

Global Value Numbering in Factor

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DEPARTMENT	COURSE	DESCRIPTION	PREREQS
COMPUTER SCIENCE	CPSC 432	INTERMEDIATE COMPILER DESIGN, WITH A FOCUS ON DEPENDENCY RESOLUTION.	CPSC 432

Factor

Factor (<http://factorcode.org/>)

- Started development September 2003—a baby among languages
- **Stack-based**
- Object-oriented
- Dynamically typed
- Extensive standard library
- High-level, yet fully **compiled**

Won't really have time to discuss the language in depth

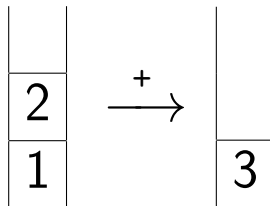
Stacks as an Evaluation Model

Example (Code)

1 2 +

Example (Execution)

```
push(1);  
push(2);  
y = pop();    // y = 2;  
x = pop();    // x = 1;  
push(x + y);  // push(3);
```



- 1 Compiler
 - Structure
 - Optimizations
- 2 Value Numbering
- 3 Results

Organization

Non-optimizing base compiler

- VM written in C++
- Responsible for basic runtime services
 - Garbage collection
 - Method dispatch
 - Polymorphic inline caches
 - ...
- Single pass—outputs assembly stubs for primitives

Optimizing compiler

- Written in Factor code
 - Possible by *bootstrapping*
- Optimizes in passes across two **intermediate representations** (IRs)
 - High-level IR (`compiler.tree`)
 - Low-level IR (`compiler.cfg`)

High-level IR

- Tree of node objects
- Very simple virtual instruction set
 - `#introduce`, `#return`
 - `#push` & `#call`
 - `#renaming`—`#copy` & `#shuffle`
 - `#declare` & `#terminate`
 - `#branch`—`#if` & `#dispatch`
 - `#phi`
 - `#recursive`, `#enter-recursive`, `#call-recursive`,
`#return-recursive`
 - `#alien-node`, `#alien-invoke`, `#alien-indirect`,
`#alien-assembly`, `#alien-callback`
- Input/output values of stack given unique names

High-level IR

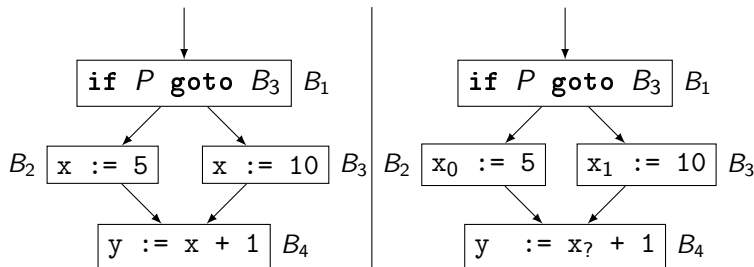
1 2 +

Example

```
V{  
  T{ #push { literal 1 } { out-d { 6256273 } } }  
  T{ #push { literal 2 } { out-d { 6256274 } } }  
  T{ #call  
    { word + }  
    { in-d V{ 6256274 6256273 } }  
    { out-d { 6256275 } }  
  }  
  T{ #return { in-d V{ 6256275 } } }  
}
```

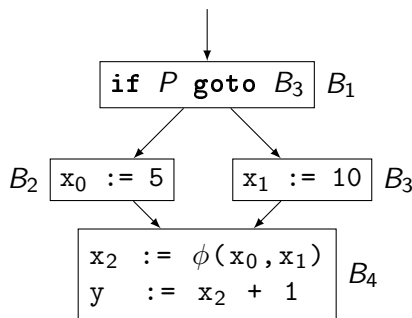
Low-level IR

- Control flow graph (CFG)
 - Basic blocks = maximal sequence of “straight-line” code
 - Directed edges = transfer of control flow
- `insn` objects modeled closely after assembly-like instructions
- **Static single assignment** (SSA) form



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Optimizations—High-level IR

```
: optimize-tree ( nodes -- nodes' )  
[  
    analyze-recursive  
    normalize  
    propagate  
    cleanup  
    dup run-escape-analysis? [  
        escape-analysis  
        unbox-tuples  
    ] when  
    apply-identities  
    compute-def-use  
    remove-dead-code  
    ?check  
    compute-def-use  
    optimize-modular-arithmetic  
    finalize  
] with-scope ;
```

Optimizations—Low-level IR

```
      : optimize-cfg ( cfg -- cfg' )  
        optimize-tail-calls  
        delete-useless-conditionals  
        split-branches  
        join-blocks  
        normalize-height  
        construct-ssa  
        alias-analysis  
→     value-numbering  
        copy-propagation  
        eliminate-dead-code ;
```

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