# Comparison of two theorem provers: Isabelle & Coq

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Comparison of two theorem provers: Isabelle & Coq

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Foundations of Formal Approach

Properties of a Formal

Classical and Intuitionistic Logics

wo Theorem

Isabelle

Coq

Comparison of the Theorem Provers

Common Features

Major Difference

## Outline

### Foundations of Formal Approach

A Formal System
Properties of a Formal System
Classical and Intuitionistic Logics

#### Two Theorem Provers

Isabelle

Coq

#### Comparison of the Theorem Provers

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## Elements of a Formal System

▶ A formula (judgement, statement) φ ∈ Φ:

$$\begin{array}{l} \phi := p \mid q \mid ... \\ \mid \phi_1 \wedge \phi_2 \mid \phi_1 \vee \phi_2 \mid \neg \phi_1 \mid \phi_1 \rightarrow \phi_2 \\ \mid \textit{true} \mid \textit{false} \\ \mid ... \end{array}$$

Propositional variables:

-  $p,q,... \in V$ 

An axiom

-  $\phi_A \in A$ 

ightharpoonup An inference rule au

- a transition function  $\tau:\Phi\to\Phi$
- ▶ A formula  $\phi$  provable from  $\Phi$   $\Phi \vdash \phi$
- ► A tautology ⊤

 $- \vdash \phi$ 

▶ A contradiction ⊥

-  $\vdash \neg \phi$ 

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Foundations of Formal Approach

## A Formal System Properties of a Forma

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A formal system is a quadruple  $\Gamma = \langle A, V, \Omega, R \rangle$ , where

- ► A set of axioms
- ► *V* set of propositional variables
- $ightharpoonup \Omega$  set of logical operators
- $\triangleright$  R set of inference rules

A formal proof of the formula  $\phi$  is a finite sequence of judgements

$$\psi_1 \xrightarrow{\tau_1} \psi_2 \xrightarrow{\tau_2} \dots \xrightarrow{\tau_n} \psi_n$$

where each  $\psi_i$  is either an axiom  $\phi_{A_i}$ , or a formula inferred from the set of previously derived formulas according the rules of inference.

# Properties of a Formal System

#### A formal system $\Gamma$ is called:

- ▶ consistent, if  $\nexists \phi \in \Gamma$ :  $\Gamma \vdash \phi \land \Gamma \vdash \neg \phi \Leftrightarrow \Gamma \nvdash \bot$ ;
- ▶ *complete*, if  $\forall \phi \in U : A \vdash \phi \lor A \vdash \neg \phi$ ;
- ▶ independent, if  $\exists a \in A : A \vdash a$ .

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# Classical Logic

Example: The Hilbert System

Set of axioms:

$$A \rightarrow (B \rightarrow A)$$

$$(A \to (B \to C)) \to ((A \to B) \to (A \to C))$$

Single inference rule:

$$\llbracket A, A \to B \rrbracket \longrightarrow B$$

Some tautologies:

$$A \lor \neg A$$

$$\neg\neg(A\vee\neg A)$$

$$A\to\neg\neg A$$

$$A \rightarrow \neg \neg A$$

$$\neg \neg A \rightarrow A$$

(DNe)

(MP)

(A1)

(A2)

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Classical and Intuitionistic Logics

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## Isabelle: First Acquaintance

- a generic proof assistant
- a successor of HOL theorem prover //TODO: cite
- created in 1986 by
  - Larry Paulson @ University of Cambridge, and
  - ► Tobias Nipkow @ Technische Universität München
- based on classical higher-order logic
- ▶ uses powerful functional language HOL
- ▶ has large collection of formalised theories //TODO: HOL, ZF, CCL, ...

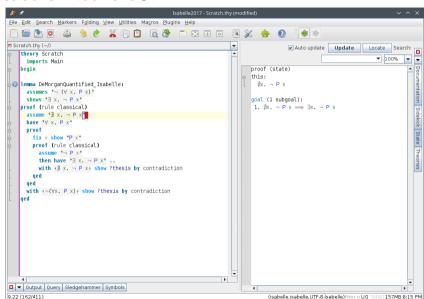
```
Example 2: ???
Example 1: Definition of basic datatypes
datatype bool =
                                          ???
  True | False
datatype nat =
  zero ("0") | Suc nat
```

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Isabelle

## Isabelle: Native GUI



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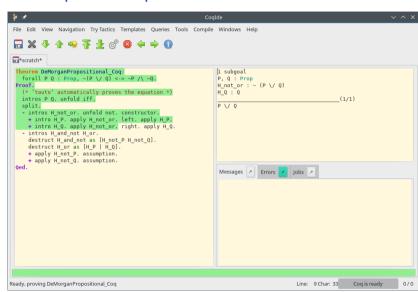
Major Differences

- a formal proof management system
- created at INRIA (Paris, France) in 1984
- based on Calculus of Inductive Constructions theory (an implementation of intuitionistic logic)
- uses powerful functional language Gallina
- ▶ has large collection of formalised theories //TODO
- widely used in software verification (proof code extraction)

```
Example 3: Definition of basic datatypes
Inductive False : Prop := . ???
Inductive True : Prop := I : True.

Inductive nat : Type :=
    | 0 : nat
    | S : nat -> nat.
```

## The Cog theorem prover



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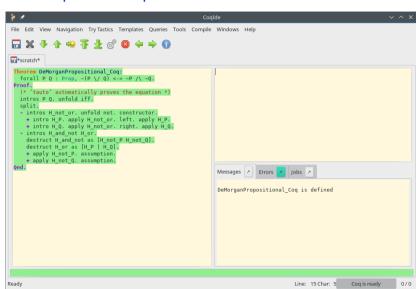
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## Summary

- ▶ The first main message of your talk in one or two lines.
- ► The second main message of your talk in one or two lines.
- ▶ Perhaps a third message, but not more than that.
- Outlook
  - Something you haven't solved.
  - Something else vou haven't solved.

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