

# TryHackMe Room Write-up

Title: 0x41haz

Description: Simple Reversing Challenge

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## Find the password!

Description: *In this challenge, you are asked to solve a simple reversing solution. Download and analyze the binary to discover the password.*

*There may be anti-reversing measures in place!*

**Objective:** What is the password?

**Tools:** Kali Linux, hexeditor, ghidra, hexdump, strings.

Before we start with anything, let's gather more information about our downloaded file.

```
> file 0x41haz-1640335532346.0x41haz
0x41haz-1640335532346.0x41haz: ELF 64-bit MSB *unknown arch 0x3e00* (SYSV)
```

Picture 1: File information.

What can we tell right now?

ELF is Executable and Linkable Format – means Linux.

MSB is Most Significant Bit, indicates Big Endian encoding type.

There might be something wrong with this file - *\*unknown arch 0x3e00\* (SYSV)*, it might be the anti-reversing measure mentioned in description. Let's now see what strings are hidden within the file.

```
> strings 0x41haz-1640335532346.0x41haz
/lib64/ld-linux-x86-64.so.2
gets
exit
puts
strlen
__cxa_finalize
__libc_start_main
libc.so.6
GLIBC_2.2.5
_ITM_deregisterTMCloneTable
__gmon_start__
_ITM_registerTMCloneTable
u/UH
2@@25$gfH
sT&@f
[]A\A]A^A_
=====
Hey , Can You Crackme ?
=====
It's jus a simple binary
Tell Me the Password :
Is it correct , I don't think so.
Nope
Well Done !!
;*3$"
```

Picture 2: Selected strings output.

If you know your architectures and file formats, you can notice that this file was build for x86-64 architecture, which uses *Little Endian encoding*.

[Link on Endianess](#)

endianness is the dominant ordering for processor architectures (x86, most ARM implementations, base RISC-V implementations) and their associated memory. File formats can

Picture 3: Wiki page selected paragraph on which endianess is used on x86 architecture.

There's our problem! We get contradictory information: Big Endian used for x86 architecture? We must change it, but how? In order to do that, we should learn more about [ELF header](#).

#### ELF header<sup>[5]</sup>

Offset		Size (bytes)		Field	Purpose
32-bit	64-bit	32-bit	64-bit		
0x00		4		e_ident[EI_MAG0] through e_ident[EI_MAG3]	0x7F followed by ELF ( 45 4c 46 ) in ASCII; these four bytes constitute the magic number.
0x04		1		e_ident[EI_CLASS]	This byte is set to either 1 or 2 to signify 32- or 64-bit format, respectively.
0x05		1		e_ident[EI_DATA]	This byte is set to either 1 or 2 to signify little or big endianness, respectively. This affects interpretation of multi-byte fields starting with offset 0x10.

Picture 3: Wiki page on ELF header contents.

First four bytes are used to determine the *magic number*, fifth byte is used to signify 32- or 64-bit format. Sixth byte is our goal. It should be changed from big endian (2) to little endian (1).

Let's use hexeditor to make changes.

```
File: ./0x41haz-1640335532346.0x4 ASCII Offset: 0-00000000 / 0-0000385F (%00)
00000000 7F 45 4C 46 02 02 01 00 00 00 00 00 00 00 00 00 .ELF.....
```

Picture 4: ELF header before change.

```
File: ./0x41haz-1640335532346.0x4 ASCII Offset: 0-00000000 / 0-0000385F (%00)
00000000 7F 45 4C 46 02 01 01 00 00 00 00 00 00 00 00 00 .ELF.....
```

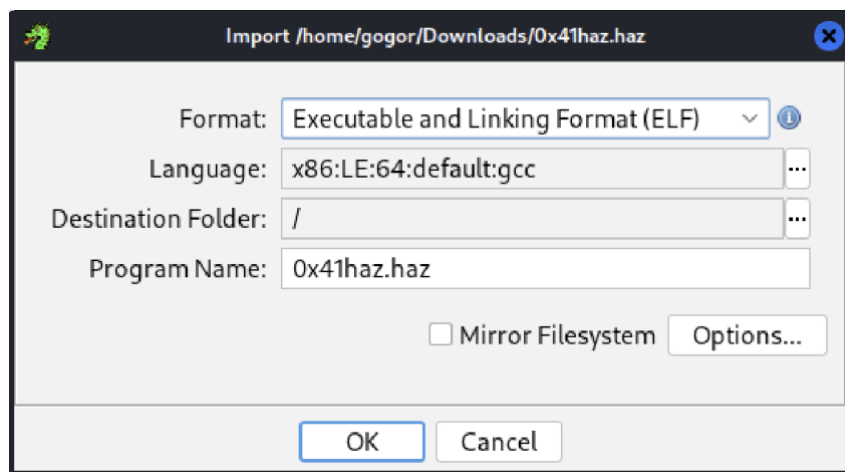
Picture 5: ELF header after change.

Now we can see if everything is correct now.

```
(gogor@kali)-[~/Downloads]
$ file 0x41haz.haz
0x41haz.haz: ELF 64-bit LSB pie executable, x86-64, version 1 (SYSV), dynamically
linked, interpreter /lib64/ld-linux-x86-64.so.2, BuildID[sha1]=6c9f2e85b64d4f12b91
136ffb8e4c038f1dc6dcd, for GNU/Linux 3.2.0, stripped
```

Picture 6: file information after changing the endian byte.

Now we can reverse engineer this bad boy. I used Ghidra.



Picture 7: Importing the file to Ghidra.

I went through the functions and in function FUN\_00101165 I found something juicy,

```
C# Decompile: FUN_00101165 - (0x41haz.haz)
1
2 undefined8 FUN_00101165(void)
3
4 {
5     size_t sVar1;
6     char local_48 [42];
7     char local_1e [14];
8     int local_10;
9     int local_c;
10
11     builtin_strncpy(local_1e,"2@@25$gfsT&@L",0xe);
12     puts("=====\nHey , Can You Crackme ?\n=====");
13     puts("It's jus a simple binary \n");
14     puts("Tell Me the Password :");
15     gets(local_48);
16     sVar1 = strlen(local_48);
17     local_10 = (int)sVar1;
18     if ((int)sVar1 != 0xd) {
19         puts("Is it correct , I don't think so.");
20         /* WARNING: Subroutine does not return */
21         exit(0);
22     }
23     local_c = 0;
24     while( true ) {
25         if (0xc < local_c) {
26             puts("Well Done !!");
27             return 0;
28         }
29         if (local_1e[local_c] != local_48[local_c]) break;
30         local_c = local_c + 1;
31     }
32     puts("Nope");
33     /* WARNING: Subroutine does not return */
34     exit(0);
35 }
36
```

Picture 7: FUN\_00101165 function contents.

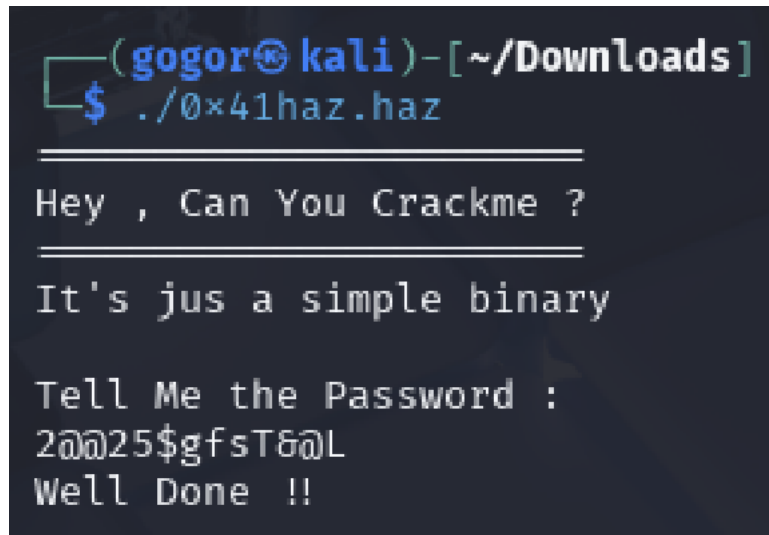
Let's focus on where the program gets our input, where it compares it and to what string.

```
builtin_strncpy(local_1e,"2@@25$gfsT&@L",0xe);
```

Picture 8: Fragment of interesting code.

Why is line on Picture 8 so interesting? Because it has something what looks like flag. After we try it, it is indeed what we're looking for, but why?

Program checks the length of our input (while ( true) loop) and the contents (if condition on line 29) character by character. If any of the sign in our input is different than the one in local\_le (variable to which the flag is copied to), the function breaks.



```
(gogor@kali)-[~/Downloads]
$ ./0x41haz.haz

=====
Hey , Can You Crackme ?
=====
It's jus a simple binary

Tell Me the Password :
2@@25$gfsT&@L
Well Done !!
```

Picture 9: Executing the code with wanted flag.