

CPT-281 - Introduction to Data Structures with C++

Module 12

Sorting Algorithms

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Sorting Algorithms

- → A sorting algorithm sorts the elements in an vector in <u>non-decreasing</u> (increasing) order.
- → Some sorting algorithms use <u>comparison</u> to sort the vectors; others do not use comparison.
- → In this class, you are only required to understand comparison sorting algorithms.

Comparison Sort

→ Only the comparison result can be used to sort the vector.

For example, for two elements a and b in the input vector, you cannot use the magnitude of a or b, but you can use their <u>comparison result</u>, either a < b, a > b, or a == b.

- **→** Comparison properties:
 - Totalness: for all a and b, either $a \le b$ or $a \ge b$.
 - Transitivity: if $a \le b$ and $b \le c$, then $a \le c$.

· In-Place Sort

- → If a sorting algorithm can directly manipulate the elements in the input vector without making additional copies of the input vector, the sorting algorithm is in-place.
- \rightarrow In-place sorting algorithms have space complexity of O(1); they only use <u>constant extra memory</u> to sort the input vector.

· Stable Sort

→ If a == b in the input vector, theoretically, either a or b may appear before the other in the sorted vector.

If a sorting algorithm can <u>guarantee</u> that for any a == b in the input vector, the relative order of a and b in the sorted vector is the same as in the input vector, the sorting algorithm is <u>stable</u>.

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Selection Sort

→ In the first iteration, you put the minimum value in the vector in the first place; in the second iteration, you put the second minimum value in the vector in the second place; and so on.

Demo of Selection Sort



• C++ Implementation of Selection Sort

```
void selection sort(vector<int>& vec) {
1
        for (size_t i = 0; i < vec.size(); i++) {</pre>
2
            // Stores the index of the min value in the rest of the vector.
3
4
            size_t min = i;
5
            // Find the min value in the rest of the vector.
6
            for (size_t j = i + 1; j < vec.size(); j++) {</pre>
7
8
                if (vec.at(j) < vec.at(min)) { min = j; }</pre>
9
            }
10
            // Swap vec.at(min) with vec.at(i) if they are not the same.
11
            if (min != i) { swap(vec.at(i), vec.at(min)); }
12
        }
13
14
   }
```

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Performance of Selection Sort

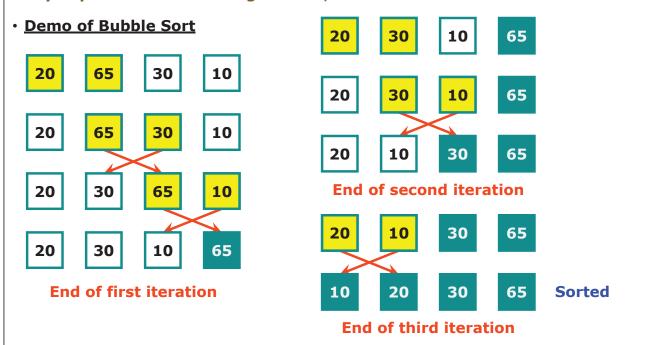
| Number of Comparisons | | | |
|-----------------------------------|----------|----------|--|
| Best Case Worst Case Average Case | | | |
| $O(n^2)$ | $O(n^2)$ | $O(n^2)$ | |

| Number of Swaps | | | |
|-----------------|------------|--|--|
| Best Case | Worst Case | | |
| O(1) | O(n) | | |

- → Selection sort is an <u>in-place</u> sorting algorithm.
- → Selection sort is a <u>stable</u> sorting algorithm.

- Bubble Sort
 - → Compare the adjacent pairs (e.g., vec[0] and vec[1], vec[1] and vec[2], vec[2] and vec[3], and so on) of elements. If they are out of order, swap them.

In the first iteration, you place the largest item; in the second iteration, you place the second largest item; and so on.



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C++ Implementation of Bubble Sort

```
void bubble sort I(vector<int>& vec) {
1
2
        for (size t i = 0; i < vec.size(); i++) {</pre>
             for (size_t j = 1; j < vec.size(); j++) {</pre>
3
                 if (vec.at(j) < vec.at(j - 1)) { // Out of order</pre>
4
5
                      swap(vec.at(j), vec.at(j - 1));
6
7
             }
8
        }
9
    }
```

- → Can we improve this solution (code)?
 - 1) We do not need to compare values that are already placed (at the end of the vector).
 - 2) If in any iteration, no swaps were ever occurred, it means _____?

 It means that the vector is already sorted.

Therefore, at the end of an iteration, if there were no swaps occurred in the iteration, then we can stop any further processing of the vector, since it is already sorted.

C++ Implementation of Bubble Sort (Improved)

```
void bubble_sort_II(vector<int>& vec) {
1
        for (size t i = 0; i < vec.size(); i++) {</pre>
2
            // Stores whether a swap occurs in this iteration.
3
4
            bool swapped = false;
5
6
            for (size t j = 1; j < vec.size(); j++) {</pre>
7
                if (vec.at(j) < vec.at(j - 1)) { // Out of order</pre>
8
                     swap(vec.at(j), vec.at(j - 1));
9
                     swapped = true;
                }
10
            }
11
12
13
            // If no swap occurred in this iteration,
            // then the vector is already sorted.
14
15
            if (!swapped) { return; }
        }
16
17
```

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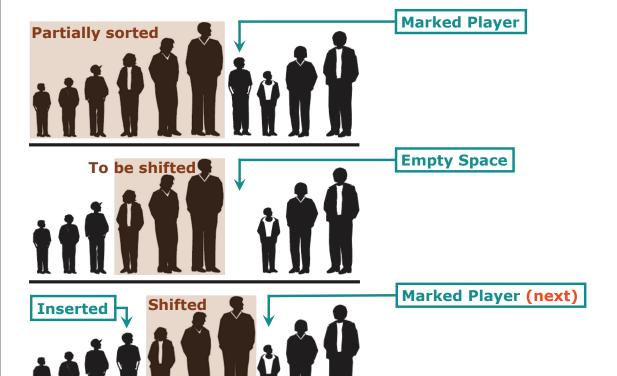
Performance of Bubble Sort

| Number of Comparisons | | | |
|-----------------------|-------------------|--------------|--|
| Best Case | Worst Case | Average Case | |
| O(n) | $O(n^2)$ | $O(n^2)$ | |

| Number of Swaps | | | |
|-----------------|------------|--|--|
| Best Case | Worst Case | | |
| O(1) | $O(n^2)$ | | |

- → Bubble sort is an <u>in-place</u> sorting algorithm.
- → Bubble sort is a <u>stable</u> sorting algorithm.
- → In general, bubble sort is a bad sorting algorithm (slow).
 - It uses up to $O(n^2)$ swaps.
 - Swap is a <u>slow operation</u>.

· Insertion Sort



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C++ Implementation of Insertion Sort

```
void insertion_sort(vector<int>& vec) {
1
2
        for (size_t mark = 1; mark < vec.size(); mark++) {</pre>
3
            int key = vec.at(mark), j;
            for (j = mark - 1; j >= 0 && vec.at(j) > key; j--) {
4
5
                vec.at(j + 1) = vec.at(j); // Data shift
6
7
            vec.at(j + 1) = key;
8
        }
9
```

• Performance of Insertion Sort

| Number of Comparisons | | | |
|-----------------------------------|----------|----------|--|
| Best Case Worst Case Average Case | | | |
| O(n) | $O(n^2)$ | $O(n^2)$ | |

| Number of Shifts | | | |
|------------------|------------|--|--|
| Best Case | Worst Case | | |
| O(1) | $O(n^2)$ | | |

- → Insertion sort is an <u>in-place</u> sorting algorithm.
- → Insertion sort is a stable sorting algorithm.

Summary of Quadratic Sorting Algorithms

- → Selection sort, bubble sort, and insertion sort are <u>in-place</u> sorting algorithms.
- → Selection sort, bubble sort, and insertion sort are <u>stable</u> sorting algorithms.
- \rightarrow Quadratic: $O(n^2)$ time complexity

External Learning Resource

- https://www.toptal.com/developers/sorting-algorithms
- https://www.geeksforgeeks.org/sorting-algorithms

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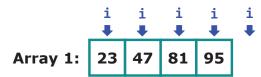
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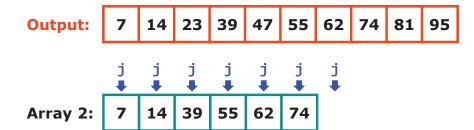
Merge Sort

- → Merge sort contains two parts:
 - 1) Merge operation
 - 2) Merge sort algorithm

Merge Operation

→ How to merge two already sorted vectors into a single sorted vector?





• C++ Implementation of the Merge Operation

```
/** Merges two sorted vectors into a single sorted vector.
1
2
        @param vec 1: first sorted vector to merge
3
        @param vec 2: second sorted vector to merge
4
        @param vec_3: merged vector
    */
5
6
   void merge(const vector<int>& vec 1, const vector<int>& vec 2,
7
               vector<int>& vec_3) {
8
        size_t i = 0, j = 0, k = 0;
9
        while (i < vec_1.size() && j < vec_2.size()) {</pre>
            if (vec_1.at(i) <= vec_2.at(j)) {</pre>
10
                vec_3.at(k++) = vec_1.at(i++);
11
            } else {
12
13
                vec 3.at(k++) = vec 2.at(j++);
14
            }
15
16
        while (i < vec 1.size()) { vec 3.at(k++) = vec 1.at(i++); }
17
        while (j < vec_2.size()) { vec_3.at(k++) = vec_2.at(j++); }</pre>
   }
18
```

 \rightarrow Time complexity of the merge operator: O(m+n)

m, n are the size of the two input vectors.

If m and n are (almost) equal, then the time complexity can be simply written as O(n).

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Merge Sort

- → Basic Idea (Algorithm):
 - 1) Divide the entire vector into two halves, the left half and right half.
 - 2) Sort the left half.
 - 3) Sort the right half.
 - 4) Merge the sorted left half and right half to form sorted whole vector.
- → [Important] How to "sort the left half"? How to "sort the right half"? Merge sort is a recursive algorithm.
- → Algorithm
 - 1) If the size of the vector is less than 2, then return (base case).
 - 2) Copy the left half of the vector into another vector (denoted as left_half).
 - 3) Copy the right half of the vector into another vector (denoted as right_half).
 - 4) Recursively sort left half.
 - 5) Recursively sort right_half.
 - 6) Merge left_half and right_half.

• C++ Implementation of Merge Sort

```
void merge sort(vector<int>& vec) {
1
2
        // Base case
3
        if (vec.size() < 2) { return; }</pre>
4
5
        // Copy the left and right half of the vector into 2 other vectors.
6
        vector<int> left, right;
7
        for (size_t i = 0; i < vec.size(); i++) {</pre>
8
            if (i < vec.size() / 2) { left.push_back(vec.at(i)); }</pre>
            else { right.push back(vec.at(i)); }
9
        }
10
11
        // Sort "left" and "right" recursively.
12
13
        merge sort(left);
        merge_sort(right);
14
15
        // Merge the sorted left and right half.
16
17
        merge(left, right, vec);
18
   }
```

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Performance of Merge Sort

 \rightarrow Merge: O(n)

ightharpoonup Merge sort: $T(n) = 2T\left(\frac{n}{2}\right) + O(n) = O(n\log n)$ (Master's theorem)

| Number of Comparisons | | | |
|-----------------------|---------------|---------------|--|
| Best Case | Average Case | | |
| $O(n \log n)$ | $O(n \log n)$ | $O(n \log n)$ | |

→ Merge sort is **not** an in-place sorting algorithm.

We need <u>extra memory space</u> to <u>store the left and right halves</u> of the input vector.

→ Merge sort is **not** a quadratic sorting algorithm.

Time complexity of merge sort is $O(n \log n)$, instead of $O(n^2)$.

→ Merge sort is a <u>stable</u> sorting algorithm.

- Best time complexity for comparison sorting algorithms
 - \rightarrow For comparison sorting algorithms, $O(n \log n)$ is already the <u>optimal time</u> <u>complexity</u>.

In other words, no comparison sorting algorithm can do better than $O(n \log n)$ (can be mathematically proved).

- → Merge sort is <u>very fast</u> since it <u>reaches the optimal time complexity</u>.
- Example of Non-Comparison Sorting Algorithm
 - → Radix Sort

| 3487 | 109 <mark>0</mark> | 11 <mark>2</mark> 8 | 1 <mark>0</mark> 90 | 1 090 |
|------|--------------------|---------------------|---------------------|--------------------|
| 2873 | 293 <mark>2</mark> | 29 <mark>3</mark> 2 | 1 <mark>1</mark> 28 | <mark>1</mark> 128 |
| 4378 | 287 <mark>3</mark> | 34 <mark>3</mark> 9 | 3 <mark>2</mark> 87 | <mark>2</mark> 873 |
| 1090 | 874 <mark>3</mark> | 87 <mark>4</mark> 3 | 4 <mark>3</mark> 78 | <mark>2</mark> 932 |
| 8743 | 348 <mark>7</mark> | 28 <mark>7</mark> 3 | 3 <mark>4</mark> 39 | <mark>3</mark> 287 |
| 1128 | 328 <mark>7</mark> | 43 <mark>7</mark> 8 | 3 <mark>4</mark> 87 | <mark>3</mark> 439 |
| 3439 | 437 <mark>8</mark> | 34 <mark>8</mark> 7 | 8 <mark>7</mark> 43 | <mark>3</mark> 487 |
| 3287 | 112 <mark>8</mark> | 32 <mark>8</mark> 7 | 2 <mark>8</mark> 73 | <mark>3</mark> 888 |
| 2932 | 388 <mark>8</mark> | 38 <mark>8</mark> 8 | 3 <mark>8</mark> 88 | <mark>4</mark> 378 |
| 3888 | 343 <mark>9</mark> | 10 <mark>9</mark> 0 | 2 <mark>9</mark> 32 | <mark>8</mark> 743 |

- → Other non-comparison sorting algorithms (not required):
 - Counting sort
 - Bucket sort