# Lecture 21: Problem Solving with Recursion

**CS 0445: Data Structures** 

#### **Constantinos Costa**

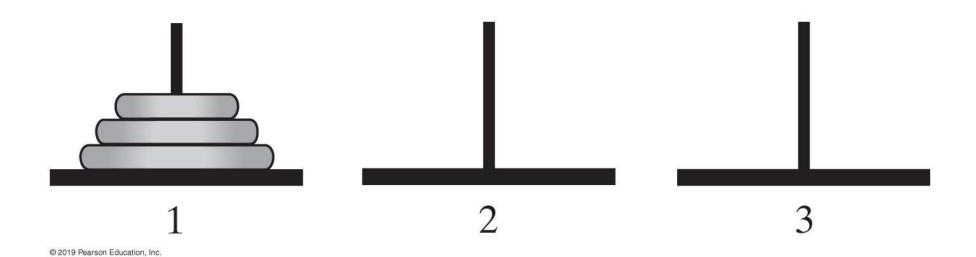
http://db.cs.pitt.edu/courses/cs0445/current.term/

Oct 23, 2019, 8:00-9:15 University of Pittsburgh, Pittsburgh, PA



# Simple Solution to a Difficult Problem

The initial configuration of the Towers of Hanoi for three disks





#### Simple Solution to a Difficult Problem

#### Rules:

- Move one disk at a time. Each disk moved must be the topmost disk.
- No disk may rest on top of a disk smaller than itself.
- You can store disks on the second (extra) pole temporarily, as long as you observe the previous two rules.

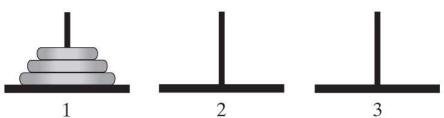


#### Simple Solution to a Difficult Problem (Part 1)

Sequence of moves for solving Towers of Hanoi problem with 3 disks

(a) The beginning configuration

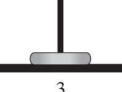
© 2019 Pearson Education, Inc.



(b) After moving a disk from pole 1 to pole 3

3





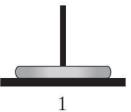
© 2019 Pearson Education, Inc.

(c) After moving a disk from pole 1 to pole 2

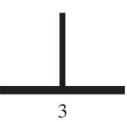
© 2019 Pearson Education, Inc.



(d) After moving a disk from pole 3 to pole 2







© 2019 Pearson Education, Inc.

#### Simple Solution to a Difficult Problem (Part 2)

Sequence of moves for solving Towers of Hanoi problem with 3 disks

(e) After moving a disk from pole 1 to pole 3

1 2

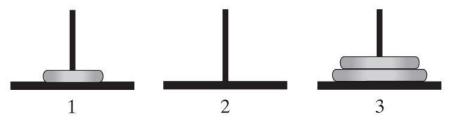
© 2019 Pearson Education, Inc.

(f) After moving a disk from pole 2 to pole 1

1 2 3

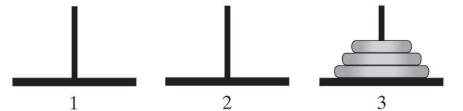
© 2019 Pearson Education, Inc.

(g) After moving a disk from pole 2 to pole 3



© 2019 Pearson Education, Inc.

© 2019 Pearson Education, Inc.

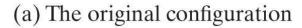


(h) After moving a disk from pole 1 to pole 3



#### A Smaller Problem

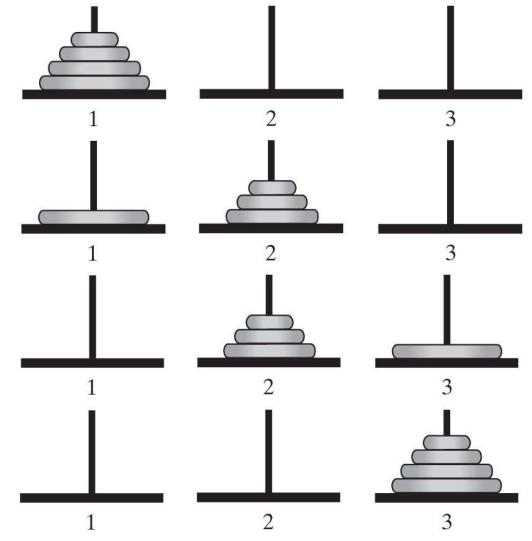
The smaller problems in a recursive solution for four disks



(b) After your friend moves three disks from pole 1 to pole 2

(c) After you move one disk from pole 1 to pole 3

(d) After your friend moves three disks from pole 2 to pole 3





#### Solutions

 Recursive algorithm to solve any number of disks.

Note: for n disks, solution will be  $2^n - 1$  moves

```
Algorithm to move numberOfDisks disks from startPole to endPole using tempPole as a spare according to the rules of the Towers of Hanoi problem if (numberOfDisks == 1)

Move disk from startPole to endPole else {

Move all but the bottom disk from startPole to tempPole

Move disk from startPole to endPole

Move all disks from tempPole to endPole
```



#### Solution

```
// Java recursive program to solve tower of hanoi puzzle
class HanoiPuzzle {
    static String toColor(int i) {
        switch (i) {...}
   // Java recursive function to solve tower of hanoi puzzle
  static void towerOfHanoi(int n, char startPole, char endPole, char tempPole) {
   if (n == 1) {
      System.out.println("Move disk (1) from pole " + startPole + " to pole " + endPole);
      return;
   towerOfHanoi(n - 1, startPole, tempPole, endPole);
   System.out.println("Move disk (" + n + ") from pole " + startPole + " to pole " + endPole);
   towerOfHanoi(n - 1, tempPole, endPole, startPole);
   // Driver method
   public static void main(String args[]) {
        int n = 4; // Number of disks
        towerOfHanoi(n, 'A', 'C', 'B'); // A, B and C are names of poles
```



# Poor Solution to a Simple Problem

- Algorithm to generate Fibonacci numbers.
- Why is this inefficient?

```
F_0 = 1

F_1 = 1

F_n = F_{n-1} + F_{n-2} when n \ge 2
```

```
Algorithm Fibonacci(n) if (n <= 1)
return 1
else
return Fibonacci(n - 1) + Fibonacci(n - 2)</pre>
```



# Poor Solution to a Simple Problem

The computation of the Fibonacci number F6

(a) Recursively  $F_2$  is computed 5 times  $F_3$  is computed 3 times  $F_4$  is computed once  $F_5$  is computed once  $F_6$  is computed once  $F_6$  is computed once  $F_6$  is  $F_7$   $F_8$   $F_9$   $F_9$ 

© 2019 Pearson Education, Inc.



## Poor Solution to a Simple Problem

The computation of the Fibonacci number F6

(a) Recursively  $F_2$  is computed 5 times  $F_3$  is computed 3 times  $F_4$  is computed 2 times  $F_5$  is computed once  $F_6$  is computed once

(b) Iteratively

$$F_0 = 1$$

$$F_1 = 1$$

$$F_2 = F_1 + F_0 = 2$$

$$F_3 = F_2 + F_1 = 3$$

$$F_4 = F_3 + F_2 = 5$$

$$F_5 = F_4 + F_3 = 8$$

$$F_6 = F_5 + F_4 = 13$$



#### Indirect Recursion

- Example
  - Method A calls Method B
  - Method B calls Method C
  - Method C calls Method A
- Difficult to understand and trace
  - But does happen occasionally



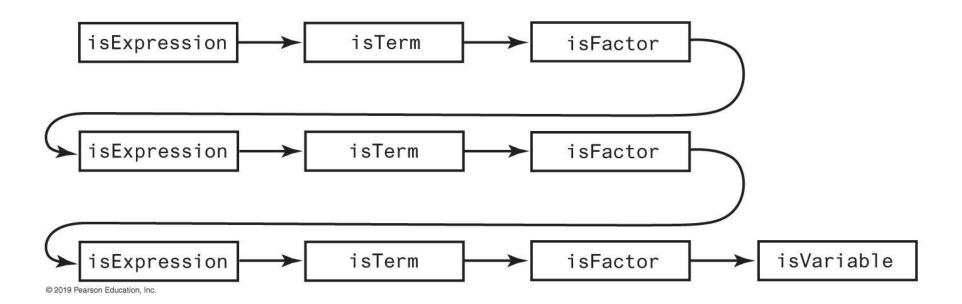
#### Indirect Recursion

- Consider evaluation of validity of an algebraic expression
  - Algebraic expression is either a term or two terms separated by a + or – operator
  - Term is either a factor or two factors separated by a \* or / operator
  - Factor is either a variable or an algebraic expression enclosed in parentheses
  - Variable is a single letter



#### Indirect Recursion

An example of indirect recursion

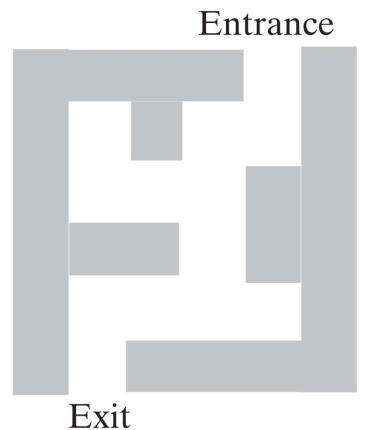


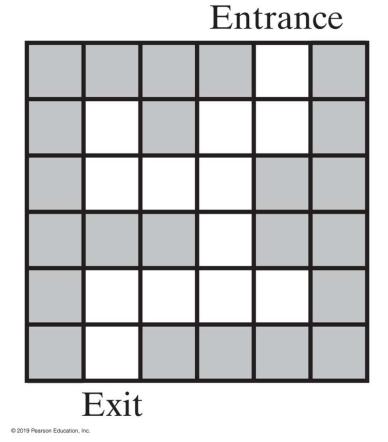


# Backtracking

· A two-dimensional maze with one entrance and one exit

(a) (b)

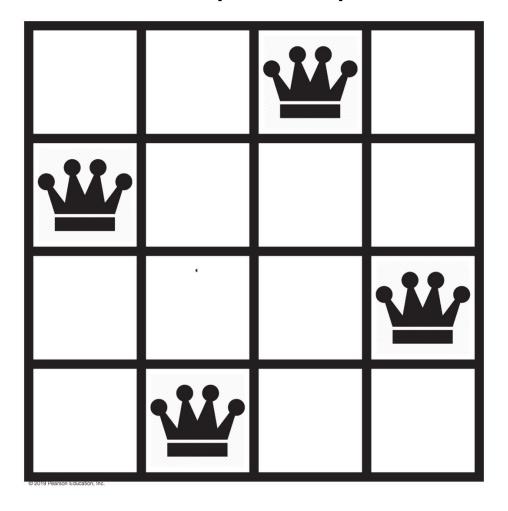






# Backtracking

A solution to the four-queens problem



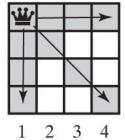


## Backtracking - Queens Solution (Part 1)

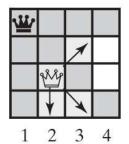
 Solving the four-queens problem by placing one queen at a time in each column

= Can be attacked by existing queens = Can be attacked by the newly placed queen = Rejected during backtracking

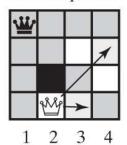
(a) The first queen in column 1.



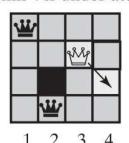
(b) The second queen in column 2. All of column 3 is under attack.



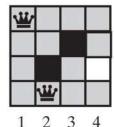
(c) Backtrack to column 2 and try another square for the queen.



(d) The third queen in column 3. All of column 4 is under attack.



(e) Backtrack to column 3, but the queen has no other move.







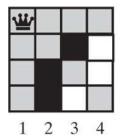
## Backtracking - Queens Solution (Part 2)

 Solving the four-queens problem by placing one queen at a time in each column

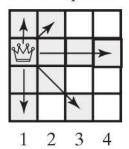
= Can be attacked by existing queens = Can be attacked by the newly placed queen = Rejected during backtracking

(f) Backtrack to column 2, but the queen has

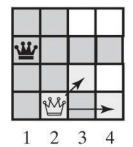
no other move.



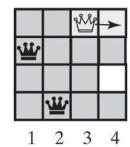
(g) Backtrack to column 1 and try another square for the queen.



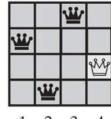
(h) The second queen in column 2.



(i) The third queen in column 3.



(j) The fourth queen in column 4. Solution!





# Happy Halloween!

