# Leinster - Basic Category Theory - Selected problem solutions for Chapter 6

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May 18, 2025

6.2.20

(a)

$$\begin{array}{ccc}
X & \xrightarrow{1} & X \\
\downarrow_{1} & & \downarrow_{\alpha} \\
X & \xrightarrow{\alpha} & Y
\end{array} \tag{1}$$

where X, Y are functors in the category [A, S]

By Lemma 5.1.32  $\alpha$  is monic in [A, S] if and only if the above square is a pullback. Using Theorem 6.2.5 there is a pullback whose image under the evaluation functor  $ev_A$  is a pullback for each  $A \in \mathbf{A}$  in S. So by the lemma  $\alpha_A$  is monic for all  $A \in \mathcal{A}$ . The other direction holds by virtue of the same theorem, that there is only one way to extend the pullback on S for each S, to a pullback on S a pullback on S for each S and S is monic for all S in S is monic for all S is monic for all S is monic for all S in S in S is monic for all S is monic for all S in S in S in S in S is monic for all S in S in

(b) Monics in  $[A^{op}, Set]$  are epics in [A, Set] and vice versa.

#### 6.2.21

(a) Use a cardinality argument as follows. There is only one identity map represented by the left hand side of the following expresssion.  $H_A(A) \cong X(A) + Y(A)$ , for all A in A. Which means that either X(A) or Y(A) must be the empty set for all A in A.

#### 6.2.22

The category of elements can be represented by  $(1 \to X)$ , where 1 is a single element set. The comma category commuting diagram becomes

$$\begin{array}{ccc}
1 & \longrightarrow & x \\
\downarrow & \downarrow & \chi_f \\
x'
\end{array} \tag{2}$$

where  $x \in X(A)$ , and  $x' \in X(A')$ , and  $f : A' \to A$ . The above diagram shows under our choice of comma category that Xf(x) = x' as required.

#### 6.2.23

A category of elements with a terminal object by definition is equivalent to the definition of a representation as a universal element in (4.6).

#### 6.2.24

Let E be a functor in the functor category  $[\mathbf{A}^{op}, \mathbf{Set}]$  and  $E \to X$  be an object of the slice category, where X is a presheaf on  $\mathbf{A}$ . We need an equivalence functor to map  $E \to X$  to some  $[\mathbf{B}^{op}, \mathbf{Set}]$ . For a given A, and consider  $\alpha_A : E(A) \to X(A)$ . For a  $x \in X(A)$  back out the definition of E with

$$\beta_A(x) = \{e : \alpha_A(e) = x\} \tag{3}$$

where  $e \in E(A), x \in X(A)$ .

So now we have (A, x) pairs as in the definition of the category of elements in Definition 6.2.16, and can construct a functor using them informally as  $(A, x) \to \beta_A(x)$ . However we do need to show that  $\beta_A(x)$  and  $\beta_{A'}(x)$  induce an f such that (Xf)(x') = x. Because the morphism of E to X is a natural transformation we know that with  $f: A \to A'$ , that  $\alpha_A(e) = (Xf)(\alpha_{A'}e')$  taken with (3) means (Xf)(x') = x as required.

In the other direction, if E is a presheaf on the category of elements,

$$E(a) = \bigsqcup_{x \in X(a)} E(a, x) \tag{4}$$

and for  $e \in E(a), x \in X(a)$  define f(e) = x.

Source of ideas for this proof.<sup>1</sup>

#### 6.2.25

### (a) i. Functoriality of $Lan_F X$

 $<sup>^{1}</sup> https://math.stackexchange.com/questions/3633646/every-slice-of-a-presheaf-category-is-again-a-presheaf-category$ 

Let the diagram given for our colimit be  $D_B := X$ , with  $(A, FA \to B)$  in  $(F \Rightarrow B)$ . To prove  $\operatorname{Lan}_F X$  is a functor we need to consider  $\operatorname{Lan}_F X$  on morphisms  $f : B \to B'$ , and  $f' : B' \to B''$ . f and f' induce the maps presented below:

$$D_{B} \xrightarrow{p_{I}} Lan_{F}(B)$$

$$\downarrow^{D(f)} \qquad \downarrow^{L(f)}$$

$$D_{B'} \xrightarrow{p_{I'}} Lan_{F}(B')$$

$$\downarrow^{D(f')} \qquad \downarrow^{L(f')}$$

$$D_{B''} \xrightarrow{p_{I''}} Lan_{F}(B'')$$

$$(5)$$

The map from  $\operatorname{Lan}_F B \to \operatorname{Lan}_F B''$  is a unique map by the colimit property of  $\operatorname{Lan}_F B$  and hence L(f')L(f) = L(f'f) as required.

## (a) ii. Bijection between $\operatorname{Lan}_F X \to Y$ and $X \to YF$ .

Consider the cocone  $(X(A) \xrightarrow{p_I} \operatorname{Lan}_F X(B))_{\{FA \to B\}}$ , for all  $(A, FA \to B)$  in  $(F \Rightarrow B)$ . To form the bijection required, make the canonical choice of  $1_{F(A)}$  in  $(F \Rightarrow F(A))$  and reevaluate  $p_I$  above which now becomes  $X(A) \to \operatorname{Lan}_F XF(A)$ . With this choice there is a single base of the cocone X(A) for every  $\operatorname{Lan}_F XF(A)$  so we can form the required bijection between  $\operatorname{Lan}_F X \to Y$  and  $X \to YF$ . The task remaining is to prove naturality between  $\operatorname{Lan}_F X$  and X.