



HW3

Abhishek Johri



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PARALLEL PROGRAMMING

University of Illinois at Urbana-Champaign

1a.

Computation:

$$A * N^2 / P$$

Communication:

$$KN$$

Comm-to-Comp Ratio

$$K * P / N$$

Efficiency:

$$1 / (1 + (K * P / N))$$

Isoefficiency:

$$\text{Let } W = N^2$$

We know that $N = \sqrt{W}$

$$P / \sqrt{W} = k; W = (P/k)^2$$

We know that the allgather on a hypercube has a depth of $\log(p)$

Communication Cost:

We start with N/P data and that amount will increase with every level of the gather.

This gives us a sum of $(\text{Alpha} + \text{Beta} * N/P * 2i)$ for when $i = 0$ to when $i = \log(p) - 1$

b.

Computation:

$$A * N / P$$

Communication:

$$KN$$

Comm-to-Comp Ratio

$$K * P / A$$

Efficiency:

$$1 / (1 + (K * P / A))$$

Isoefficiency:

Problem size is W

2. a.

$$\begin{aligned} \text{i) Diameter} &= (x-1) + (y-1) + (z-1) \\ &= 29 + 39 + 49 \\ &= 117 \end{aligned}$$

Degree = [3, 6] (range of 3 to 6, inclusive)

The diameter of this given mesh is 117 and the degree for each node is between 3 and 6.

ii) If we had a 2D mesh of 30x50, the bisection width of this would be 30. If we have a 3D mesh of 30x40x50 we just have the 40 layers of the 30x50 stacked one top of one another. Because of this we have a bisection width of $30 * 40 = 1200$. In order to find the bisection bandwidth of the mesh we will need to multiply the bisection by the channel bandwidth. We are given a channel bandwidth as 4 GB/s

$$\begin{aligned}\text{bisection bandwidth} &= \text{channel bandwidth} * \text{bisection width} \\ &= 4(\text{GB/s}) * 1200 = 4800(\text{GB/s})\end{aligned}$$

So we have a resulting bisection bandwidth of the network will be 4800 GB/s

iii) Since we need to minimize bandwidth usage we need to take advantage of all the 2GB/s channel bandwidths. Because of this we need to cut parallels to the x axis apart from the z axis.

$$\begin{aligned}\text{Bisection width} &= 40 * 50 \\ &= 2000\end{aligned}$$

This give us 2 GB/s

$$\begin{aligned}\text{Bisection bandwidth} &= \text{channel bandwidth} * \text{bisection width} \\ &= 2 * 2000 \\ &= 4000 \text{ GB/s}\end{aligned}$$

b.

Assumptions:

1. The original 2D grid of processes are i, j , and k where $i = [0, 199]$, $j = [0, 299]$ and $k = [0, 0]$
2. The 3D grid of the processes defined above are labeled with i', j', k' where $i' = [0, 29]$, $j' = [0, 39]$, $k' = [0, 49]$
3. The total # of nodes in all of the cases are the same:
 - a. $(200 * 300 \Rightarrow 60000 \Rightarrow 30 * 40 * 50)$ is true

There are various mappings that can be used in such a situation.

The simplest mapping would take advantage of this:

$$\begin{aligned}30 * 10 &= 300 \\ 40 * 50 / 10 &= 200\end{aligned}$$

From this we get the following mapping

$$f(i', j', k') = (j/10, i/5, k/4)$$

From this mapping from the y-axis from the original grid to the x-axis of the 3D mesh.

This uses the x axis then from the original grid. This finally creates the y and z axes of the 3D meshes.