Austin Keelin Yilmaz COMP 3270 4/5/2018

The code for this assignment was written and compiled using the jGRASP IDE, version 2.04\_04

I certify that I wrote the code I am submitting. I did not copy whole or parts of it from another student or have another person write the code for me. Any code I am reusing in my program is clearly marked as such with its source clearly identified in comment.

## Algorithm-1

Step	Cost of each execution	Total # of times executed
1	1	1
2	1	n+1
3	1	$\sum i=1$ to n (i+1)
4	1	$\sum_{i=1}^{n} i=1 \text{ to n } (i)$
5	1	$\sum i=1$ to $n \sum j=1$ to $i (j+1)$
6	6	$\sum i=1$ to $n \sum j=1$ to $i(j)$
7	7	$\sum_{i=1}^{\infty} i=1 \text{ to n } (i)$
8	2	1

Multiply col.1 with col.2, add across rows and simplify

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T(n) = 2 + n + n(n+1)/2 + n + n(n+1)/2 + (n^3 + 6n^2 + 5n)/6 + (n^3 + 3n^2 + 2n) + 7(n(n+1)/2) + 2
= (n^3/6 + n^3) + (n^2/2 + n^2/2 + n^2 + 3n^2 + 7n^2/2) + (n + n/2 + n + n/2 + 5n/6 + 2n + 7n/2) + (2 + 2)
= 7n^3/6 + 17n^2/2 + 28n/3 + 4
= O(n^3)
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## Algorithm-2

Step	Cost of each execution	Total # of times executed
1	1	1
2	1	n+1
3	1	n
4	1	$\sum i=1$ to n (i+1)
5	6	$\sum$ i=1 to n (i)
6	7	$\sum$ i=1 to n (i)
7	2	1

Multiply col.1 with col.2, add across rows and simplify

$$T(n) = 1 + (n+1) + n + \sum_{i=1}^{n} to \ n \ (i+1) + 6\sum_{i=1}^{n} to \ n \ (i) + 7\sum_{i=1}^{n} to \ n \ (i) + 2$$

$$= 3 + 2n + 14\sum_{i=1}^{n} to \ n \ (i) + \sum_{i=1}^{n} to \ n \ (i)$$

$$= 3 + 3n + 7(n(n+1)) = 2 + 10n + 7n^{2}$$

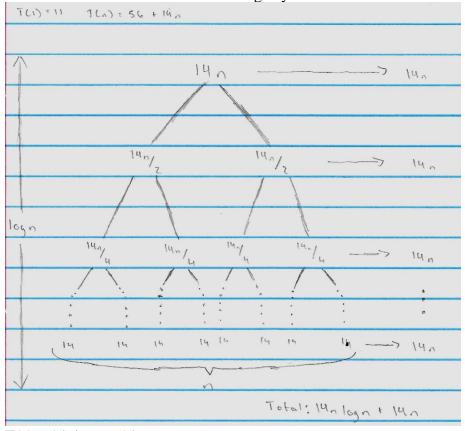
$$= O(n^{2})$$

## Algorithm-3

Step	Cost of each execution	Total # of times executed in any single
		recursive call

1	4	1			
2	11	1			
Steps exe	Steps executed when the input is a base case: 1 or 2				
First recurrence relation: T(n=1 or n=0) = 4 for n=0; 11 for n=1					
3	5	1			
4	2	1			
5	1	(n/2) + 1			
6	6	n / 2			
7	7	n / 2			
8	2	1			
9	1	(n/2) + 1			
10	6	n / 2			
11	7	n / 2			
12	4	1			
13	4	(cost excluding the recursive call) 1			
14	5	(cost excluding the recursive call) 1			
15	17	1			
Steps executed when input is NOT a base case:1 – 15					
Second recurrence relation: $T(n>1) = 56 + 14n$					
Simplified second recurrence relation (ignore the constant term): $T(n>1) = 14n$					

Solve the two recurrence relations using any method:



T(n) = 14nlogn + 14n= O(n log n)

Algorithm-4

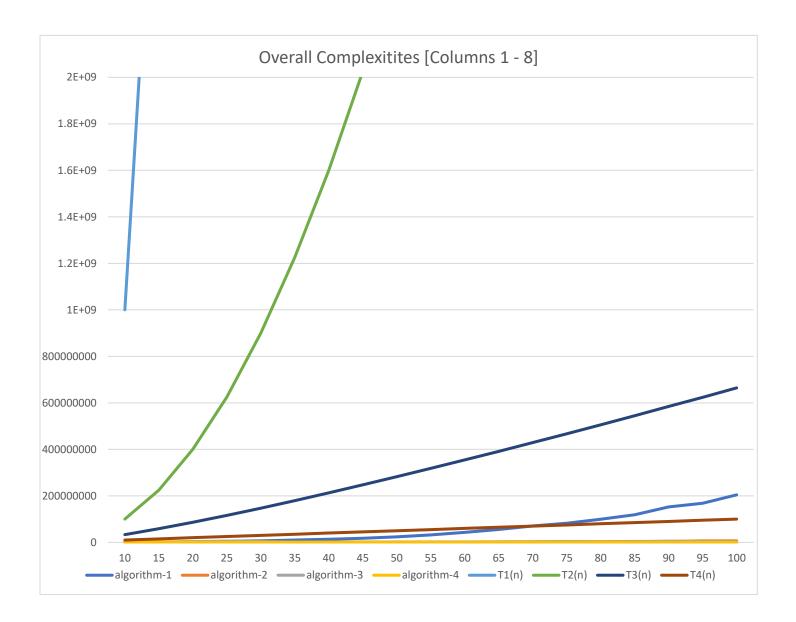
Step	Cost of each execution	Total # of times executed
1	1	1
2	1	1
3	1	n + 1
4	10	n
5	7	n
6	2	1

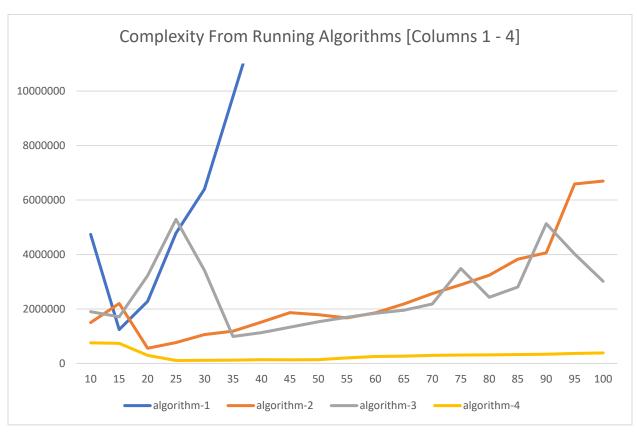
Multiply col.1 with col.2, add across rows and simplify T(n) = 1 + 1 + n + 1 + 10n + 7n + 2 = 18n + 5

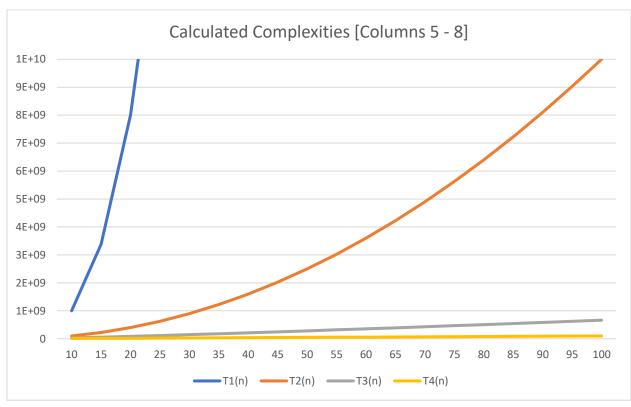
$$T(n) = 1 + 1 + n + 1 + 10n + 7n + 2$$

$$= 18n + 5$$

$$= O(n)$$







The actual time taken matched the predicted complexity for each algorithm as the input size increased. In the early runs with the smaller input sizes, where the algorithms would reach large complexity early and then dip, the data appears to suggest that the predictions were incorrect but as the algorithms continued running and reached larger input sizes, the data for the complexity of these runs each match with the respective calculated complexities.