UIU QUESTION BANK

MID-TERM QUESTION SOLUTIONS

THEORY OF COMPUTATION

CSE 2233

SOLUTION BY

NURUL ALAM ADOR

UPDATED TILL ***FALL 2023***

nurulalamador.github.io/**UIUQuestionBank**

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1. Design DFAs that accept the following languages:

   a) L = accepts any string that has an even number of 'a' or odd number of 'b' over alphabet {a, b}

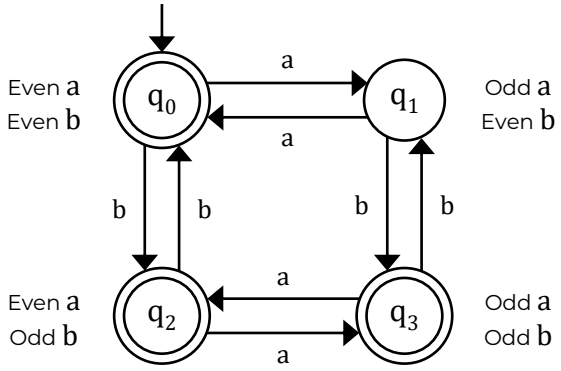
   b) L = accepts any string which starts with an odd number of 'r' and ends with 'brb' over the alphabet {b, r}

   c) L = accepts any string where the last two symbols are different over the alphabet {a, b}

   Solution:

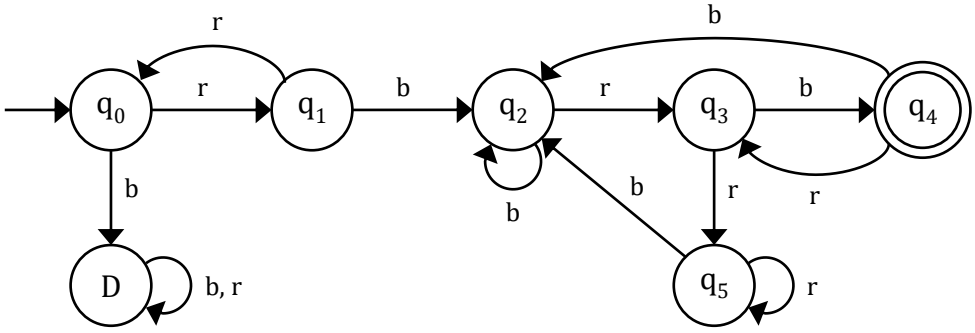
   a) L = accepts any string that has an even number of 'a' or odd number of 'b' over alphabet {a, b}

   The DFA has been designed below:



   b) L = accepts any string which starts with an odd number of 'r' and ends with 'brb' over the alphabet {b, r}

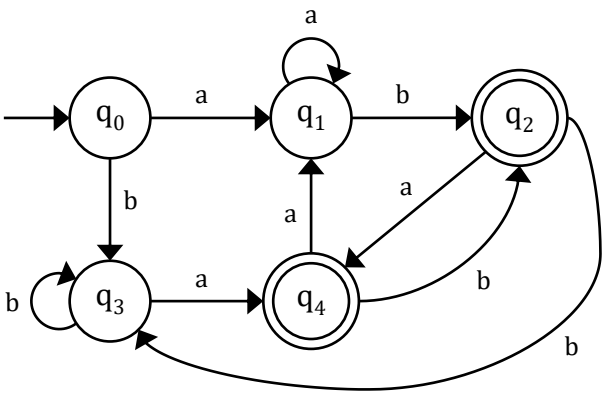
   The DFA has been designed below:



   c) L = accepts any string where the last two symbols are different over the alphabet {a, b}

   The DFA has been designed below:

   [P.T.O]



2. Design NFAs that accept the following languages:

a) L = { w | w starts with 'a' or 'b' and contains 'cca' and ends with 'b' or 'c' } | ∑ = {a, b, c}

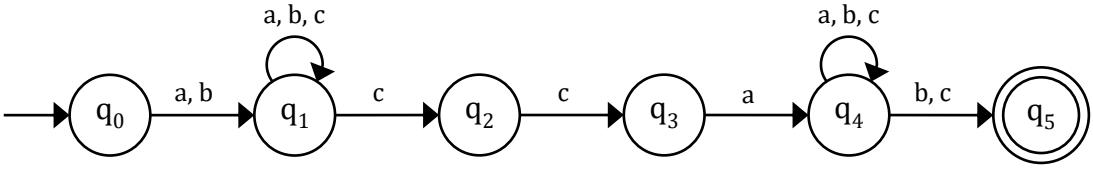
b) L = { w | w starts and ends with different symbols with total length of at least 2 } | ∑ = {a, b, c}

c) L = { w | w contains 'xyz' or 'yzx' or 'zxx' and ends with 'yz' } | ∑ = {x, y, z}

Solution:

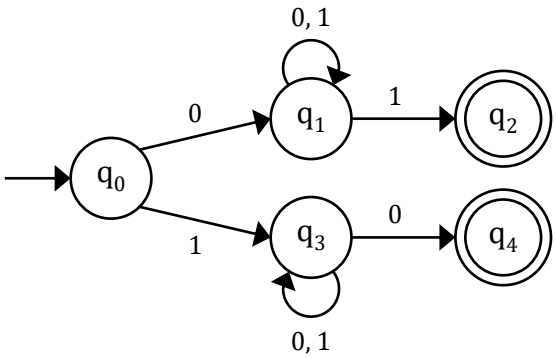
a) L = { w | w starts with 'a' or 'b' and contains 'cca' and ends with 'b' or 'c' } | ∑ = {a, b, c}

The NFA has been designed below:



b) L = { w | w starts and ends with different symbols with total length of at least 2 } | ∑ = {a, b, c}

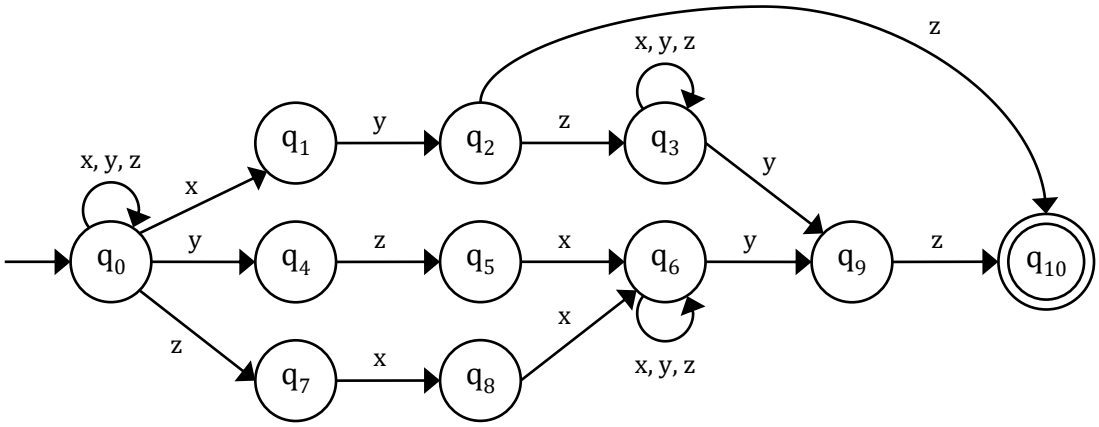
The NFA has been designed below:



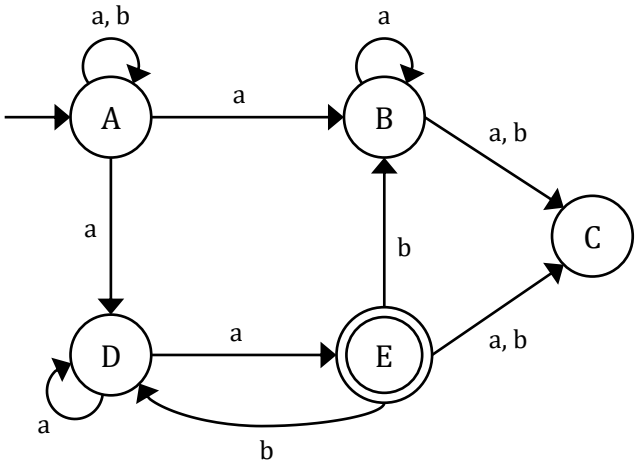
c) L = { w | w contains 'xyz' or 'yzx' or 'zxx' and ends with 'yz' } | ∑ = {x, y, z}

The NFA has been designed below:

[P.T.O]



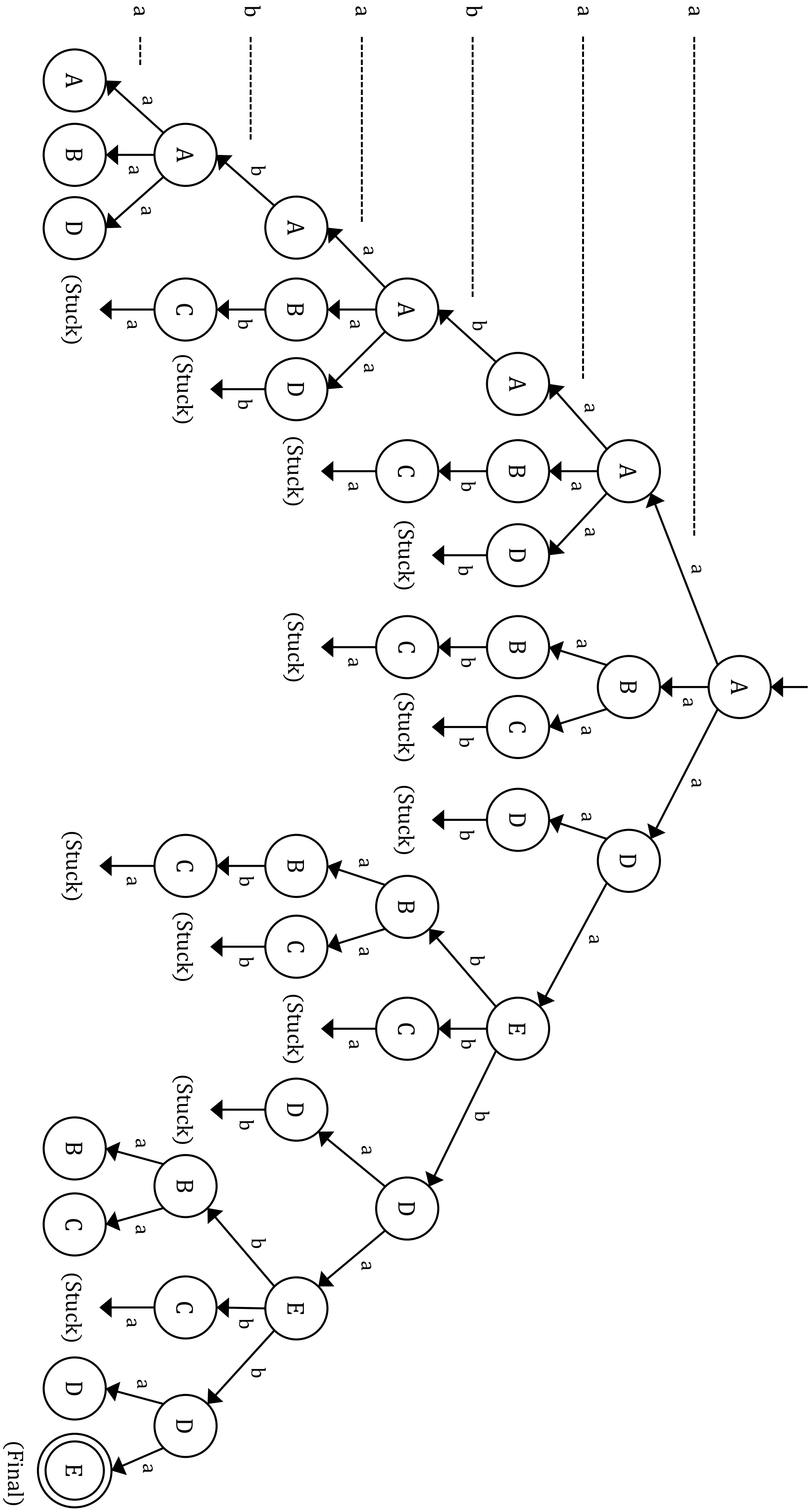
3. Consider the following NFA, and show with help of NFA-tree whether the string "aababa" is accepted.



Solution:

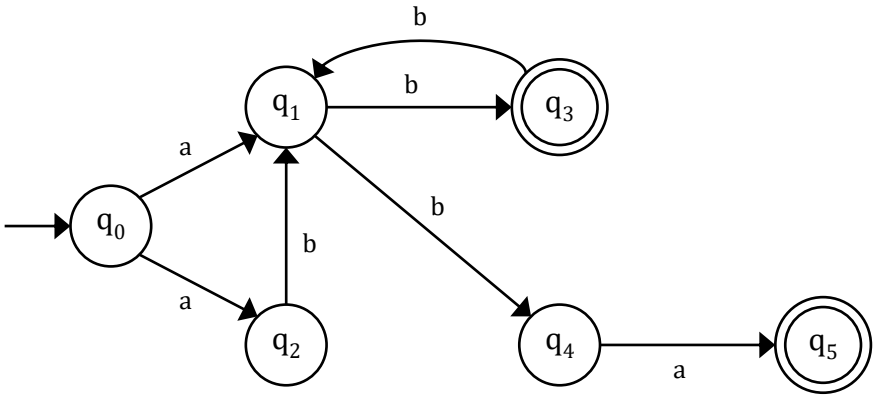
NFA Tree:

[P.T.O]


(Stuck)
(Final)

With help of NFA-tree, we can see the string "aababa" reach the final state.
∴ The string "aababa" is accepted.

4. Convert the following NFA over the alphabet $\sum = \{a, b\}$ to an equivalent DFA



Solution:

Transition Table of the given NFA:

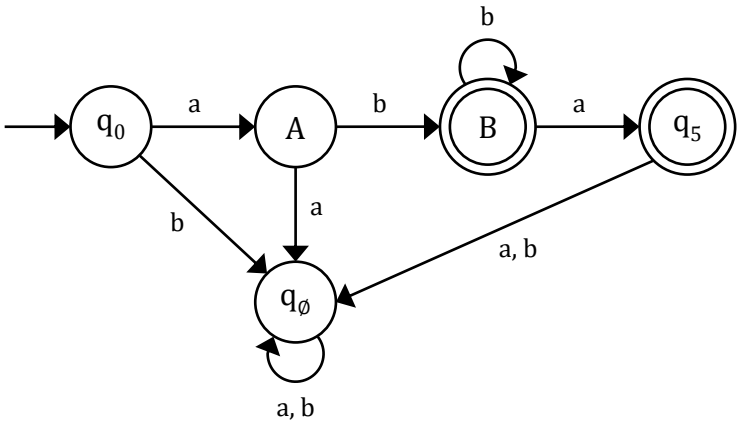
	a	**b**
→ **q0**	q1, q2	∅
q1	∅	q3, q4
q2	∅	q1
***** **q3**	∅	q1
q4	q5	∅
***** **q5**	∅	∅

Transition Table for the Equivalent DFA:

	a	**b**
→ **q0**	{ q1, q2 } = A	{ ∅ } = q∅
A = { q1, q2 }	{ ∅ } = q∅	{ q1, q3, q4 } = B
***** **B = { q1, q3, q4 }**	q5	{ q1, q3, q4 } = B
***** **q5**	{ ∅ } = q∅	{ ∅ } = q∅
q∅ = { ∅ }	{ ∅ } = q∅	{ ∅ } = q∅

Equivalent DFA Diagram:

[P.T.O]



5. a) Convert the following regular expressions to finite automata:

i) (ab)* + (a + ab)* b* (a + b)*

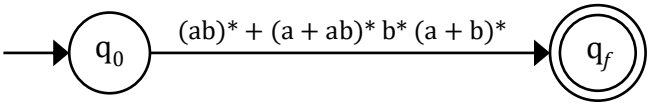
ii) [a + ba (a + b)]* a (ba)* b*

Solution:

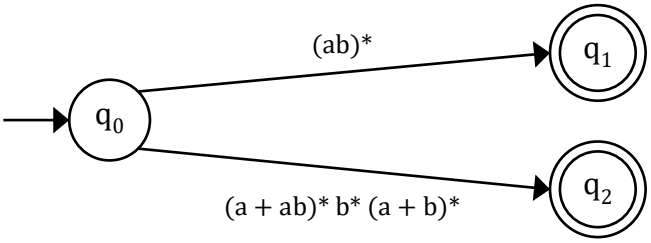
i) (ab)* + (a + ab)* b* (a + b)*

The equivalent finite automata for the following regular expression has been constructed below:

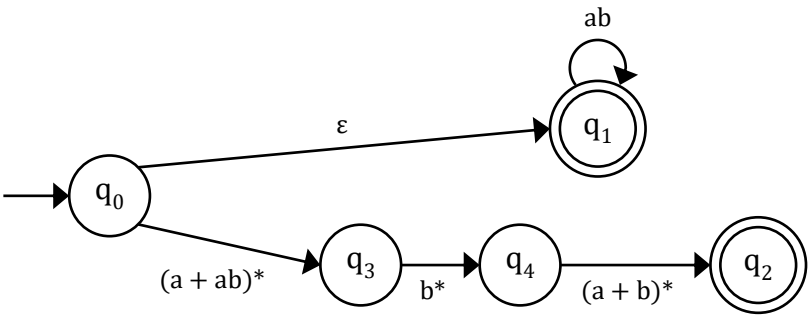
Step 1:



Step 2:

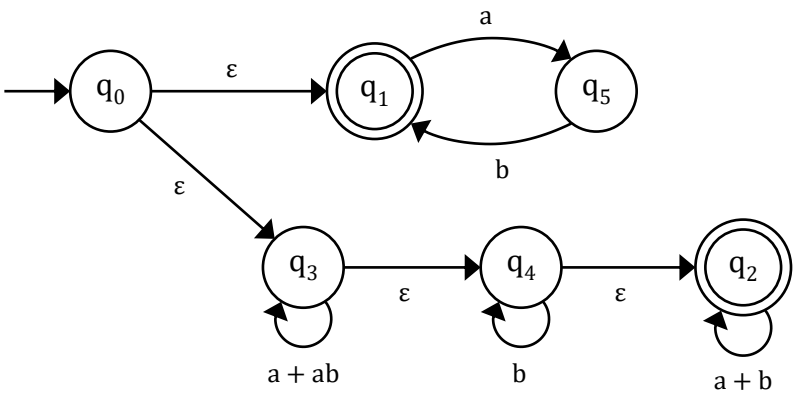


Step 3:

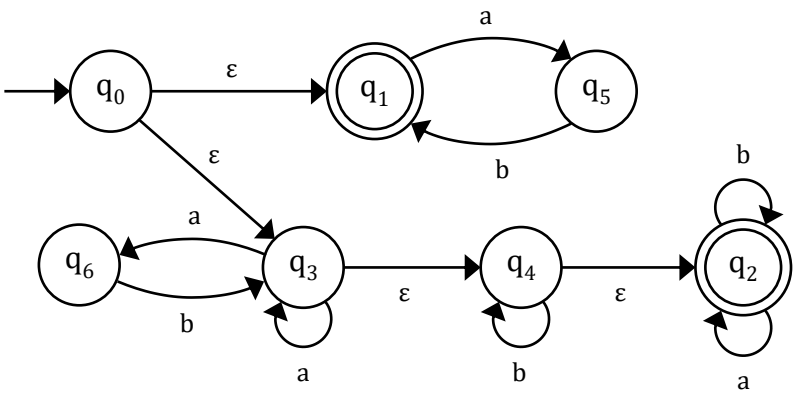


[P.T.O]

Step 4:



Step 5:

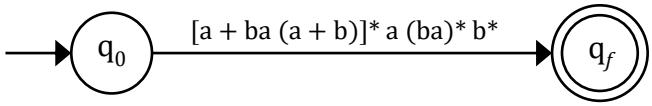


This is our final finite automata for following regular expression.

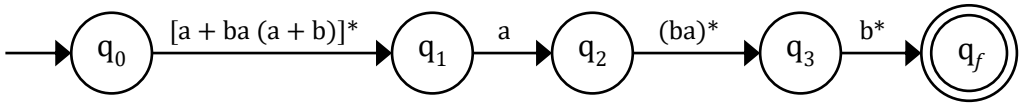
ii) [a + ba (a + b)]* a (ba)* b*

The equivalent finite automata for the following regular expression has been constructed below:

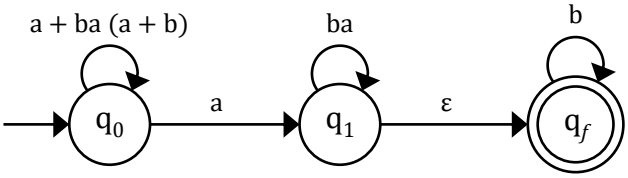
Step 1:



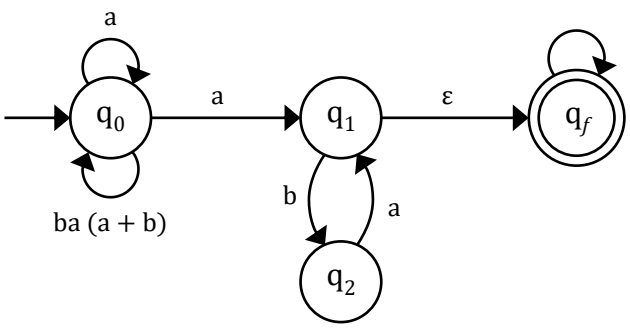
Step 2:



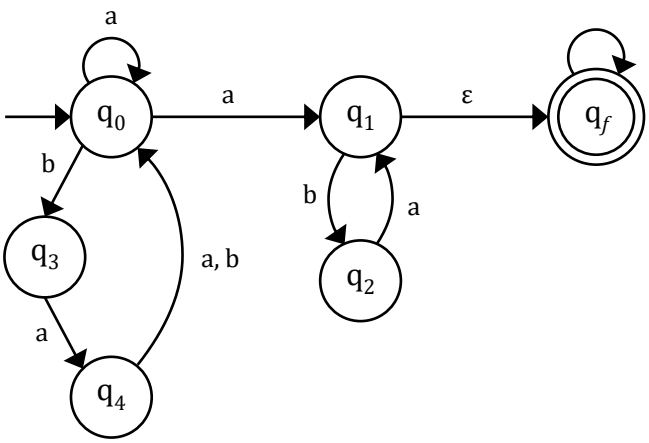
Step 3:



Step 4:



Step 5:



This is our final finite automata for following regular expression.

b) Convert the following regular expressions to finite automata:

i) L = { strings such that the 4th symbol from the right is b over the alphabet {a, b} }

ii) L = { strings such that they start and end with 'a' over the alphabet {a, b, c} }

Solution:

i) L = { strings such that the 4th symbol from the right is b over the alphabet {a, b} }

Regular Expression: (a | b)* a (a | b) (a | b) (a | b)

ii) L = { strings such that they start and end with 'a' over the alphabet {a, b, c} }

Regular Expression: a (a | b | c)* a | a

Summer 2023

1. Design **DFA**'s that accepts the following languages:

a) $L = \{ w \mid w \text{ does not contain '101'} \} \mid \sum = \{0,1\}$

b) $L = \{ w \mid w \text{ starts with an even number of 'a', contains 'ba' and ends with 'baa'} \} \mid \sum = \{a, b\}$

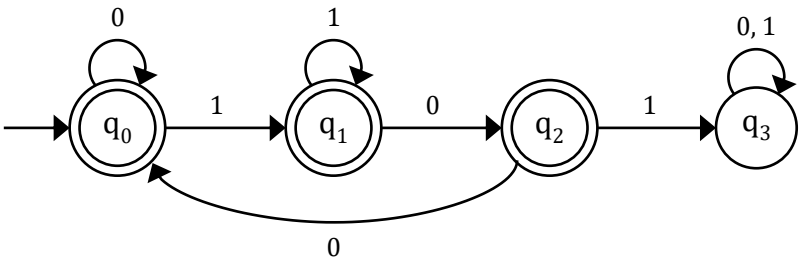
c) $L = \{ w \mid w \text{ is a palindrome with a max length of 3} \} \mid \sum = \{0,1\}$

d) $L = \{ a^i b^j \mid i \geq 0, j \geq 0, \ i + j \text{ is an odd number} \} \mid \sum = \{a, b\}$

Solution:

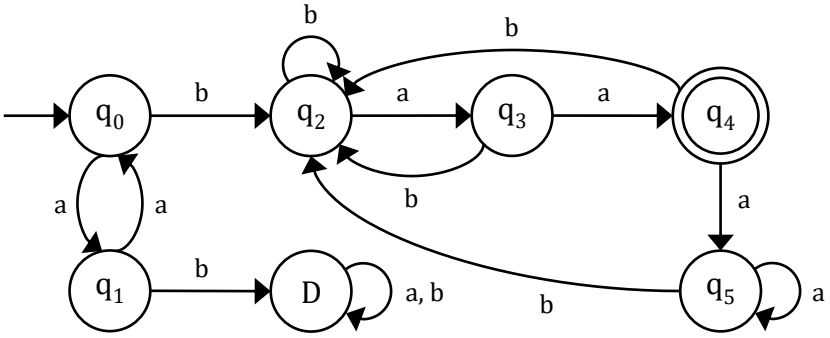
a) $L = \{ w \mid w \text{ does not contain '101'} \} \mid \sum = \{0,1\}$

The DFA has been designed below:



b) $L = \{ w \mid w \text{ starts with an even number of 'a', contains 'ba' and ends with 'baa'} \} \mid \sum = \{a, b\}$

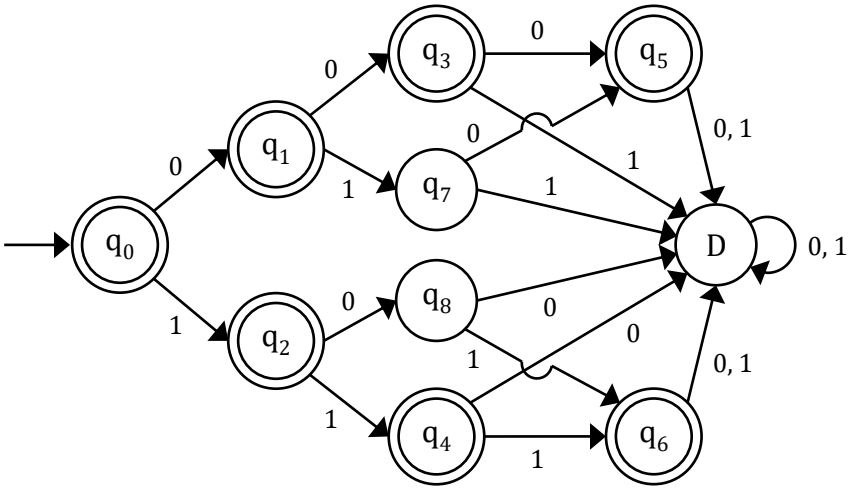
The DFA has been designed below:



c) $L = \{ w \mid w \text{ is a palindrome with a max length of 3} \} \mid \sum = \{0,1\}$

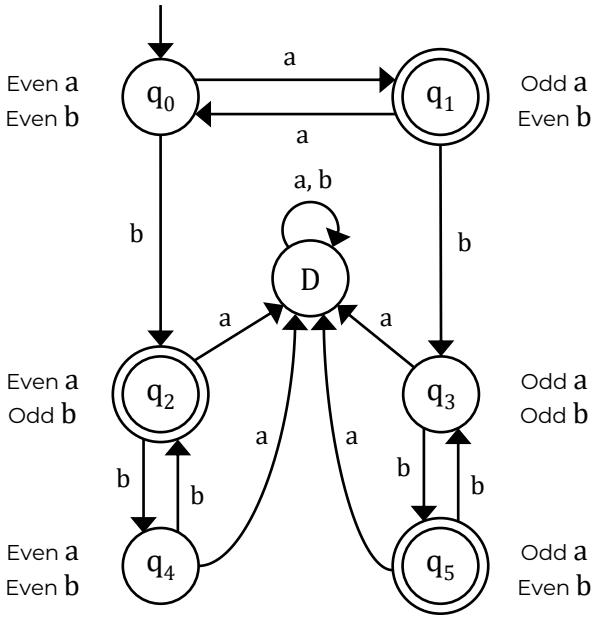
The DFA has been designed below:

[P.T.O]



d) $L = \{ a^i b^j \mid i \geq 0, j \geq 0,\ i + j \text{ is an odd number} \} \mid \sum = \{a, b\}$

The DFA has been designed below:



We Know,

Even + Even = Even
Odd + Odd = Even
Even + Odd = Odd

2. Design **NFA**'s that accepts the following languages:

a) $L = \{ w \mid w$ doesn't start with 'a' or 'b' and contains 'bbc' and ends with 'ab' $\} \mid \sum = \{a, b, c\}$

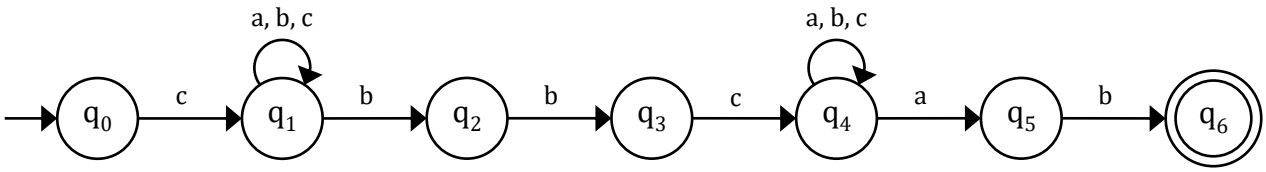
b) $L = \{ w \mid w$ starts with '10' or '21' and contains '220' and ends with '112' $\} \mid \sum = \{0,1,2\}$

c) $L = \{ w \mid w$ starts and ends with either 'xzy' or 'xy' $\} \mid \sum = \{x, y, z\}$

Solution:

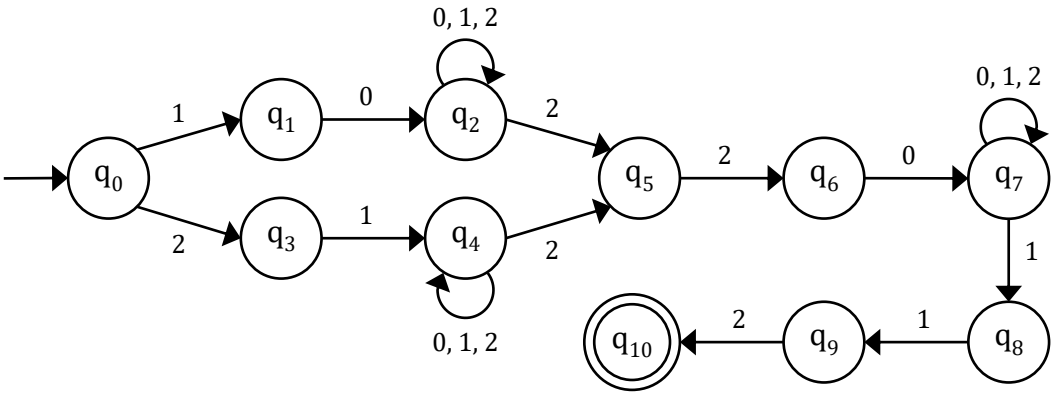
a) $L = \{ w \mid w$ doesn't start with 'a' or 'b' and contains 'bbc' and ends with 'ab' $\} \mid \sum = \{a, b, c\}$

The NFA has been designed below:



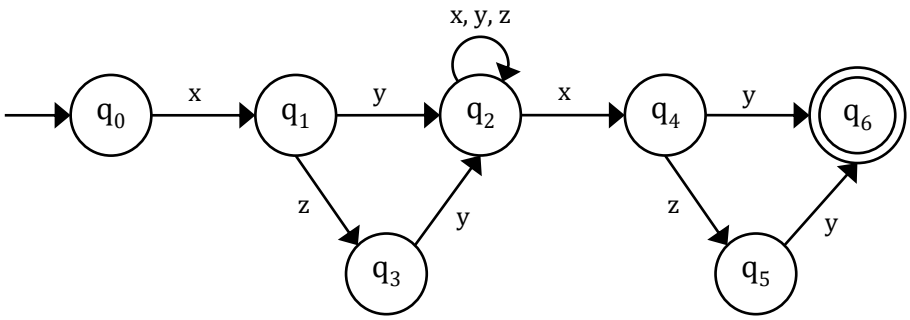
b) L = { w | w starts with '10' or '21' and contains '220' and ends with '112' } | ∑ = {0,1,2}

The NFA has been designed below:

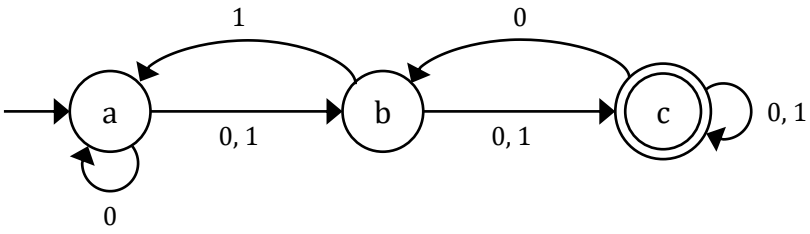


c) L = { w | w starts and ends with either 'xzy' or 'xy' } | ∑ = {x, y, z}

The NFA has been designed below:



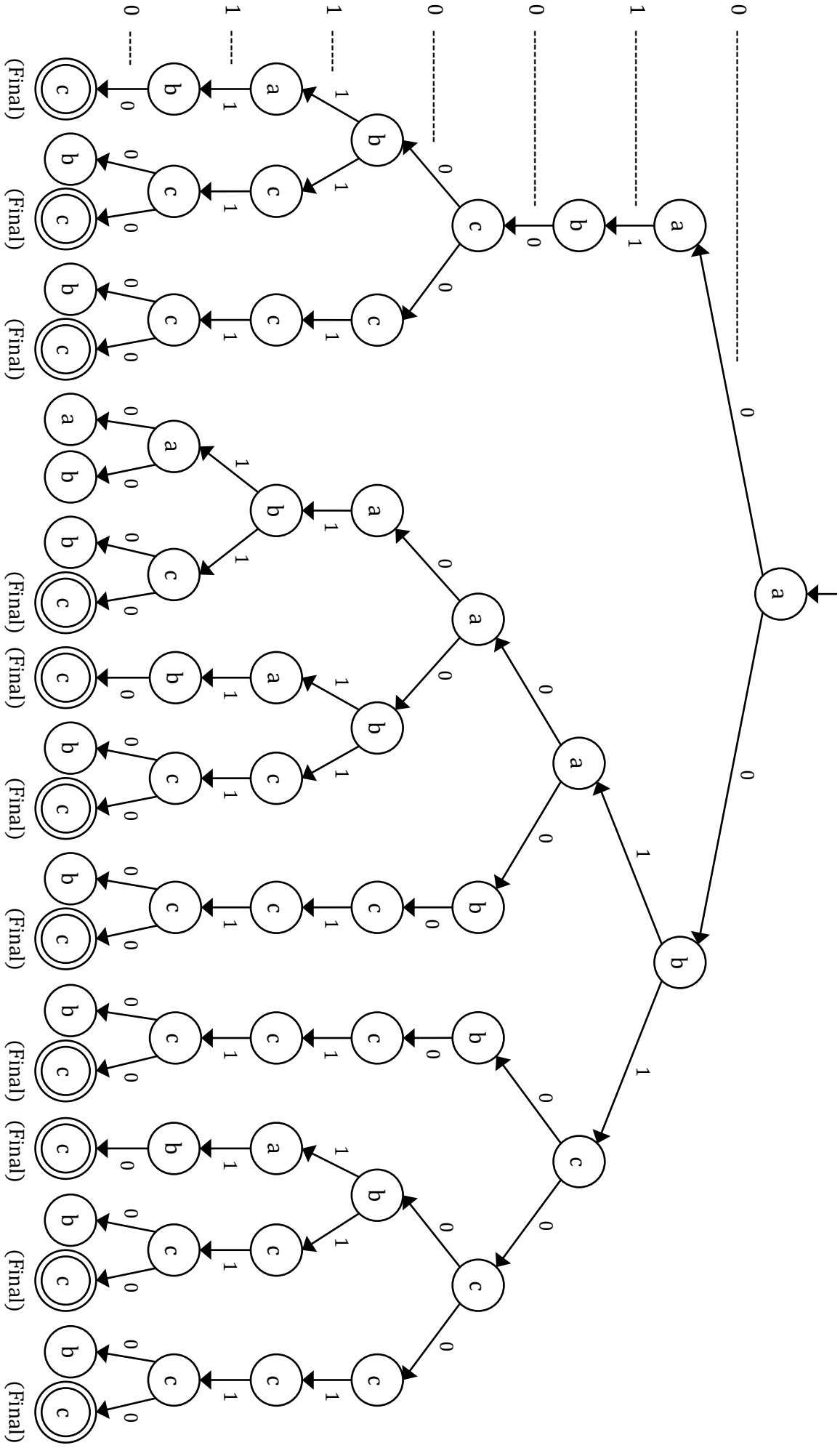
3. Consider the following NFA, and show with help of **NFA-tree** whether the string **"0100110"** is accepted.



Solution:

NFA Tree:

[P.T.O]


a
b
c
(Final)
0
1

With help of NFA-tree, we can see the string "0100110" reach the final state.
∴ The string "0100110" is accepted.

4. Convert the following NFA over the alphabet $\sum = \{1,2,3\}$ to an equivalent DFA including the diagram.

	1	**2**	**3**
→ a	{ a, b, d }	{ a, c }	{ d }
b	∅	{ a, d }	{ a, e }
*** c**	{ a, b, c, d, e }	∅	{ b, c }
d	{ d, e }	{ d }	{ a, e }
*** e**	{ b, e }	∅	∅

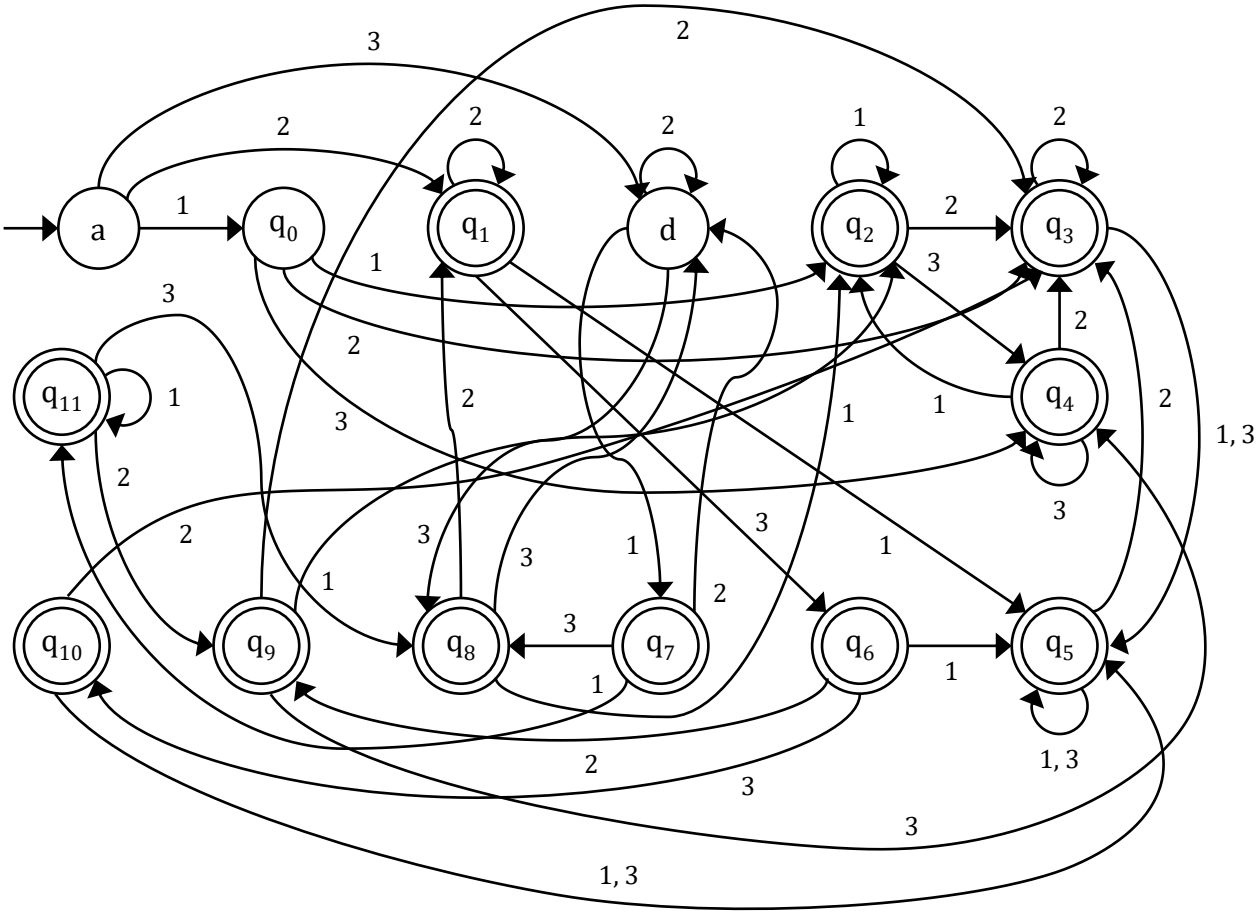
Solution:

Transition Table for the Equivalent DFA:

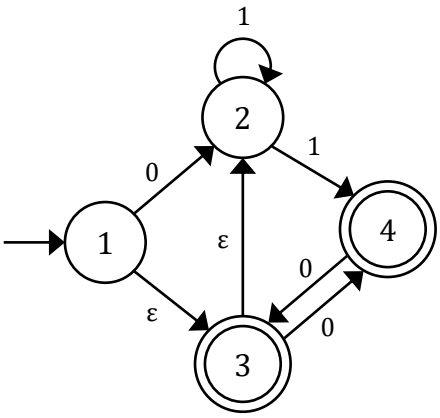
	1	**2**	**3**
→ a	{ a, b, d } = q_0	{ a, c } = q_1	d
q_0 = { a, b, d }	{ a, b, d, e } = q_2	{ a, c, d } = q_3	{ a, d, e } = q_4
*** q_1 = { a, c }**	{ a, b, c, d, e } = q_5	{ a, c } = q_1	{ b, c, d } = q_6
d	{ d, e } = q_7	d	{ a, e } = q_8
*** q_2 = { a, b, d, e }**	{ a, b, d, e } = q_2	{ a, c, d } = q_3	{ a, d, e } = q_4
*** q_3 = { a, c, d }**	{ a, b, c, d, e } = q_5	{ a, c, d } = q_3	{ a, b, c, d, e } = q_5
*** q_4 = { a, d, e }**	{ a, b, d, e } = q_2	{ a, c, d } = q_3	{ a, d, e } = q_4
*** q_5 = { a, b, c, d, e }**	{ a, b, c, d, e } = q_5	{ a, c, d } = q_3	{ a, b, c, d, e } = q_5
*** q_6 = { b, c, d }**	{ a, b, c, d, e } = q_5	{ a, d } = q_9	{ a, b, c, e } = q_{10}
*** q_7 = { d, e }**	{ b, d, e } = q_{11}	d	{ a, e } = q_8
*** q_8 = { a, e }**	{ a, b, d, e } = q_2	{ a, c } = q_1	d
q_9 = { a, d }	{ a, b, d, e } = q_2	{ a, c, d } = q_3	{ a, d, e } = q_4
*** q_{10} = { a, b, c, e }**	{ a, b, c, d, e } = q_5	{ a, c, d } = q_3	{ a, b, c, d, e } = q_5
*** q_{11} = { b, d, e }**	{ b, d, e } = q_{11}	{ a, d } = q_9	{ a, e } = q_8

Equivalent DFA Diagram:

[P.T.O]



5. Convert the following ε-NFA over the alphabet $\sum = \{0,1\}$ to an equivalent DFA.



Solution:

Transition Table of the given ε-NFA:

	0	**1**
→ 1	2	∅
2	∅	2, 4
*** 3**	4	∅
*** 4**	3	∅

ε-Closure of all state of the given ε-NFA:

ε-Closure (1) = { 1, 2, 3 }

ε-Closure (2) = { 2 }

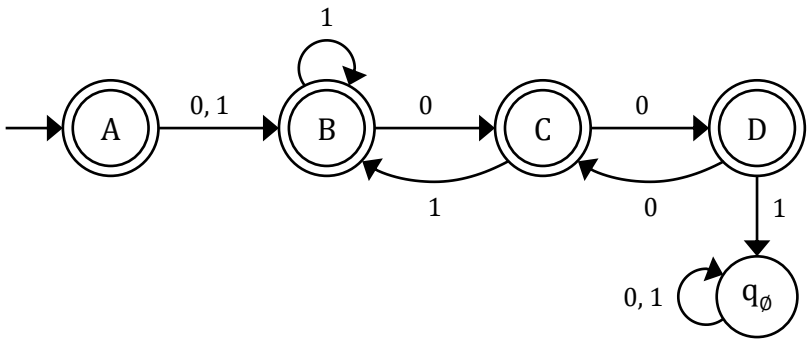
ε-Closure (3) = { 3, 2 }

ε-Closure (4) = { 4 }

Transition Table for the Equivalent DFA:

	0	**1**
→ ∗ **A = { 1, 2, 3 }**	{ 2, 4 } = B	{ 2, 4 } = B
∗ **B = { 2, 4 }**	{ 2, 3 } = C	{ 2, 4 } = B
∗ **C = { 2, 3 }**	{ 4 } = D	{ 2, 4 } = B
∗ **D = { 4 }**	{ 2, 3 } = C	{ Ø } = $q_ø$
$q_ø$ = { Ø }	{ Ø } = $q_ø$	{ Ø } = $q_ø$

Equivalent DFA Diagram:



1. Design DFAs that accepts the following languages:

   a) L = contains even number of 'a' and ends with 'bc' | $\sum = \{a, b, c\}$

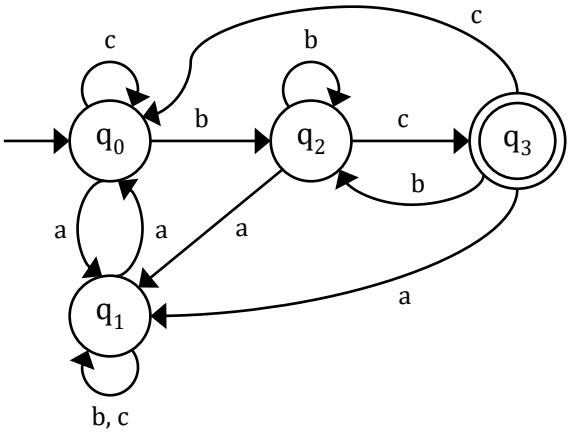
   b) L = does not contain 'mnm' | $\sum = \{m, n, w\}$

   c) L = starts with 'gh' and contains 'kgh' and ends with 'gh' | $\sum = \{g, h, k\}$

   Solution:

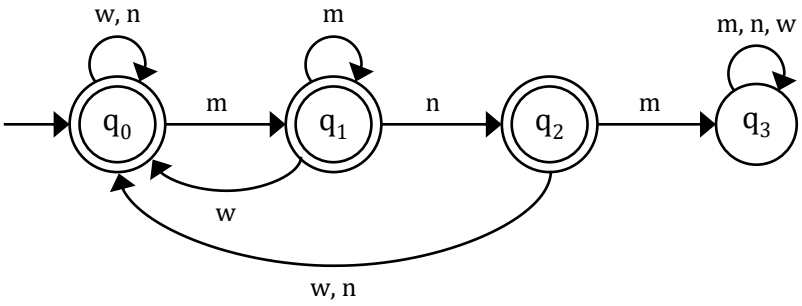
   a) L = contains even number of 'a' and ends with 'bc' | $\sum = \{a, b, c\}$

   The DFA has been designed below:



   b) L = does not contain 'mnm' | $\sum = \{m, n, w\}$

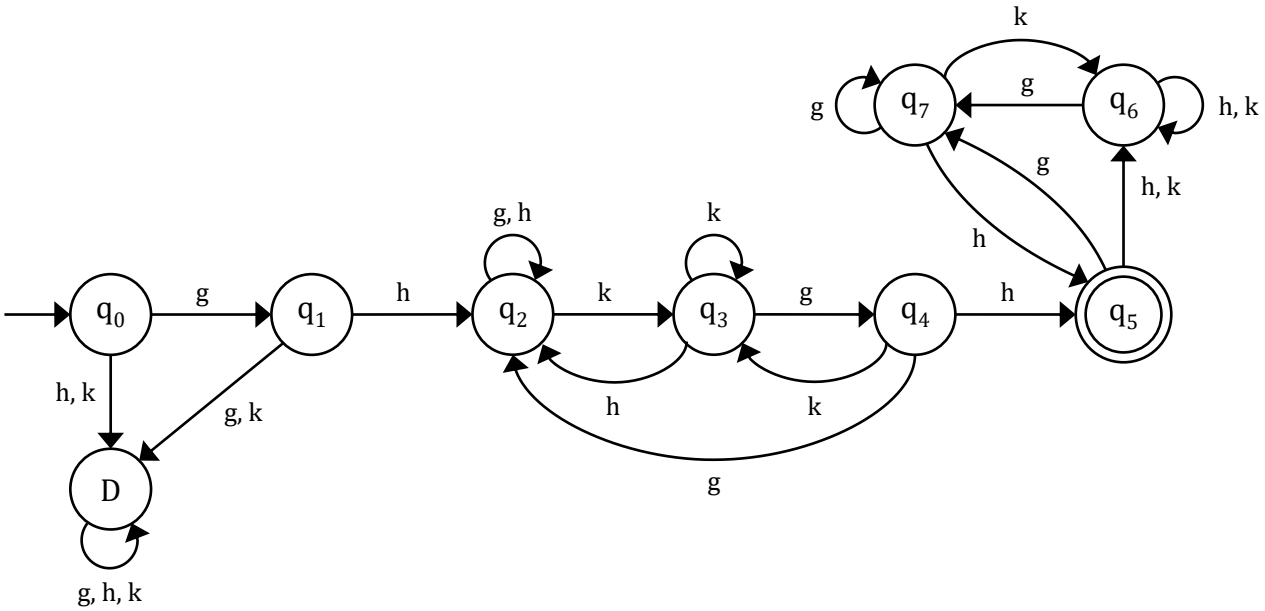
   The DFA has been designed below:



   c) L = starts with 'gh' and contains 'kgh' and ends with 'gh' | $\sum = \{g, h, k\}$

   The DFA has been designed below:

   [P.T.O]



2. Design NFAs that accepts the following languages:

a) L = starts with 'p', and contains 'rqp', and ends with 'qr' | $\sum = \{p, q, r\}$

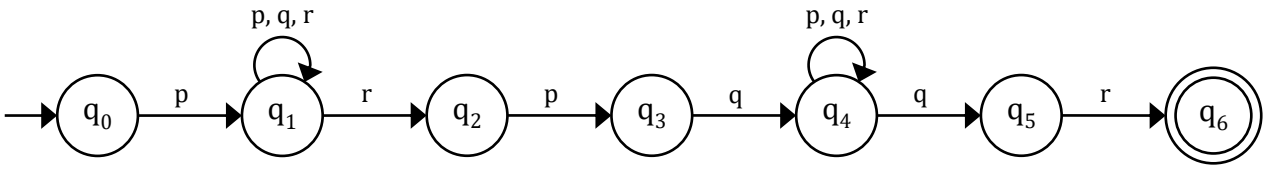
b) L = starts with '11' or '21' and contains '210' and ends with '101' | $\sum = \{0,1,2\}$

c) L = starts with 'xyz' and contains 'yyz' or 'zyx' and ends with 'zy' | $\sum = \{x, y, z\}$

Solution:

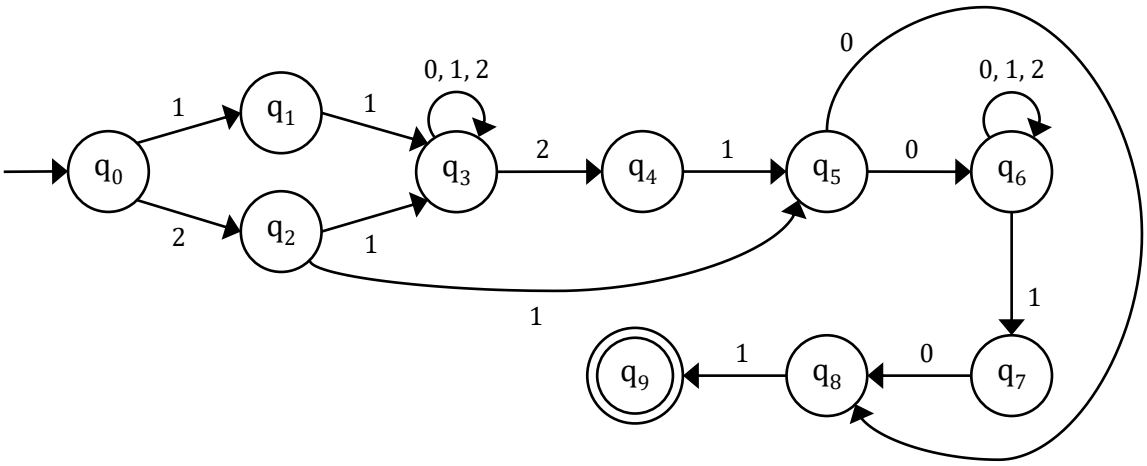
a) L = starts with 'p', and contains 'rqp', and ends with 'qr' | $\sum = \{p, q, r\}$

The NFA has been designed below:



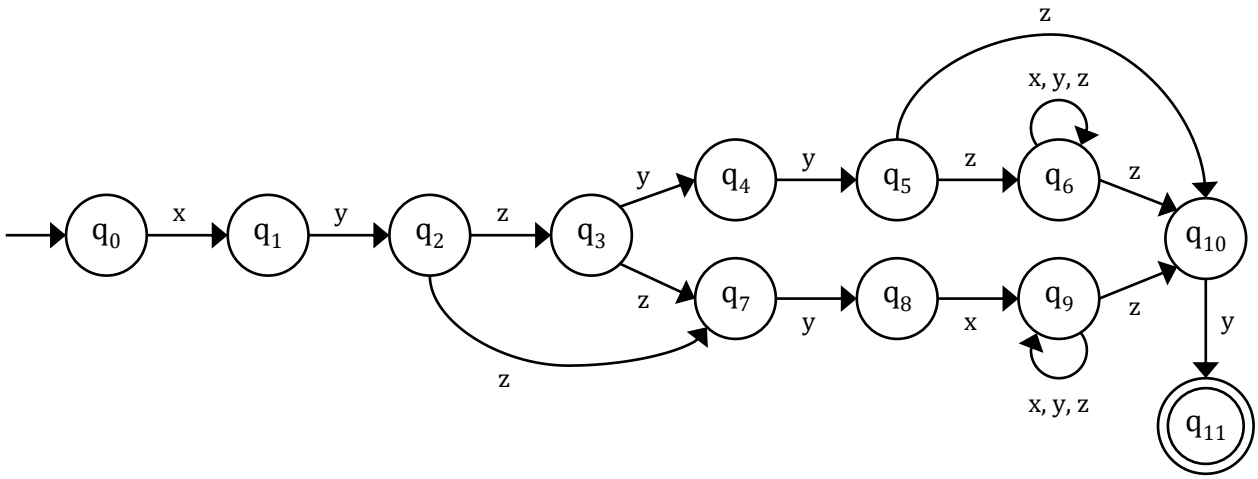
b) L = starts with '11' or '21' and contains '210' and ends with '101' | $\sum = \{0,1,2\}$

The NFA has been designed below:

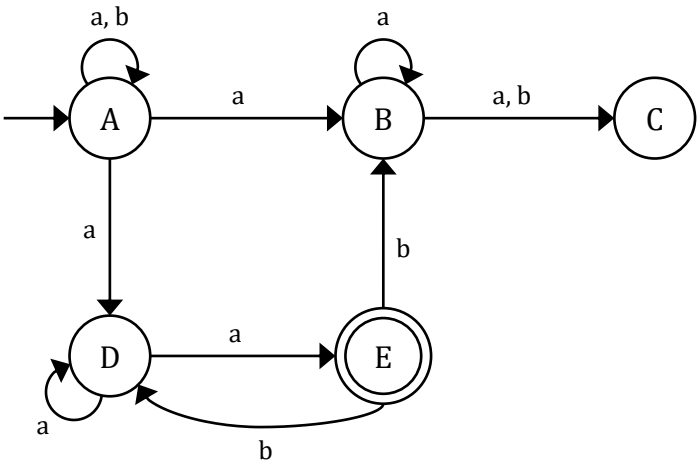


c) L = starts with 'xyz' and contains 'yyz' or 'zyx' and ends with 'zy' | $\sum = \{x, y, z\}$

The NFA has been designed below:



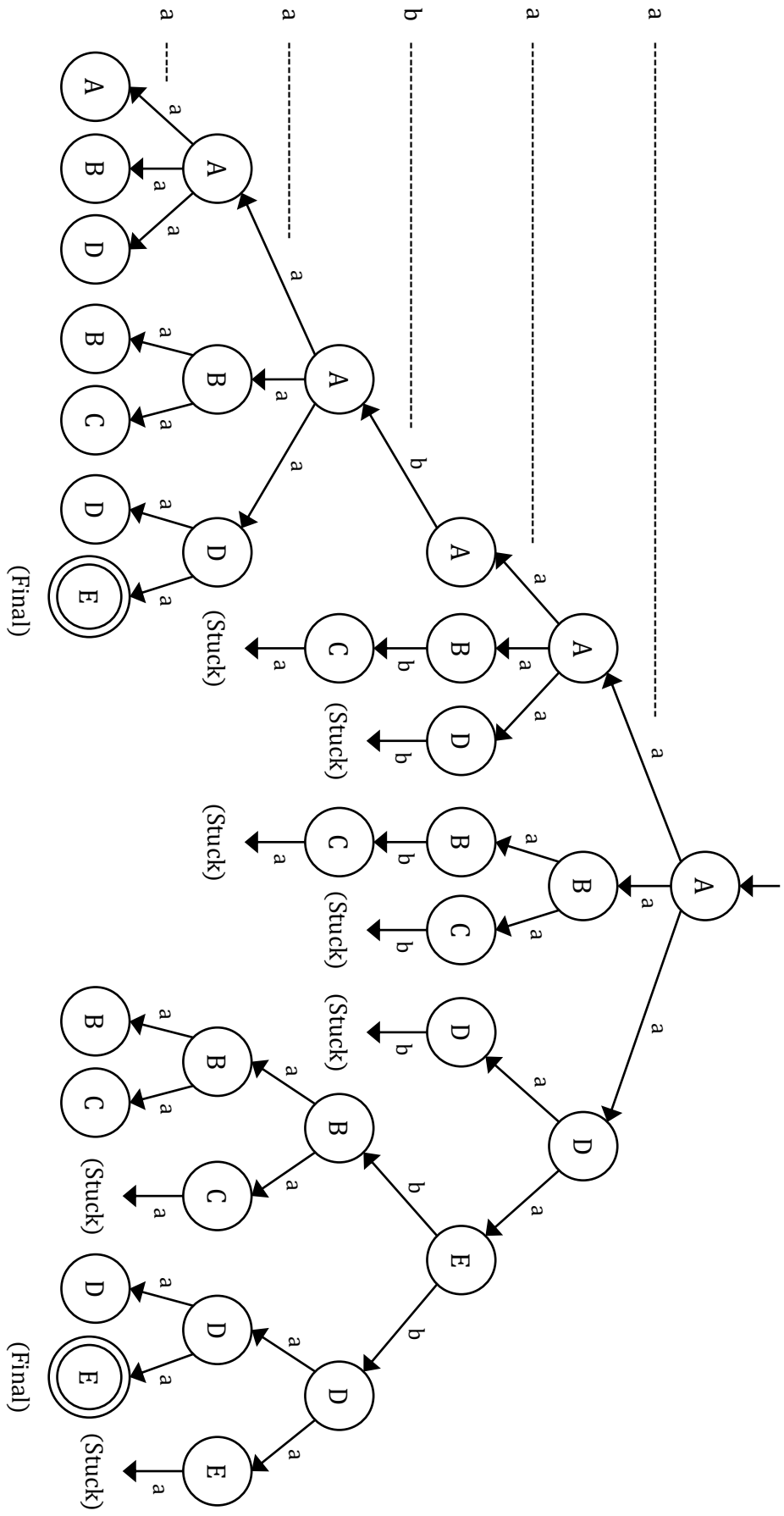
3. Consider the following NFA, and show with help of NFA-tree whether the string "aabaa" is accepted or not.



Solution:

NFA Tree:

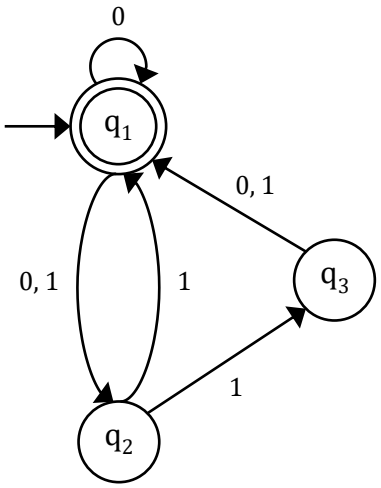
[P.T.O]



With help of NFA-tree, we can see the string “aabaa” reach the final state.

∴ The string “aabaa” is accepted.

4. Convert the following **NFA** over the alphabet $\sum = \{0,1\}$ to an equivalent **DFA**.



Solution:

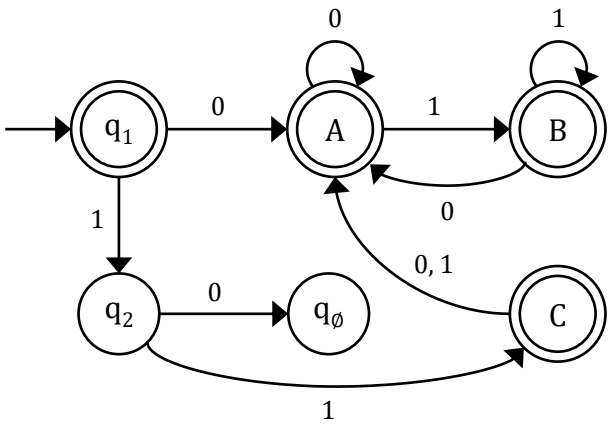
Transition Table of the given NFA:

	0	**1**
→ * **q1**	q1, q2	q2
q2	∅	q1, q3
q3	q1	q1

Transition Table for the Equivalent DFA:

	0	**1**
→ * **q1**	{ q1, q2 } = A	q2
* **A = { q1, q2 }**	{ q1, q2 } = A	{ q1, q2, q3 } = B
q2	{ ∅ } = q∅	{ q1, q3 } = C
* **B = { q1, q2, q3 }**	{ q1, q2 } = A	{ q1, q2, q3 } = B
* **C = { q1, q3 }**	{ q1, q2 } = A	{ q1, q2 } = A

Equivalent DFA Diagram:



5. Design Regular Expression for the following languages where $\sum = \{a, b\}$

a) All strings w having even length strings and starting with **a** or odd length strings starting with **b**.

b) All strings w which begins and ends with **b**.

c) All strings w where every **a** is followed by at least one **b**.

Solution:

a) All strings w having even length strings and starting with **a** or odd length strings starting with **b**.

Regular Expression: $a\ (a \mid b)\ ((a \mid b)\ (a \mid b))^*\ \mid\ b\ ((a \mid b)\ (a \mid b))^*$

b) All strings w which begins and ends with **b**.

Regular Expression: $b\ (a \mid b)^*\ b\ \mid\ b$

c) All strings w where every **a** is followed by at least one **b**.

Regular Expression: $b^*\ (ab\ b^*)^*$

1. Design DFAs that accepts the following languages:

a) L = { w | w starts with 'ab' and contains 'bba' and ends with 'bb' } | ∑ = {a, b}

b) L = { w | w contains the set of all strings that has length exactly 3 and its third symbol is from the left side is 'a' } | ∑ = {a, b}

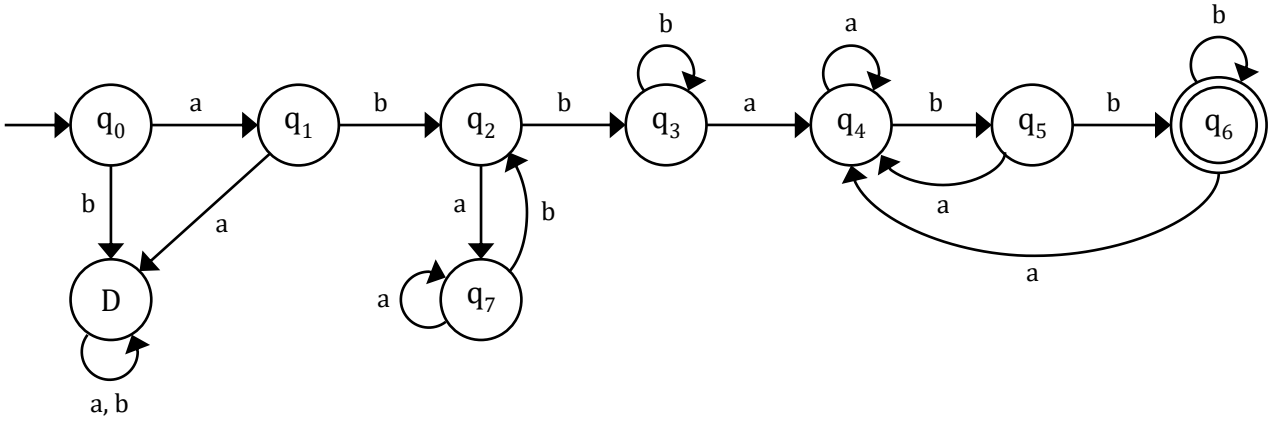
c) L = { w | w contains the set of all strings that has neither '00' nor '11' as substring } | ∑ = {0,1,2}

d) L = { w | w contains the set of all strings whose length always returns remainder 2 when divided by 4 } | ∑ = {0,1}

Solution:

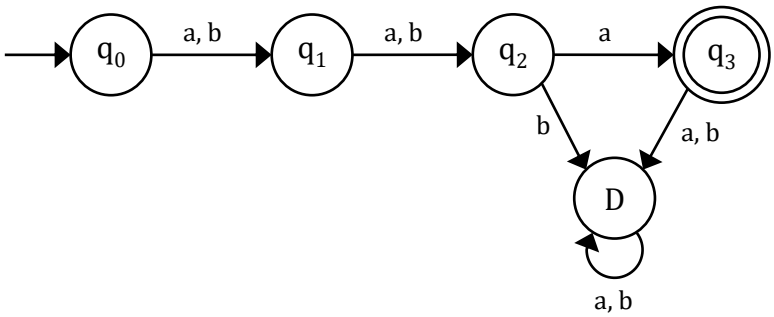
a) L = { w | w starts with 'ab' and contains 'bba' and ends with 'bb' } | ∑ = {a, b}

The DFA has been designed below:



b) L = { w | w contains the set of all strings that has length exactly 3 and its third symbol is from the left side is 'a' } | ∑ = {a, b}

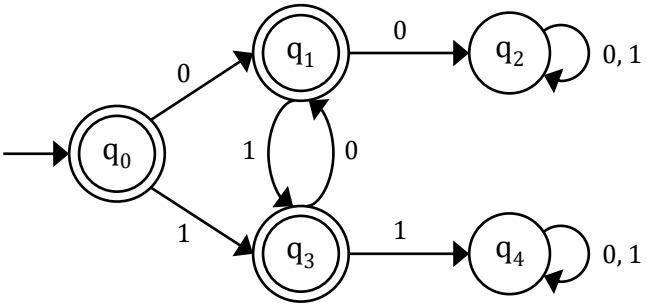
The DFA has been designed below:



c) L = { w | w contains the set of all strings that has neither '00' nor '11' as substring } | ∑ = {0,1,2}

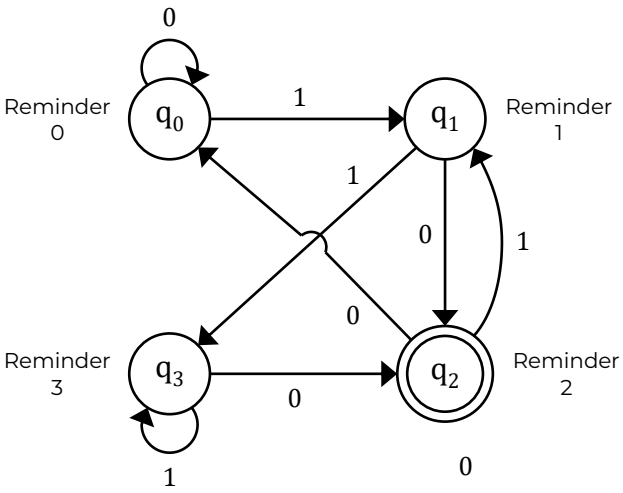
[P.T.O]

The DFA has been designed below:



d) L = { w | w contains the set of all strings whose length always returns remainder 2 when divided by 4 } | $\sum = \{0,1,2\}$

The DFA has been designed below:



Decimal	Binary	Reminder
0	0	0
1	1	1
2	10	2
3	11	3
4	100	0
5	101	1
6	110	2
7	111	3
8	1000	0

2. Design NFAs that accepts the following languages:

a) L = ends with 'x' and contains 'yxz' and starts with 'xy' | $\sum = \{x, y, z\}$

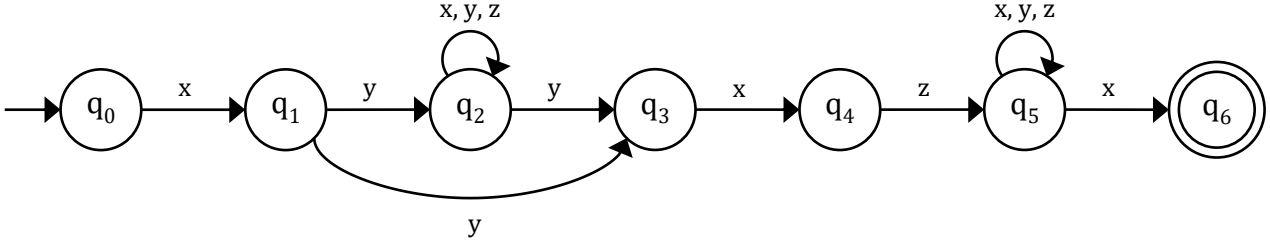
b) L = starts with 'pq' or 'qr' and contains 'pqp' or 'qrr' and ends with 'qqr' | $\sum = \{p, q, r\}$

c) L = starts with '211' and contains '112' or '321' and ends with '1' | $\sum = \{1,2,3\}$

Solution:

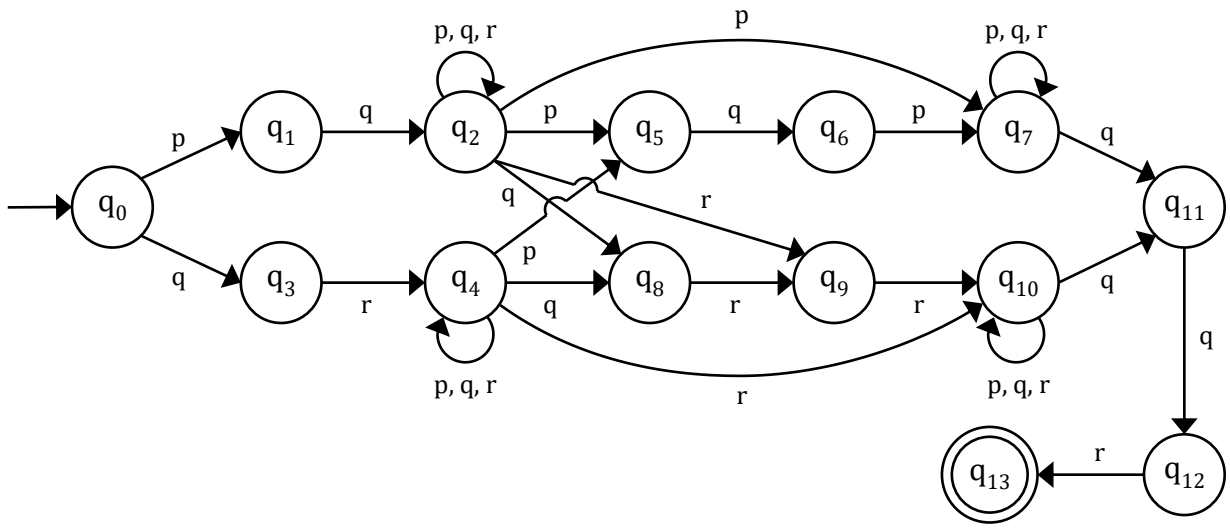
a) L = ends with 'x' and contains 'yxz' and starts with 'xy' | $\sum = \{x, y, z\}$

The NFA has been designed below:



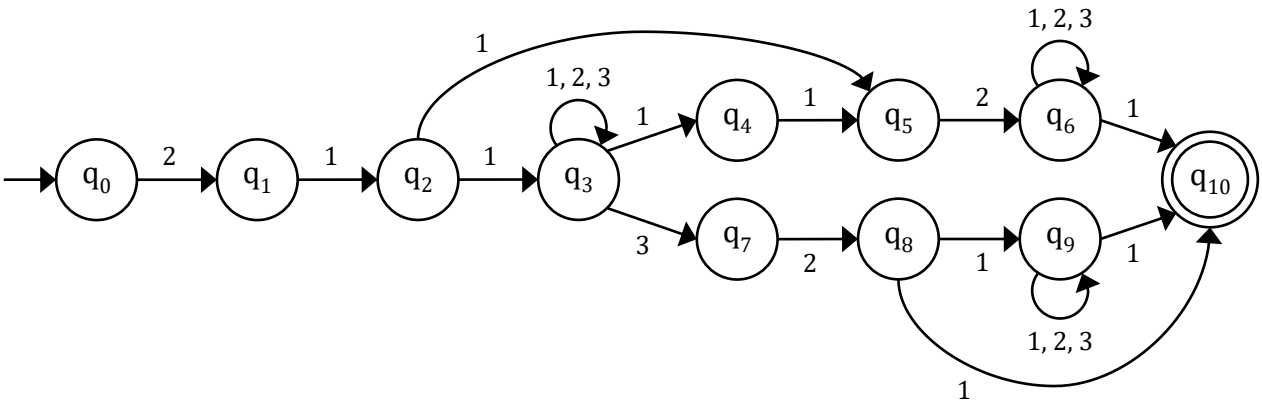
b) L = starts with 'pq' or 'qr' and contains 'pqp' or 'qrr' and ends with 'qqr' | $\sum = \{p, q, r\}$

The NFA has been designed below:

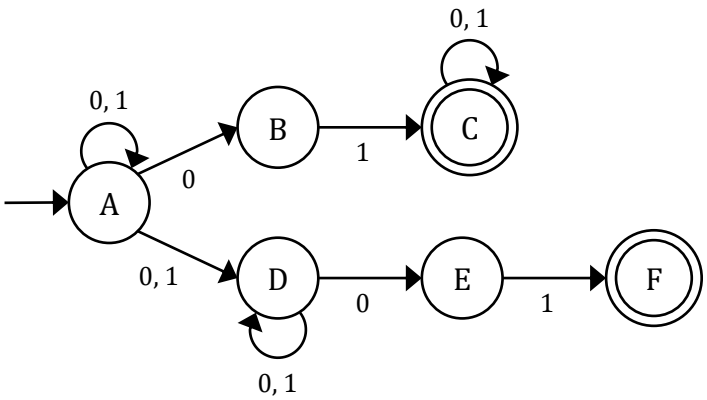


c) L = starts with '211' and contains '112' or '321' and ends with '1' | $\sum = \{1,2,3\}$

The NFA has been designed below:



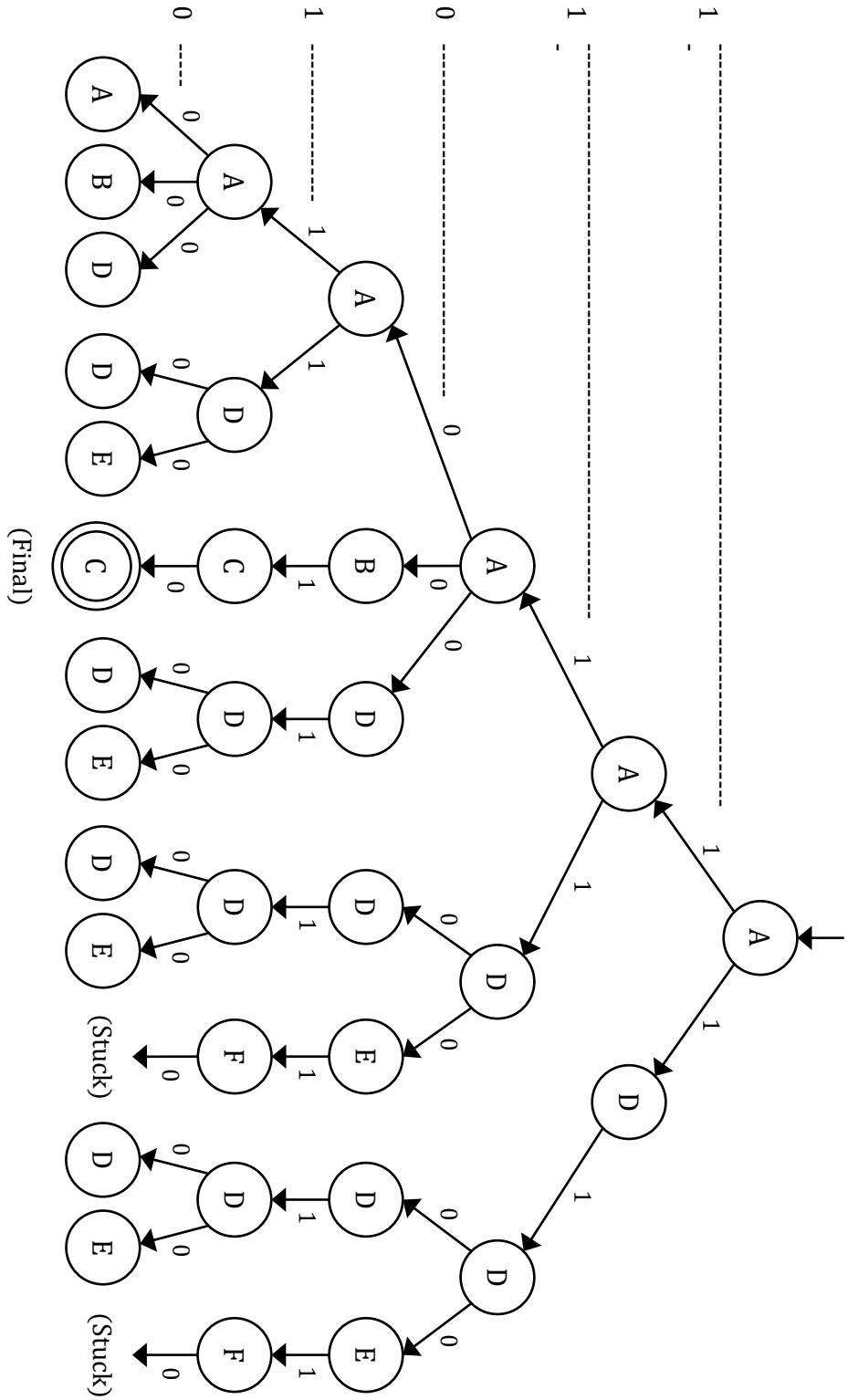
3. Consider the following NFA, and show with help of NFA-tree whether the string "11010" is accepted or not.



Solution:

NFA Tree:

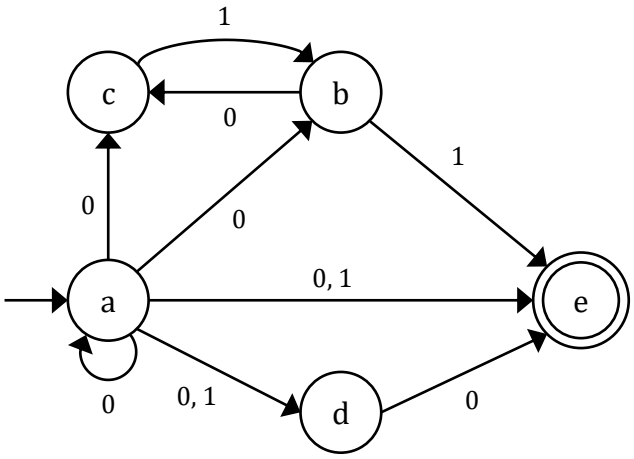
[P.T.O]



With help of NFA-tree, we can see the string "aabaa" reach the final state.
∴ The string "11010" is accepted.

4. Convert the following NFA over the alphabet $\sum = \{0,1\}$ to an equivalent DFA.

[P.T.O]



Solution:

Transition Table of the given NFA:

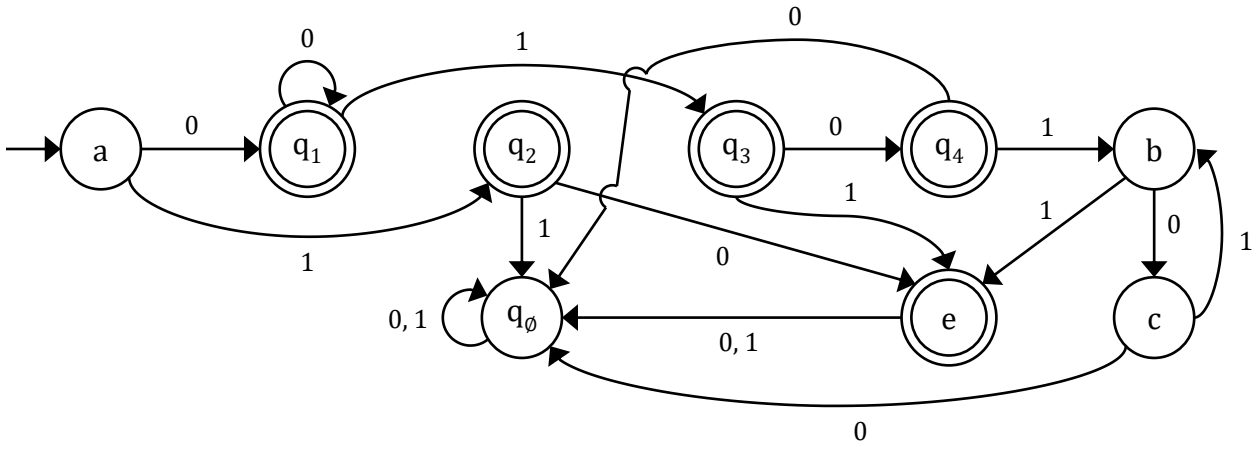
	0	**1**
→ a	a, b, c, d, e	d, e
b	c	e
c	∅	b
d	e	∅
*** e**	∅	∅

Transition Table for the Equivalent DFA:

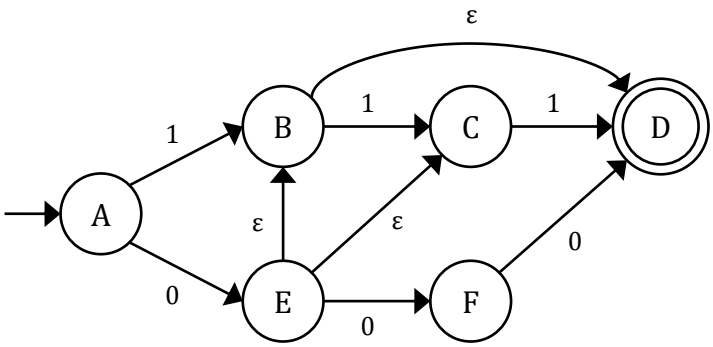
	0	**1**
→ a	{ a, b, c, d, e } = q_1	{ d, e } = q_2
*** q_1 = { a, b, c, d, e }**	{ a, b, c, d, e } = q_1	{ b, d, e } = q_3
*** q_2 = { d, e }**	e	{ ∅ } = $q_\emptyset$
*** q_3 = { b, d, e }**	{ c, e } = q_4	e
*** e**	{ ∅ } = $q_\emptyset$	{ ∅ } = $q_\emptyset$
*** q_4 = { c, e }**	{ ∅ } = $q_\emptyset$	b
b	c	e
c	{ ∅ } = $q_\emptyset$	b
*** $q_\emptyset$ = { ∅ }**	{ ∅ } = $q_\emptyset$	{ ∅ } = $q_\emptyset$

Equivalent DFA Diagram:

[P.T.O]



5. Convert the following ε-NFA over the alphabet $\sum = \{0,1\}$ to an equivalent DFA.



Solution:

Transition Table of the given ε-NFA:

	0	**1**
→ A	E	B
B	Ø	C
C	Ø	D
*** D**	Ø	Ø
E	F	Ø
F	D	Ø

ε-Closure of all state of the given ε-NFA:

$\varepsilon\text{-Closure}(A) = \{ A \}$

$\varepsilon\text{-Closure}(B) = \{ B, D \}$

$\varepsilon\text{-Closure}(C) = \{ C \}$

$\varepsilon\text{-Closure}(D) = \{ D \}$

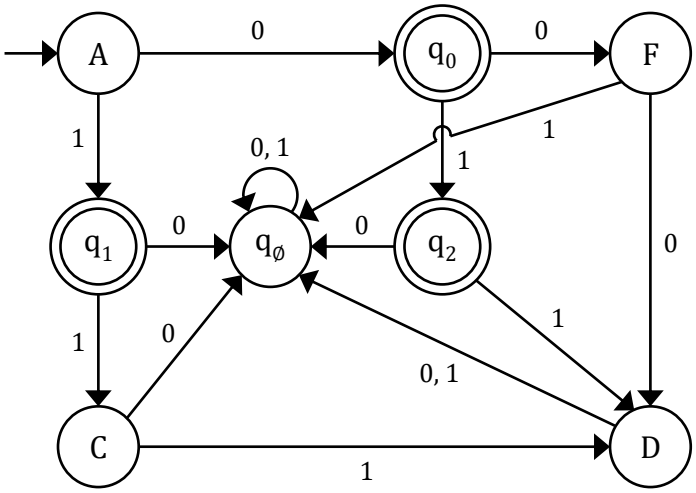
$\varepsilon\text{-Closure}(E) = \{ E, B, C, D \}$

$\varepsilon\text{-Closure}(F) = \{ F \}$

Transition Table for the Equivalent DFA:

	0	**1**
→ A	{ B, C, D, E } = q_0	{ B, D } = q_1
*** q_0 = { B, C, D, E }**	F	{ C, D } = q_2
*** q_1 = { B, D }**	{ ∅ } = $q_\emptyset$	C
F	D	{ ∅ } = $q_\emptyset$
*** q_2 = { C, D }**	{ ∅ } = $q_\emptyset$	D
C	{ ∅ } = $q_\emptyset$	D
*** D**	{ ∅ } = $q_\emptyset$	{ ∅ } = $q_\emptyset$
$q_\emptyset$ = { ∅ }	{ ∅ } = $q_\emptyset$	{ ∅ } = $q_\emptyset$

Equivalent DFA Diagram:



Summer 2022

1. Design DFAs that accepts the following languages:

a) L = { w | w starts and ends with different symbols and the length of w is even } | $\sum = \{0,1\}$

b) L = { w | w contains at least two 'a's and at most one 'b' } | $\sum = \{a, b\}$

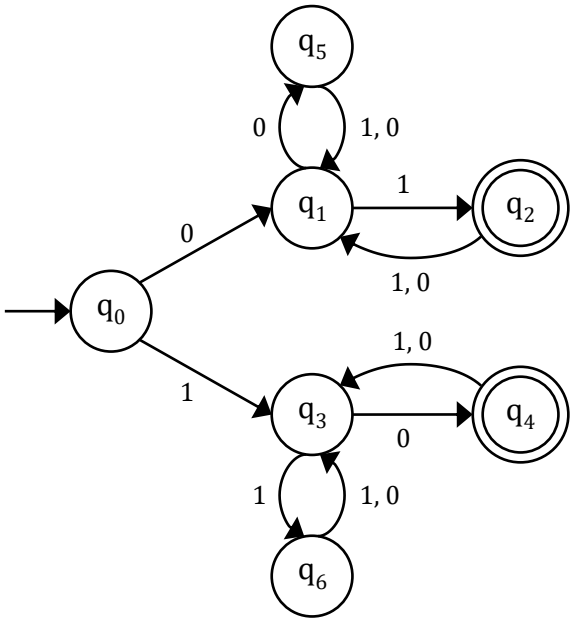
c) L = { w | w contains even number of 0's or odd number of 2's. } over $\sum = \{0,1,2\}$

d) L = { w | w contains all the binary number which is divisible by 3 or ends with with '011' } | $\sum = \{0,1\}$

Solution:

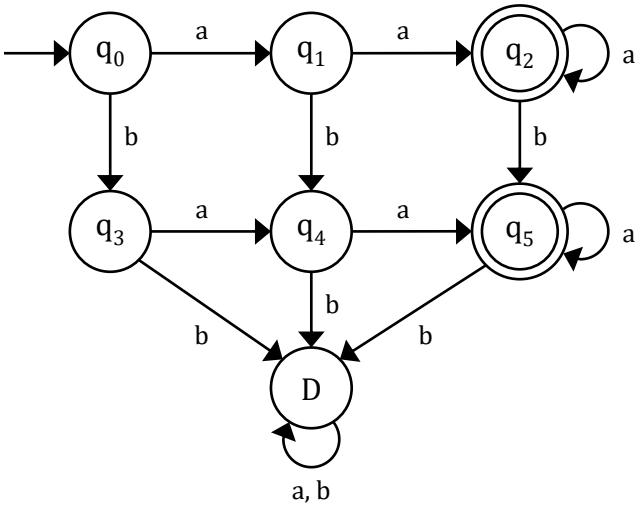
a) L = { w | w starts and ends with different symbols and the length of w is even } | $\sum = \{0,1\}$

The DFA has been designed below:



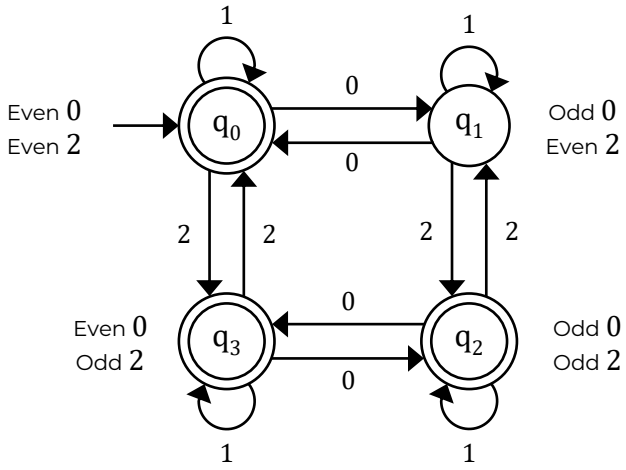
b) L = { w | w contains at least two 'a's and at most one 'b' } | $\sum = \{a, b\}$

The DFA has been designed below:



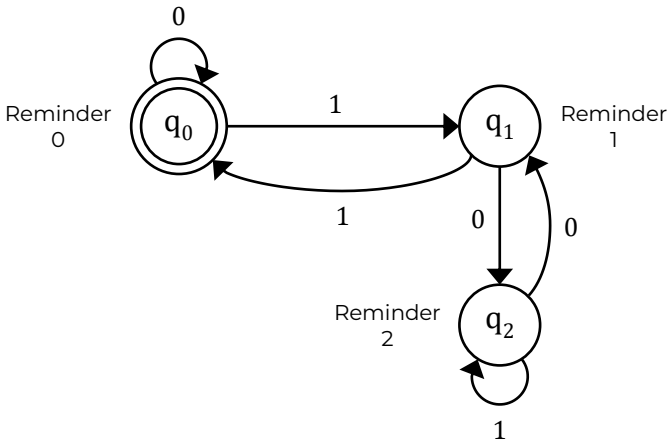
c) L = { w | w contains even number of 0's or odd number of 2's. } over $\sum = \{0,1,2\}$

The DFA has been designed below:



d) L = { w | w contains all the binary number which is divisible by 3 or ends with with '011' }

The DFA has been designed below:



Decimal	Binary	Reminder
0	0	0
1	1	1
2	10	2
3	11	0
4	100	1
5	101	2
6	110	0

2. Design NFAs that accepts the following languages:

a) L = ends with 'b' and contains 'bbcb' and starts with 'aacd' | $\sum = \{a, b, c, d\}$

b) L = contains 'bba' or 'abb' or 'acc' and starts with 'ab' or 'bc' | $\sum = \{a, b, c\}$

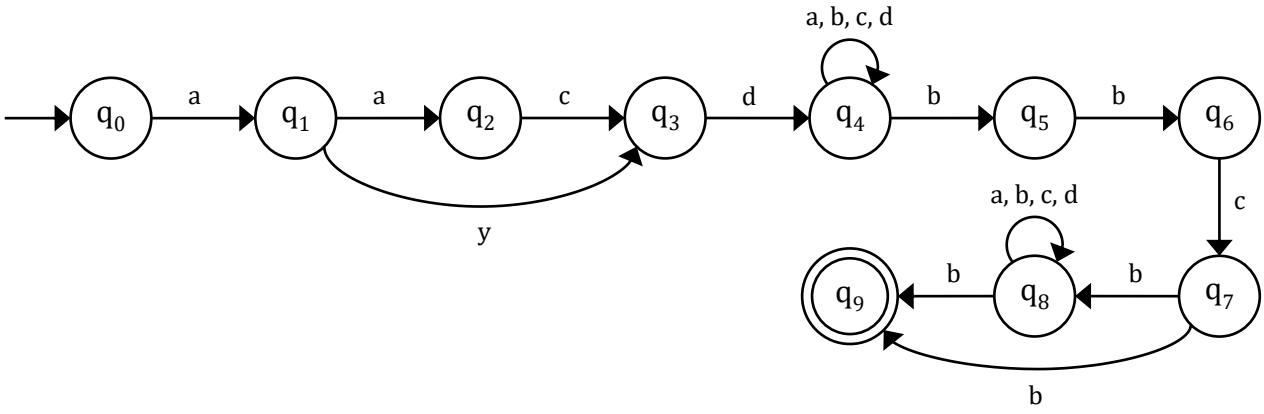
c) L = starts with '121' and contains '212' or '312' and ends with '2' | $\sum = \{1,2,3\}$

Solution:

a) L = ends with 'b' and contains 'bbcb' and starts with 'aacd' | $\sum = \{a, b, c, d\}$

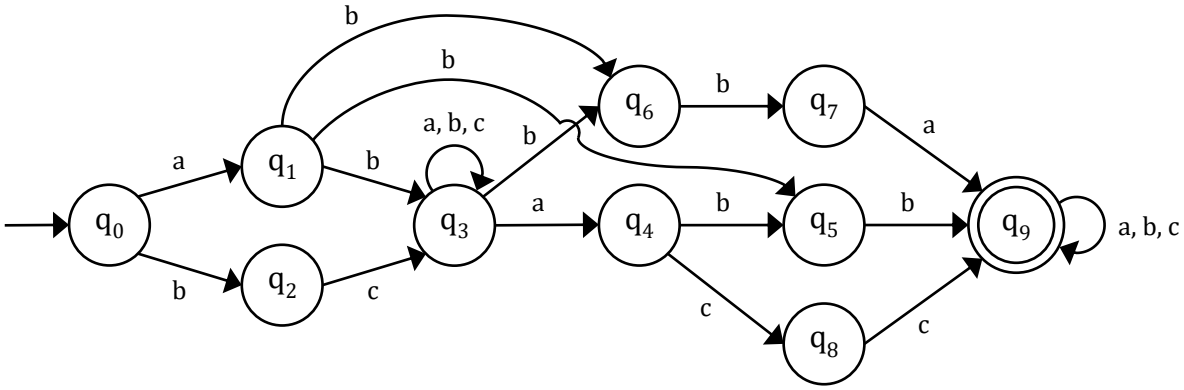
The NFA has been designed below:

[P.T.O]



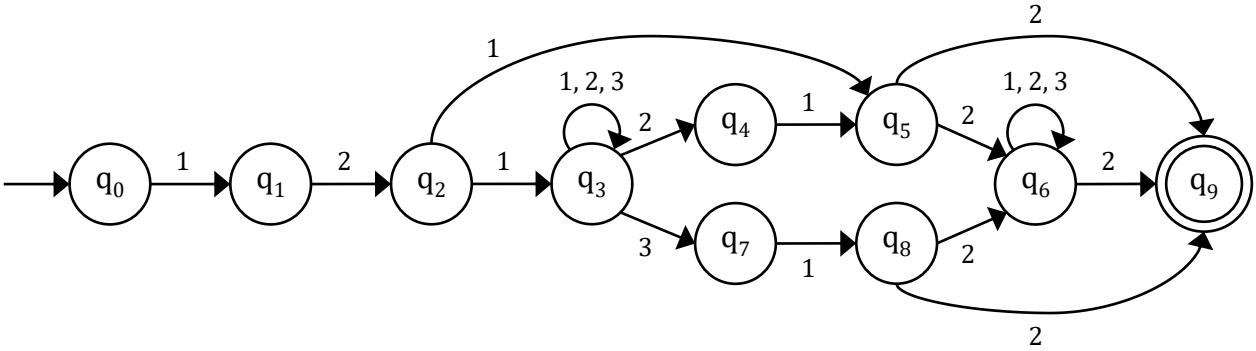
b) L = contains 'bba' or 'abb' or 'acc' and starts with 'ab' or 'bc' | $\sum = \{a, b, c\}$

The NFA has been designed below:

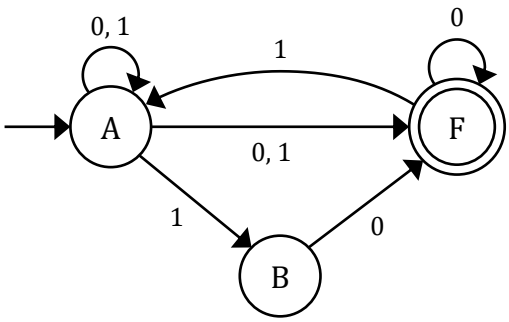


c) L = starts with '121' and contains '212' or '312' and ends with '2' | $\sum = \{1,2,3\}$

The NFA has been designed below:

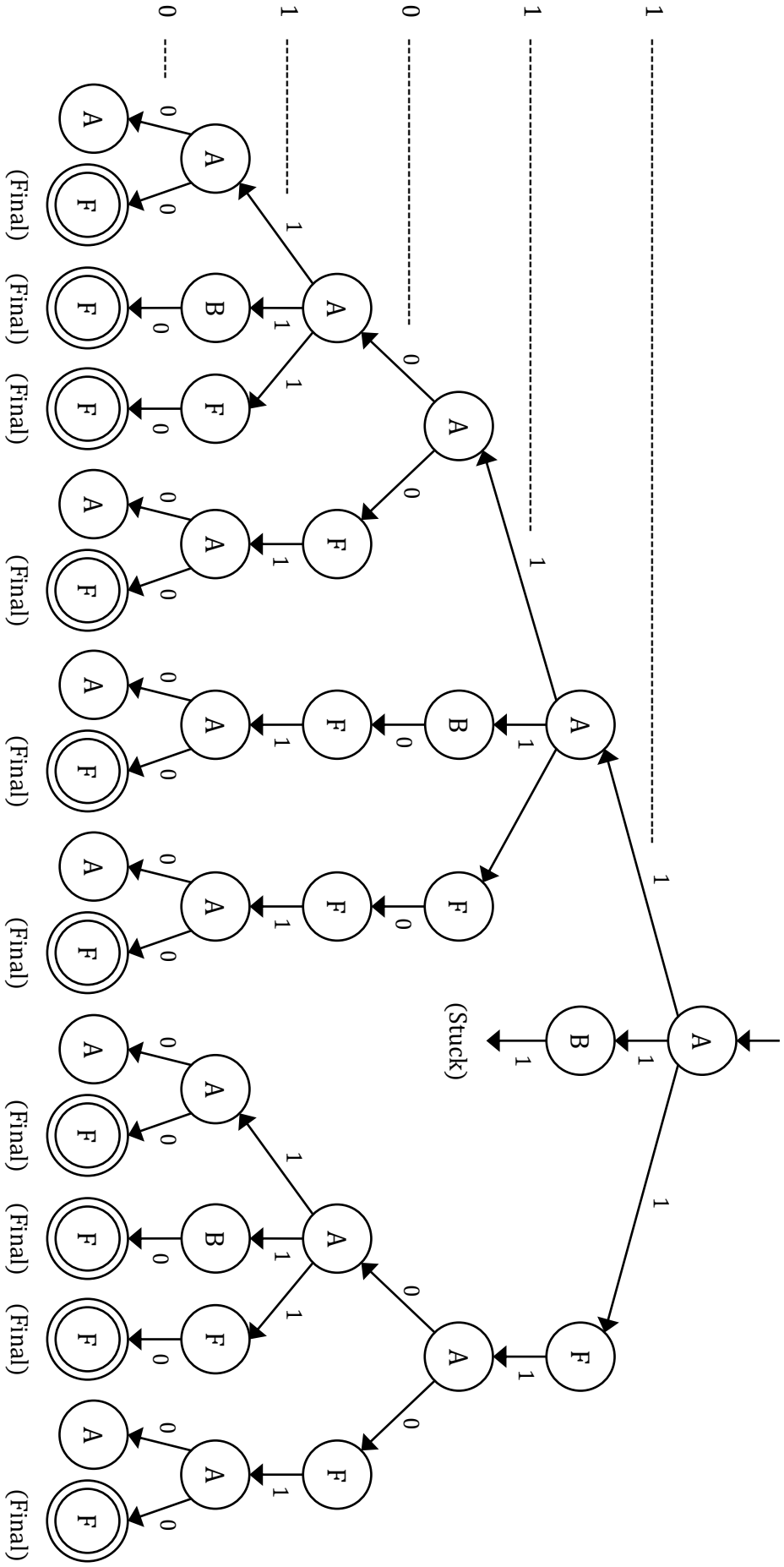


3. Consider the following NFA, and show with help of NFA-tree whether the string "11010" is accepted or not.



Solution:

NFA Tree:



With help of NFA-tree, we can see the string “11010” reach the final state.
∴ The string “11010” is accepted.