

# UNIT-III: Declarative Programming

## Paradigm: Functional Programming



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# OUTLINE OF UNIT-3

## Sub-Unit

## Contents

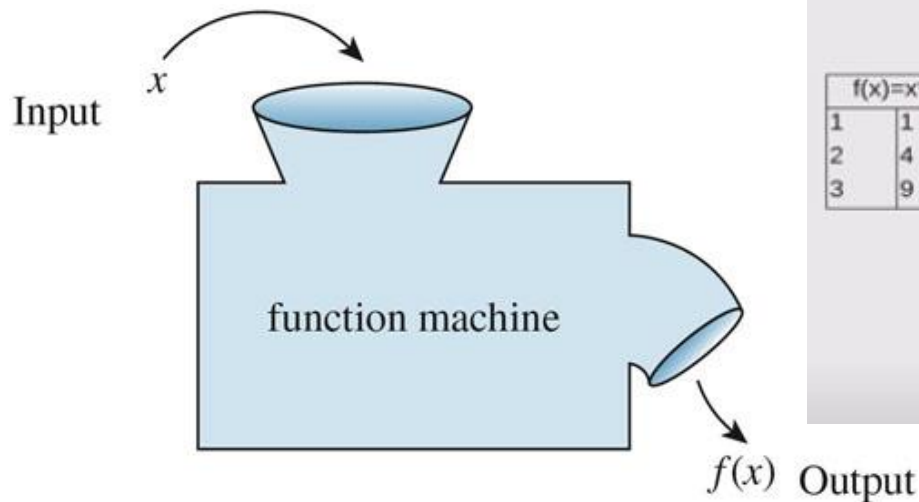
- 3.0 Introduction to functional programming
  - 3.1 Introduction to Lambda Calculus
  - 3.2 Functional Programming Concepts
  - 3.3 Evaluation order
  - 3.4 Higher order functions
  - 3.5 I/O- Streams and Monads
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## 3.0: Introduction to Functional Programming



# INTRODUCTION TO FUNCTIONAL PROGRAMMING

1. Functional programming starts with a point of view that a program is a function
2. Function can be thought of as a black box that takes input and produces output
3. A program is a function that transforms input to output



$f(x) = x \times x$	
1	1
2	4
3	9

$$f(x) = x^2$$
$$f(x) = \{x \times x | x \in \text{Naturals}\}$$
$$f(\bar{x}) = \bar{x}^T \cdot \bar{x}$$

- Maps Inputs To Outputs
- Computationally agnostic

4. So when we are writing a program in a functional programming language, we are
  - ✓ Specifying the rules on how to generate the output from a given input
  - ✓ In computation we apply the rules to generate the expected output



**Definition : Functional programming** is a programming paradigm — a style of building the structure and elements of computer programs — that treats computation as the evaluation of mathematical functions and avoids changing-state and mutable data

## Advantages of functional programming

- Functional programming is based on mathematical functions
- Easier to determine inputs
- Easier to determine outputs
- Easier to demonstrate prove that you have a correct program
- Easier to test programs that are too difficult to prove
- Examples of functional programming paradigm
- Haskell, Lisp, Python, Erlang, Racket, F#,



# Functional Programming

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- Functional programming is the process of building software by **composing pure functions**, avoiding shared state, mutable data, and side-effects.
  - Functional programming is **declarative** (telling the computer what you want to do) rather than **imperative** (telling the computer exactly how to do that), and application state flows through pure functions.
  - Functional programming is **based on mathematical functions**.
  - Functions are first class and can be higher order
  - It allows us to handle functions as if they were normal data types
  - Functions can either accept another function as argument or can return a function themselves
  - Some of the popular functional programming languages include: Lisp, Python, Erlang, Haskell, Clojure, etc.
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# BUILDING UP PROGRAMs

How do we build up these programs.....

Assume that

- We have some built in functions and values
- Use these to build more complex functions
- Example
- We have

Whole Numbers

Set of whole numbers:

$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \dots\}$

Successor function, *succ*

$\text{succ } 0 = 1$

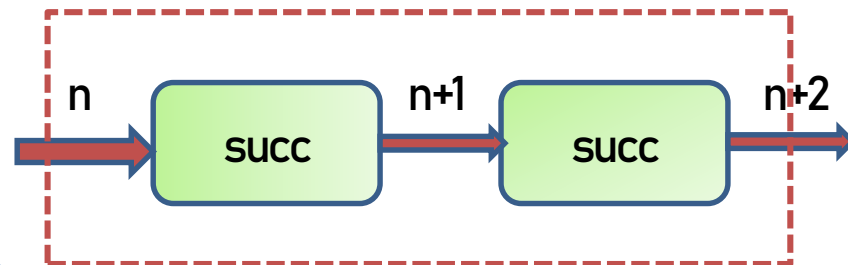
$\text{succ } 1 = 2$

$\text{succ } 2 = 3$



We can **compose** *succ* twice to build a new function

$\text{plusTwo } n = \text{succ}(\text{succ } n)$





## Whole Numbers

Set of whole numbers:

$\{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \dots \}$

Successor function, *succ*

$\text{succ } 0=1$

$\text{succ } 1=2$

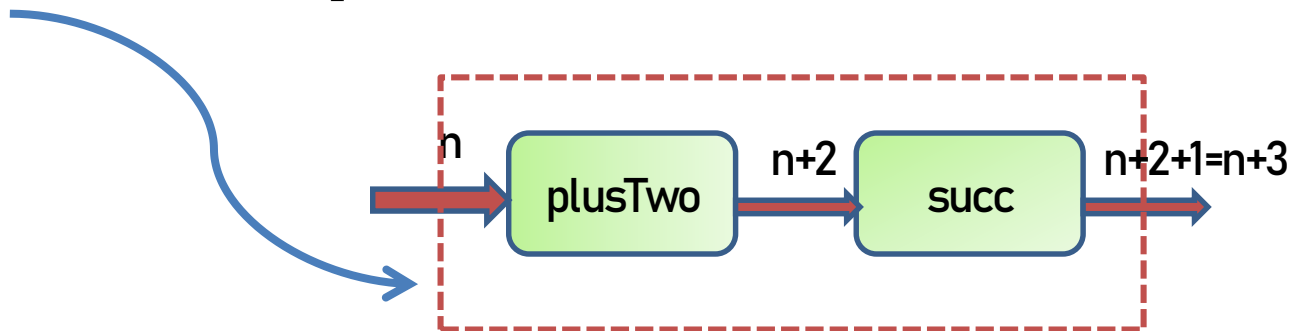
$\text{succ } 2=3$

function, plusTwo n

$\text{succ}(\text{succ } n)$

We can **compose** **succ** twice to build a new function

$\text{plusThree } n = \text{succ}(\text{plusTwo } n)$



We can combine functions to form new compositions/function





# CORE PROGRAMMING CHARACTERISTICS/CONCEPTS

1. PURE Functions: Functions are pure in Functional Programming
  - A function called multiple times with the same arguments will always return the same value. Always.
2. Functions are first class and can be higher order
  - It allows you to handle functions as if they were normal data types
  - Functions can either accept another function as argument or can return a function themselves
3. Variables are Immutable
  - You can't modify a variable after it has been initialized
  - You can create new variables, but you can't modify the existing variables
4. Functional programming is based on Lambda Calculus

## Definition:

**Lambda calculus** (also written as  **$\lambda$ -calculus**) is a formal system in mathematical logic for expressing computation based on function abstraction and application using variable binding and substitution.



# INTRODUCTION TO LAMBDA CALCULUS

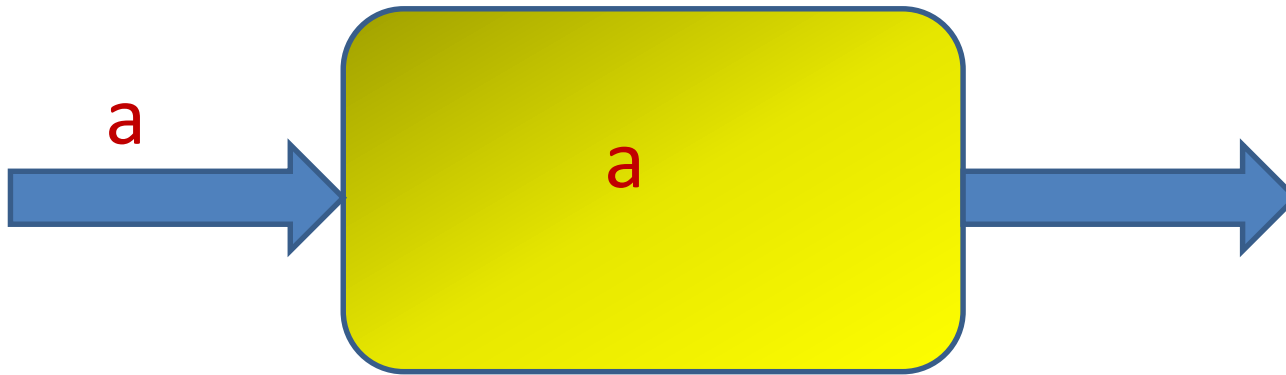
$\lambda a. a$   
IDENTITY

FUNCTION  
SIGNIFIER  $\rightarrow$   $\lambda a. a$

PARAMETER VARIABLE  
 $\downarrow$   
FUNCTION  
SIGNIFIER  $\rightarrow$   $\lambda a. a$

PARAMETER VARIABLE  
 $\downarrow$   
FUNCTION  
SIGNIFIER  $\rightarrow$   $\lambda a. a$   $\leftarrow$  RETURN  
EXPRESSION

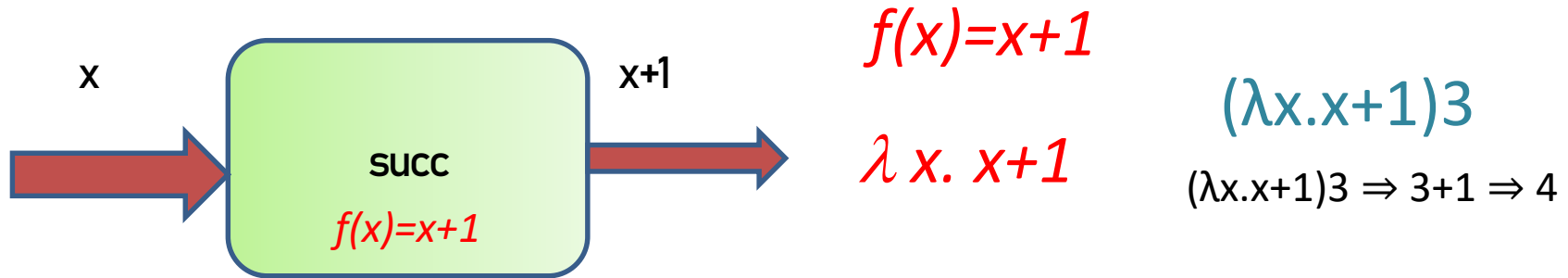
PARAMETER VARIABLE  
 $\downarrow$   
FUNCTION  
SIGNIFIER  $\rightarrow$   $\lambda a. a$   $\leftarrow$  RETURN  
EXPRESSION  
 $\underbrace{\hspace{10em}}$   
LAMBDA ABSTRACTION



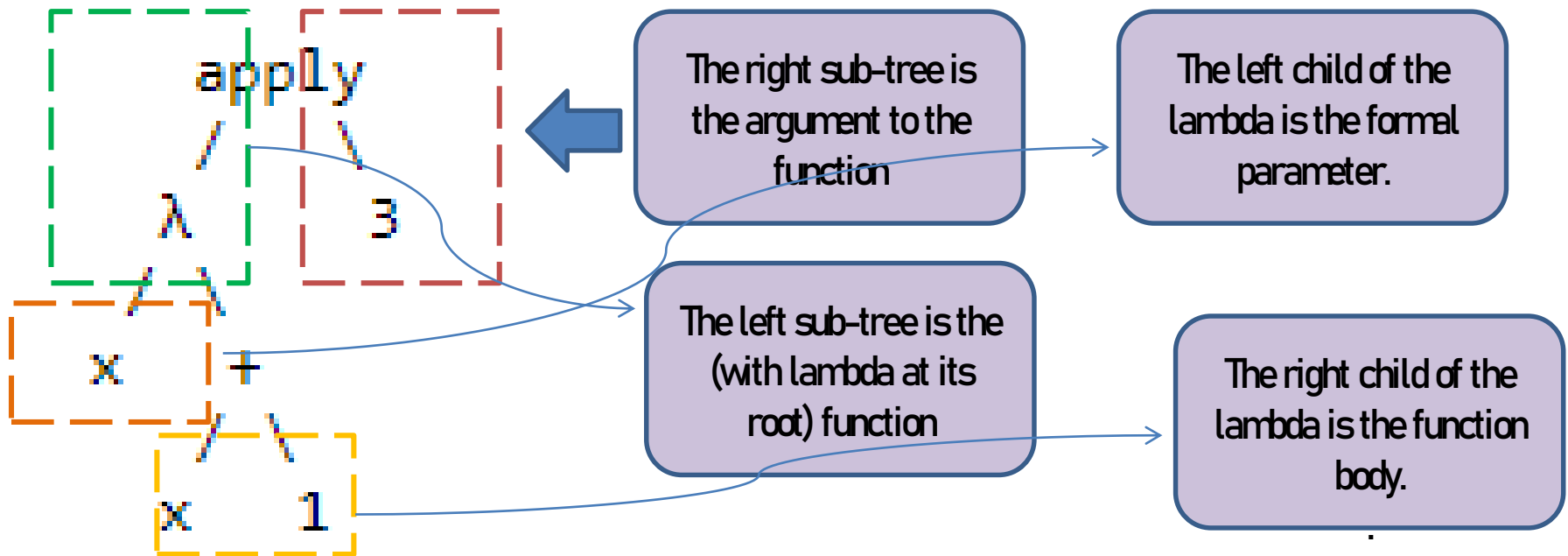
$\lambda a.a$



# FUNCTIONS USING LAMBDA CALCULUS-SUCC FUNCTION

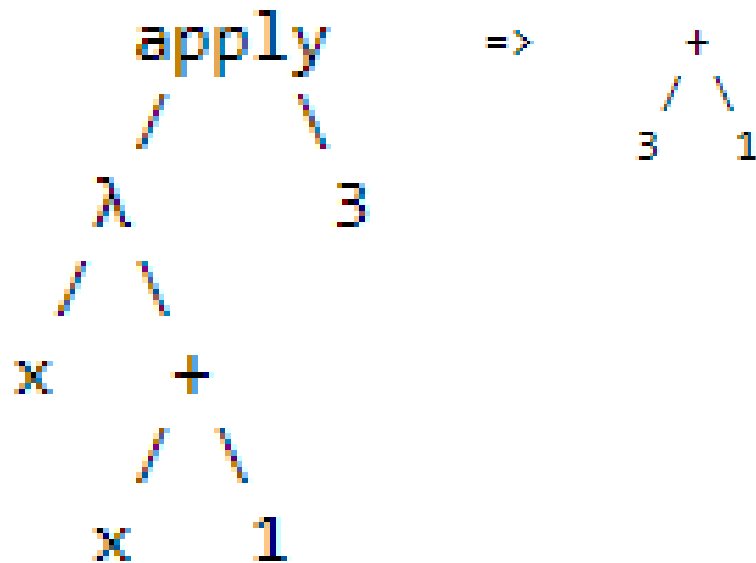


abstract-syntax tree (where  $\lambda$  is the abstraction operator, and  $\text{apply}$  is the application operator)



There is only one apply node in our example;

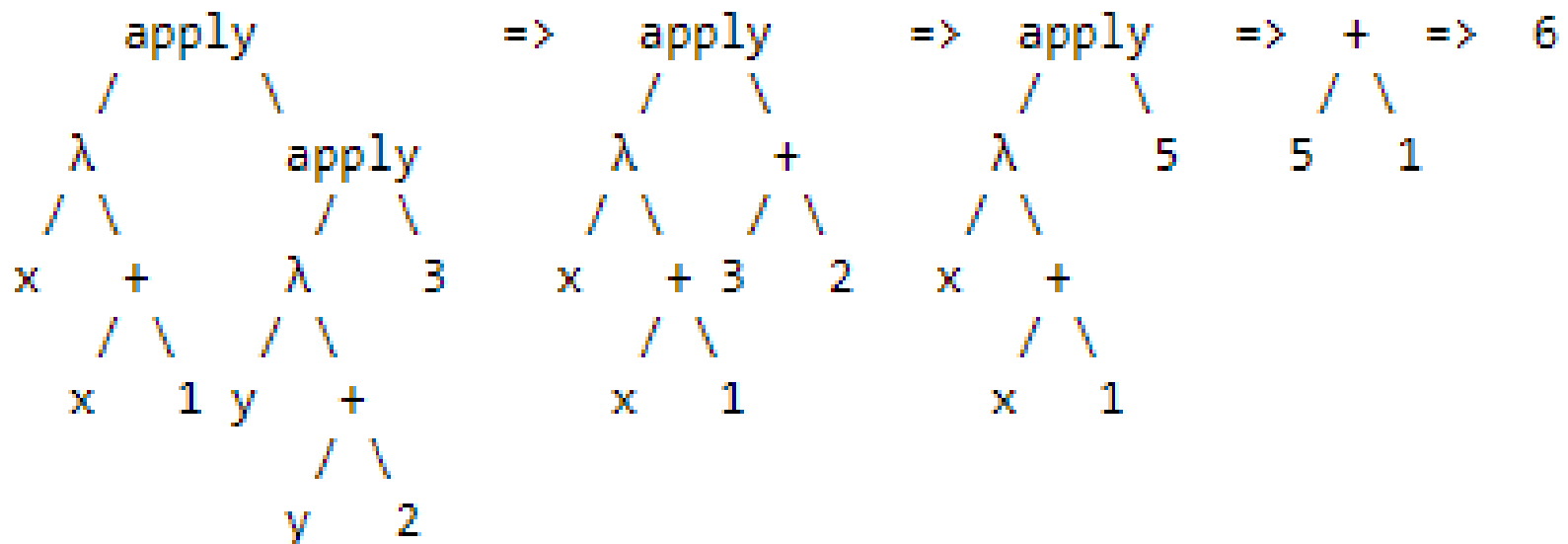
- The argument is 3
- The function is  $\lambda x.x+1$ ;
- The formal parameter is x
- The function body is  $x+1$ . Here's the rewriting step:



## FUNCTIONS USING LAMBDA CALCULUS-PLUSONE, PLUSTWO

$(\lambda x.x+1)((\lambda y.y+2)3)$

- ✓ The first lambda expression defines the "plus-one" function
- ✓ The argument to that function is itself an application
- ✓ which applies the "plus-two" function to the value 3



## FREE and BOUND VARIABLES

In an expression, each appearance of a variable is either "free" (to  $\lambda$ ) or "bound" (to a  $\lambda$ ).

✓ **Bound Variable:** a variable that is associated with some lambda.

✓ **Free Variable:** a var that is *not* associated with any lambda.

$(\lambda x. xy)$



✓ The variable  $x$  is bound and  $y$  is free.

✓ The  $x$  in the body of the first expression from the left is bound to the first  $\lambda$ .

$(\lambda x. x)(\lambda y. yx)$



✓ The  $y$  in the body of the second expression is bound to the second  $\lambda$  and the  $x$  is free

It is very important to notice that the  $x$  in the second expression is totally independent of the  $x$  in the first expression.





## Function Application

- **Function application** –
- A function application, often called a lambda application, consists of an expression followed by an expression:  $\text{expr expr}$ .
  - The notation  $E1.E2$  to denote the **application of function  $E1$  to actual argument  $E2$** .
- The **first** expression is a **function abstraction** and the **second** expression is the **argument** to which the function is applied.
- **Expressions** can be thought of as **programs** in the language of lambda calculus.
- **All functions** in lambda calculus **have exactly one argument**.
- **Multiple-argument functions** are represented by **currying**
  - For example, the lambda expression  $\lambda x. (+ x 1) 2$  is an application of the function  $\lambda x. (+ x 1)$  to the argument 2.
  - This function application  $\lambda x. (+ x 1) 2$  can be evaluated by substituting the argument 2 for the formal parameter  $x$  in the body  $(+ x 1)$ .
  - Doing this we get  $(+ 2 1)$ . **This substitution is called a beta reduction.**
  - Beta reductions are like **macro substitutions** in C. To do beta reductions correctly we may need to rename bound variables in lambda expressions to avoid name clashes.

## Function Application

- Function application associates left-to-right; thus,  $f\ g\ h = (f\ g)h$ .
- Function application binds more tightly than  $\lambda$ ; thus,  $\lambda x. f\ g\ x = (\lambda x. (f\ g))x$ .
  - Multiple expressions:  $E_1E_2E_3 \dots E_n$       $(\dots((E_1E_2)E_3)\dots E_n)$
- Functions in the lambda calculus are first-class citizens; that is to say, **functions can be used as arguments to functions** and **functions can return functions as results**.

- Evaluating Lambda Calculus:
- Ex1:  $(+ (* 5 6) (* 8 3))$
- Here, we can't start with '+' because it only operates on numbers. There are two reducible expressions:  $(* 5 6)$  and  $(* 8 3)$ .

- We can reduce either one first. For example –

$(+ (* 5 6) (* 8 3))$

$(+ 30 (* 8 3))$

$(+ 30 24)$

$= 54$

# DATA TYPES

- Bool
- Char
- Int
- Float
- Double
- List
- Tuple
- Function

- In Haskell all computations are done via the evaluation of expressions
- Examples of expressions include **atomic** values (built-in) such as
  - the integer 5,
  - the character 'a', and
  - the function  $\backslash x \rightarrow x+1$ , as well as structured values such as
  - the list  $[1,2,3]$  and
  - the pair  $('b',4)$ .





## Types (set of Values)

- Bool
- Char
- Int (64 bit)
- Integer (Superset of Int)
- Float
- Double
- List
- Tuple
- Function

## Examples

- :type True
- :type "hi"
- :type 5
- :type 5.34
- :type (True, False)

## Type Class

- **EQ**
  - Type class is an interface which provides the functionality to test the **equality** of an expression.
- **Num and Fractional**
  - This type class is used for numeric operations. Types such as **Int, Integer, Float, and Double** come under this Type class.
- **Integral**
  - sub-class of the Num Type Class.
  - Int and Integer are the types under this Type class.
- **Floating**
  - sub-class of the Num Type Class.
  - Float and Double come under this type class.

# ARITHMETIC AND LOGICAL OPERATORS

- $2 + 3$
- $2 - 3$
- $2 * 3$
- $2 * (-3)$
- $2 / 3$
- `it (result)`
- $50 * 100 - 4999$
- $50 * (100 - 4999)$
- $40 * 100 - 3000 + 50 / 5$
- $(40 * 100 - 3000 + 50) / 5$
- $2 + \text{"hi"}$

- `False`
- `True`
- `True && False`
- `True && True`
- `False || True`
- `not False`
- `not (True && False)`
- `not (True || False)`

```
Prelude> :t "a"
```

```
Prelude> '\97\  
'a'
```

```
Prelude> '\67' 'C'
```

```
Prelude> :t "mrinmoyee.in"
```

```
Prelude> [1,2,3,4,5]
```

```
Prelude> (1,1,'a')
```



# COMPARATIVE OPERATORS

- $2 == 3$
- $2 == 0$
- $2 /= 2$
- $2 /= 0$
- $2 < 3$
- $2 > 3$
- $2 \wedge 3$
- `not (2 < 3)`
- `"hi" == "hi"`
- `"hi" == "Hi"`





## INBUILT FUNCTIONs

1. succ 6
2. succ (succ 5)
3. min 5 6
4. max 5 6
5. max 101 101
6. succ 9 + max 5 4 + 1
7. (max 5 4)+(succ 9)+1
8. (succ 9) + (max 5 4) + 1
9. We wanted to get the successor of the product of numbers 9 and 10. we couldn't write **succ 9 \* 10** because that would get the successor of 9, which would then be multiplied by 10
10. succ 9\*10
11. succ (9\*10)
- 12. div 92 10**
13. div 3 4
14. div 4 3
15. 4/3
16. mod 7 5
17. mod 3 1
18. mod 7 2
19. reverse "hello"
20. x=45
21. print x
22. return True
23. return False
24. x <- return 35
25. print x
26. putStrLn "hello"

