

Dijkstra's Algorithm

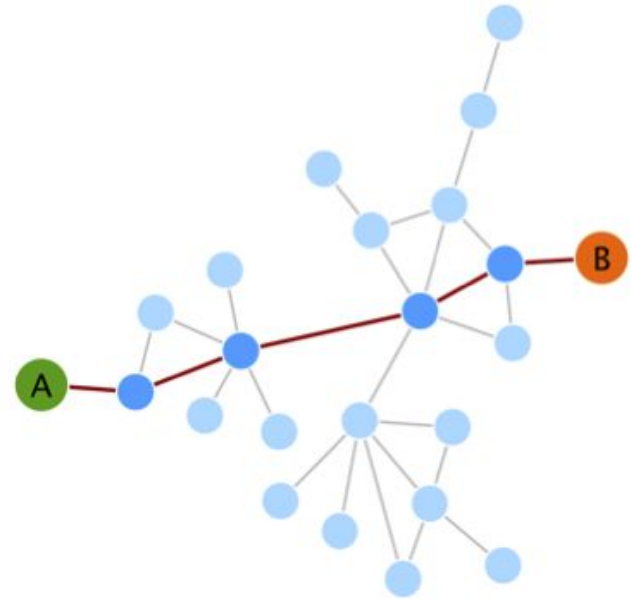


SCTG Lunch & Learn

What is Dijkstra's Algorithm?

Dijkstra's algorithm is used to find the shortest path between 2 nodes in a graph

Named after the Dutch computer scientist Edsger Dijkstra (1930-2002)

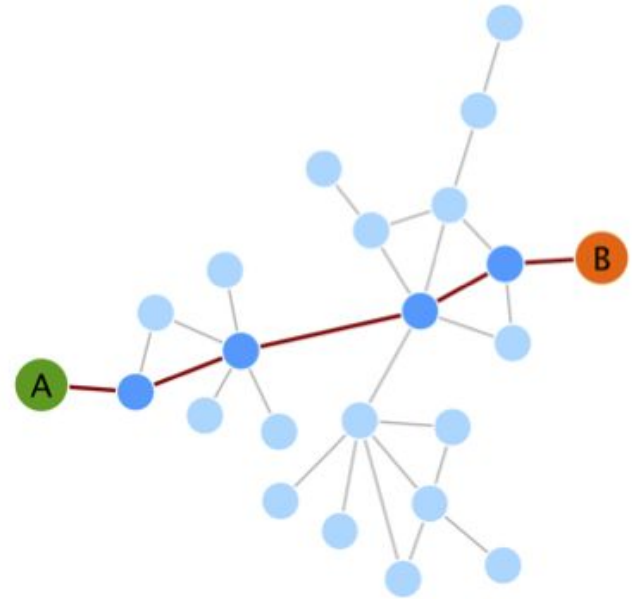


What is a graph?

A set of vertices (nodes) and connecting edges (arcs)

Example:

- Road/highway system
- Facebook/social network
- The internet

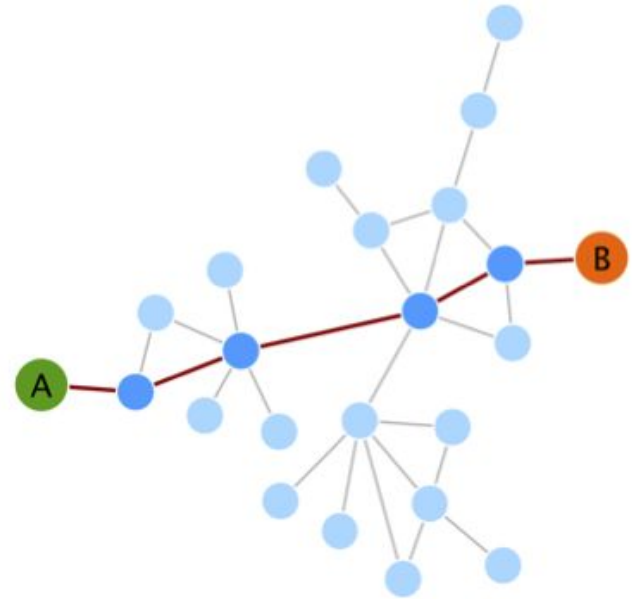


What is Dijkstra's Algorithm?

Algorithm to find the shortest path between 2 nodes in a directed graph.

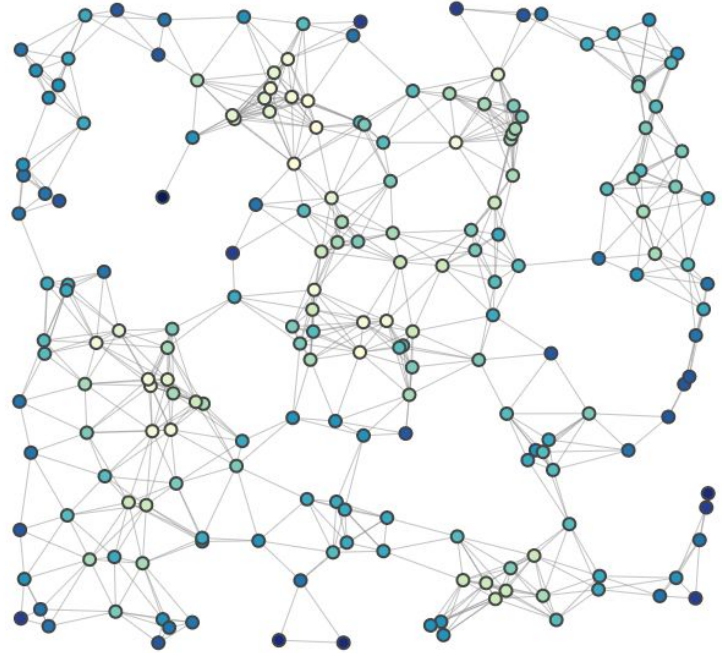
Shortest path is not the fewest number of edges traversed. It is the shortest path along weighted edges.

Note: Path is not *necessarily* distance. Could be number of connections, time of traversal, etc



Alternatives & Motivation

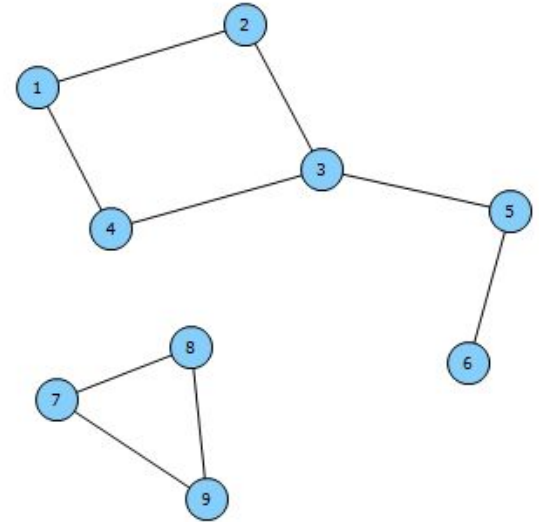
- BFS or DFS
- Find all paths
- Time Efficiency (Big O)



Dijkstra's Algorithm

Conditions:

- All nodes are known in advance (or discover all in pre-processing step)
- Graph is directed
- Edges have non-negative weight
- The nodes are connected*



Computation

S: Source/starting node

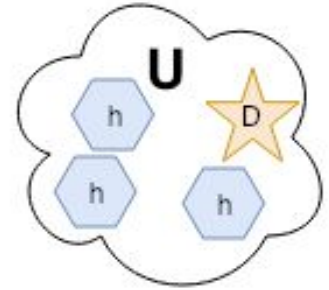
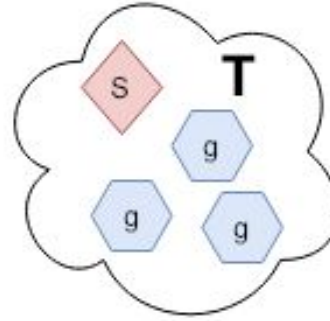
D: Destination/ending node

T: Set of nodes we have visited (traversed to)

U: Set of nodes we have not visited (un-visited)

g: A node in the set T

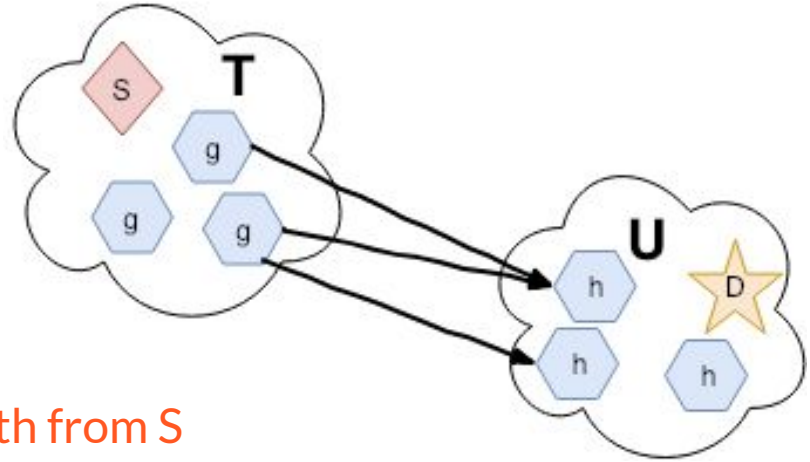
h: A node in the set U

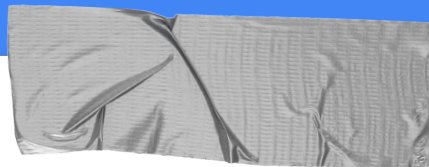


Computation

While D is not in T

- Find all edges (e) that cross from T to U
- Determine the edge, e^* , with shortest path from S [through g] to h along e^*
 - Dijkstra's greedy criterion
- Add h to T, remove from U
- Note: g does not need to be h from last iteration





Bookkeeping

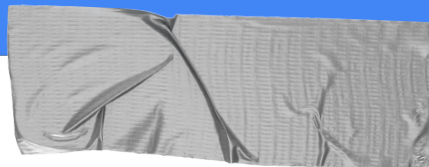
Whenever node is added to T , need to keep track of:

→ **Total distance traveled**

Shortest distance from S to g (so we know how far we will need to go)

→ **Path traveled**

Path from S to g (so we know how to traverse the graph)

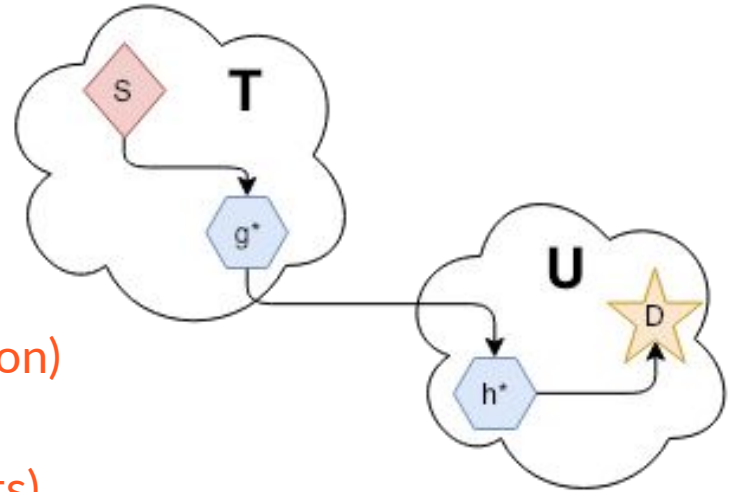


Demo

What is the fastest way to
get a free lunch?

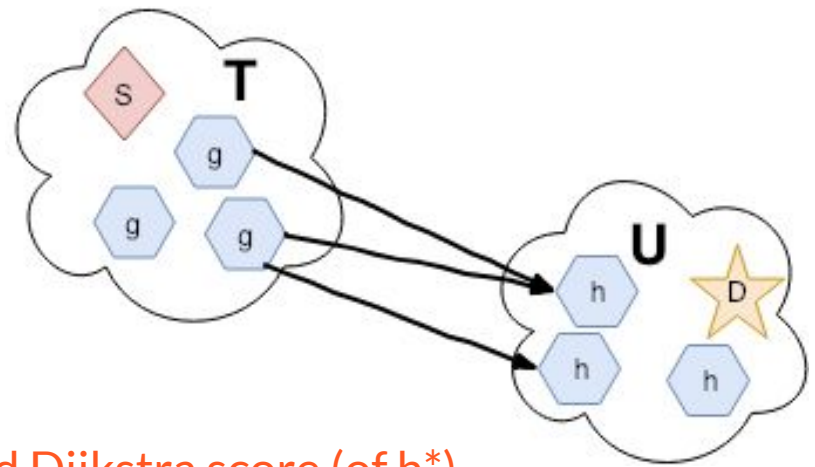
Proof of Correctness

- Overall: determine some path from S to D
- Base: add S to V with length 0
- Iteration: add node h^* to V
- Given $S \rightarrow g^* \rightarrow h^* \rightarrow D$,
 - $S \rightarrow g^*$ is shortest path to g^* (by induction)
 - $g^* \rightarrow h^*$ is some distance from g^* to h^*
 - $h^* \rightarrow D \geq 0$ (non-negative edge weights)
- But $g^* \rightarrow h^*$ is shortest path (Dijkstra's greedy criterion)
- Therefore, chosen path $S \rightarrow D$ is not longer than any other path from $S \rightarrow D$

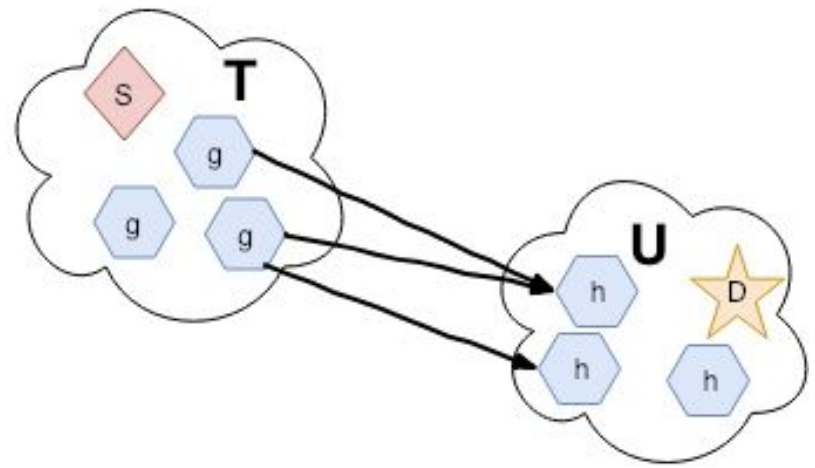


Run Time

- Heap: $O(\log(x))$ inserts and extract min
- Heap of nodes in U
 - Value of node in U is smallest greed Dijkstra score (of h^*)
 - If no edge from T to h, key of h is infinity (should never happen)
- Two rounds:
 - For each node h, determine shortest path from T to h
 - Determine shortest of all paths to h (result of 1st round)
- Every time node h^* is removed from U, check every edge from h^* into U and recalculate shortest path for that node



Run Time



- $m = \# \text{ edges}; n = \# \text{ nodes}$
- If path from S to every node, $m \geq n - 1$
 - $m + n \leq 2m + 1 = O(m)$
- We remove at most $n-1$ nodes from heap of U $\rightarrow O(n)$
- Any edge will require at most 1 key recalculation (insert/delete) in heap of U
- Heap operations:
 - $O(n)$ [extract min from heap] + $O(m)$ [heap key recalculation]
 - $O(m + n)$ operations = $O(m)$ [from above]
 - Each heap operation takes log time $\rightarrow O(m \cdot \log(n))$

Further Enhancements

A* (A star) algorithm

- Extension of Dijkstra's algorithm
- “Best first search”
- Uses heuristic function to determine estimate of minimum distance from h to D
- Prevents going down paths in the “wrong direction”

Resources

Wikipedia:

https://en.wikipedia.org/wiki/Dijkstra%27s_algorithm

Stanford Graph Search and Data Structures Class

<https://www.coursera.org/learn/algorithms-graphs-data-structures>

Demo Code:

<https://startupnextdoor.com/dijkstras-algorithm-in-python-3/>

Code from the demo

<https://github.com/ajstocchetti/dijkstra-lunch-and-learn>

Questions?

If you were inspired and want to give your own Lunch & Learn,
contact me