Exam 4 Report

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${\bf Abstract}$

This report answers the three questions in the CIS 634 Exam 4. There are theories proved for questions 1 and 2 and an opinioned response for policy spaces in question 3.

Acknowledgments: I received no assistance with this exercise.			

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1 Executive Summary

All requirements for this assignment have been satisfied. The first question is answered in the code in the HOL/exam4Script.sml source file. There are four derived inference rules for all the access control scheme of the large distributed network The simple opener state machine theory described in problem 2 can be found in simpleOpener.Script. Question 3 involves my opinion on policy spaces. A formulated proof is given for the derived inference rule for the accessc control scheme described in problem 1. In addition, responses are given for the encryption scheme for question 2 and the question on theroem provers and aritificial intelligence of question 3.

2 Question 1

2.1 Problem Statement

The following problem involves an access control scheme controlling authorization and access to a non-local filesystem on a very large distributed network. The proof below is for a user Ursala requesting to have a file printed. Access control to the filesystem on the network is controlled by three servers AS(Authentication Server), TGS(Ticket Granting Server), and the FS (File Server).

- 1. Request a ticket bearing ticket from the AS, which the AS will respond with the Ticket Bearing Ticket and encrypted session key
- 2. The user's client program will receive the session key and ticket, and make a request to get a ticket for the FS server using these items.
- 3. Ursala will send a message, encrypted with the TGS session key, to the TGS server as well as the Ticket Bearing Ticket, and then receive a service granting ticket to be sent to the FS server with the new session key to be used with FS as well..
- 4. Ursala will send a new message to the FS encrypted with the latest session key that requests access to print the file along with the service granting ticket from before. If Ursala is authorized to print the file, the FS will honor the request.

HOL datatypes for probelm one below.

```
access = print | read | write
keys = KTGS | KFS | KU | KUTGS | KUFS
people = Ursala
princs = PR servers | User people | Key keys
props = USE serviceKey | PK people keys | AC access
servers = AS | TGS | FS
serviceKey = SERV servers | K keys
```

2.2 Relevant Code

Part A. Inference rule for the AS authenticating Ursala's request to use the TGS service. The ticket granting ticket is represented here by the Name (Key KTGS) says prop (PK Ursala KUTGS)) formula. The ticket is the (PK Ursala KUTGS) proposition, and the (Key KTGS) says prefix means that it is encrypted with the TGS key, which is only available to the AS and TGS servers.

```
\vdash (M,Oi,Os) sat Name (User Ursala) says prop (USE (SERV TGS)) \Rightarrow (M,Oi,Os) sat prop (USE (SERV TGS)) impf Name (Key KU) says prop (USE (K KUTGS)) andf Name (Key KTGS) says prop (PK Ursala KUTGS) \Rightarrow (M,Oi,Os) sat Name (User Ursala) controls prop (USE (SERV TGS)) \Rightarrow (M,Oi,Os) sat Name (Key KU) says prop (USE (K KUTGS)) andf Name (Key KTGS) says prop (PK Ursala KUTGS)
```

Part B: The inference for part B deals with the network client on Ursala's machine making the decision to use the session key that it obtains from the message from AS that was encrypted with Ursala's and the AS's shared key. After deciding to use the session key (Key KUTGS), Ursala will send the encrypted request to use the FS server, symbolized by Name (Key KUTGS) says prop (USE (SERV FS)), and also the Ticket Bearing Ticket encrypted with (Key KTGS) from before.

```
\vdash (M, Oi, Os) sat Name (Key KU) speaks_for Name (PR AS) \Rightarrow
   (M, Oi, Os) sat
   prop (USE (K KUTGS)) impf
   Name (Key KUTGS) says prop (USE (SERV FS)) \Rightarrow
   (M,Oi,Os) sat Name (PR AS) controls prop (USE (K KUTGS)) \Rightarrow
   (M, Oi, Os) sat
   Name (Key KU) says prop (USE (K KUTGS)) andf
   Name (Key KTGS) says prop (PK Ursala KUTGS) \Rightarrow
   (M,Oi,Os) sat
   Name (Key KUTGS) says prop (USE (SERV FS)) andf
   Name (Key KTGS) says prop (PK Ursala KUTGS)
 val formb1 = ''(Name (Key KU) says prop (USE (K KUTGS)) andf
                 Name (Key KTGS) says prop (PK Ursala KUTGS)):(props, princs, 'd, 'e)
val formb2 = ''Name (PR AS) controls prop (USE (K KUTGS)): (props, princs, 'd, 'e) For
val formb3 = ''(Name (Key KU) speaks_for Name (PR AS)):(props, princs, 'd, 'e)Form'
val formb4 = ''prop (USE (K KUTGS)) impf Name (Key KUTGS) says prop (USE (SERV F
val [formb1_thm, formb2_thm, formb3_thm, formb4_thm] = map ACL_ASSUM[formb1, formb2,
val formb5_thm = ACL_SIMP1 formb1_thm
val formb6\_thm = ACL\_SIMP2 formb1\_thm
val formb7_thm = SPEAKS_FOR formb3_thm formb5_thm
val formb8_thm = CONTROLS formb2_thm formb7_thm
val formb9\_thm = ACLMP formb8\_thm formb4\_thm
val session_key_receipt_thm = ACL_CONJ formb9_thm formb6_thm
```

Part C: The next inference rule proved is the decision to grant the request for services ticket, symbolized by formula Name (Key KFS) says ((Name (Key KUFS)) speaks_for (Name (User Ursala))). This message is encyrpted with the key only known by the TGS and FS servers, and it tells the FS server that the session key KUFS belongs to the user Ursala. The TGS wil also generate a session key KUFS, which is encrypted with KUTGS then sent.

The decision to grant these items is dependent on the TGS able to decrypt the ticket (PK Ursala KUTGS) which says that the key KUTGS is valid. Once that key is decided to be valid, the request USE SERV FS which was also sent can be deemed valid since it was encrypted with KUTGS, which should only be known to Ursala and AS.

Note, the formula (Name (Key KUFS) says prop (AC read) is derived from a valid ticket because the User Ursala will have access to the session key KUFS generated and can make this request to the final server.

```
\vdash (M, Oi, Os) sat Name (Key KTGS) speaks_for Name (PR AS) \Rightarrow
   (M,Oi,Os) sat
   prop (PK Ursala KUTGS) impf
   Name (Key KUFS) says prop (AC read) \Rightarrow
   (M, Oi, Os) sat
   prop (PK Ursala KUTGS) impf
   Name (Key KFS) says
   Name (Key KUFS) speaks_for Name (User Ursala) \Rightarrow
   (M,Oi,Os) sat Name (PR AS) controls prop (PK Ursala KUTGS) \Rightarrow
   (M,Oi,Os) sat
   Name (Key KUTGS) says prop (USE (SERV FS)) andf
   Name (Key KTGS) says prop (PK Ursala KUTGS) \Rightarrow
   (M,Oi,Os) sat
   Name (Key KFS) says
   Name (Key KUFS) speaks_for Name (User Ursala) andf
   Name (Key KUFS) says prop (AC read)
val formc1 = 'Name (Key KUTGS) says prop (USE (SERV FS)) andf
                 Name (Key KTGS) says prop (PK Ursala KUTGS): (props, princs, 'd, 'e) Form'
val formc2 = ''(Name (PR AS)) controls (prop (PK Ursala KUTGS)):(props, princs, 'd, 'e)F
val formc3 = ''(Name (Key KTGS) speaks_for (Name (PR AS))):(props, princs, 'd, 'e)Form''
val formc4 = ''prop (PK Ursala KUTGS) impf (Name (Key KFS) says ((Name (Key KUFS)) sp
val formc5 = ''prop (PK Ursala KUTGS) impf (Name (Key KUFS) says prop (AC read)):(pro
val [formc1_thm, formc2_thm, formc3_thm, formc4_thm, formc5_thm] = map ACL_ASSUM[formc1, formc1, formc1, formc1, formc2_thm]
val formc6_thm = ACL_SIMP1 formc1_thm
val formc7_thm = ACL_SIMP2 formc1_thm
val formc8_thm = SPEAKS_FOR formc3_thm formc7_thm
val formc9_thm = CONTROLS formc2_thm formc8_thm
val formc10_thm = ACL_MP formc9_thm formc4_thm
```

val request_for_services_thm = ACL_CONJ formc10_thm formc11_thm

val formc11_thm = ACLMP formc9_thm formc5_thm

Part D: The final inference rule proof involes the FS server granting Ursala's request to print. The FS authenticated the request using the request ticket sent to it by Ursala. It encovers the ticket by decrypting with the shared key KFS, then knows that the session key KUFS is valid to represent Ursala. Since the message requesting printing was encyrpted with this key KUFS, it knows Alice must have made the request.

```
\vdash (M, Oi, Os) sat Name (Key KFS) speaks_for Name (PR TGS) \Rightarrow
   (M,Oi,Os) sat Name (User Ursala) controls prop (AC print) \Rightarrow
   (M, Oi, Os) sat
   Name (PR TGS) controls
   Name (Key KUFS) speaks_for Name (User Ursala) \Rightarrow
   (M, Oi, Os) sat
   Name (Key KFS) says
   Name (Key KUFS) speaks_for Name (User Ursala) andf
   Name (Key KUFS) says prop (AC print) \Rightarrow
   (M, Oi, Os) sat prop (AC print)
val formd1 = ''(Name (Key KFS) says (Name (Key KUFS) speaks_for Name (User Ursals
                              (Name (Key KUFS) says prop (AC print)):(props, princs,
val formd2 = ''(Name (Key KFS) speaks_for Name (PR TGS)):(props, princs, 'd, 'e)Form
val formd3 = ''(Name (PR TGS) controls (Name (Key KUFS) speaks_for Name (User Ur
val formd4 = ''(Name (User Ursala) controls prop(AC print)):(props, princs, 'd, 'e)
val [formd1_thm, formd2_thm, formd3_thm, formd4_thm] = map ACL_ASSUM[formd1, formd2,
val formd5_thm = ACL_SIMP1 formd1_thm
val formd6_thm = ACL_SIMP2 formd1_thm
val formd7_thm = SPEAKS_FOR formd2_thm formd5_thm
val formd8_thm = CONTROLS formd3_thm formd7_thm
val formd9_thm = SPEAKS_FOR formd8_thm formd6_thm
val service_request_thm = CONTROLS formd4_thm formd9_thm
     Execution Transcript
      HOL-4 [Kananaskis 11 (stdknl, built Sat Aug 19 09:30:06 2017)]
      For introductory HOL help, type: help "hol";
       To exit type <Control>-D
[extending loadPath with Holmakefile INCLUDES variable]
>>>>>>> Loading acl_infRules
>>>>>> > < HOL message: Created theory "exam4">>
<<HOL message: Defined type: "servers">>
<<HOL message: Defined type: "people">>
<<HOL message: Defined type: "keys">>
<<HOL message: Defined type: "princs">>
```

<<HOL message: Defined type: "access">>

```
<<HOL message: Defined type: "serviceKey">>
<<HOL message: Defined type: "props">>
Theory: exam4
Parents:
    aclDrules
Type constants:
    access 0
   keys 0
    people 0
    princs 0
    props 0
    servers 0
    serviceKey 0
Term constants:
    AC
                      :access -> props
    AS
                      :servers
    FS
                      :servers
    K
                      :keys -> serviceKey
    KFS
                      :keys
    KTGS
                      :keys
    KU
                      :keys
    KUFS
                      :keys
    KUTGS
                      :keys
                      :keys -> princs
    Key
    PΚ
                      :people -> keys -> props
    PR
                      :servers -> princs
    SERV
                      :servers -> serviceKey
    TGS
                      :servers
    USE
                      :serviceKey -> props
    Ursala
                      :people
    User
                      :people -> princs
    access2num
                      :access -> num
    access_CASE
                      :access -> -> -> ->
    access_size
                      :access -> num
    keys2num
                      :keys -> num
    keys_CASE
                      :keys -> -> -> -> ->
    keys_size
                      :keys -> num
    num2access
                      :num -> access
    num2keys
                      :num -> keys
    num2people
                      :num -> people
    num2servers
                      :num -> servers
                      :people -> num
    people2num
    people_CASE
                      :people -> ->
                      :people -> num
    people_size
```

```
princs_CASE
                      :princs ->
                       (servers -> ) ->
                       (people -> ) -> (keys -> ) ->
                      :princs -> num
    princs_size
    print
                       :access
    props_CASE
                      :props ->
                       (serviceKey -> ) ->
                       (people -> keys -> ) -> (access -> ) ->
    props_size
                      :props -> num
    read
                      :access
    servers2num
                      :servers -> num
    servers_CASE
                      :servers -> -> -> ->
    servers_size
                      :servers -> num
    serviceKey_CASE
                      :serviceKey -> (servers -> ) -> (keys -> ) ->
    serviceKey_size
                      :serviceKey -> num
    write
                      :access
Definitions:
    @tempAS_def
      |-AS = num2servers 0
    @tempFS_def
      |- FS = num2servers 2
    @tempKFS_def
      |-KFS = num2keys 1
    @tempKTGS_def
      |-KTGS = num2keys 0
    @tempKUFS_def
      |-KUFS = num2keys 4
    @tempKUTGS_def
      |-KUTGS = num2keys 3
    @tempKU_def
      |-KU = num2keys 2
    @tempTGS_def
      |- TGS = num2servers 1
    @tempUrsala_def
      |- Ursala = num2people 0
    @tempprint_def
      |- print = num2access 0
    @tempread_def
      |- read = num2access 1
    @tempwrite_def
      |- write = num2access 2
    access_BIJ
      |-(a. num2access (access2num a) = a)
         r. (n. n < 3) r (access2num (num2access r) = r)
    access_CASE
      |- x v0 v1 v2.
```

```
(case x of print \Rightarrow v0 | read \Rightarrow v1 | write \Rightarrow v2) =
       (m. if m < 1 then v0 else if m = 1 then v1 else v2)
         (access2num x)
access_TY_DEF
  |- rep. TYPE_DEFINITION (n. n < 3) rep
access_size_def
  |-x. access_size x = 0
keys_BIJ
  |-(a. num2keys (keys2num a) = a)
     r. (n. n < 5) r (keys2num (num2keys r) = r)
keys_CASE
  |- x v0 v1 v2 v3 v4.
       (case x of
          KTGS => v0
        | KFS => v1
        | KU => v2
        | KUTGS => v3
        | KUFS => v4) =
       (m.
          if m < 2 then if m = 0 then v0 else v1
          else if m < 3 then v2
          else if m = 3 then v3
          else v4) (keys2num x)
keys_TY_DEF
  |- rep. TYPE_DEFINITION (n. n < 5) rep</pre>
keys_size_def
  |-x. \text{ keys\_size } x = 0
people_BIJ
  |-(a. num2people (people2num a) = a)
     r. (n. n < 1) r (people2num (num2people r) = r)
people_CASE
  |-x v0. (case x of Ursala => v0) = (m. v0) (people2num x)
people_TY_DEF
  |- rep. TYPE_DEFINITION (n. n < 1) rep</pre>
people_size_def
  |-x. people_size x = 0
princs_TY_DEF
  |- rep.
       TYPE_DEFINITION
         (a0.
             'princs'.
               (a0.
                  (a.
                     a0 =
                     (a.
                        ind_type$CONSTR 0 (a,ARB,ARB)
                           (n. ind_type$BOTTOM)) a)
```

```
(a.
                    a0 =
                     (a.
                        ind_type$CONSTR (SUC 0) (ARB,a,ARB)
                          (n. ind_type$BOTTOM)) a)
                  (a.
                    a0 =
                     (a.
                        ind_type$CONSTR (SUC (SUC 0)) (ARB,ARB,a)
                          (n. ind_type$BOTTOM)) a)
                  'princs' a0)
              'princs' a0) rep
princs_case_def
  |- (a f f1 f2. princs_CASE (PR a) f f1 f2 = f a)
     (a f f1 f2. princs_CASE (User a) f f1 f2 = f1 a)
     a f f1 f2. princs_CASE (Key a) f f1 f2 = f2 a
princs_size_def
  |- (a. princs_size (PR a) = 1 + servers_size a)
     (a. princs_size (User a) = 1 + people_size a)
     a. princs_size (Key a) = 1 + keys_size a
props_TY_DEF
  |- rep.
       TYPE_DEFINITION
         (a0'.
            'props' .
              (a0'.
                 (a.
                    a0' =
                     (a.
                        ind_type$CONSTR 0 (a,ARB,ARB,ARB)
                          (n. ind_type$BOTTOM)) a)
                  (a0 a1.
                    a0' =
                     (a0 a1.
                        ind_type$CONSTR (SUC 0) (ARB,a0,a1,ARB)
                          (n. ind_type$BOTTOM)) a0 a1)
                  (a.
                    a0' =
                     (a.
                        ind_type$CONSTR (SUC (SUC 0)) (ARB,ARB,ARB,a)
                          (n. ind_type$BOTTOM)) a)
                  'props' a0')
              'props' a0') rep
props_case_def
  |- (a f f1 f2. props_CASE (USE a) f f1 f2 = f a)
     (a0 a1 f f1 f2. props_CASE (PK a0 a1) f f1 f2 = f1 a0 a1)
     a f f1 f2. props_CASE (AC a) f f1 f2 = f2 a
```

```
props_size_def
      |- (a. props_size (USE a) = 1 + serviceKey_size a)
         (a0 a1.
            props_size (PK a0 a1) =
            1 + (people_size a0 + keys_size a1))
         a. props_size (AC a) = 1 + access_size a
    servers_BIJ
      |- (a. num2servers (servers2num a) = a)
         r. (n. n < 3) r (servers2num (num2servers r) = r)
    servers_CASE
      |- x v0 v1 v2.
           (case x of AS \Rightarrow v0 | TGS \Rightarrow v1 | FS \Rightarrow v2) =
           (m. if m < 1 then v0 else if m = 1 then v1 else v2)
             (servers2num x)
    servers_TY_DEF
      |- rep. TYPE_DEFINITION (n. n < 3) rep
    servers_size_def
      |-x. servers_size x = 0
    serviceKey_TY_DEF
      |- rep.
           TYPE_DEFINITION
             (a0.
                 'serviceKey' .
                   (a0.
                      (a.
                         a0 =
                         (a.
                            ind_type$CONSTR 0 (a,ARB)
                              (n. ind_type$BOTTOM)) a)
                      (a.
                         a0 =
                         (a.
                            ind_type$CONSTR (SUC 0) (ARB,a)
                              (n. ind_type$BOTTOM)) a)
                      'serviceKey' a0)
                   'serviceKey' a0) rep
    serviceKey_case_def
      |- (a f f1. serviceKey_CASE (SERV a) f f1 = f a)
         a f f1. serviceKey_CASE (K a) f f1 = f1 a
    serviceKey_size_def
      |- (a. serviceKey_size (SERV a) = 1 + servers_size a)
         a. serviceKey_size (K a) = 1 + keys_size a
Theorems:
    access2num_11
      |-aa'. (access2num a = access2num a') (a = a')
    access2num_ONTO
```

```
|-r.r < 3 a. r = access2num a
access2num_num2access
  |-r.r < 3 \quad (access2num \quad (num2access r) = r)
access2num thm
  |- (access2num print = 0) (access2num read = 1)
     (access2num write = 2)
access Axiom
  |-x0 x1 x2. f. (f print = x0) (f read = x1) (f write = x2)
access_EQ_access
  |- a a'. (a = a') (access2num a = access2num a')
access_case_cong
  |- M M' v0 v1 v2.
       (M = M') ((M' = print) (v0 = v0'))
       ((M' = read) (v1 = v1')) ((M' = write) (v2 = v2'))
       ((case M of print \Rightarrow v0 | read \Rightarrow v1 | write \Rightarrow v2) =
        case M' of print => v0' | read => v1' | write => v2')
access_case_def
  |- (v0 v1 v2.
        (case print of print => v0 | read => v1 | write => v2) =
        v0)
     (v0 v1 v2.
        (case read of print => v0 | read => v1 | write => v2) =
     v0 v1 v2.
       (case write of print => v0 | read => v1 | write => v2) = v2
access_distinct
  |- print read print write read write
access_induction
  |- P. P print P read P write a. P a
access_nchotomy
  |-a. (a = print) (a = read) (a = write)
datatype_access
  |- DATATYPE (access print read write)
datatype_keys
  |- DATATYPE (keys KTGS KFS KU KUTGS KUFS)
datatype_people
  |- DATATYPE (people Ursala)
datatype_princs
  |- DATATYPE (princs PR User Key)
datatype_props
  |- DATATYPE (props USE PK AC)
datatype_servers
  |- DATATYPE (servers AS TGS FS)
datatype_serviceKey
  |- DATATYPE (serviceKey SERV K)
init_auth_thm
  |- (M,Oi,Os) sat Name (User Ursala) says prop (USE (SERV TGS))
```

```
(M,Oi,Os) sat
     prop (USE (SERV TGS)) impf
     Name (Key KU) says prop (USE (K KUTGS)) andf
     Name (Key KTGS) says prop (PK Ursala KUTGS)
     (M,Oi,Os) sat
     Name (User Ursala) controls prop (USE (SERV TGS))
     (M,Oi,Os) sat
     Name (Key KU) says prop (USE (K KUTGS)) andf
     Name (Key KTGS) says prop (PK Ursala KUTGS)
keys2num_11
  |-a a'. (keys2num a = keys2num a') (a = a')
keys2num_ONTO
  |-r.r < 5 a. r = keys2num a
keys2num_num2keys
  |-r. r < 5  (keys2num (num2keys r) = r)
keys2num_thm
  |-(keys2num KTGS = 0) (keys2num KFS = 1) (keys2num KU = 2)
     (keys2num KUTGS = 3) (keys2num KUFS = 4)
keys_Axiom
  |- x0 x1 x2 x3 x4.
      f.
         (f KTGS = x0) (f KFS = x1) (f KU = x2)
         (f KUTGS = x3) (f KUFS = x4)
keys_EQ_keys
  |-aa'. (a=a') (keys2num a=keys2num a')
keys_case_cong
  |- M M' v0 v1 v2 v3 v4.
       (M = M') ((M' = KTGS) (v0 = v0'))
       ((M' = KFS) (v1 = v1')) ((M' = KU) (v2 = v2'))
       ((M' = KUTGS) (v3 = v3')) ((M' = KUFS) (v4 = v4'))
       ((case M of
           KTGS => v0
         | KFS => v1
         | KU => v2
         | KUTGS => v3
         | KUFS => v4) =
        case M' of
         KTGS => v0'
        | KFS => v1'
        | KU => v2'
        | KUTGS => v3'
        | KUFS => v4')
keys_case_def
  |- (v0 v1 v2 v3 v4.
        (case KTGS of
           KTGS => v0
         | KFS => v1
```

```
| KU => v2
         | KUTGS => v3
         | KUFS => v4) =
        v0)
     (v0 v1 v2 v3 v4.
        (case KFS of
           KTGS => v0
         | KFS => v1
         | KU => v2
         | KUTGS => v3
         | KUFS => v4) =
        v1)
     (v0 v1 v2 v3 v4.
        (case KU of
           KTGS => v0
         | KFS => v1
         | KU => v2
         | KUTGS => v3
         | KUFS => v4) =
        v2)
     (v0 v1 v2 v3 v4.
        (case KUTGS of
           KTGS => v0
         | KFS => v1
         | KU => v2
         | KUTGS => v3
         | KUFS => v4) =
        v3)
     v0 v1 v2 v3 v4.
       (case KUFS of
         KTGS => v0
        | KFS => v1
        | KU => v2
        | KUTGS => v3
        | KUFS => v4) =
       v4
keys_distinct
  |- KTGS KFS KTGS KU KTGS KUTGS KTGS KUFS
     KFS KU KFS KUTGS KFS KUFS KU KUTGS KU KUFS
     KUTGS KUFS
keys_induction
  |- P. P KFS P KTGS P KU P KUFS P KUTGS a. P a
keys_nchotomy
  |- a.
       (a = KTGS) (a = KFS) (a = KU) (a = KUTGS) (a = KUFS)
num2access_11
  |- r r'.
```

```
r < 3 r' < 3 ((num2access r = num2access r') (r = r'))
num2access_ONTO
  |-a.r. (a = num2access r) r < 3
num2access access2num
  |-a. num2access (access2num a) = a
num2access_thm
  |- (num2access 0 = print) (num2access 1 = read)
     (num2access 2 = write)
num2keys_11
  |-rr'.r' < 5r' < 5 ((num2keys r = num2keys r') (r = r'))
num2keys_ONTO
  |-a.r. (a = num2keys r) r < 5
num2keys_keys2num
  |-a. num2keys (keys2num a) = a
num2keys_thm
  |- (num2keys 0 = KTGS) (num2keys 1 = KFS) (num2keys 2 = KU)
     (num2keys 3 = KUTGS) (num2keys 4 = KUFS)
num2people_11
  |- r r'.
       r < 1 r' < 1 ((num2people r = num2people r') (r = r'))
num2people_ONTO
  |-a.r. (a = num2people r) r < 1
num2people_people2num
  |-a. num2people (people2num a) = a
num2people_thm
  |- num2people 0 = Ursala
num2servers_11
  |- r r'.
      r' < 3
       ((num2servers r = num2servers r') (r = r'))
num2servers_ONTO
  |-a.r. (a = num2servers r) r < 3
num2servers_servers2num
  |- a. num2servers (servers2num a) = a
num2servers_thm
  |- (num2servers 0 = AS) (num2servers 1 = TGS)
     (num2servers 2 = FS)
people2num_11
  |- a a'. (people2num a = people2num a') (a = a')
people2num_ONTO
  |-r.r < 1 a. r = people2num a
people2num_num2people
  |-r.r < 1 \pmod{people2num (num2people r)} = r
people2num_thm
  |- people2num Ursala = 0
people_Axiom
```

```
|-x0.f.f Ursala = x0
people_EQ_people
  |- a a'. (a = a') (people2num a = people2num a')
people_case_cong
  |- M M' vO.
       (M = M') ((M' = Ursala) (v0 = v0'))
       ((case M of Ursala => v0) = case M' of Ursala => v0')
people_case_def
  |- v0. (case Ursala of Ursala => v0) = v0
people_induction
  |- P. P Ursala a. P a
people_nchotomy
  |-a.a=Ursala
princs_11
  |-(a a'. (PR a = PR a') (a = a'))
     (a a'. (User a = User a') (a = a'))
     a a'. (Key a = Key a') (a = a')
princs_Axiom
  |- f0 f1 f2.
       fn.
         (a. fn (PR a) = f0 a) (a. fn (User a) = f1 a)
         a. fn (Key a) = f2 a
princs_case_cong
  |- M M' f f1 f2.
       (M = M') (a. (M' = PR a) (f a = f' a))
       (a. (M' = User a) (f1 a = f1' a))
       (a. (M' = Key a) (f2 a = f2' a))
       (princs_CASE M f f1 f2 = princs_CASE M' f' f1' f2')
princs_distinct
  |- (a' a. PR a User a') (a' a. PR a Key a')
    a' a. User a Key a'
princs_induction
  |- P.
       (s. P (PR s)) (p. P (User p)) (k. P (Key k)) p. P p
princs_nchotomy
  |-pp. (s. pp = PR s) (p. pp = User p) k. pp = Key k
props_11
  |-(a a'. (USE a = USE a') (a = a'))
     (a0 a1 a0' a1'.
        (PK a0 a1 = PK a0' a1') (a0 = a0') (a1 = a1'))
    a a'. (AC a = AC a') (a = a')
props_Axiom
  |- f0 f1 f2.
       fn.
         (a. fn (USE a) = f0 a)
         (a0 a1. fn (PK a0 a1) = f1 a0 a1) a. fn (AC a) = f2 a
props_case_cong
```

```
|- M M' f f1 f2.
       (M = M') (a. (M' = USE a) (f a = f' a))
       (a0 a1. (M' = PK a0 a1) (f1 a0 a1 = f1' a0 a1))
       (a. (M' = AC a) (f2 a = f2' a))
       (props_CASE M f f1 f2 = props_CASE M' f' f1' f2')
props_distinct
  |- (a1 a0 a. USE a PK a0 a1) (a' a. USE a AC a')
     a1 a0 a. PK a0 a1 AC a
props_induction
  |- P.
       (s. P (USE s)) (p k. P (PK p k)) (a. P (AC a))
props_nchotomy
  |-pp. (s. pp = USE s) (p k. pp = PK p k) a. pp = AC a
request_for_services_thm
  |- (M,Oi,Os) sat Name (Key KTGS) speaks_for Name (PR AS)
     (M,Oi,Os) sat
     prop (PK Ursala KUTGS) impf
     Name (Key KUFS) says prop (AC read)
     (M,Oi,Os) sat
     prop (PK Ursala KUTGS) impf
     Name (Key KFS) says
     Name (Key KUFS) speaks_for Name (User Ursala)
     (M,Oi,Os) sat Name (PR AS) controls prop (PK Ursala KUTGS)
     (M,Oi,Os) sat
     Name (Key KUTGS) says prop (USE (SERV FS)) andf
     Name (Key KTGS) says prop (PK Ursala KUTGS)
     (M,Oi,Os) sat
     Name (Key KFS) says
     Name (Key KUFS) speaks_for Name (User Ursala) andf
     Name (Key KUFS) says prop (AC read)
servers2num_11
  |- a a'. (servers2num a = servers2num a') (a = a')
servers2num_ONTO
  |-r.r < 3 a. r = servers2num a
servers2num_num2servers
  |-r. r < 3  (servers2num (num2servers r) = r)
servers2num_thm
  |- (servers2num AS = 0) (servers2num TGS = 1)
     (servers2num FS = 2)
servers_Axiom
  |-x0 x1 x2. f. (f AS = x0) (f TGS = x1) (f FS = x2)
servers_EQ_servers
  |- a a'. (a = a') (servers2num a = servers2num a')
servers_case_cong
  |- M M' v0 v1 v2.
       (M = M') ((M' = AS) (v0 = v0'))
```

```
((M' = TGS) (v1 = v1')) ((M' = FS) (v2 = v2'))
       ((case M of AS => v0 | TGS => v1 | FS => v2) =
        case M' of AS => v0' | TGS => v1' | FS => v2')
servers_case_def
  |- (v0 v1 v2.
        (case AS of AS \Rightarrow v0 | TGS \Rightarrow v1 | FS \Rightarrow v2) \Rightarrow v0)
     (v0 v1 v2.
        (case TGS of AS \Rightarrow v0 | TGS \Rightarrow v1 | FS \Rightarrow v2) = v1)
     v0 v1 v2. (case FS of AS => v0 | TGS => v1 | FS => v2) = v2
servers_distinct
  |- AS TGS AS FS TGS FS
servers_induction
  |- P. PAS PFS PTGS a. Pa
servers_nchotomy
  |-a. (a = AS) (a = TGS) (a = FS)
serviceKey_11
  |-(a a'. (SERV a = SERV a') (a = a'))
     a a'. (K a = K a') (a = a')
serviceKey_Axiom
  |-f0 f1. fn. (a. fn (SERV a) = f0 a) a. fn (K a) = f1 a
serviceKey_case_cong
  I- M M' f f1.
       (M = M') (a. (M' = SERV a) (f a = f' a))
       (a. (M' = K a) (f1 a = f1' a))
       (serviceKey_CASE M f f1 = serviceKey_CASE M' f' f1')
serviceKey_distinct
  |- a' a. SERV a K a'
serviceKey_induction
  |- P. (s. P (SERV s)) (k. P (K k)) s. P s
serviceKey_nchotomy
  |-ss. (s. ss = SERV s) k. ss = K k
service_request_thm
  |- (M,Oi,Os) sat Name (Key KFS) speaks_for Name (PR TGS)
     (M,Oi,Os) sat Name (User Ursala) controls prop (AC print)
     (M,Oi,Os) sat
     Name (PR TGS) controls
     Name (Key KUFS) speaks_for Name (User Ursala)
     (M,Oi,Os) sat
     Name (Key KFS) says
     Name (Key KUFS) speaks_for Name (User Ursala) andf
     Name (Key KUFS) says prop (AC print)
     (M,Oi,Os) sat prop (AC print)
session_key_receipt_thm
  |- (M,Oi,Os) sat Name (Key KU) speaks_for Name (PR AS)
     (M,Oi,Os) sat
     prop (USE (K KUTGS)) impf
     Name (Key KUTGS) says prop (USE (SERV FS))
```

Process HOL finished

3 Question 2

3.1 Problem Statement

This problem explores defining a simple state machine in the form of simpleOpener Theory. The dataytpes are defined below as well as the next state, next output, and other relevant theorems for the state machine. Source for the state machine can be found in appendix B.

3.2 Datatypes/Theorems

```
command = i0 | i1
output = 00 | o1
state = SO \mid S1
 \vdash (simpleOpenerns S0 i1 = S1) \land (simpleOpenerns S1 i0 = S0)
 ⊢ (simpleOpenerout SO i1 = o1) ∧ (simpleOpenerout S1 i0 = o0)
 \vdash i0 \neq i1
 \vdash o0 \neq o1
 \vdash (\forall ins outs.
        TR i1 (CFG (i1::ins) SO outs) (CFG ins S1 (o1::outs))) \land
    \forall ins outs.
       TR iO (CFG (iO::ins) S1 outs) (CFG ins SO (oO::outs))
\vdash \ \forall \, P. \ P \ \mathtt{SO} \ \mathtt{i1} \ \land \ P \ \mathtt{S1} \ \mathtt{i0} \ \land \ P \ \mathtt{S0} \ \mathtt{i0} \ \land \ P \ \mathtt{S1} \ \mathtt{i1} \ \Rightarrow \ \forall \, v \ v_1. \ P \ v \ v_1
\vdash \forall P. P \text{ SO il} \land P \text{ Sl io} \land P \text{ SO io} \land P \text{ Sl il} \Rightarrow \forall v \ v_1. P \ v \ v_1
 \vdash (\forall x \ x1s \ s_1 \ out1s \ x2s \ out2s \ s_2.
        TR x (CFG x1s s_1 out1s) (CFG x2s s_2 out2s) \iff
        \exists NS \ Out \ ins.
            (x1s = x::ins) \land (x2s = ins) \land (s_2 = NS \ s_1 \ x) \land
            (out2s = Out s_1 x::out1s)) \land
    \forall x \ x1s \ s_1 \ out1s \ x2s \ out2s.
       TR x (CFG x1s s_1 out1s)
          (CFG x2s (simpleOpenerns s_1 x)
               (simpleOpenerout s_1 \ x :: out2s)) \iff
       \exists ins. (x1s = x::ins) \land (x2s = ins) \land (out2s = out1s)
 \vdash \forall s \ x \ ins \ outs.
       TR x (CFG (x::ins) s outs)
          (CFG ins (simpleOpenerns s x)
              (simpleOpenerout s \ x :: outs))
 \vdash TR x (CFG (x::ins) s outs)
       (CFG ins (simpleOpenerns s x)
            (simpleOpenerout \ s \ x::outs)) \iff
    Trans x s (simpleOpenerns s x)
 \vdash S0 \neq S1
```

3.3 Execution Transcript

```
HOL-4 [Kananaskis 11 (stdknl, built Sat Aug 19 09:30:06 2017)]
      For introductory HOL help, type: help "hol";
      To exit type <Control>-D
          _____
[extending loadPath with Holmakefile INCLUDES variable]
>>>>>> > (<HOL message: Created theory "simpleOpener">>
<<HOL message: Defined type: "command">>
<<HOL message: Defined type: "state">>
<<HOL message: Defined type: "output">>
<< HOL message: mk_functional:
 pattern completion has added 2 clauses to the original specification.>>
<<HOL warning: GrammarDeltas.revise_data:
 Grammar-deltas:
   overload_on("simpleOpenerns_tupled")
  invalidated by DelConstant(simpleOpener$simpleOpenerns_tupled)>>
Equations stored under "simpleOpenerns_def".
Induction stored under "simpleOpenerns_ind".
<< HOL message: mk_functional:
 pattern completion has added 2 clauses to the original specification.>>
<< HOL warning: GrammarDeltas.revise_data:
 Grammar-deltas:
   overload_on("simpleOpenerout_tupled")
  invalidated by DelConstant(simpleOpener$simpleOpenerout_tupled)>>
Equations stored under "simpleOpenerout_def".
Induction stored under "simpleOpenerout_ind".
Exporting theory "simpleOpener" ... done.
Theory "simpleOpener" took 0.31254s to build
Theory: simpleOpener
Parents:
   Sm
Type constants:
   command 0
   output 0
   state 0
Term constants:
   S0
                     :state
   S1
                     :state
                  :command -> -> ->
   command2num
   command_CASE
```

```
command_size
                      :command -> num
    i0
                      :command
    i1
                      :command
    num2command
                      :num -> command
                      :num -> output
    num2output
    num2state
                      :num -> state
                      :output
    01
                      :output
    output2num
                      :output -> num
    output_CASE
                      :output -> -> ->
    output_size
                      :output -> num
    simpleOpenerns
                      :state -> command -> state
    simpleOpenerout
                      :state -> command -> output
    state2num
                      :state -> num
    state_CASE
                      :state -> ->
    state_size
                      :state -> num
Definitions:
    @tempSO_def
      |-S0 = num2state 0
    @tempS1_def
      |-S1 = num2state 1
    @tempi0_def
      |-i0 = num2command 0
    @tempi1_def
      |-i1 = num2command 1
    @tempo0_def
      |- 00 = num2output 0
    @tempo1_def
      |- o1 = num2output 1
    command_BIJ
      |-(a. num2command (command2num a) = a)
         r. (n. n < 2) r (command2num (num2command r) = r)
    command_CASE
      I- x v0 v1.
           (case x of i0 => v0 | i1 => v1) =
           (m. if m = 0 then v0 else v1) (command2num x)
    command_TY_DEF
      |- rep. TYPE_DEFINITION (n. n < 2) rep</pre>
    command_size_def
      |-x. command_size x = 0
    output_BIJ
      |- (a. num2output (output2num a) = a)
         r. (n. n < 2) r (output2num (num2output r) = r)
    output_CASE
      |- x v0 v1.
           (case x of o0 => v0 | o1 => v1) =
```

```
(m. if m = 0 then v0 else v1) (output2num x)
   output_TY_DEF
      |- rep. TYPE_DEFINITION (n. n < 2) rep
    output_size_def
      |-x. output_size x = 0
    state_BIJ
      |-(a. num2state (state2num a) = a)
        r. (n. n < 2) r (state2num (num2state r) = r)
   state_CASE
      |- x v0 v1.
           (case x of S0 => v0 | S1 => v1) =
           (m. if m = 0 then v0 else v1) (state2num x)
   state_TY_DEF
      |- rep. TYPE_DEFINITION (n. n < 2) rep</pre>
   state_size_def
      |-x. state_size x = 0
Theorems:
   command2num_11
      |- a a'. (command2num a = command2num a') (a = a')
    command2num_ONTO
      1-r. r < 2 a. r = command2num a
    command2num_num2command
      |-r. r < 2 \pmod{num2command r} = r
    command2num_thm
      |-(command2num i0 = 0) (command2num i1 = 1)
    command_Axiom
      |-x0 x1. f. (f i0 = x0) (f i1 = x1)
    command_EQ_command
      |- a a'. (a = a') (command2num a = command2num a')
   command_case_cong
      |- M M' v0 v1.
           (M = M') ((M' = i0) (v0 = v0'))
           ((M' = i1) (v1 = v1'))
           ((case M of i0 => v0 | i1 => v1) =
           case M' of i0 => v0' | i1 => v1')
    command_case_def
      |-(v0 v1. (case i0 of i0 => v0 | i1 => v1) = v0)
         v0 v1. (case i1 of i0 => v0 | i1 => v1) = v1
    command_distinct
      |- i0 i1
   command_distinct_clauses
      l- i0 i1
    command_induction
      |- P. P i0 P i1 a. P a
    command_nchotomy
      |-a. (a = i0) (a = i1)
```

```
datatype_command
  |- DATATYPE (command i0 i1)
datatype_output
  |- DATATYPE (output o0 o1)
datatype_state
  |- DATATYPE (state SO S1)
num2command 11
  |- r r'.
       r < 2
       r' < 2
       ((num2command r = num2command r') (r = r'))
num2command_ONTO
  |-a.r. (a = num2command r) r < 2
num2command_command2num
  |-a. num2command (command2num a) = a
num2command_thm
  |- (num2command 0 = i0) (num2command 1 = i1)
num2output_11
  |- r r'.
       r < 2 r' < 2 ((num2output r = num2output r') (r = r'))
num2output_ONTO
  |-a.r. (a = num2output r) r < 2
num2output_output2num
  |- a. num2output (output2num a) = a
num2output_thm
  |-(num2output 0 = o0) (num2output 1 = o1)
num2state_11
  |- r r'.
       r < 2 r' < 2 ((num2state r = num2state r') (r = r'))
num2state_ONTO
  |-a.r. (a = num2state r) r < 2
num2state_state2num
  |- a. num2state (state2num a) = a
num2state_thm
  |-(num2state 0 = S0) (num2state 1 = S1)
output2num_11
  |- a a'. (output2num a = output2num a') (a = a')
output2num_ONTO
  |-r.r < 2 a. r = output2num a
output2num_num2output
  |-r. r < 2  (output2num (num2output r) = r)
output2num_thm
  |- (output2num o0 = 0) (output2num o1 = 1)
output_Axiom
  |-x0 x1. f. (f o0 = x0) (f o1 = x1)
output_EQ_output
  |- a a'. (a = a') (output2num a = output2num a')
```

```
output_case_cong
  |- M M' v0 v1.
       (M = M') ((M' = 00) (v0 = v0'))
       ((M' = o1) (v1 = v1'))
       ((case M of o0 => v0 | o1 => v1) =
       case M' of o0 => v0' | o1 => v1')
output_case_def
  |-(v0 v1. (case o0 of o0 => v0 | o1 => v1) = v0)
     v0 v1. (case o1 of o0 => v0 | o1 => v1) = v1
output_distinct
  |- o0 o1
output_distinct_clauses
  |- o0 o1
output_induction
  |- P. P o 0 P o 1 a. P a
output_nchotomy
  |-a. (a = o0) (a = o1)
simpleCounter_rules
  |- (ins outs.
        TR i1 (CFG (i1::ins) SO outs) (CFG ins S1 (o1::outs)))
     ins outs.
       TR iO (CFG (iO::ins) S1 outs) (CFG ins SO (oO::outs))
simpleOpenerTR_clauses
  |- (x x1s s1 out1s x2s out2s s2.
        TR x (CFG x1s s1 out1s) (CFG x2s s2 out2s)
        NS Out ins.
          (x1s = x::ins) (x2s = ins) (s2 = NS s1 x)
          (out2s = Out s1 x::out1s))
     x x1s s1 out1s x2s out2s.
       TR x (CFG x1s s1 out1s)
         (CFG x2s (simpleOpenerns s1 x)
            (simpleOpenerout s1 x::out2s))
       ins. (x1s = x::ins) (x2s = ins) (out2s = out1s)
simpleOpenerTR_rules
  |- s x ins outs.
       TR x (CFG (x::ins) s outs)
         (CFG ins (simpleOpenerns s x) (simpleOpenerout s x::outs))
simpleOpenerTrans_Equiv_TR
  |- TR x (CFG (x::ins) s outs)
       (CFG ins (simpleOpenerns s x) (simpleOpenerout s x::outs))
     Trans x s (simpleOpenerns s x)
simpleOpenerns_def
  |- (simpleOpenerns S0 i1 = S1) (simpleOpenerns S1 i0 = S0)
simpleOpenerns_ind
  |- P. P SO i1 P S1 i0 P S0 i0 P S1 i1 v v1. P v v1
simpleOpenerout_def
  |- (simpleOpenerout SO i1 = o1) (simpleOpenerout S1 i0 = o0)
```

```
simpleOpenerout_ind
      state2num_11
      |- a a'. (state2num a = state2num a') (a = a')
   state2num_ONTO
      |-r.r < 2 a. r = state2num a
   state2num num2state
     |-r.r < 2  (state2num (num2state r) = r)
   state2num_thm
     |-(state2num S0 = 0) (state2num S1 = 1)
   state_Axiom
      |-x0 x1. f. (f S0 = x0) (f S1 = x1)
   state_EQ_state
      |- a a'. (a = a') (state2num a = state2num a')
   state_case_cong
      |- M M' v0 v1.
          (M = M') ((M' = SO) (vO = vO'))
          ((M' = S1) (v1 = v1'))
          ((case M of SO => vO | S1 => v1) =
           case M' of SO \Rightarrow vO' | S1 \Rightarrow v1')
   state_case_def
      |-(v0 v1. (case S0 of S0 => v0 | S1 => v1) = v0)|
        v0 v1. (case S1 of S0 => v0 | S1 => v1) = v1
   state_distinct
      |- S0 S1
   state_distinct_clauses
      I- S0 S1
   state_induction
      |- P. P SO P S1 a. P a
   state_nchotomy
     |-a. (a = S0) (a = S1)
structure simpleOpenerScript:
  sig
  end
val it = (): unit
*** Emacs/HOL command completed ***
```

4 Question 3

In the healthcare industry, access control polocies can be overuled in cases of emergency in order to swiftly get patients the care that they need. Unfortunately, this also opens the door to people who want to take advantage of these "break the glass" policies. The policy spaces described in the paper may offer a good way to provide a methodology to deciding when to break the polocies. The paper describes a good way of splitting up the polocies that should never be broken into their own category, which is good in keeping malicious people from resources that should never be exploited. The more interesting cases are the policis in the unknown space. The enviornmental effects that govern these polocies may help healthcare providers react to disastrous cases are situations more quickly, since they have a methodology in breaking their poloices instead of leaving them up to interpretation by individual gatekeepers. On the flip side, the complications added to the access control logic to incorporate this dynamic behavior may perhaps lead to more unforseen vulnerabilities, as the environment could behave in unpredictable ways. Therefore, much care and effort must be made when decision these policy spaces.

A Source For exam4Script.sml

```
The following code is from HOL/exam4Script.sml
Exam4 Theory: Access Control Mechanisms
*)
(*
   Author: Alfred Murabito
*)
   Date: 3/26/2020
(*
* )
structure exam4Script = struct
(* only necessary when working interactively
app load ["acl_infRules", "aclrulesTheory", "aclDrulesTheory", "exam4Theory"]
open \ acl_{-}infRules \ aclrules Theory \ aclDrules Theory \ exam4 Theory
* )
open HolKernel boolLib Parse bossLib
open acl_infRules aclrulesTheory aclDrulesTheory
val _ = new_theory "exam4"
(* Theory Definitions *)
val =
Datatype
'servers = AS | TGS | FS'
val_{-} =
Datatype
'people = Ursala'
val_{-} =
Datatype
'keys = KTGS | KFS | KU | KUTGS | KUFS'
val_{-} =
Datatype
'princs = PR servers | User people | Key keys'
val =
Datatype
'access = print | read | write'
\mathbf{val}_{-} =
Datatype
```

```
'serviceKey = SERV servers | K keys'
val _ =
Datatype
'props = USE serviceKey | PK people keys | AC access'
val formal = 'Name (User Ursala) says prop (USE (SERV TGS)): (props, princs, 'd, 'e
val forma2 = ''Name (User Ursala) controls prop (USE (SERV TGS)):(props,princs,'o
val forma3 = ''prop (USE (SERV TGS)) impf (Name (Key KU) says prop (USE (K KUTGS
\mathbf{val} \ [ \mathbf{forma1\_thm} \ , \mathbf{forma2\_thm} \ , \mathbf{forma3\_thm} \ ] \ = \ \mathbf{map} \ \mathbf{ACL\_ASSUM} \ [ \mathbf{forma1} \ , \mathbf{forma2} \ , \mathbf{forma3} \ ]
val forma4_thm = CONTROLS forma2_thm forma1_thm
val forma_thm = ACL_MP forma4_thm forma3_thm
val init_auth_thm = DISCH_ALL forma_thm
val _ = save_thm("init_auth_thm", init_auth_thm)
val formb1 = ''(Name (Key KU) says prop (USE (K KUTGS)) andf
                Name (Key KTGS) says prop (PK Ursala KUTGS)):(props, princs, 'd, 'e)
val formb2 = 'Name (PR AS) controls prop (USE (K KUTGS)): (props, princs, 'd, 'e) For
val formb3 = ''(Name (Key KU) speaks_for Name (PR AS)):(props,princs,'d,'e)Form'
val formb4 = ''prop (USE (K KUTGS)) impf Name (Key KUTGS) says prop (USE (SERV F
val [formb1_thm, formb2_thm, formb3_thm, formb4_thm] = map ACL_ASSUM[formb1, formb2,
val formb5_thm = ACL_SIMP1 formb1_thm
val formb6\_thm = ACL\_SIMP2 formb1\_thm
val formb7_thm = SPEAKS_FOR formb3_thm formb5_thm
val formb8_thm = CONTROLS formb2_thm formb7_thm
val formb9_thm = ACL_MP formb8_thm formb4_thm
val formb_thm = ACL_CONJ formb9_thm formb6_thm
val session_key_receipt_thm = DISCH_ALL formb_thm
val _ = save_thm("session_key_receipt_thm", session_key_receipt_thm)
val formc1 = ''Name (Key KUTGS) says prop (USE (SERV FS)) andf
                Name (Key KTGS) says prop (PK Ursala KUTGS): (props, princs, 'd, 'e) Fe
val formc2 = ''(Name (PR AS)) controls (prop (PK Ursala KUTGS)):(props, princs, 'd
val formc3 = ''(Name (Key KTGS) speaks_for (Name (PR AS))):(props, princs, 'd, 'e)Fe
val formc4 = "prop (PK Ursala KUTGS) impf (Name (Key KFS) says ((Name (Key KUFS
val formc5 = ''prop (PK Ursala KUTGS) impf (Name (Key KUFS) says prop (AC read))
```

```
val [formc1_thm, formc2_thm, formc3_thm, formc4_thm, formc5_thm] = map ACL_ASSUM[formc1, formc1, formc1, formc1, formc2_thm]
val formc6_thm = ACL_SIMP1 formc1_thm
val formc7_thm = ACL_SIMP2 formc1_thm
val formc8_thm = SPEAKS_FOR formc3_thm formc7_thm
val formc9_thm = CONTROLS formc2_thm formc8_thm
val formc10_thm = ACLMP formc9_thm formc4_thm
val formc11_thm = ACL_MP formc9_thm formc5_thm
val formc_thm = ACL_CONJ formc10_thm formc11_thm
val request_for_services_thm = DISCH_ALL formc_thm
val _ = save_thm("request_for_services_thm", request_for_services_thm)
val formd1 = ''(Name (Key KFS) says (Name (Key KUFS) speaks_for Name (User Ursala)))
                           (Name (Key KUFS) says prop (AC print)): (props, princs, 'd, 'e
val formd2 = ''(Name (Key KFS) speaks_for Name (PR TGS)):(props, princs, 'd, 'e)Form''
val formd3 = ''(Name (PR TGS) controls (Name (Key KUFS) speaks_for Name (User Ursala)
val formd4 = ''(Name (User Ursala) controls prop(AC print)):(props, princs, 'd, 'e)Form'
val [formd1_thm, formd2_thm, formd3_thm, formd4_thm] = map ACL_ASSUM[formd1, formd2, formd
val formd5_thm = ACL_SIMP1 formd1_thm
val formd6\_thm = ACL\_SIMP2 formd1\_thm
val formd7_thm = SPEAKS_FOR formd2_thm formd5_thm
val formd8_thm = CONTROLS formd3_thm formd7_thm
val formd9_thm = SPEAKS_FOR formd8_thm formd6_thm
val formd_thm = CONTROLS formd4_thm formd9_thm
val service_request_thm = DISCH_ALL formd_thm
val _ = save_thm("service_request_thm", service_request_thm)
val _ = print_theory "-";
val = export_theory();
end
    Source For simpleOpenerScript.sml
  The following code is from HOL/simpleOpenerScript.sml
```

(* Exam 4: Simple Opener Theory

(*Alfred Murabito

```
(* March 26 2020
* )
structure simpleOpenerScript = struct
(* interactive mode
app load ["TypeBase", "smTheory", "sminfRules", "simpleOpenerTheory"];
open TypeBase smTheory sminfRules simpleOpenerTheory
*)
open HolKernel boolLib Parse bossLib
open TypeBase smTheory sminfRules
val _ = new_theory "simpleOpener"
(* Define the inputs
*)
(* -----
val _ =
Datatype
'command = i0 | i1'
val command_distinct_clauses = distinct_of ':command''
val _ = save_thm("command_distinct_clauses", command_distinct_clauses)
(* Define the states
*)
(* --
val _ =
Datatype
'state = S0 \mid S1'
val state_distinct_clauses = distinct_of '':state''
val _ = save_thm("state_distinct_clauses", state_distinct_clauses)
(* Define the outputs
*)
(* --
val_{-} =
Datatype
'output = 00 \mid 01'
val output_distinct_clauses = distinct_of ':output''
```

```
val _ = save_thm("output_distinct_clauses", output_distinct_clauses)
* )
val simpleOpenerns_def =
Define '(simpleOpenerns S0 i1 = S1) /\ (simpleOpenerns S1 i0 = S0)'
(* Define next-output function for machine simpleOpener
* )
val simpleOpenerout_def =
Define '(simpleOpenerout S0 i1 = o1) /\ (simpleOpenerout S1 i0 = o0)'
(*\ Specialize\ TR\_rules\ to\ simpleOpener
* )
*/
(*

val simpleOpenerTR_rules = SPEC_TR ''simpleOpenerns'' ''simpleOpenerout''
val _ = save_thm("simpleOpenerTR_rules", simpleOpenerTR_rules)
*)
*)
(*

val simpleOpenerTR_clauses = SPEC_TR_clauses ''simpleOpenerns'' ''simpleOpenerout''
val _ = save_thm("simpleOpenerTR_clauses", simpleOpenerTR_clauses)
*)
(*
val simpleOpenerTrans_Equiv_TR = SPEC_Trans_Equiv_TR ''simpleOpenerns'' ''simpleOpener
val _ = save_thm("simpleOpenerTrans_Equiv_TR", simpleOpenerTrans_Equiv_TR)
(* Theorems corresponding to the tabular specification of simpleOpener
(* -----
val th1 = REWRITE_RULE[simpleOpenerns_def, simpleOpenerout_def](SPECL[''S0'', ''i1''']
val th2 = REWRITERULE[simpleOpenerns_def, simpleOpenerout_def](SPECL[''S1'', ''i0'']
```

```
val simpleCounter_rules = LIST_CONJ [th1,th2]
val _ = save_thm("simpleCounter_rules", simpleCounter_rules)
val _ = export_theory ()
val _ = print_theory "-"
end
```