

Lecture 27: BST Implementation

CS 0445: Data Structures

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<http://db.cs.pitt.edu/courses/cs0445/current.term/>

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Binary Search Tree

- For each node in a binary search tree
 - Node's data is greater than all data in node's left subtree
 - Node's data is less than all data in node's right subtree
- Every node in a binary search tree is the root of a binary search tree

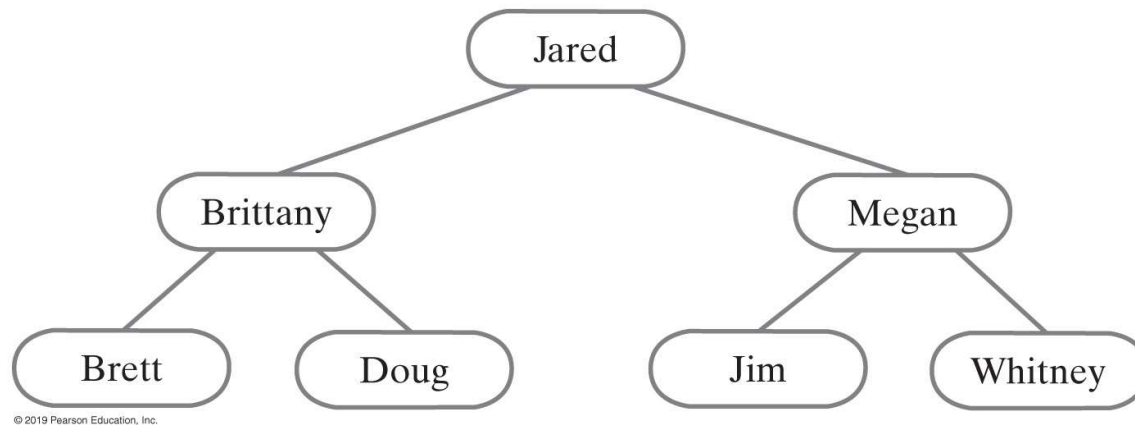


FIGURE 26-1 A binary search tree of names

Interface for the Binary Search Tree (Part 1)

- LISTING 26-1 An interface for a search tree

```
/** An interface for a search tree. */
public interface SearchTreeInterface<T extends Comparable<? super T>>
    extends TreeInterface<T>
{
    /** Searches for a specific entry in this tree.
     * @param anEntry An object to be found.
     * @return True if the object was found in the tree. */
    public boolean contains(T anEntry);

    /** Retrieves a specific entry in this tree.
     * @param anEntry An object to be found.
     * @return Either the object that was found in the tree or
     *         null if no such object exists. */
    public T getEntry(T anEntry);

    /** Adds a new entry to this tree, if it does not match an existing
     * object in the tree. Otherwise, replaces the existing object with
     * the new entry.
     * @param anEntry An object to be added to the tree.
     * @return Either null if anEntry was not in the tree but has been added, or
     *         the existing entry that matched the parameter anEntry
     *         and has been replaced in the tree. */
    public T add(T anEntry);
}
```



Interface for the Binary Search Tree (Part 2)

- LISTING 26-1 An interface for a search tree

```
/** Removes a specific entry from this tree.  
    @param anEntry An object to be removed.  
    @return Either the object that was removed from the tree or  
            null if no such object exists. */  
public T remove(T anEntry);  
  
/** Creates an iterator that traverses all entries in this tree.  
    @return An iterator that provides sequential and ordered access  
            to the entries in the tree. */  
public Iterator<T> getInorderIterator();  
} // end SearchTreeInterface
```



Understanding the Specifications

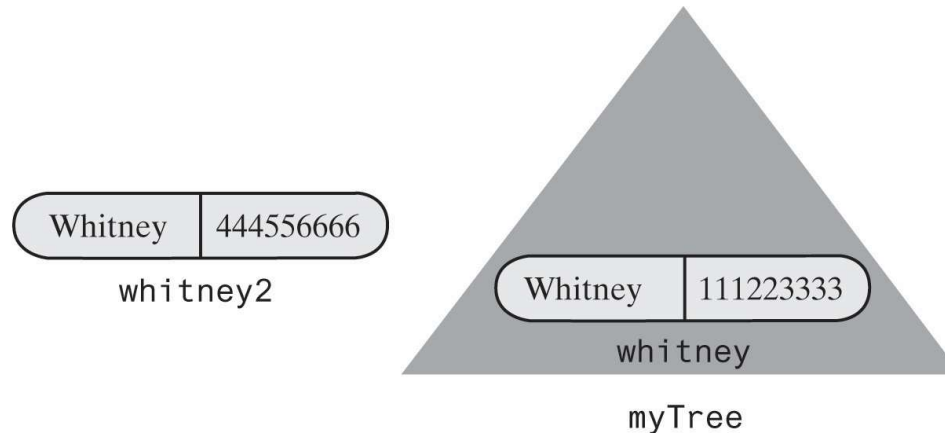
- Methods will use return values instead of exceptions to indicate whether an operation has failed
 - **getEntry**, returns same entry it is given to find
 - **getEntry** returns an object in tree and matches given entry according to the entry's **compareTo** method



Adding to a Binary Search Tree

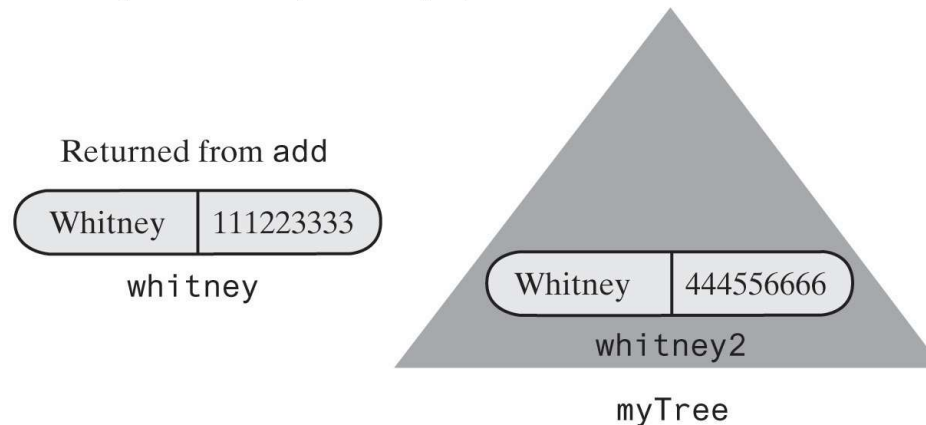
- FIGURE 26-2 Adding an entry that matches an entry already in a binary search tree

(a) Before `myTree.add(whitney2)` executes



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(b) After `myTree.add(whitney2)` executes



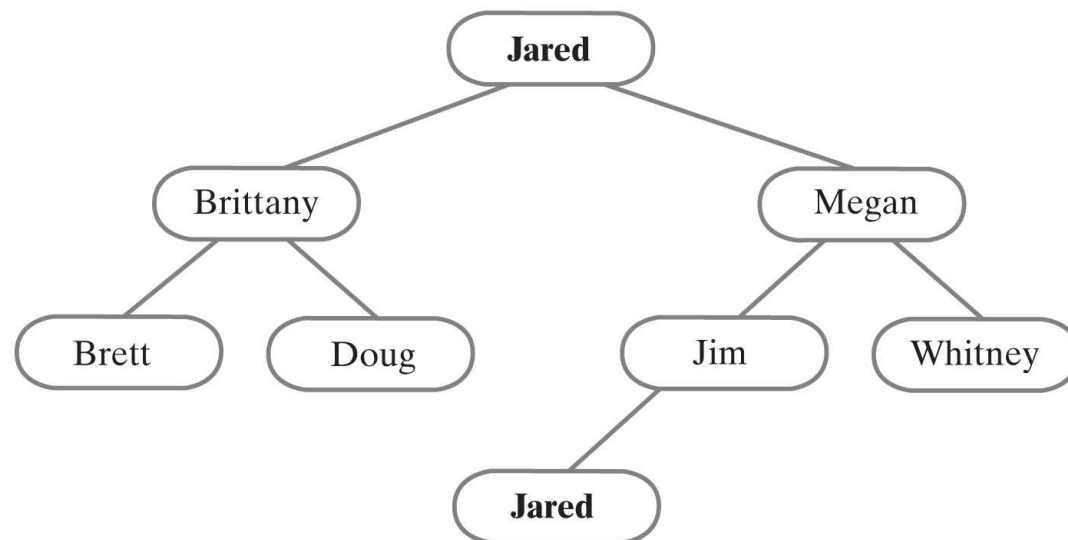
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Duplicate Entries

- If any entry e has a duplicate entry d , we arbitrarily require that d occur in the right subtree of e 's node
- For each node in a binary search tree:
 - Data in a node is greater than data in node's left subtree
 - Data in a node is less than or equal to data in node's right subtree

FIGURE 26-3
A binary search
tree with duplicate
entries



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Beginning the Class Definition

- LISTING 26-2 An outline of the class `BinarySearchTree`

```
package TreePackage;
import java.util.Iterator;
/** A class that implements ADT binary search tree by extending BinaryTree. */
public class BinarySearchTree<T extends Comparable<? super T>>
    extends BinaryTree<T> implements SearchTreeInterface<T>
{
    public BinarySearchTree()
    {
        super();
    } // end default constructor

    public BinarySearchTree(T rootEntry)
    {
        super();
        setRootNode(new BinaryNode<T>(rootEntry));
    } // end constructor

    // Disable setTree (see Segment 26.6)
    public void setTree(T rootData, BinaryTreeInterface<T> leftTree,
        BinaryTreeInterface<T> rightTree)
    {
        throw new UnsupportedOperationException();
    } // end setTree

    /* Implementations of other methods goes here. */
} // end BinarySearchTree
```



Searching and Retrieving

- Recursive algorithm to search a binary search tree

Algorithm **bstSearch(binarySearchTree, desiredObject)**

// Searches a binary search tree for a given object.

// Returns true if the object is found.

if (binarySearchTree is empty)

return false

else if (desiredObject == *object in the root of binarySearchTree*)

return true

else if (desiredObject < *object in the root of binarySearchTree*)

return bstSearch(*left subtree of binarySearchTree*, desiredObject)

else

return bstSearch(*right subtree of binarySearchTree*, desiredObject)



Searching and Retrieving

- Algorithm that describes actual implementation more closely

Algorithm **bstSearch(binarySearchTreeRoot, desiredObject)**

// Searches a binary search tree for a given object.

// Returns true if the object is found.

if (binarySearchTreeRoot is null)

return false

else if (desiredObject == *object in* binarySearchTreeRoot)

return true

else if (desiredObject < *object in* binarySearchTreeRoot)

return bstSearch(*left child of* binarySearchTreeRoot, desiredObject)

else

return bstSearch(*right child of* binarySearchTreeRoot, desiredObject)



Searching and Retrieving

- The method `getEntry` uses `findEntry`

```
public T getEntry(T anEntry)
{
    return findEntry(getRootNode(), anEntry);
} // end getEntry

private T findEntry(BinaryNode<T> rootNode, T anEntry)
{
    T result = null;

    if (rootNode != null)
    {
        T rootEntry = rootNode.getData();

        if (anEntry.equals(rootEntry))
            result = rootEntry;
        else if (anEntry.compareTo(rootEntry) < 0)
            result = findEntry(rootNode.getLeftChild(), anEntry);
        else
            result = findEntry(rootNode.getRightChild(), anEntry);
    } // end if

    return result;
} // end findEntry
```



Searching and Retrieving

- Method `contains` can simply call `getEntry` to see whether a given entry is in the tree

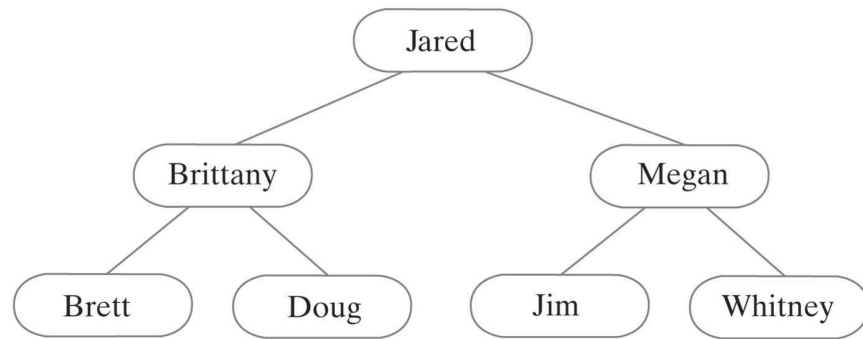
```
public boolean contains(T anEntry)
{
    return getEntry(anEntry) != null;
} // end contains
```



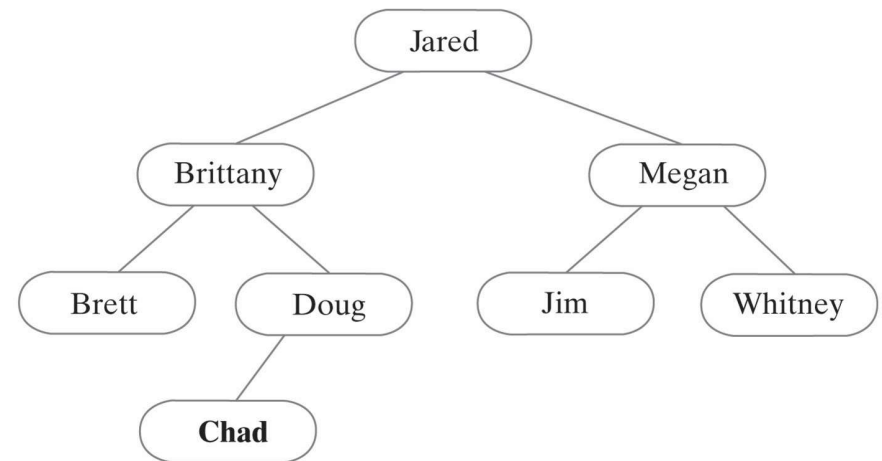
Adding to a Binary Search Tree

- FIGURE 26-4 A binary search tree before and after adding Chad

(a) A binary search tree

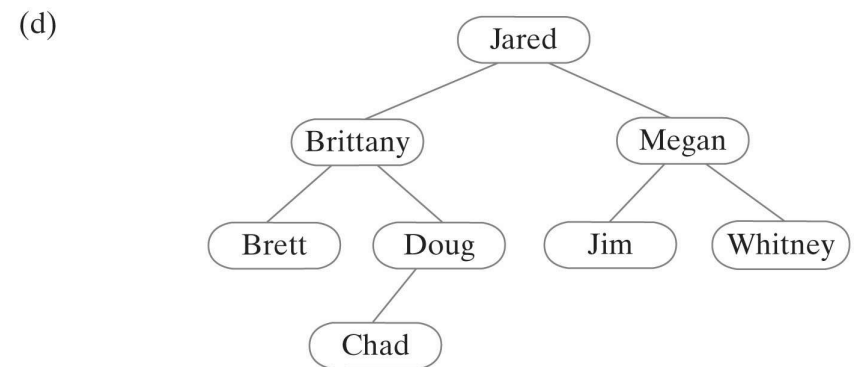
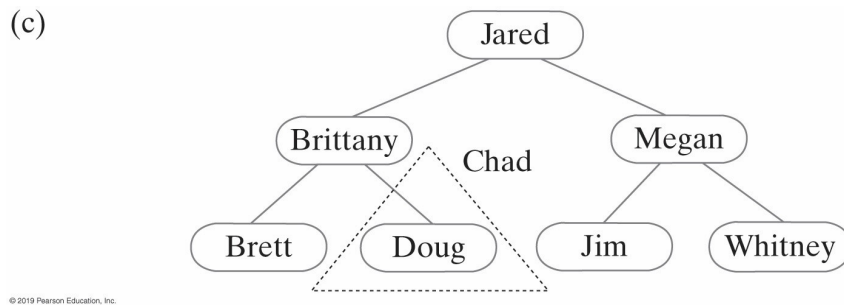
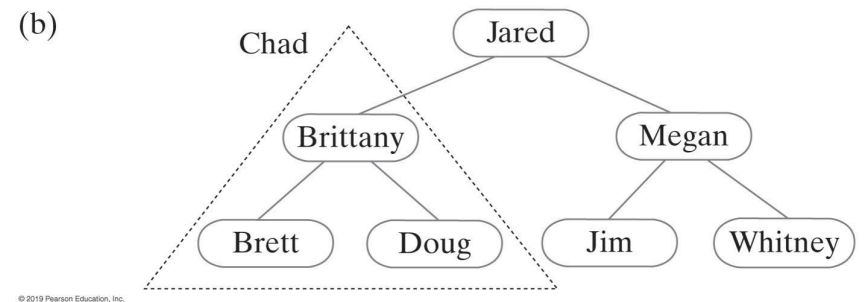
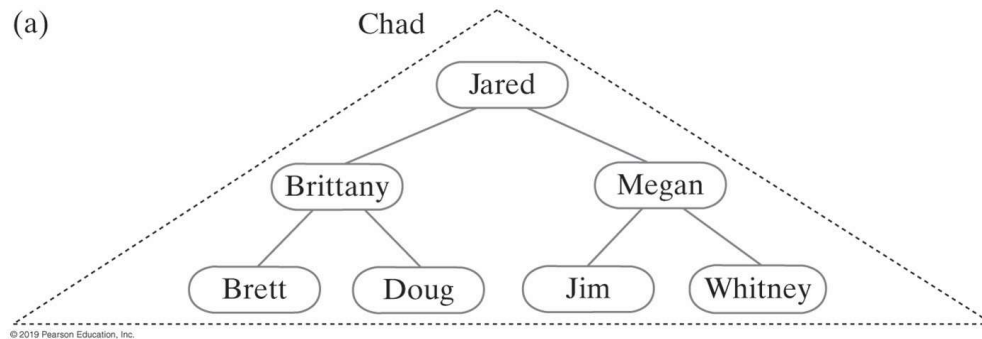


(b) The same tree after adding *Chad*



Adding to a Binary Search Tree

- FIGURE 26-5 Recursively adding Chad to smaller subtrees of a binary search tree



Recursive Add Implementation

- Recursive algorithm for adding a new entry

Algorithm **addEntry(binarySearchTree, anEntry)**

// Adds an entry to a binary search tree that is not empty.

*// Returns null if anEntry did not exist already in the tree. Otherwise, returns the
// tree entry that matched and was replaced by anEntry.*

result = null

if (anEntry matches the entry in the root of binarySearchTree)

{

result = entry in the root

Replace entry in the root with anEntry

}

else if (anEntry < entry in the root of binarySearchTree)

{

if (the root of binarySearchTree has a left child)

result = addEntry(left subtree of binarySearchTree, anEntry)

else

Give the root a left child containing anEntry

}

else // anEntry > entry in the root of binarySearchTree

{

if (the root of binarySearchTree has a right child)

result = addEntry(right subtree of binarySearchTree, anEntry)

else

Give the root a right child containing anEntry

}

return result



Recursive Implementation

- Handle the addition to an empty binary search tree as a special case

Algorithm **add(binarySearchTree, anEntry)**

// Adds an entry to a binary search tree.

// Returns null if anEntry did not exist already in the tree. Otherwise, returns the

// tree entry that matched and was replaced by anEntry.

result = null

if (binarySearchTree is empty)

Create a node containing anEntry and make it the root of binarySearchTree

else

result = addEntry(binarySearchTree, anEntry)

return result;



Recursive Implementation

- The public method `add`

```
public T add(T anEntry)
{
    T result = null;

    if (isEmpty())
        setRootNode(new BinaryNode<>(anEntry));
    else
        result = addEntry(getRootNode(), anEntry);

    return result;
} // end add
```



Recursive Implementation — method `addEntry`

// Adds anEntry to the nonempty subtree rooted at rootNode.

```
private T addEntry(BinaryNode<T> rootNode, T anEntry)
{
    // Assertion: rootNode != null
    T result = null;
    int comparison = anEntry.compareTo(rootNode.getData());

    if (comparison == 0)
    {
        result = rootNode.getData();
        rootNode.setData(anEntry);
    }
    else if (comparison < 0)
    {
        if (rootNode.hasLeftChild())
            result = addEntry(rootNode.getLeftChild(), anEntry);
        else
            rootNode.setLeftChild(new BinaryNode<>(anEntry));
    }
    else
    {
        if (rootNode.hasRightChild())
            result = addEntry(rootNode.getRightChild(), anEntry);
        else
            rootNode.setRightChild(new BinaryNode<>(anEntry));
    } // end if

    return result;
} // end addEntry
```



Iterative Implementation (Part 1)

- Iterative algorithm for adding a new entry

Algorithm **addEntry(binarySearchTree, anEntry)**

// Adds a new entry to a binary search tree that is not empty.

*// Returns null if anEntry did not exist already in the tree. Otherwise, returns the
// tree entry that matched and was replaced by anEntry.*

result = null

currentNode = root node of binarySearchTree **found = false**

while (found is false)

{

if (anEntry matches the entry in currentNode)

{

found = true

result = entry in currentNode

Replace entry in currentNode with anEntry

}

else if (newEntry < entry in currentNode)

{



Iterative Implementation (Part 2)

- Iterative algorithm for adding a new entry

```
    if (currentNode has a left child)
        currentNode = left child of currentNode
    else
    {
        found = true
        Give currentNode a left child containing anEntry
    }
} // end if-else
else // anEntry > entry in currentNode
{
    if (currentNode has a right child)
        currentNode = right child of currentNode
    else
    {
        found = true
        Give currentNode a right child containing anEntry
    }
} // end if
} // end while
return result
```



Iterative Implementation— method `addEntry` (Part 1)

```
private T addEntry(T anEntry) {
    BinaryNode<T> currentNode = getRootNode();
    // Assertion: currentNode != null
    T result = null;
    boolean found = false;

    while (!found)
    {
        T currentEntry = currentNode.getData();
        int comparison = anEntry.compareTo(currentEntry);

        if (comparison == 0)
        { // anEntry matches currentEntry;
          // return and replace currentEntry
            found = true;
            result = currentEntry;
            currentNode.setData(anEntry);
        }
        else if (comparison < 0)
        {
            if (currentNode.hasLeftChild())
                currentNode = currentNode.getLeftChild();
            else
            {
                found = true;
                currentNode.setLeftChild(new BinaryNode<>(anEntry));
            } // end if
        }
    }
}
```

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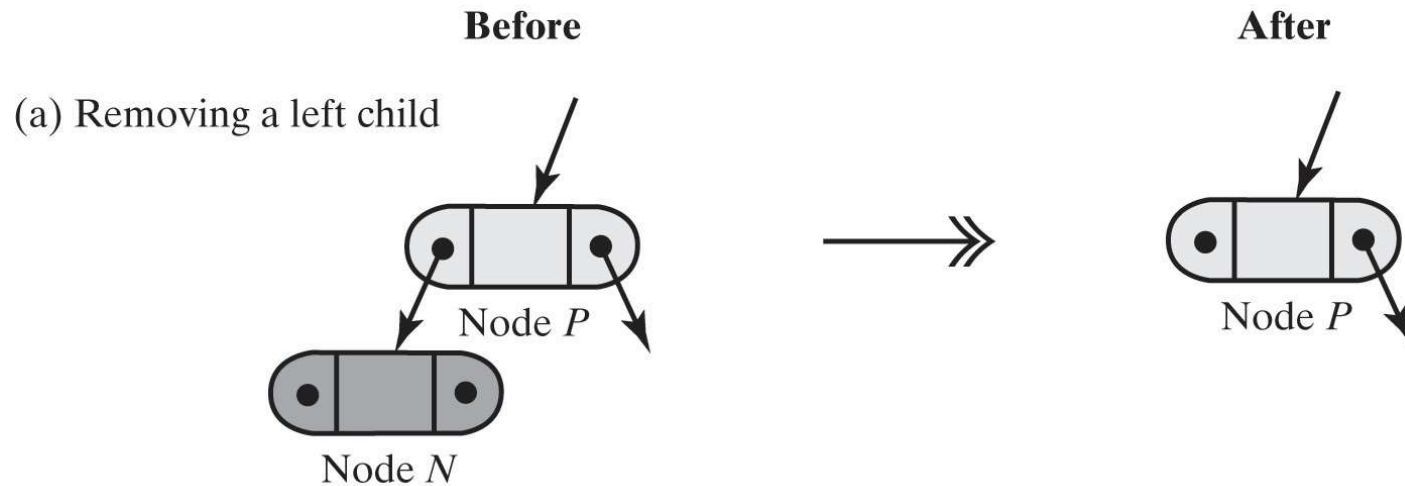
Iterative Implementation— method `addEntry` (Part 2)

```
else {  
    // Assertion: comparison > 0  
  
    if (currentNode.hasRightChild())  
        currentNode = currentNode.getRightChild();  
    else  
    {  
        found = true;  
        currentNode.setRightChild(new BinaryNode<>(anEntry));  
    } // end if  
} // end if  
} // end while  
  
return result;  
} // end addEntry
```

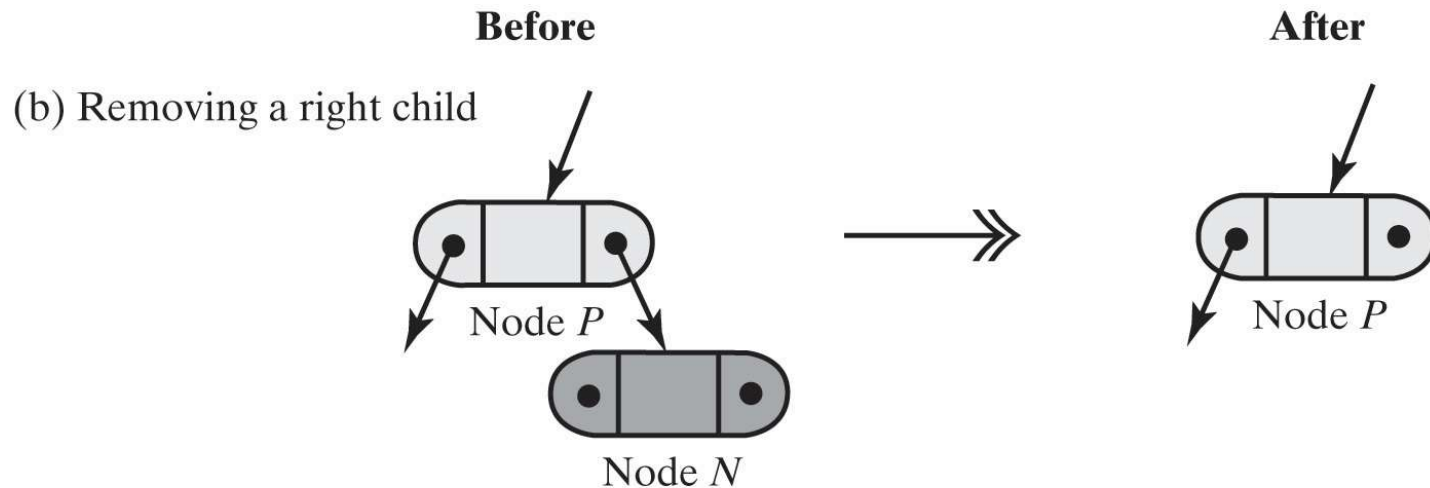


Removing a Value

- FIGURE 26-6 Removing a leaf node *N* from its parent node *P*



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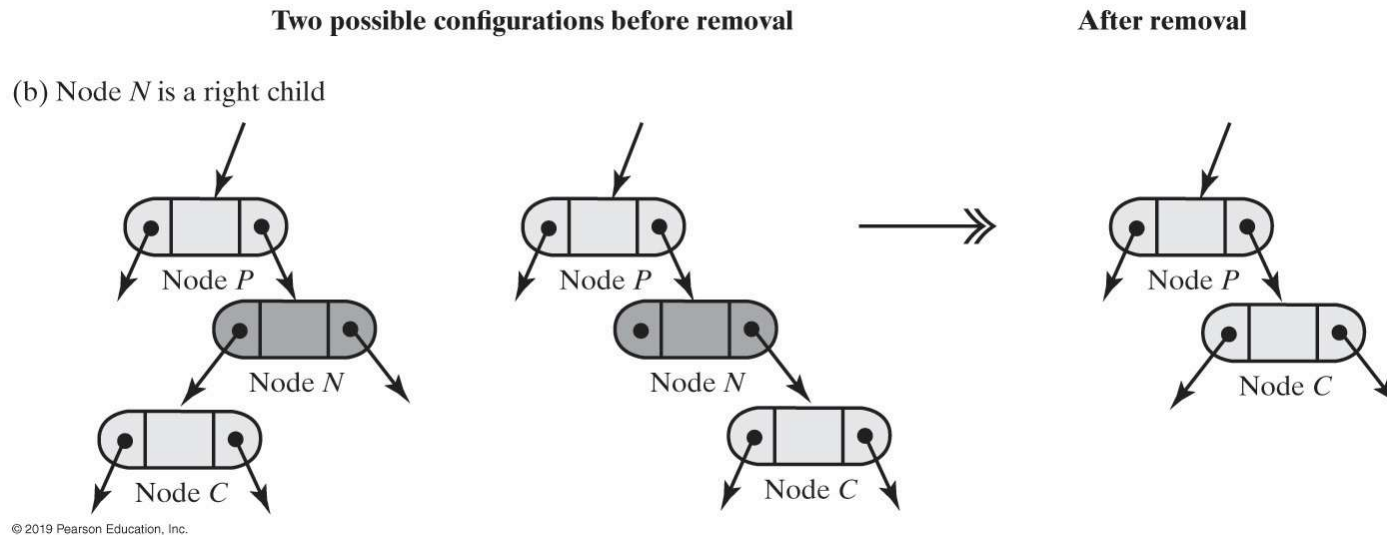
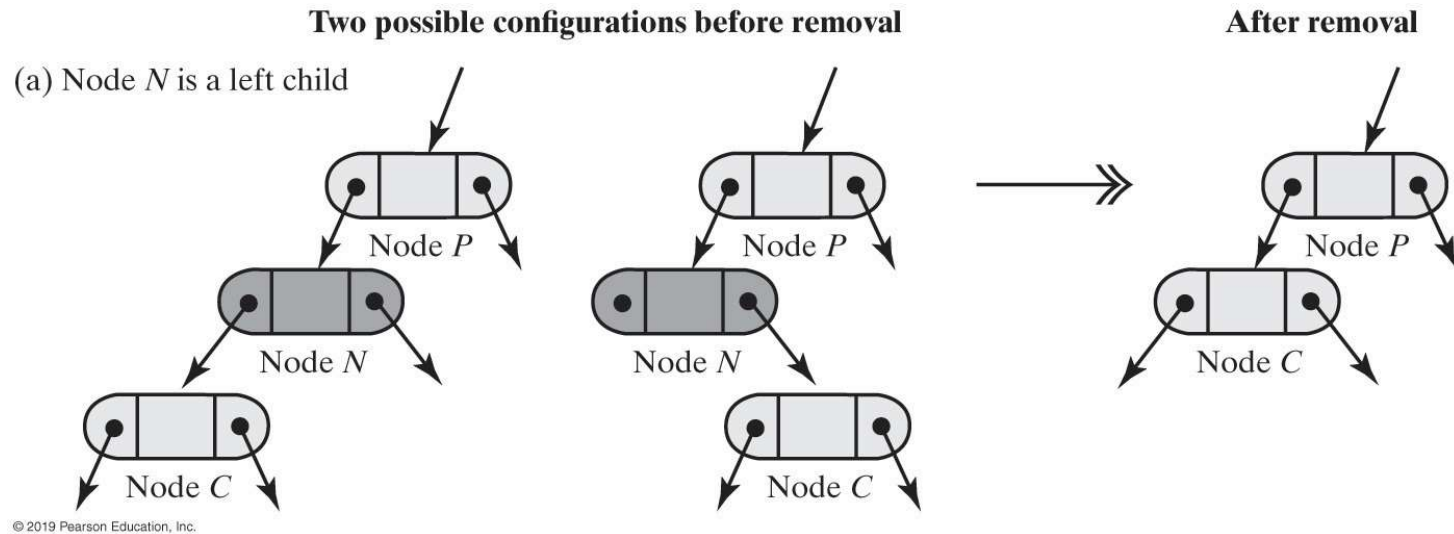


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Removing a Value

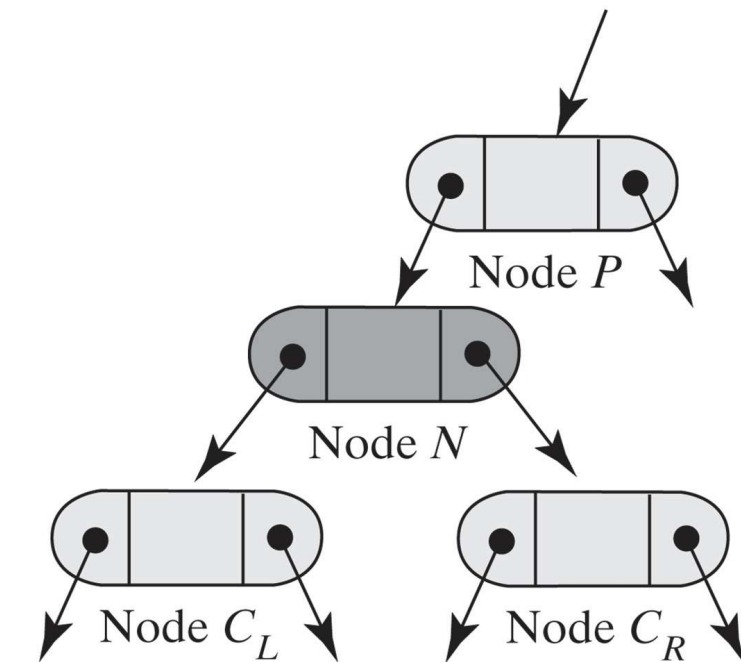
- FIGURE 26-7 Removing a node N from its parent node P when N has one child



Removing a Value

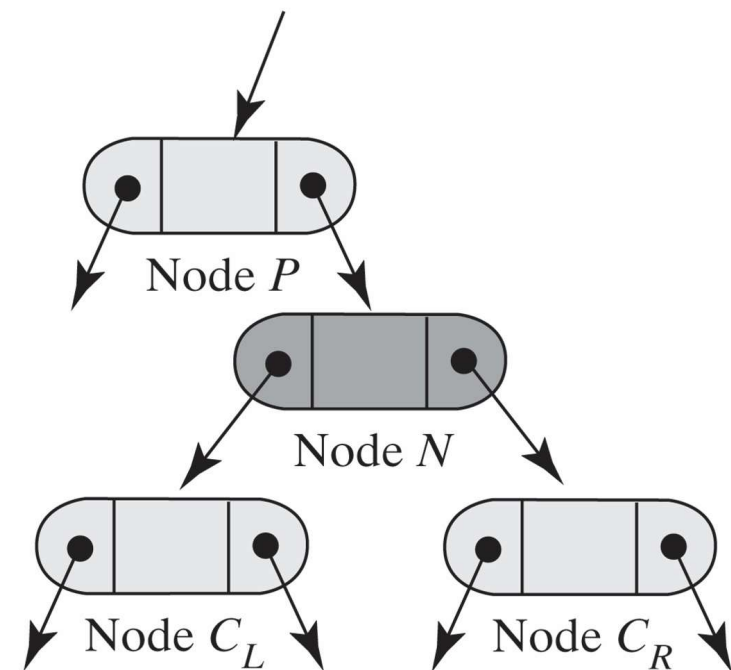
- FIGURE 26-8 Two possible configurations of a node N that has two children

(a) Node N is a left child



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(b) Node N is a right child



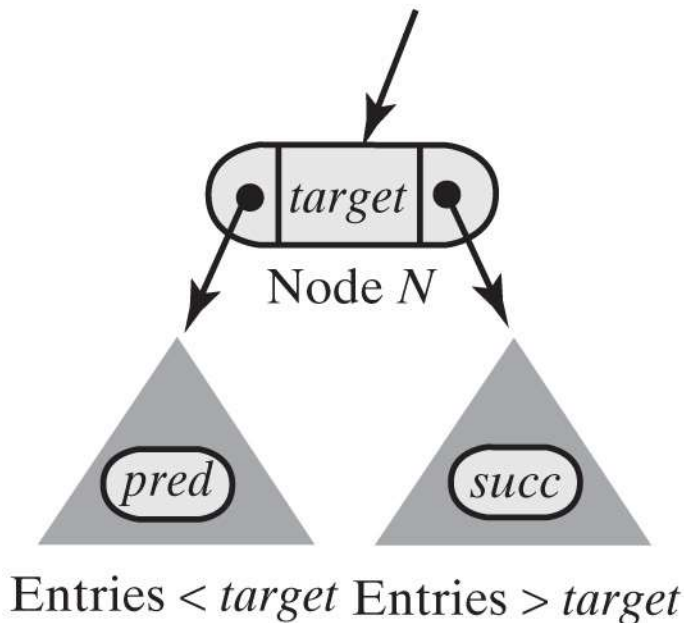
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Removing a Value

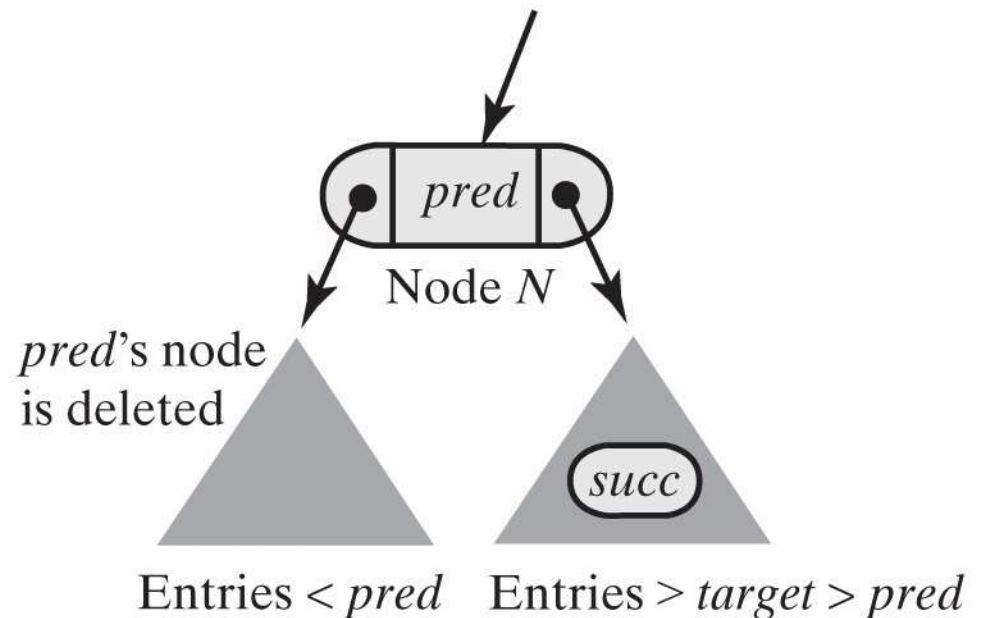
- FIGURE 26-9 Node N and its subtrees before and after removing target

(a) $pred$ is immediately before $target$,
 $succ$ is immediately after $target$



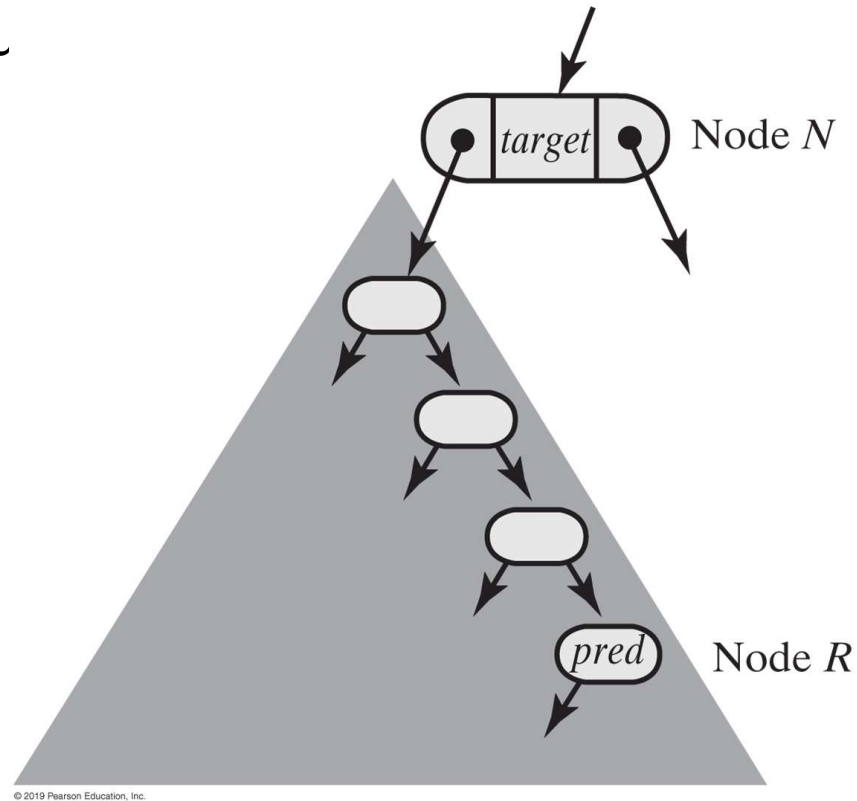
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(b) $pred$ replaces $target$,
effectively removing it



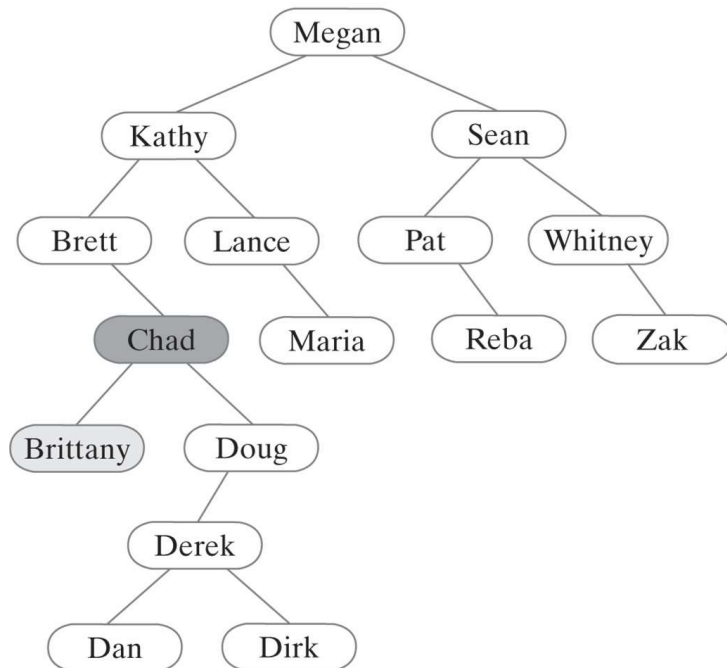
Removing a Value

- FIGURE 26-10 The largest entry *pred* in node *N*'s left subtree occurs in the *sl*



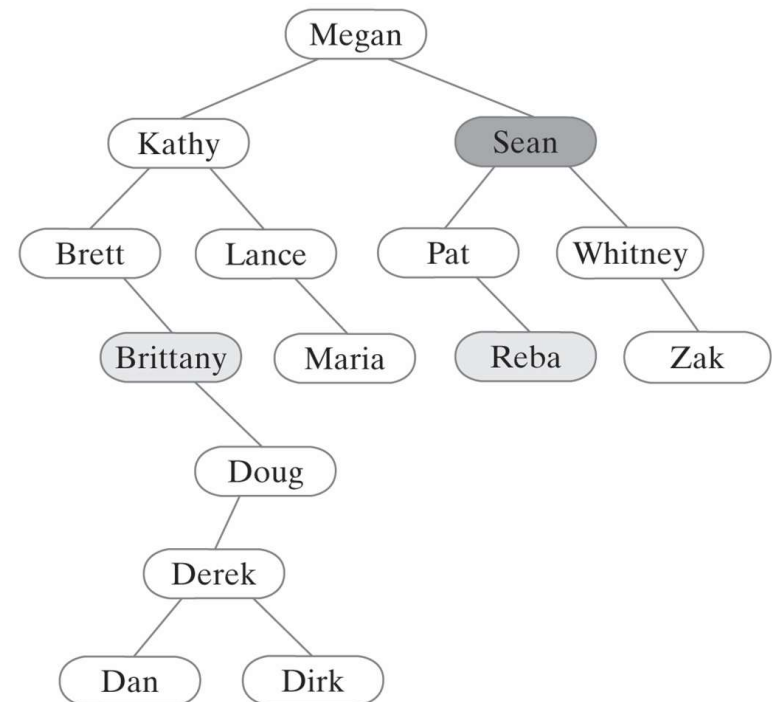
Successive removals from a binary search tree (Part 1)

(a) A binary search tree



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(b) The tree after removing *Chad*

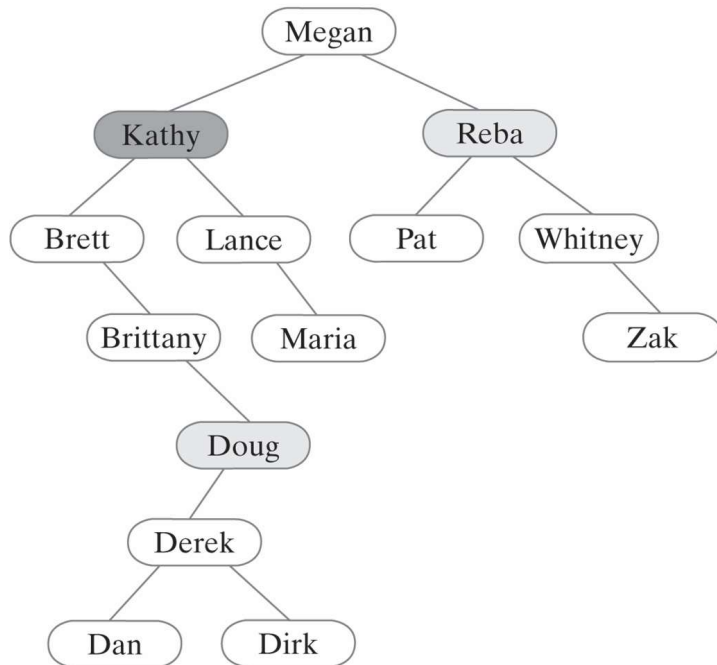


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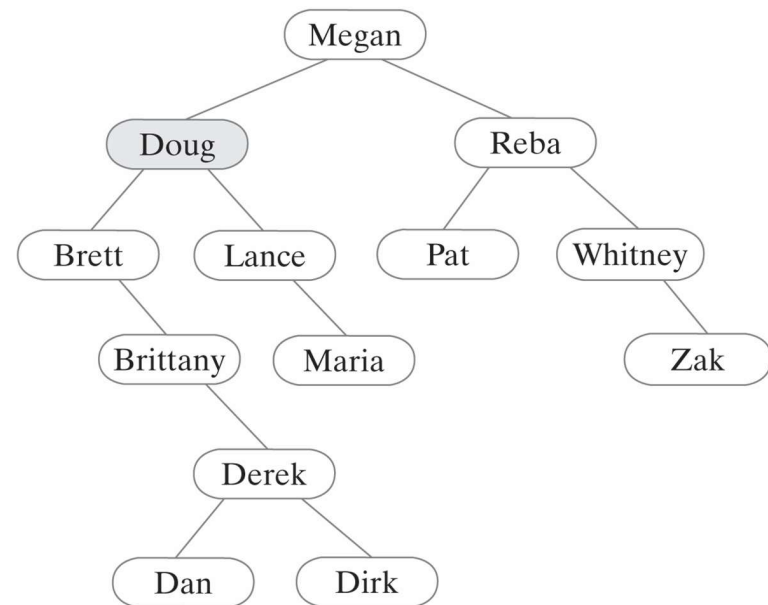
Successive removals from a binary search tree (Part 2)

(c) The tree after removing *Sean*



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(d) The tree after removing *Kathy*



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Removing a Value

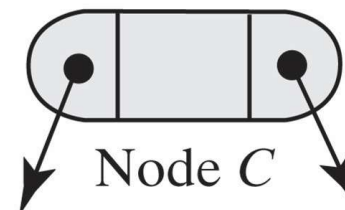
- FIGURE 26-12 Removing the root when it has one child

(a) Two possible configurations of a tree's root with one child



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(b) The tree after removing its root



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Recursive Implementation

- Recursive algorithm describes the method's logic at a high level

Algorithm remove(binarySearchTree, anEntry)

oldEntry = null

if (binarySearchTree *is not empty*)

{

 if (anEntry *matches the entry in the root of* binarySearchTree)

 {

 oldEntry = *entry in root*

 removeFromRoot(*root of* binarySearchTree)

 }

 else if (anEntry < *entry in root*)

 oldEntry = remove(*left subtree of* binarySearchTree, anEntry)

 else // anEntry > *entry in root*

 oldEntry = remove(*right subtree of* binarySearchTree, anEntry)

}

return oldEntry



Recursive Implementation

- The public method `remove`

```
public T remove(T anEntry)
{
    ReturnObject oldEntry = new ReturnObject(null);
    BinaryNode<T> newRoot = removeEntry(getRootNode(), anEntry, oldEntry);
    setRootNode(newRoot);

    return oldEntry.get();
} // end remove
```



Recursive Implementation

- The private method `removeEntry`

// Removes an entry from the tree rooted at a given node.

```
private BinaryNode<T> removeEntry(BinaryNode<T> rootNode, T anEntry,
                                   ReturnObject oldEntry)
{
    if (rootNode != null)
    {
        T rootData = rootNode.getData();
        int comparison = entry.compareTo(rootData);

        if (comparison == 0)    // anEntry == root entry
        {
            oldEntry.set(rootData);
            rootNode = removeFromRoot(rootNode);
        }
        else if (comparison < 0) // anEntry < root entry
        {
            BinaryNode<T> leftChild = rootNode.getLeftChild();
            BinaryNode<T> subtreeRoot = removeEntry(leftChild, anEntry, oldEntry);
            rootNode.setLeftChild(subtreeRoot);
        }
        else                    // anEntry > root entry
        {
            BinaryNode<T> rightChild = rootNode.getRightChild();
            // A different way of coding than for left child:
            rootNode.setRightChild(removeEntry(rightChild, anEntry, oldEntry));
        } // end if
    } // end if

    return rootNode;
} // end removeEntry
```



Recursive Implementation

- The algorithm `removeFromRoot`

Algorithm removeFromRoot(rootNode)

// Removes the entry in a given root node of a subtree.

if (rootNode has two children)

{

largestNode = node with the largest entry in the left subtree of rootNode

Replace the entry in rootNode with the entry in largestNode

Remove largestNode from the tree

}

else if (rootNode has a right child)

rootNode = rootNode's right child

else

rootNode = rootNode's left child // Possibly null

// Assertion: If rootNode was a leaf, it is now null

return rootNode



Recursive Implementation

- The private method `removeFromRoot`

// Removes the entry in a given root node of a subtree.

```
private BinaryNode<T> removeFromRoot(BinaryNode<T> rootNode)
```

```
{
```

// Case 1: rootNode has two children

```
if (rootNode.hasLeftChild() && rootNode.hasRightChild())
```

```
{
```

// Find node with largest entry in left subtree

```
BinaryNode<T> leftSubtreeRoot = rootNode.getLeftChild();
```

```
BinaryNode<T> largestNode = findLargest(leftSubtreeRoot);
```

// Replace entry in root

```
rootNode.setData(largestNode.getData());
```

// Remove node with largest entry in left subtree

```
rootNode.setLeftChild(removeLargest(leftSubtreeRoot));
```

```
} // end if
```

// Case 2: rootNode has at most one child

```
else if (rootNode.hasRightChild())
```

```
rootNode = rootNode.getRightChild();
```

```
else
```

```
rootNode = rootNode.getLeftChild();
```

// Assertion: If rootNode was a leaf, it is now null

```
return rootNode;
```

```
} // end removeEntry
```

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Recursive Implementation

- The private method `findLargest`

```
// Finds the node containing the largest entry in a given tree.  
// rootNode is the root node of the tree.  
// Returns the node containing the largest entry in the tree.  
private BinaryNode<T> findLargest(BinaryNode<T> rootNode)  
{  
    if (rootNode.hasRightChild())  
        rootNode = findLargest(rootNode.getRightChild());  
  
    return rootNode;  
} // end findLargest
```



Recursive Implementation

- The private method `removeLargest`

// Removes the node containing the largest entry in a given tree.

// rootNode is the root node of the tree.

// Returns the root node of the revised tree.

```
private BinaryNode<T> removeLargest(BinaryNode<T> rootNode)
{
    if (rootNode.hasRightChild())
    {
        BinaryNode<T> rightChild = rootNode.getRightChild();
        rightChild = removeLargest(rightChild);
        rootNode.setRightChild(rightChild);
    }
    else
        rootNode = rootNode.getLeftChild();

    return rootNode;
} // end removeLargest
```



Iterative Implementation

- Pseudocode that describes `remove`

Algorithm `remove(anEntry)`

result = null

currentNode = *node that contains a match for anEntry*

parentNode = **currentNode**'s parent

if (**currentNode** != null) *// That is, if entry is found*

{

result = **currentNode**'s data (*the anEntry to be removed from the tree*)

// Case 1

if (**currentNode** *has two children*)

 {

// Get node to remove and its parent

nodeToRemove = *node containing anEntry inorder predecessor; it has at most one child*

parentNode = **nodeToRemove**'s parent

Copy entry from nodeToRemove to currentNode

currentNode = **nodeToRemove**

// Assertion: currentNode is the node to be removed; it has at most one child

// Assertion: Case 1 has been transformed to Case 2

 }

// Case 2: currentNode has at most one child

Delete currentNode from the tree

}

return result

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Iterative `remove` Implementation (Part 1)

- The public method `remove`

```
public T remove(T entry)
{
    T result = null;

    // Locate node (and its parent) that contains a match for entry
    NodePair pair = findNode(entry);
    BinaryNode<T> currentNode = pair.getFirst();
    BinaryNode<T> parentNode = pair.getSecond();

    if (currentNode != null) // Entry is found
    {
        result = currentNode.getData(); // Get entry to be removed

        // Case 1: currentNode has two children
        if (currentNode.hasLeftChild() && currentNode.hasRightChild())
        {
            // Replace entry in currentNode with the entry in another node
            // that has at most one child; that node can be deleted

            // Get node to remove (contains inorder predecessor; has at
            // most one child) and its parent
            pair = getNodeToRemove(currentNode);
            BinaryNode<T> nodeToRemove = pair.getFirst();
            parentNode = pair.getSecond();
        }
    }
}
```



Iterative `remove` Implementation (Part 2)

- The public method `remove`

```
// Copy entry from nodeToRemove to currentNode
currentNode.setData(nodeToRemove.getData());

currentNode = nodeToRemove;
// Assertion: currentNode is the node to be removed; it has at
//           most one child
// Assertion: Case 1 has been transformed to Case 2
} // end if

// Case 2: currentNode has at most one child; delete it
removeNode(currentNode, parentNode);
} // end if

return result;
} // end remove
```



Iterative Implementation

- The private method `findNode`

```
private NodePair findNode(T entry)
{
    NodePair result = new NodePair();
    boolean found = false;

    // ...

    if (found)
        result = new NodePair(currentNode, parentNode);
    // Located entry is currentNode.getData()

    return result;
} // end findNode
```



Iterative Implementation

- Pseudocode for the private method `getNodeToRemove`

// Find the in-order predecessor by searching the left subtree; it will be the largest

// entry in the subtree, occurring in the node as far right as possible

leftSubtreeRoot = *left child of currentNode*

rightChild = **leftSubtreeRoot**

priorNode = **currentNode**

while (**rightChild** *has a right child*)

{

priorNode = **rightChild**

rightChild = *right child of rightChild*

}

// Assertion: rightChild is the node to be removed and has no more than one child



Iterative getNodeToRemove Implementation

- Implementation of the private method `getNodeToRemove`

```
private NodePair getNodeToRemove(BinaryNode<T> currentNode)
{
    // Find node with largest entry in left subtree by
    // moving as far right in the subtree as possible
    BinaryNode<T> leftSubtreeRoot = currentNode.getLeftChild();
    BinaryNode<T> rightChild = leftSubtreeRoot;
    BinaryNode<T> priorNode = currentNode;

    while (rightChild.hasRightChild())
    {
        priorNode = rightChild;
        rightChild = rightChild.getRightChild();
    } // end while

    // rightChild contains the inorder predecessor and is the node to
    // remove; priorNode is its parent

    return new NodePair(rightChild, priorNode);
} // end getNodeToRemove
```



Iterative `removeNode` Implementation

- The private method `removeNode`

```
private void removeNode(BinaryNode<T> nodeToRemove,  
                        BinaryNode<T> parentNode)  
{  
    BinaryNode<T> childNode;  
  
    if (nodeToRemove.hasLeftChild())  
        childNode = nodeToRemove.getLeftChild();  
    else  
        childNode = nodeToRemove.getRightChild();  
  
    // Assertion: if nodeToRemove is a leaf, childNode is null  
  
    if (nodeToRemove == getRootNode())  
        setRootNode(childNode);  
    else if (parentNode.getLeftChild() == nodeToRemove)  
        parentNode.setLeftChild(childNode);  
    else  
        parentNode.setRightChild(childNode);  
} // end removeNode
```



Efficiency of Operations

- For tree of height h
 - The operations add, remove, and **getEntry** are $O(h)$
- If tree of n nodes has height $h = n$
 - These operations are $O(n)$
- Shortest tree is full
 - Results in these operations being $O(\log n)$



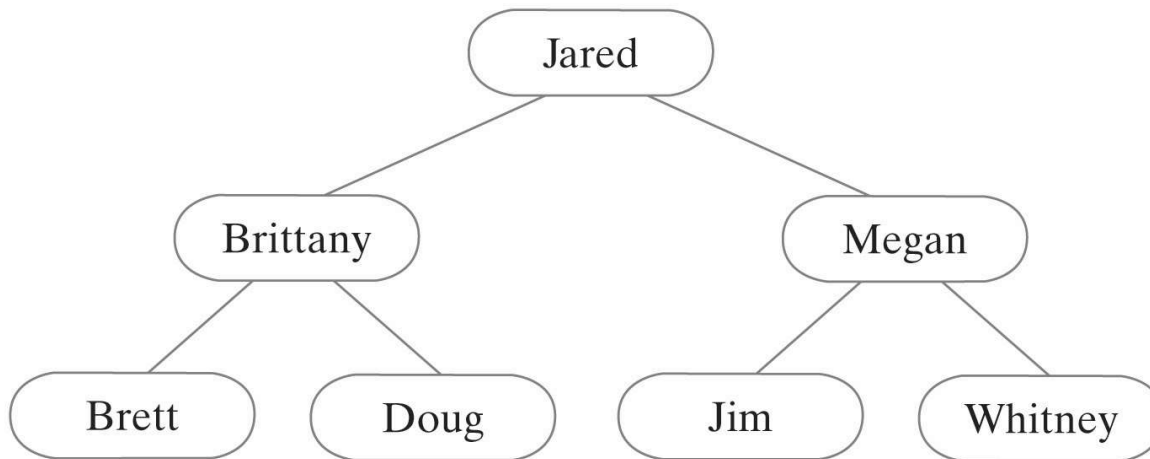
Efficiency of Operations

- FIGURE 26-13 Two binary search trees that contain the same data

(b) The tallest binary search tree having seven nodes

Operations are $O(n)$

(a) The shortest binary search tree having seven nodes



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Operations are $O(\log n)$

