

# 1D "Warm Up" - Approximation Theory

Aim: Illustrate key ideas about model accuracy in "simple" 1D context.

## Problem 1. Polynomial Interpolation.

Given samples from a continuous map,

$$(*) \quad y_i = G(x_i), \quad \text{for } i=0, 1, 2, \dots, n,$$

where  $G: [0, 1] \rightarrow \mathbb{R}$ , find the unique degree  $\leq n$  polynomial that reproduces the observed data, i.e., interpolates  $G$ .

## Algorithm 1. "Dictionary" Interpolation

1. Select basis  $\{e_0(x), \dots, e_n(x)\} \subseteq P_n$

2. Solve

$$\begin{bmatrix} e_0(x_0) & \dots & e_n(x_0) \\ \vdots & \ddots & \vdots \\ e_0(x_n) & \dots & e_n(x_n) \end{bmatrix} \begin{bmatrix} c_0 \\ \vdots \\ c_n \end{bmatrix} = \begin{bmatrix} y_0 \\ \vdots \\ y_n \end{bmatrix}$$

3. Interpolant is  $F_n(x) = c_0 e_0(x) + \dots + c_n e_n(x)$

Q: How accurate is the interpolant?

"Model  
Error"

$$E_n = \sup_{x \in [0,1]} |G(x) - F_n(x)|$$

## Analysis of Interpolation

To analyze the accuracy of the interpolant, we follow a strategy that is common in approximation theory.

### Step 1. Best Approximation in $P_n$ .

Q: How well is  $G$  approximated by polynomials in  $P_n$ ?

### Step 2. Suboptimality of $P_n$ -interpolant.

Q: How "far from best" is the interpolant?

## Best Approximation

We want to find  $p_* \in P_n$  that solves

$$(*) \quad p_* = \underset{p \in P_n}{\operatorname{argmin}} \|p - G\|,$$

where  $\|p\| = \sup_{x \in [0,1]} |p(x)|$  is the "sup" norm.

### Key Fact 1. Existence: Uniqueness

If  $G \in C[0,1]$  (continuous on  $[0,1]$ ), then  
there is a unique best approximation  
 $p_* \in P_n$  that satisfies  $(*)$ .

pf | Existence: the space  $P_n$  is a  
complete  $n$ -dimensional vector space and the subset

$$\beta = \{p \in P_n \text{ s.t. } \|p - G\| \leq \|G\|\}$$

is closed! bounded, hence, compact.

The error function  $\rho \rightarrow \|p - G\|$  is continuous w.r.t.  $\rho \in \mathbb{B}$  and, therefore, achieves its minimum on  $\mathbb{B}$ .

Uniqueness: Chebyshev Equioscillation Thm.

Key Fact 2. Estimates for Smooth Functions

The error in the best polynomial approximation to  $G$  is tightly linked to the "smoothness" or "regularity" of  $G$ .

Jackson Thm. Let  $G \in C^k[0, 1]$  ( $k \geq 1$ ) and  $n \geq k-1$ . Then,

$$\inf_{P \in P_n} \|p - G\| \leq \left(\frac{\pi}{2}\right)^k \frac{1}{(n+1) \dots (n-k+2)} \|G^{(k)}\|_b.$$

Essentially, the error decreases algebraically

$$\Rightarrow \inf_{P \in P_n} \|P - G\| \leq C_K \frac{\|G^{(K)}\|}{n^K}$$

Pf | Integration-by-parts n-times  
+ a few "tricks." See attached reading.

## Relation to Interpolant

How far is interpolant from "best"?

Given distinct points  $x_0, \dots, x_N \in [0, 1]$ , every continuous function  $f: [0, 1] \rightarrow \mathbb{R}$  has a unique polynomial interpolant of deg  $\leq N$ :

$$P = P_x f.$$

The map  $P_x: C[0, 1] \rightarrow P_n$  is a linear projection of  $C[0, 1]$  onto  $P_n$  with

$$\|P_x\| = \sup_{f \in C[0, 1]} \frac{\|P_x f\|}{\|f\|} \geq 1.$$

This amplification factor, the operator norm of  $P_x$  determines how far the interpolants can be from "best."

Claim: Given  $X = \{x_0, \dots, x_n\} \subset [-1, 1]$ ,  
 $f \in C[-1, 1]$ , and  $p_x = \underset{p \in P_n}{\operatorname{argmin}} \|p - f\|$ ,

$$\|f - P_x f\| \leq (1 + \|P_x\|) \|f - p_x\|$$

interp. error

best fit error

$$\boxed{pf} \quad \|f - P_x f\| \leq \|f - p_x\| + \|p_x - P_x f\|$$

Now,  $P_x p_x = p_x$  because  $p_x \in P_n$ , so

$$\|p_x - P_x f\| = \|P_x(p_x - f)\| \leq \|P_x\| \|f - p_x\|$$

$$\Rightarrow \|f - P_x f\| \leq (1 + \|P_x\|) \|f - p_x\|.$$

For "good" sets of interpolation nodes

(slow growth)  $\|P_x\| \sim \log(\# \text{nodes})$  as  $n \rightarrow \infty$ .

For equally-spaced nodes,

(exponential growth)  $\|P_x\| \sim 2^{n+1} (\log n)^{-1}$  as  $n \rightarrow \infty$ .

Lemma The operator norm of  $P_x$  is called the Lebesgue constant and can be calculated:

$$\|P_x\| = \max_{x \in [0,1]} \sum_{k=0}^n \left| \prod_{\substack{i=0 \\ i \neq k}}^n \frac{x - x_i}{x_k - x_i} \right|.$$

(sketch)  $Pf$  / The Lagrange form of the interpolant

$$[P_x f](x) = \sum_{k=0}^n f(x_k) l_k(x), \text{ where}$$

"nodal polynomial"  $l_k(x) = \prod_{\substack{i=0 \\ i \neq k}}^n \frac{x - x_i}{x_k - x_i} = \begin{cases} 1 & x = x_k \\ 0 & x = x_j, j \neq k \end{cases}$

$$\Rightarrow \sup_{f \in C[0,1]} \frac{\|P_x f\|}{\|f\|} \leq \max_{x \in [0,1]} \sum_{k=0}^n |l_k(x)|$$

$\Leftarrow$  For  $f(x) = 1$ , " $\leq$ " is " $=$ "

## Model Accuracy for Bly. Interp.

Collecting the results of our 2-steps analysis of interpolation yields:

Then let  $G \in C^k[0,1]$  for some  $k \geq 1$  and let  $p_n \in P_n$  interpolate  $G$  at distinct points  $x_0, \dots, x_n \in [0,1]$ . For  $n \geq k-1$ ,

$$\|G - p_n\| \leq \|P_x\| \|G - p_x\| \leq C_k \|P_x\| \|G^{(k)}\| n^{-k}$$

where  $p_x = \underset{p \in P_n}{\operatorname{argmin}} \|G - p\|$

is the best degree  $\leq n$  approximation of  $G$ .

$\Rightarrow \|G - p_x\|$  describes how well the "truth" is approximated by a degree  $\leq n$  polynomial model.

$\Rightarrow$  This is controlled by the "smoothness" of  $G$  and depends both on  $G$  and our choice of a  $P_n$ -model.

$\Rightarrow \|P_x\|$  describes the influence  
of the data distribution on  $[0, 1]$   
on the accuracy of the approximation.

## Regression: Oversampling

In data-driven settings, we may not have control over the sampling points, which can severely limit the accuracy of our model because  $\|P_x\|$  often grows rapidly as  $n$  increases.

Q! What can we do if we are stuck with "bad" sample points?

Idea! Try using a lower degree polynomial model and find the "best fit" to the available data.

Interpolation  $\Rightarrow$  Regression