Progress

Course outline

course work?

Week 0

Week 1

Week 2

Week 3

Week 4

How does an NPTEL online

Lecture 16 : Variants of EDF

Lecture 17 : Rate Monotonic

 Lecture 18 : Rate Monotonic Schedulability Analysis

 Lecture 19 : Rate Monotonic Scheduling: Miscellaneous

Quiz: Week 4 : Assignment

Feedback Form of Week 4

Schedulability Analysis

and Rate Monotonic

Scheduling

issues

Week 5

Week 6

Week 7

Week 8

Week 9

Week 10

Week 11

Week 12

Assignments Solution

Live Interactive Session

Download Videos

Lecture 20 : RMS Generalizations

Lecture Materials

NPTEL » Real-Time Systems

Week 4 : Assignment 4 The due date for submitting this assignment has passed.

As per our records you have not submitted this assignment.

Task 3

O c.

a.

O b.

O c.

O d.

O e.

Score: 0

O c.

O d.

○ e.

Score: 0

○ c.

d.

O e.

Score: 0

a.

b.

O c.

O d.

О e.

___ a.

_ b.

_ c.

d.

___ e.

Score: 0

a.

b.

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O d.

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Score: 0

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c.

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Score: 0

___ a.

■ b.

_ c.

d.

_ e.

a.

b.

○ c.

d.

О е.

f.

Score: 0

uniprocessor are false?

set.

Accepted Answers:

T1

T2

T3

c. e.

Accepted Answers:

Task

Task1

Task 2

Task 3

No, the answer is incorrect.

Accepted Answers:

the task T1?

Task Name

on a uniprocessor are true?

getting late.

Consider only Liu-Layland criterion for schedulability.

Guaranteed to be schedulable

d. Guaranteed to be unschedulable

e. Liu-Layland criterion cannot be used

20

25

40

T1

T2

T3

Phase

100

200

300

b.

d.

Accepted Answers:

Task 3

criterion for determining their schedulability. Task Set A (All times are in milliseconds)

Task Name Execution time Period Deadline Task1 20 100 100 15 150 150 Task 2

Task Name	Execution time	Period	Deadline
Task1	40	100	100
Task 2	20	80	80
Task 3	50	200	200

- e. Liu-Layland criterion cannot be applied

- criterion for determining their schedulability.

40

Task Name Execution time Period Deadline 100 Task1 40 100 Task 2 80 20 80 50 200 200 Task 3 None of the task sets is schedulable

- Accepted Answers:
 - uniprocessor. Which one of the following characteristics can guarantee schedulability of a task set even when its processor utilization approaches 100%? a. Task periods are relatively prime to each other
- No, the answer is incorrect. Accepted Answers:
 - monotonic scheduler. For the two tasks to run without missing any deadlines, what should be the maximum processor utilization due to the two tasks?

Consider the periodic real-time task shown in the following table. The task set is to be scheduled on

a uniprocessor using a rate monotonic scheduler. What would be the worst case completion time for

Execution time

Task Set (All times are in milliseconds)

30

50

20

 a. 30 milliseconds b. 40 milliseconds c. 50 milliseconds 70 milliseconds

Period

100

150

80

Deadline

100

150

80

- No, the answer is incorrect. Accepted Answers:
- uniprocessor, RMS can be as proficient as EDF (earliest deadline first) for task sets whose periods satisfy some constraints. d. For scheduling real-time tasks in practical uniprocessor based real-time systems, sub-optimal heuristic scheduling algorithms are normally used as optimal scheduling algorithms are computationally intractable.

Consider a set of periodic hard real-time tasks whose characteristics are shown in the following

Task Set A (All times are in milliseconds)

Period

100

150

400

Deadline

Deadline

100

150

80

Deadline

100

200

500

100

150

400

table. Would this task set be schedulable on a uniprocessor under rate monotonic scheduling?

Execution time

a. Not likely to be schedulable but no guarantee can be given

b. Likely to be schedulable but no guarantee can be given

- No, the answer is incorrect.

 - monotonic scheduler. The task characteristics are given in the following table. What would be the worst case completion time for the task T2? Task Set (All times are in milliseconds) Task Name Period Execution time Phase

200

150

100

 a. 30 milliseconds b. 55 milliseconds

c. 85 milliseconds

d. 110 milliseconds

e. 135 milliseconds

- No, the answer is incorrect. Score: 0 Accepted Answers:
 - By Liu-Layland's criterion it will not meet any of its deadlines c. By Liu-Lehoczky's criterion it will meet all its deadline d. By Liu-Lehoczky's criterion it will miss all its deadline e. By Liu-Lehoczky's criterion it will occasionally miss some of its deadline

Liu-Lehoczky criterion cannot be used, consequently no results can be given

Suppose a single periodic hard real-time task requires an execution time of 100 milliseconds every

100 milliseconds. This task is to be run on a uniprocessor using a rate monotonic scheduler. Which

Consider a set of three periodic real-time tasks whose characteristics have been shown in the

following table. Would this task set be schedulable on a uniprocessor under rate monotonic

500 Task 3 The task set is not schedulable The task set is schedulable Prima facie the task set is schedulable, but further analysis using the Liu-Lehoczky's

criterion needs to be made before any guarantees can be given

non-zero phasings

e. Liu-Layland criterion cannot be used

EDF as compared while using RMS

No, the answer is incorrect. Accepted Answers:

d. Schedulability cannot be determined by using Liu-Layland criterion since the tasks have

 EDF is the optimal task scheduler for running a set of periodic hard real-time tasks on a uniprocessor. EDF has better transient overload handling capability as compared to RMS. RMS is more proficient as compared to the time-sliced round-robin algorithm for

uniprocessor, RMS can be as proficient as EDF under some constraints on the task

- No, the answer is incorrect. Score: 0 Accepted Answers: a.
- Also assume that Priority(Ti) gives the priority value assigned to the task Ti. Task Period (in millisec) Deadline (in millisec) Execution Time (in millisec) T1 20 150 150 T2 40 200 200 100 100
 - d. Priority(T1)=2, Priority(T2)=1, Priority(T3)=3 e. Priority(T1)=3, Priority(T2)=1, Priority(T3)=2 f. Priority(T1)=1, Priority(T2)=3, Priority(T3)=2

No, the answer is incorrect.

Accepted Answers:

Due on 2021-09-01, 23:59 IST.

1 point

Mentor

Consider two periodic task sets A and B shown in the following two tables. Would these two task

1 point Consider two periodic task sets A and B shown in the following two tables. Would these two task sets be schedulable on uniprocessors under rate monotonic scheduling? Use the Liu-Lehoczky's Task Set A (All times are in milliseconds) Task Name **Execution time** Period Deadline Task1 20 100 100 15 150 150 Task 2

400

400

Task Set B (All times are in milliseconds)

 Task set A is schedulable but Task set B is not Task set B is schedulable but Task set A is not d. Both task sets are schedulable.

No, the answer is incorrect. Suppose a rate monotonic scheduler is used to schedule a set of periodic hard real-time tasks on a

b. Task periods are harmonically related c. Tasks have zero phasings d. Tasks have non-zero phasings e. Task execution times are harmonically related a. b.

a. 0.887 b. 0.824 c. 0.765 d. 0.736 e. 0.698 a. ○ b.

Consider two arbitrary periodic hard real-time tasks are to be run on a uniprocessor using a rate

e. 90 milliseconds

 Under RMS (rate monotonic scheduler), the achievable utilization of a set of hard real-time periodic tasks would be lower when task periods are multiples of each other compared to the case when they are not. Even if a task set passes the Liu-Layland's criterion, the schedulability of the task set under RMS cannot be guaranteed unless it passes the Liu-Lehoczky's criterion as well. While scheduling a set of independent hard real-time periodic tasks on a

When a set of tasks are scheduling using a rate monotonic scheduler, a higher

Which of the following statements concerning scheduling of a set of periodic hard real-time tasks

priority task is guaranteed to not miss its deadline on account of a lower priority task

No, the answer is incorrect.

100

150

80

Consider a set of three periodic hard real-time tasks is to be run on a uniprocessor using a rate

25

25

30

- e.
- f. No, the answer is incorrect. Score: 0 Accepted Answers:

scheduling? Use the Liu-Layland's criterion for determining schedulability.

of the following statements are true regarding the schedulability of this task?

a. By Liu-Layland's criterion it will meet all its deadlines

- Task Set A (All times are in milliseconds) Task Execution time Period Phase 50 100 Task1 40 20 200 Task 2 100 150 50

Which of the following statements concerning scheduling a set of periodic real-time tasks on a

a. While scheduling a set of independent hard real-time periodic tasks on a

- scheduling a set of soft real-time tasks on a uniprocessor. e. Resource sharing among a set of real-time tasks is easier to support while using the
- b. d.

Suppose the set of three tasks shown in the following table are to be run on a uniprocessor by

the tasks? Assume that the higher the priority value assigned to a task, the lower is its priority.

using a rate monotonic scheduler. Which one of the following is a correct priority assignment to

- a. Priority(T1)=1, Priority(T2)=2, Priority(T3)=3 b. Priority(T1)=3, Priority(T2)=2, Priority(T3)=1 Priority(T1)=2, Priority(T2)=3, Priority(T3)=1

○ d. O e. No, the answer is incorrect.

Task set A is schedulable but Task set B is not Task set B is schedulable but Task set A is not d. Both task sets are schedulable (a. (b.

40 400 400 Task Set B (All times are in milliseconds)

About the Course

1 point sets be schedulable on uniprocessors under rate monotonic scheduling? Use the Liu-Layland

Ask a Question