.top();
9          stk.pop();
10     }
11 } // Time complexity: O(n)

```

- **[Example 2] Finding palindromes**

→ A **palindrome** is a string that reads the same in either direction—left to right or right to left.

e.g., "I saw I was I"

→ Using a stack to test whether a string is palindromic

- 1) Create an empty stack.
- 2) Push the characters in the string onto the stack.
- 3) Construct a string using the popped characters from the stack.
- 4) Compare the constructed string against the original string.

```

1  /** Tests whether a string is a palindrome.
2      @param s: string to test
3      @return: {true} if the string is palindromic; {false} otherwise
4  */
5  bool is_palindromic(const string& s) {
6      List_Stack<char> stk;
7      for (string::const_iterator it = s.begin(); it != s.end(); it++) { stk.push(*it); }
8      string reversed;
9      while (!stk.empty()) {
10         reversed.push_back(stk.top());
11         stk.pop();
12     }
13     return reversed == s;
14 } // Time complexity: O(n)

```


• Arithmetic Expressions

→ Arithmetic expressions consist of operands and operators.

→ 3 types of arithmetic expressions

Prefix Expression	Infix Expression	Postfix Expression
c+ab	c(a+b)	ab+c*

→ Evaluating arithmetic expressions

Prefix Expression	Infix Expression	Postfix Expression	Value
* 4 7	4 * 7	4 7 *	28
* 4 + 7 2	4 * (7 + 2)	4 7 2 + *	36
- * 4 7 20	4 * 7 - 20	4 7 * 20 -	8
+ 3 / * 4 7 2	3 + 4 * 7 / 2	3 4 7 * 2 / +	17

→ Advantages of using postfix expressions instead of infix expressions:

- There are **no parentheses** in postfix expressions.
- User does **not** need to use **precedence rules** to evaluate postfix expressions.

• [Example 4] Evaluating postfix expressions

→ Algorithm

stack $\leftarrow \emptyset$

while there are more tokens

token \leftarrow Get next token

if token.is_operand()

stack.push(token)

else // "token" is an operator.

right_operand \leftarrow stack.pop()

left_operand \leftarrow stack.pop()

result \leftarrow Evaluate the operation

stack.push(result)

endif

endwhile

return stack.top()

[Example]

"3 4 7 * 2 / +"

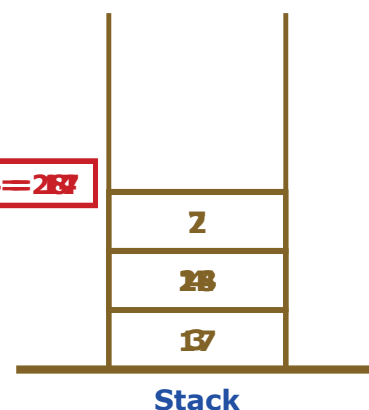


Evaluation result: 17

Evaluate $28 * 3 = 84$

right_operand = 14

left_operand = 28



```

1  /** Evaluates a postfix expression.
2      @param postfix: postfix expression to evaluate
3      @return: evaluation result
4      @throws exception: divide-by-zero
5  */
6  int eval_postfix(const string& postfix) {
7      istream iss(postfix);
8      List_Stack<int> stk;
9      string token; // Current token
10     while (iss >> token) {
11         if (isdigit(token.front())) { stk.push(stoi(token)); }
12         else {
13             int right = stk.top();
14             stk.pop();
15             int left = stk.top();
16             stk.pop();
17
18             // Supported operators
19             if (token == "+") { stk.push(left + right); }
20             if (token == "-") { stk.push(left - right); }
21             if (token == "*") { stk.push(left * right); }
22             if (token == "/" ) {
23                 if (!right) { throw exception("Divide by zero"); }
24                 stk.push(left / right);
25             }
26         }
27     }
28     return stk.top();
29 } // Time complexity: O(n)

```

• **[Example 5] Converting from infix expression to postfix expression**

→ **Algorithm**

while there are more tokens in infix expression

token ← Get next token

if token.is_operand()

 Append token to postfix expression

else if token == '('

 Push token onto the stack

else if token.is_operator()

while !stack.empty() and stack.top() != '('
 and precedence(token) ≥ precedence(stack.top())

 Pop top off stack and append to postfix expression

endwhile

 Push token onto the stack

else // "token" == ')'

while stack.top() != '('

 Pop top off stack and append to postfix expression

endwhile

 stack.pop() // Remove the opening parenthesis from the stack.

endif

endwhile

while !stack.empty()

 Append stack.top() to postfix expression

 stack.pop()

endwhile

- **[Example]** "a + b * (c + d / e)"

↑↑↑↑↑↑↑↑↑↑↑↑

Postfix: a b c d e / + * +

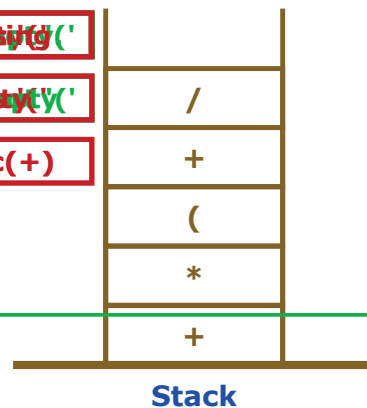
→ In the infix-to-postfix process, precedence rules must be used.

The precedence() function

Top of Stack is '('

Top of Stack is '/'

Prec(*) > Prec(+)



```

1  /** Returns the precedence of an operator.
2      @param oper: operator to return its precedence
3      @return: precedence
4      @throws exception: operator unsupported
5  */
6  int precedence(const string& oper) {
7      if (oper == "*" || oper == "/") { return 5; }
8      if (oper == "+" || oper == "-") { return 6; }
9      throw exception("Unsupported operator");
10 } // Time complexity: O(1)

```

→ In order to correctly parse the infix expression, all tokens must be surrounded by spaces on the left and on the right.

"(3 + 2) * 5" is correct.

"(3 + 2) * 5" is incorrect.

→ All the sample code assumes that the input arithmetic expression is valid, without errors.

```

1  /** Converts an infix expression to postfix expression.
2      @param infix_exp: infix expression to convert
3      @return: postfix expression converted from the infix expression
4  */
5  string infix_to_postfix(const string& infix_exp) {
6      istringstream iss(infix_exp);
7      ostringstream oss;
8      List_Stack<string> stk;
9      string token;
10     while (iss >> token) {
11         if (isdigit(token.front())) { oss << ' ' << token; }
12         else if (token == "(") { stk.push(token); }
13         else if (token == ")") {
14             while (stk.top() != "(") {
15                 oss << ' ' << stk.top();
16                 stk.pop();
17             }
18             stk.pop();
19         } else {
20             while (!stk.empty() && stk.top() != "(" && precedence(token) >= precedence(stk.top())) {
21                 oss << ' ' << stk.top();
22                 stk.pop();
23             }
24             stk.push(token);
25         }
26     }
27     while (!stk.empty()) {
28         oss << ' ' << stk.top();
29         stk.pop();
30     }
31     return oss.str();
32 } // Time complexity: O(n)

```