Lab 04: Algorithm Analysis

CS 0445: Data Structures

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http://db.cs.pitt.edu/courses/cs0445/current.term/

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- Mathematical notation to describe complexity of an algorithm compared to the size of the input
 - how many extra steps as the input grows
- Used to classify algorithms according to:
 - Runtime Main topic of discussion
 - Memory usage



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Recall the formal definition from lecture:

The function f(n) is O(g(n)) if:

- for some positive real number, c
- for some positive integer n_0
- $f(n) \le c * g(n)$ for all $n \ge n_0$

That is, c * g(n) is an upper bound on f(n), for sufficiently large n



Big O Notation – Broken Down

Our actual function

Must be a sufficiently large n (larger than n_0)

$$f(n) \le c * g(n)$$
 for all $n \ge n_0$

Some multiplicative constant c

The growth rate E.g., $n, n^2, \log n$, etc.



Big O Notation – Broken Down

$$f(n) \le c * g(n)$$
 for all $n \ge n_0$

- This boils down to...
 - Ignore multiplicative constants
 - Drop lower order terms
 - Take our worst growth rate, g(n).



- What's the Big O of:
 - \circ n^2
 - 2 + log n
 - o 2*n
 - $0 4 * n^2 + 2 * n$
 - \circ 2ⁿ + 2 * n



$$\circ n^2 \longrightarrow O(n^2)$$

$$0 4 * n^2 + 2 * n$$

$$\circ 2^n + 2 * n$$



$$\circ$$
 $n^2 \longrightarrow O(n^2)$

$$\circ$$
 2 + log n \longrightarrow O(log n)

$$0.4*n^2+2*n$$

$$\circ 2^n + 2 * n$$



$$\circ n^2 \longrightarrow O(n^2)$$

$$0 4 * n^2 + 2 * n$$

$$\circ 2^n + 2 * n$$



$$\circ$$
 $n^2 \longrightarrow O(n^2)$

$$\circ$$
 2 + log n \longrightarrow O(log n)

$$0 + n^2 + 2 * n \longrightarrow O(n^2)$$

$$\circ$$
 2ⁿ + 2 * n



$$\circ n^2 \longrightarrow O(n^2)$$

$$0 + n^2 + 2 * n \longrightarrow O(n^2)$$

$$\circ 2^n + 2 * n \longrightarrow O(2^n)$$



```
int counter = 0;
for (int i=0; i < n ; i++) {
    counter++;
}</pre>
```





```
int counter = 0;
for (int i=0; i < n ; i++) {
    counter++;
}</pre>
Will repeat n
times
```

O(n)



```
int counter = 0;
for (int i=0; i < n; i++) {
    counter++;
    int j = 6;
    int k = 10;
    int m = k*j;
}</pre>
```



```
int counter = 0;
for (int i=0; i < n ; i++) {
    counter++;
    int j = 6;
    int k = 10;
    int m = k*j;
}</pre>

O(n)
```

These additional operations do not affect how many times this loop repeats so they do not affect the growth rate



```
for (int i=0; i < n ; i++) {
    for (int j=0; j < n ; j++) {
        counter++;
    }
}</pre>
```



```
for (int i=0; i < n; i++) {
    for (int j=0; j < n; j++) {
        counter++;
    }
    Will repeat
    n times</pre>
```

 $O(n^2)$



```
for (int i=0; i < n ; i++) {
    print_mult_row(i, n);
}

public void print_mult_row(int mul, int size) {
    for ( int j = 0; j <= size; j++ ) {
        System.out.print(mul*j + "\t");
    }
}</pre>
```



```
for (int i=0; i < n; i++) {</pre>
                                       Will repeat
      print mult row(i, n);
                                       n times
                                          O(n^2)
 public void print mult row(int mul, int size) {
       for ( int j = 0; j <= size; j++ ) {</pre>
             System.out.print(mul*j + "\t");
                                               Will repeat
                                               n times
```



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How to analyze code

- Remember to look at the loops
 - Determine how many times they'll repeat compared to n
 - Multiply the runtime of inner loops by the runtime of outer loops



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Worksheet

