#### Lecture 10: Stacks

#### **CS 0445: Data Structures**

#### **Constantinos Costa**

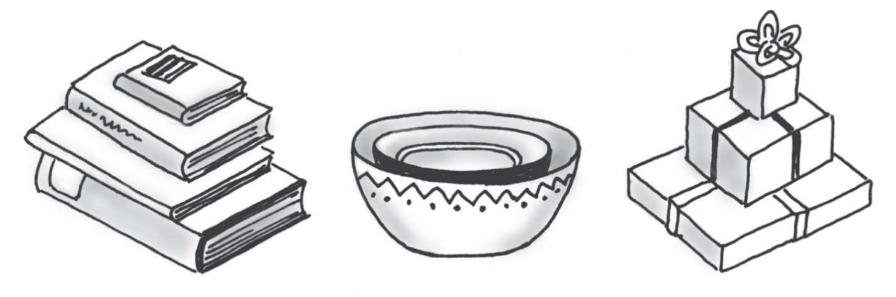
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#### Stacks

- Add item on top of stack
- Remove item that is topmost
  - Last In, First Out ... LIFO



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#### Specifications of the ADT Stack

- Data
  - A collection of objects in reverse chronological order and having the same data type

Pseudocode	UML	Description
push(newEntry)	+push(newEntry: T): void	Task: Adds a new entry to the top of the stack. Input: newEntry is the new entry. Output: None.
pop()	+pop(): T	Task: Removes and returns the stack's top entry. Input: None. Output: Returns the stack's top entry. Throws an exception if the stack is empty before the operation.
peek()	+peek(): T	Task: Retrieves the stack's top entry without changing the stack in any way.  Input: None.  Output: Returns the stack's top entry.  Throws an exception if the stack is empty.
isEmpty()	+isEmpty(): boolean	Task: Detects whether the stack is empty. Input: None. Output: Returns true if the stack is empty.
clear()	+clear(): void	Task: Removes all entries from the stack. Input: None. Output: None.

#### Design Decision

- When stack is empty
  - What to do with pop and peek?
- Possible actions
  - Assume that the ADT is not empty;
  - Return null.
  - Throw an exception (which type?).

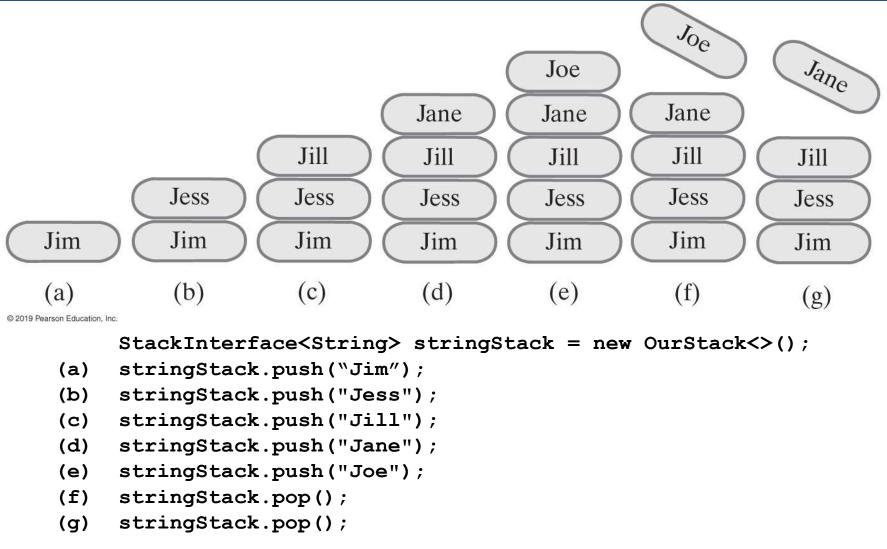


#### Interface for the ADT Stack

```
/** An interface for the ADT stack. */
public interface StackInterface<T>
 /** Adds a new entry to the top of this stack.
   @param newEntry An object to be added to the stack. */
 public void push(T newEntry);
 /** Removes and returns this stack's top entry.
   @return The object at the top of the stack.
   @throws EmptyStackException if the stack is empty before the operation. */
 public T pop();
 /** Retrieves this stack's top entry.
    @return The object at the top of the stack.
   @throws EmptyStackException if the stack is empty. */
 public T peek();
 /** Detects whether this stack is empty.
   @return True if the stack is empty. */
 public boolean isEmpty();
 /** Removes all entries from this stack. */
 public void clear();
} // end StackInterface
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```



#### Example of a Stack





#### Security Note

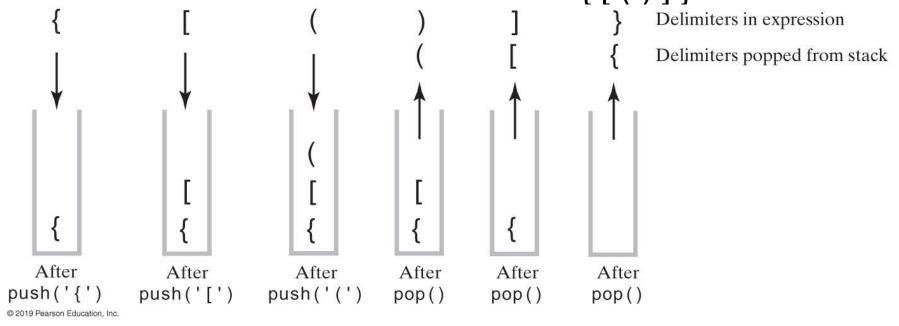
- Design guidelines
  - Use preconditions and postconditions to document assumptions.
  - Do not trust client to use public methods correctly.
  - Avoid ambiguous return values.
  - Prefer throwing exceptions instead of returning values to signal problem.



- Infix:
  - each binary operator appears between its operandsa + b
- Prefix:
  - each binary operator appears before its operands+ a b
- Postfix:
  - each binary operator appears after its operandsa b +
- Balanced expressions: delimiters paired correctly

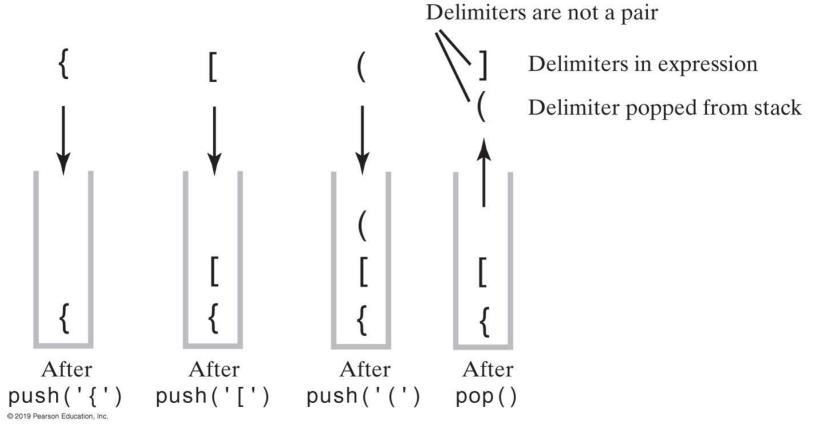


 The contents of a stack during the scan of an expression that contains the balanced delimiters{ [ ( ) ] }



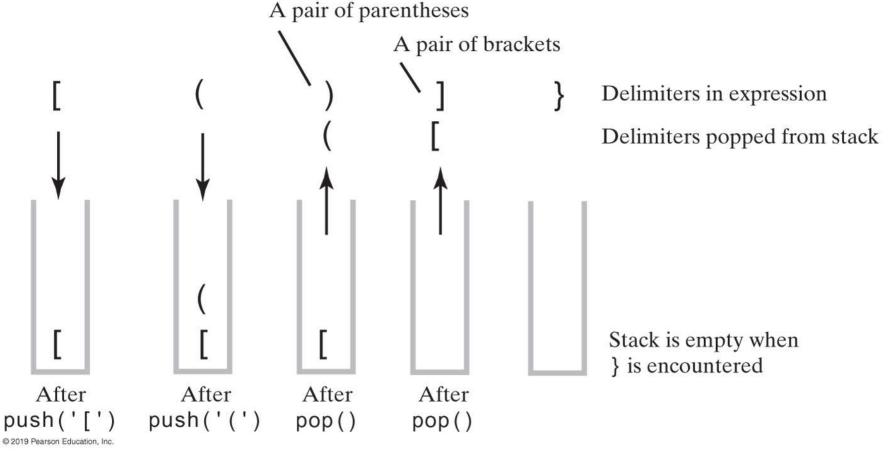


 The contents of a stack during the scan of an expression that contains the unbalanced delimiters { [ ( ] ) }



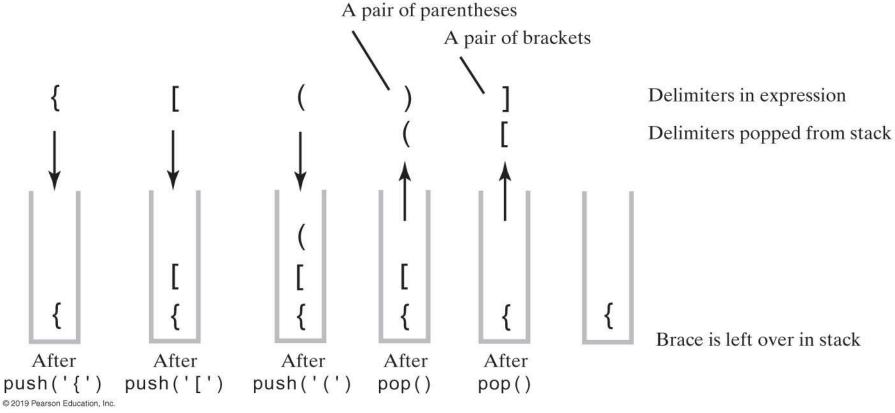


 The contents of a stack during the scan of an expression that contains the unbalanced delimiters [ ( ) ] }





 The contents of a stack during the scan of an expression that contains the unbalanced delimiters { [ ( ) ]





```
Algorithm checkBalance(expression)
//Returns true if the parentheses, brackets, and braces in an expression are paired correctly.
isBalanced = true // The absence of delimiters is balanced
while ((isBalanced == true) and not at end of expression)
   nextCharacter = next character in expression
   switch (nextCharacter)
      case '(': case '[': case '{':
                                   Push nextCharacter onto stack
                                   break
      case ')': case ']': case '}':
          if (stack is empty)
             isBalanced = false
          else
             openDelimiter = top entry of stack
             Pop stack
             isBalanced = true or false according to whether openDelimiter and
                                  nextCharacter are a pair of delimiters
          break
if (stack is not empty)
   isBalanced = false
return isBalanced
```



## Implementation of Algorithm (Part 1)

```
/** A class that checks whether the parentheses, brackets, and braces
 in a string occur in left/right pairs. */
public class BalanceChecker
 // Returns true if the given characters, open and close, form a pair
 // of parentheses, brackets, or braces.
 private static boolean isPaired(char open, char close)
   return (open == '(' && close == ')') | |
       (open == '[' && close == ']') | |
       (open == '{' && close == '}');
 } // end isPaired
 /** Decides whether the parentheses, brackets, and braces
   in a string occur in left/right pairs.
    @param expression A string to be checked.
   @return True if the delimiters are paired correctly. */
 public static boolean checkBalance(String expression)
 } // end checkBalance
} // end BalanceChecker
```



#### Implementation of Algorithm (Part 2)

```
StackInterface<Character> openDelimiterStack = new LinkedStack<>();
 int characterCount = expression.length();
 boolean isBalanced = true:
 int index = 0;
 char nextCharacter = ' ';
 while (isBalanced && (index < characterCount)) {
   nextCharacter = expression.charAt(index);
   switch (nextCharacter) {
     case '(': case '[': case '{':
      openDelimiterStack.push(nextCharacter);
      break:
     case ')': case ']': case '}':
      if (openDelimiterStack.isEmpty())
        isBalanced = false;
      else {
        char openDelimiter = openDelimiterStack.pop();
        isBalanced = isPaired(openDelimiter, nextCharacter);
      } // end if
      break:
     default: break; // Ignore unexpected characters
   } // end switch
   index++;
 } // end while
 if (!openDelimiterStack.isEmpty())
   isBalanced = false;
 return isBalanced;
} // end checkBalance
```



# Converting Infix to Postfix

a + b \* c

Next Character in Infix Expression	Postfix Form	Operator Stack (bottom to top)
а	a	
+	a	+
b	a b	+
*	a b	+ *
c	abc	+ *
	a b c *	+
	abc*+	

a ^ b ^ c

Next Character in Infix Expression	Postfix Form	Operator Stack (bottom to top)
а	а	
٨	a	٨
b	a b	٨
٨	a b	^^
c	a b c	^ ^
	a b c ^ a b c ^ ^	٨
	a b c ^ ^	

a - b + c

Next Character in Infix Expression	Postfix Form	Operator Stack (bottom to top)
а	a	
_	a	<u> </u>
b	a b	_
+	a b - a b -	
	a b -	+
C	ab-c	+
	ab-c+	



#### Converting Infix to Postfix

Steps in converting an infix expression to postfix form

Next Character from Infix Expression	Postfix Form	Operator Stack (bottom to top)	
a	a		
/	a	/	
b	a b	/	
*	a b /		
20	a b /	*	
(	a b /	* (	
c	a b / c	* (	
+	a b / c	* (+	
(	a b / c	* (+ (	
d	a b / c d	* (+ (	
_	a b / c d	* (+ (-	
e	ab/cde	* (+ (-	
)	a b / c d e -	* (+ (	
	a b / c d e -	*(+	
)	a b / c d e − +	* (	
32	a b / c d e - +	*	
	ab/cde-+*		



#### Infix-to-postfix Conversion

 To convert an infix expression to postfix form, you take the following actions, according to the symbols you encounter, as you process the infix expression from left to right:

Operand	Append each operand to the end of the output expression.	
Operator ^	Push ^ onto the stack.	
Operator +, -, *, or /	Pop operators from the stack, appending them to the output expression, until either the stack is empty or its top entry has a lower precedence than the newly encountered operator. Then push the new operator onto the stack.	
Open parenthesis	Push ( onto the stack.	
Close parenthesis	Pop operators from the stack and append them to the output expression until an open parenthesis is popped. Discard both parentheses.	



## Infix-to-postfix Algorithm (Part 1)

```
Algorithm convertToPostfix(infix)
// Converts an infix expression to an equivalent postfix expression.
operatorStack = a new empty stack
postfix = a new empty string
while (infix has characters left to parse)
   nextCharacter = next nonblank character of infix
   switch (nextCharacter)
      case variable:
         Append nextCharacter to postfix
         break
      case '^' :
         operatorStack.push(nextCharacter)
         break
      case '+' : case '-' : case '*' : case '/' :
         while (!operatorStack.isEmpty() and
           precedence of nextCharacter <= precedence of operatorStack.peek())</pre>
            Append operatorStack.peek() to postfix
             operatorStack.pop()
        operatorStack.push(nextCharacter)
        break
```



### Infix-to-postfix Algorithm (Part 2)

```
case '(':
         operatorStack.push(nextCharacter)
         break
      case ') ' : // Stack is not empty if infix expression is valid
          topOperator = operatorStack.pop()
         while (topOperator != '(')
             Append topOperator to postfix
             topOperator = operatorStack.pop()
         break
      default:
         break // Ignore unexpected characters
while (!operatorStack.isEmpty())
   topOperator = operatorStack.pop()
   Append topOperator to postfix
return postfix
```

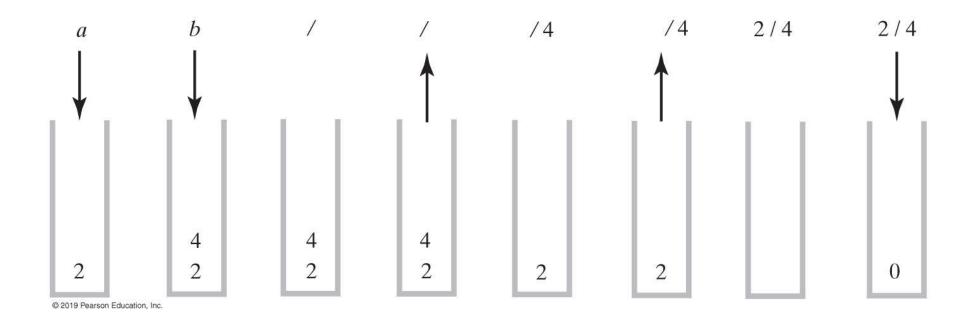


#### **Evaluating Postfix Expressions**

```
Algorithm evaluatePostfix(postfix)
// Evaluates a postfix expression.
valueStack = a new empty stack
while (postfix has characters left to parse)
   nextCharacter = next nonblank character of postfix
   switch (nextCharacter)
      case variable:
         valueStack.push(value of the variable nextCharacter)
         break
      case '+' : case '-' : case '*' : case '/' : case '^' :
         operandTwo = valueStack.pop()
         operandOne = valueStack.pop()
         result = the result of the operation in nextCharacter and
                                  its operands operandOne and operandTwo
         valueStack.push(result)
         break
      default: break // Ignore unexpected characters
return valueStack.peek()
```

## **Evaluating Infix Expressions**

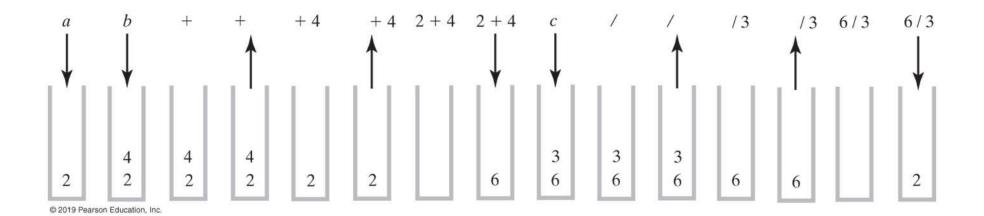
The stack during the evaluation of the postfix expression
 a b / when a is 2 and b is 4





#### **Evaluating Infix Expressions**

The stack during the evaluation of the postfix expression a
 b + c / when a is 2, b is 4, and c is 3





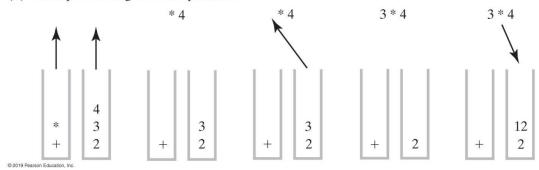
## Evaluating Infix Expressions

(a) After reaching the end of the expression

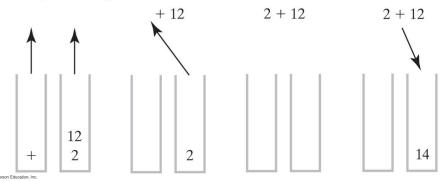


- Two stacks during the evaluation of
- a + b \* c when a is 2, b is 3, and c is 4

(b) While performing the multiplication



(c) While performing the addition



# Evaluating Infix Expressions (Part 1)

```
Algorithm evaluateInfix(infix) // Evaluates an infix expression.
operatorStack = a new empty stack
valueStack = a new empty stack
while (infix has characters left to process)
{
   nextCharacter = next nonblank character of infix
   switch (nextCharacter)
      case variable:
         valueStack.push(value of the variable nextCharacter)
         break
      case '^':
         operatorStack.push(nextCharacter)
         break
      case '+' : case '-' : case '*' : case '/' :
         while (!operatorStack.isEmpty() and
                   precedence of nextCharacter <= precedence of operatorStack.peek())</pre>
            // Execute operator at top of operatorStack
            topOperator = operatorStack.pop()
            operandTwo = valueStack.pop()
            operandOne = valueStack.pop()
            result = the result of the operation in
                            topOperator and its operands operandOne and operandTwo
            valueStack.push(result)
         operatorStack.push(nextCharacter)
         break
```



# Evaluating Infix Expressions (Part 2)

```
case '(':
         operatorStack.push (nextCharacter)
         break
      case ') ' : // Stack is not empty if infix expression is valid
         topOperator = operatorStack.pop()
         while (topOperator != '(')
            operandTwo = valueStack.pop()
            operandOne = valueStack.pop()
            result = the result of the operation in
                             topOperator and its operands operandOne and operandTwo
            valueStack.push(result)
            topOperator = operatorStack.pop()
         break
      default: break // Ignore unexpected characters
   }
while (!operatorStack.isEmpty())
   topOperator = operatorStack.pop()
   operandTwo = valueStack.pop()
   operandOne = valueStack.pop()
   result = the result of the operation in
                   topOperator and its operands operandOne and operandTwo
   valueStack.push(result)
return valueStack.peek()
```



#### The Application Program Stack

The program stack as a program executes

```
public static
     void main(string[] arg)
        int x = 5;
        int y = methodA(x);
     } // end main
    public static
100
     int methodA(int a)
        int z = 2;
        methodB(z);
120
        . . .
        return z;
     } // end methodA
    public static
     void methodB(int b)
     } // end methodB
```

```
(a) When main
                       (b) When methodA
                                              (c) When methodB
  begins execution
                          begins execution
                                                 begins execution
                                                  methodB
                          methodA
                                                  methodA
                            PC = 100
                             a = 5
   main
                          main
                                                 main
                            PC = 50
     PC = 1
                            arg = \dots
     arg = \dots
                            x = 5
                             y = 0
```

Program

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Program stack at three points in time (PC is the program counter)

PC = 150

b = 2

PC = 120

a = 5

z = 2

PC = 50

x = 5

v = 0

 $arg = \dots$ 



#### Java Class Library: The Class Stack

- Found in java.util
- Methods

```
- A constructor - creates an empty stack
- public T push(T item);
- public T pop();
- public T peek();
- public boolean empty();
```

