

# COM SCI 132 Week 2

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## The Java Compiler Compiler (JCC)

- Can be thought of as "Lex and Yacc for Java"
- It is based on LL(k) rather than LALR(1).
- Grammars are written in EBNF.
- The Java Compiler Compiler transforms an EBNF grammar into an LL(k) parser.
- The JavaCC grammar can have embedded action code written in Java, just like a Yacc grammar can have embedded action code written in C.
- The lookahead can be changed by writing LOOKAHEAD(...).
- The whole input is given in just one file (not two).

**Input format:** one file containing a header, token specifications for lexical analysis, and grammar.

**Example of a token specification:**

```
TOKEN:
{
    < INTEGER_LITERAL: ( ["1"-"9"] (["0"-"9"])* | "0" ) >
}
```

**Example of a production:**

```
void StatementListReturn():
{}
{
    ( Statement() )* "return" Expression() ";"
}
```

**Generating a parser with JavaCC**

```
javacc fortran.jj // generates a parser with a specified name
javac Main.java // Main.java contains a call of the parser
java Main < prog.f // parses the program prog.f
```

## The Visitor Pattern

According to Gamma, Helm, Johnson, Vlissides: **Design Patterns**, 1995:

For **object-oriented programming**, the Visitor pattern **enables** the definition of a **new operation** on an **object structure without changing the classes** of the objects.

## First Approach: Instanceof and Type Casts

The running Java example: summing an integer list.

```
interface List {}
```

```
class Nil implements List {}
```

```
class Cons implements List {  
    int head;  
    List tail;  
}
```

```
List l;
```

Suppose we have a `List l` in the code to parse. Is `l` a `Nil` or a `Cons`?

One way to check is to do type casting:

```
List l; // The List-object  
int sum = 0;  
boolean proceed = true;  
while(proceed) {  
    if (l instanceof Nil) proceed = false;  
    else if (l instanceof Cons) {  
        sum += ((Cons) l).head; // Notice these two type casts!  
        l = ((Cons) l).tail;  
    }  
}
```

**Advantage:** The code is written without touching the classes `Nil` and `Cons`.

**Drawback:** The code constantly uses type casts and `instanceof` to determine what class of object it is considering.

## Second Approach: Dedicated Methods

The first approach is **not** object-oriented! To access parts of an object, the classical approach is to use dedicated methods which both access and act on the subobjects.

Instead, of writing a `sum()` method that operates on list types, just simply add a `.sum()` method to the `List` interface. All classes that implement the interface can then override the method.

```
class Nil implements List {  
    public int sum() { return 0; }  
}
```

```
class Cons implements List {  
    int head;  
    List tail;  
    public int sum() {  
        return head + tail.sum();  
    }  
}
```

**Advantage:** The type casts and `instanceof` operations have disappeared, and the code can be written in a systematic way.

**Drawback:** For each new operation on `List`-objects, write new dedicated methods and recompile all classes.

## Third Approach: The Visitor Pattern

### The Idea:

- Divide the code into an object structure and a Visitor (akin to Functional Programming!)
- Insert an `accept` method in each class. Each `accept` method takes a Visitor as an argument.
- A Visitor contains a `visit` method for each class (overloading!) A method for a class *C* takes an argument of type *C*.

```
interface List {  
    void accept(Visitor v);  
}
```

```
interface Visitor {  
    void visit(Obj x);  
    void visit(Cons x);  
}
```

The purpose of the `accept` methods is to invoke the `visit` method in the Visitor which can handle the current object.

```
class Nil implements List {  
    public void accept(Visitor v) {  
        v.visit(this);  
    }  
}
```

```
class Cons implements List {  
    int head;  
    List tail;  
    public void accept(Visitor v) {  
        v.visit(this);  
    }  
}
```

**Notice:** the `accept` method in both classes is the same! Due to the way the Visitor is implemented, the method always calls `visit`. This is beneficial because it allows the compiler to easily parse classes by appending the same code to each class. In fact, the `accept` method is always implemented the way it is shown above, for any class!

The control flow goes back and forth between the `visit` methods in the Visitor and the `accept` methods in the object structure.

```
class SumVisitor implements Visitor {  
    int sum = 0;  
    public void visit(Obj x) {}  
    public void visit(Cons x) {  
        sum += x.head;  
        x.tail.accept(this);  
    }  
}  
.....
```

```
SumVisitor sv = new SumVisitor();  
l.accept(sv);  
System.out.println(sv.sum);
```

**Also notice:** since the `accept` method is titled `accept` in every class, the `visit` method in the `Visitor` can just call the `accept` method, and it automatically calls the method for the correct class!

## Visitors: Summary

- Visitor makes adding new operations easy. Simply write a new visitor.
- A visitor gathers related operations. It also separates unrelated ones.
- Adding new classes to the object structure is hard. Key consideration: are you most likely to change the algorithm applied over an object structure, or are you most likely to change the classes of objects that make up the structure.
- Visitors can accumulate state.
- Visitor can break encapsulation. Visitor's approach assumes that the interface of the data structure classes is powerful enough to let visitors do their job. As a result, the pattern often forces you to provide public operations that access internal state, which may compromise its encapsulation.

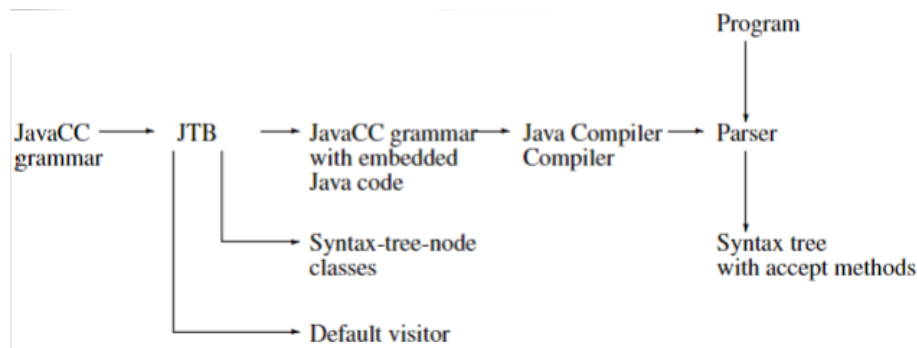
## The Java Tree Builder

The Java Tree Builder (JTB) is a front for The Java Compiler Compiler. It supports the building of syntax trees which can be traversed using visitors.

JTB transforms a bare JavaCC grammar into three components:

1. a JavaCC grammar with embedded Java code for building a syntax tree;
2. one class for every form of syntax tree node; and
3. a default visitor which can do a depth-first traversal of a syntax tree.

The produced JavaCC grammar can then be processed by the JCC to give a parser which produces syntax trees. The produced syntax trees can now be traversed by a Java program by writing subclasses of the default visitor.



### Example (simplified)

JTB produces a syntax-tree-node class for `Assignment`:

```
public class Assignment implements Node {
    PrimaryExpression f0; AssignmentOperator f1;
    Expression f2;

    public Assignment (PrimaryExpression n0, AssignmentOperator n1, Expression n2) {
        f0 = n0;
    }
}
```

```

        f1 = n1;
        f2 = n2;
    }

    public void accept(visitor.Visitor v) {
        v.visit(this);
    }
}

```

Notice the `accept` method; it invokes the method `visit` for `Assignment` in the default visitor.

The default visitor looks like this:

```

public class DepthFirstVisitor implements Visitor {
    ...
    //
    // f0 -> PrimaryExpression()
    // f1 -> AssignmentOperator()
    // f2 -> Expression()
    //
    public void visit(Assignment n) {
        n.f0.accept(this);
        n.f1.accept(this);
        n.f2.accept(this);
    }
}

```

Notice the body of the method which visits each of the three subtrees of the `Assignment` node.