

# Notes for Week 8/1/19 - 8/8/19

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## 1 Bound on Eigenvalue Tail Decay of Neighborhood Graph

Let  $\mathcal{D} = [0, 1]^d$ . We consider two graphs over data on  $\mathcal{D}$ .

Suppose we observe the random design  $X = x_1, \dots, x_n$  independently sampled from probability measure  $P$  supported on  $\mathcal{D}$ . For any  $r \geq 0$ , let the kernel function  $\eta_r : \mathcal{D} \times \mathcal{D} \rightarrow \mathbb{R}_{\geq 0}$  be given by  $\eta_r(x, y) = \mathbf{1}(\|x - y\| \leq r)$ . Then, let the  $r$ -neighborhood graph over  $x_1, \dots, x_n$  be the undirected, unweighted graph  $G = (V, E)$ , where  $V = [n]$  and for  $i, j \in [n]$ ,  $(i, j) \in E$  if  $\eta_r(x_i, x_j) = 1$ .

Let  $B$  be the incidence matrix associated with  $G$ , and let  $L = B^T B$  be the corresponding Laplacian. Write  $L = V \Lambda V^T$  for the eigen decomposition of  $L$ , where  $V = (v_1 \dots v_n)$  is an orthonormal matrix with the eigenvectors of  $L$  as its columns, and  $\Lambda$  is a diagonal matrix with entries  $\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$ . For a given  $s > 0$  and  $C > 0$ , let  $N(C) = \#\{k : \lambda_k^s \leq C^2\}$ . We wish to prove the following result.

**Theorem 1.** *Let  $r \rightarrow 0$  as  $n \rightarrow \infty$  sufficiently slowly so that  $r(\log n/n)^{1/d} \rightarrow \infty$ . Then, for each  $s > 0$  and  $C \leq \sqrt{2}$ ,*

$$N(C) \leq nC^{d/s}. \quad (1)$$

*with probability tending to one as  $n \rightarrow \infty$ .*

## 2 Theory

To prove Theorem 1, we will make use of another graph on points within  $\mathcal{D}$ , whose spectral properties are very well-understood. Let  $\xi$  be the set of evenly spaced lattice points over  $\mathcal{D}$ ; formally  $\xi = \{k/n : k \in [\ell]^d\}$  where  $\ell = n^{1/d}$ , and we define  $[\ell]^d = \{(\ell_1, \dots, \ell_d) : \ell_k \in [\ell] \text{ for each } \ell_k\}$ . Then, let the grid graph over  $\xi$  be given by  $\tilde{G} = (\tilde{V}, \tilde{E})$ , where  $\tilde{V} = \xi$  and  $(\xi_k, \xi_{k'})$  is in  $\tilde{E}$  if  $\|\xi_k - \xi_{k'}\|_1 = 1/\ell$ .

Let  $\tilde{L}$  be the Laplacian of  $\tilde{G}$ , and let  $\tilde{\lambda}_1 \leq \dots \leq \tilde{\lambda}_n$  be the ordered eigenvalues of  $\tilde{L}$ . Write  $\tilde{N}(C) = \#\{k : \tilde{\lambda}_k^s \leq C^2\}$ . We have that the desired scaling rate (1) holds with respect to the grid graph  $\tilde{G}$ .

**Lemma 1.** *For each  $s$  and every  $C \leq \sqrt{2}$ , we have*

$$\tilde{N}(C) \leq 2^d (nC^{d/s} + 1) \quad (2)$$

*Proof.* It will be sufficient to show that

$$\lambda_k^s \leq C^2 \Rightarrow \left\lfloor k^{1/d} \right\rfloor^d \leq C^{d/s} n \quad (3)$$

To show this, observe that for any  $\tau \in \mathbb{N}$  and  $k = \tau^d$ , we have

$$\lambda_k \geq 4 \sin^2 \left( \frac{\pi k^{1/d}}{2n^{1/d}} \right) \geq \frac{\pi^2 k^{2/d}}{4n^{2/d}} \wedge 2$$

Therefore, if  $\lambda_k^s \leq C^2$  and  $C \leq \sqrt{2}$ , this implies

$$\frac{\pi^{2s} k^{2s/d}}{4^s n^{2s/d}} \leq C^2$$

and rearranging, we obtain

$$k \leq \frac{C^{d/s} 2^d n}{\pi^d} \leq C^{d/s} n$$

If  $k^{1/d}$  is not a natural number, applying the same argument to  $k' = \lfloor k^{1/d} \rfloor^d$  yields (3).  $\square$

In light of Lemma 1, to prove Theorem 1 it is sufficient to show that  $\tilde{L} \preceq L$ , since by the Courant-Fischer min-max theorem, the ordering  $\tilde{L} \preceq L$  implies that  $\tilde{\lambda}_k \leq \lambda_k$  for all  $k \in [n]$ . The next result details the conditions under which this ordering holds. This condition will be stated with respect to the min-max matching distance between  $\xi$  and  $X$ , i.e. the minimum over all bijections  $T : \xi \rightarrow X$  such that

$$\min_T \max_{i \in [n]} |T^{-1}(x_i) - x_i|$$

**Lemma 2.** *For any radius  $r$  satisfying*

$$r \geq 2 \min_T \max_{i \in [n]} |T^{-1}(x_i) - x_i| + n^{-1/d} \quad (4)$$

*we have that  $\tilde{L} \preceq L$ .*

*Proof.* Let  $T_\star$  achieve the min-max matching distance, i.e

$$\max_{i \in [n]} |T_\star^{-1}(x_i) - x_i| = \min_T \max_{i \in [n]} |T^{-1}(x_i) - x_i|.$$

It will be sufficient to prove that for every pair  $(T_\star^{-1}(x_i), T_\star^{-1}(x_j)) \in \tilde{E}$ , the corresponding edge  $(i, j)$  is in  $E$ . To see this, let  $A$  denote the adjacency matrix associated with the neighborhood graph  $G$ , and  $\tilde{A}$  the adjacency matrix associated with the grid  $\tilde{G}$ . Precisely  $A$  is the  $n \times n$  matrix with entries  $A_{ij} = \eta_r(x_i, x_j)$ , and  $\tilde{A}$  is the  $n \times n$  matrix with entries  $\tilde{A}_{ij} = \mathbf{1}\{T_\star^{-1}(x_i), T_\star^{-1}(x_j) \in \tilde{E}\}$ . Our goal is to show that, for every  $z \in \mathbb{R}^d$ , we have

$$z^T L z = \frac{1}{2} \sum_{i,j=1}^n (z_i - z_j)^2 A_{ij} \leq \frac{1}{2} \sum_{i,j=1}^n (z_i - z_j)^2 \tilde{A}_{ij} = z^T \tilde{L} z.$$

which certainly holds if  $\tilde{A}_{ij} \leq A_{ij}$  for all  $i, j \in [n]$ .

Now, assume  $(T_\star^{-1}(x_i), T_\star^{-1}(x_j)) \in \tilde{E}$ . This implies

$$\begin{aligned} \|x_i - x_j\|_2 &\leq \|x_i - T_\star^{-1}(x_i)\|_2 + \|T_\star^{-1}(x_j) - T_\star^{-1}(x_i)\|_2 + \|x_j - T_\star^{-1}(x_j)\|_2 \\ &\leq 2 \max_{i \in [n]} |T_\star^{-1}(x_i) - x_i| + n^{-1/d} \\ &\leq r. \end{aligned}$$

so  $\eta_r(x_i, x_j) = 1$  and therefore  $(i, j) \in E$ .  $\square$

Finally, the following Lemma demonstrates that, for sufficiently large  $r$ , the condition (4) will hold with high probability.

**Lemma 3.** *Assume  $P$  has density  $p$  which is bounded above and below uniformly over  $\mathcal{D}$ ; that is, there exist constants  $p_{\min}$  and  $p_{\max}$  such that*

$$0 < p_{\min} < p(x) < p_{\max} < \infty, \quad \text{for all } x \in \mathcal{D}.$$

*Then, for any  $r = r_n$  such that  $r(\log(n)/n)^{1/d} \rightarrow \infty$ , we have that the event*

$$r < 2 \min_T \max_{i \in [n]} |T^{-1}(x_i) - x_i| + n^{-1/d}$$

*occurs with probability tending to 0 as  $n \rightarrow \infty$ .*

Together Lemmas 1, 2, and 3 imply Theorem 1.