$\land$	/\ or \land	and, conjunction
$\lor$	\/ or \lor	or, disjunction
$\neg$	or \lnot or \neg	not
$\in$	\in	in
$\notin$	\notin	not in
$\langle x, y \rangle$	<< x, y>>	a tuple containing some x, y
$<$	<	less than
$\leq$	\leq or =<	less than or equal
$\ll$	\ll	much less?
$\equiv$	<=> or \equiv	is equivalent to
$>$	>	greater
$\geq$	\geq or >=	greater or equal
$\gg$	\gg	much greater?
$\prec$	\prec	precedes
$\preceq$	\preceq	precedes or equals
$\succ$	\succ	succeeds
$\succeq$	\succeq	succeeds or equals
$\subset$	\subset	subset
$\subseteq$	\subseteq	subset or equal
$\sqsubset$	\sqsubset	bag subset/is a refinement?
$\sqsubseteq$	\sqsubseteq	bag subset or equal/is a refinement or equal?
$A \vdash B$	|-	B can be derived from A?
$[S \to T]$	[S -> T]	set of functions
$\to$	->	step
$\cap$	\cap or \intersect	intersection
$\sqcap$	\sqcap	
$\oplus$	(+) or \oplus	bag union
$\ominus$	(-) or \ominus	bag difference
$\odot$	(.) or \odot	
$\otimes$	(\X) \otimes	Cartesian product
$\oslash$	(/) or \oslash	
$\exists$	\E	existential quantification (there exists)
$\exists!$	\exists!	there exists exactly one
$\mathbf{\exists}$	\EE	temporal existential quantification, 'hiding'
$f[e]$	f[e]	function application
$[A]_v$	[A]_v	action operator, 'square A sub v', A happens or v is unchanged, [A \/ v' = v], allows stuttering
WF_v	WF_v	weak fairness variables
SF_v	SF_v	strong fairness variables
$\supseteq$	\supseteq	superset
$\supset$	\supset	superset
$\sqsupset$	\sqsupset	bag superset
$\sqsupseteq$	\sqsupseteq	bag superset or equal
$\dashv$	-|	
$\models$	|=	models/satisfies a temporal formula
$\leftarrow$	<-	substitution
$\cup$	\cup or \union	union
$\sqcup$	\sqcup	
$\uplus$	\uplus	
$\times$	\X or \times	multiply
$\wr$	\wr	
$\propto$	\propto	propositional something?

$\forall$	\A	universal quantification (for all)
$\mathbf{\forall}$	\AA	temporal universal quantification
$\langle A \rangle_v$	<<A>>_v	action operator, 'angle A sub v', A happens and v changes, [A /\ v' \# v]
$\Rightarrow$	=>	implies
$\triangleq$	==	is equivalent
$\neq$	\neq or #	not equal
$\Box$	[]	always in the future/henceforth
$\Diamond$	<>	sometime(s) in the future/eventually
$n..m$	n .. m	integer interval, n to m inclusive (Naturals module)
x	x	operator of arity 2
$SUBSETS$	SUBSET S	set of all subsets of S
$CHOOSEx \in S : p$	CHOOSE x \in S : p	Choose x such that x is in S, and p is TRUE
$\leadsto$	>	leads to
$E \stackrel{+}{\triangleright} M$	-+->	E guarantees M: M remains true at least one step longer than E does
$[h_1 > e_1, ..., h_n \mapsto e_n]$	[h_1 |-> e_1, ..., h_n |-> e_n]	function/record constructor
$[x \in S \mid -> e]$	[x \in S |-> e]	function constructor
$[h_1 : S_1, ..., h_n : S_n]$	[h1: s1, ..., hn: sn]	set of records
		empty set
$e_1, ..., e_n$	e1, ..., en	set
$x \in S : p$	x \in S : p	set constructor
$e : x \in S$	e: x \in S	set constructor
$\div$	\div	integer division
$A \cdot B$	A \cdot B	composition of actions, executing A then B as one step
$\circ$	\o or \circ	concatenate sequences
$\bullet$	\doteq	
$\star$	\star	
$\bigcirc$	\bigcirc	
$\simeq$	\sim	stuttering equivalent
$\asymp$	\asymp	
$\approx$	\approx	
$\cong$	\cong	
$\doteq$	\doteq	
x^y	x^y	exponentiation
$'$	'	prime
$\sim$	\sim	stuttering equivalent
!	!	new record (in EXCEPT expression)
@	@	previous record field value (in EXCEPT expression)
:>	:>	One key-value mapping in a function (TLC module)
@@	@@	Function composition (TLC module)
α	\alpha	alpha
β	\beta	beta
γ	\gamma	gamma
Γ	\Gamma	Gamma
δ	\delta	delta
Δ	\Delta	Delta
ϵ	\epsilon	epsilon
ε	\varepsilon	variant epsilon
ζ	\zeta	zeta
η	\eta	eta
θ	\theta	theta

ϑ	\vartheta	variant theta
Θ	\Theta	Theta
ι	\iota	iota
κ	\kappa	kappa
λ	\lambda	lambda
Λ	\Lambda	Lambda
μ	\mu	mu
ν	\nu	nu
o	o	omicron
π	\pi	pi
Π	\Pi	Pi
ρ	\rho	rho
ϱ	\varrho	variant rho
σ	\sigma	sigma
ς	\varsigma	variant sigma
Σ	\Sigma	Sigma
τ	\tau	tau
υ	\upsilon	upsilon
Υ	\Upsilon	Upsilon
ϕ	\phi	phi
φ	\varphi	variant phi
Φ	\Phi	Phi
χ	\chi	chi
ψ	\psi	psi
Ψ	\Psi	Psi
ω	\omega	omega
Ω	\Omega	Omega
∂	\partial	partial