### Global Value Numbering in Factor

### Alex Vondrak

ajvondrak@csupomona.edu

September 1, 2011

PAGE 3			
DEPARTMENT	COURSE	DESCRIPTION	PREREQS
COMPUTER SCIENCE		INTERMEDIATE COMPILER DESIGN, WITH A FOCUS ON DEPENDENCY RESOLUTION.	CPSC 432
		Harrison Contraction Company Contraction	D. 100

### **Factor**

Factor (http://factorcode.org/)

- Started development September 2003—a baby among languages
- Stack-based
- Object-oriented
- Dynamically typed
- Extensive standard library
- High-level, yet fully compiled

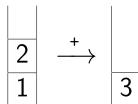
Won't really have time to discuss the language in depth

### Stacks as an Evaluation Model

```
Example (Code)
1 2 +
```

Example (Execution)

```
push(1);
push(2);
y = pop(); // y = 2;
x = pop(); // x = 1;
push(x + y); // push(3);
```



- Compiler
  - Structure
  - Optimizations
- Value Numbering
  - Local Value Numbering
  - Global Value Numbering

## Organization

### Non-optimizing base compiler

- VM written in C++
- Responsible for basic runtime services
  - Garbage collection
  - Method dispatch
  - Polymorphic inline caches
  - . . .
- Single pass—outputs assembly stubs for primitives

### Optimizing compiler

- Written in Factor code
  - Possible by bootstrapping
- Optimizes in passes across two intermediate representations (IRs)
  - High-level IR (compiler.tree)
  - Low-level IR (compiler.cfg)

### High-level IR

- Tree of node objects
- Very simple virtual instruction set
  - #introduce, #return
  - #push & #call
  - #renaming—#copy & #shuffle
  - #declare & #terminate
  - #branch—#if & #dispatch
  - #phi
  - #recursive, #enter-recursive, #call-recursive, #return-recursive
  - #alien-node, #alien-invoke, #alien-indirect, #alien-assembly, #alien-callback
- Input/output values of stack given unique names

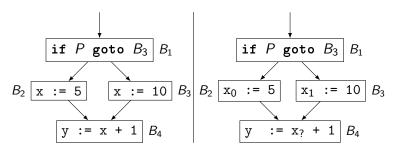
### High-level IR

1 2 +

```
Example
₩.
    T{ #push { literal 1 } { out-d { 6256273 } } }
    T{ #push { literal 2 } { out-d { 6256274 } } }
    T{ #call
        { word + }
        { in-d V{ 6256273 6256274 } }
        { out-d { 6256275 } }
    }
    T{ #return { in-d V{ 6256275 } } }
```

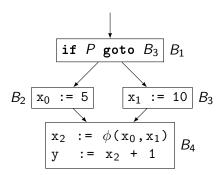
### Low-level IR

- Control flow graph (CFG)
  - Basic blocks = maximal sequence of "straight-line" code
  - Directed edges = transfer of control flow
- insn objects modeled closely after assembly-like instructions
- Static single assignment (SSA) form



### Low-level IR

- Control flow graph (CFG)
  - Basic blocks = maximal sequence of "straight-line" code
  - Directed edges = transfer of control flow
- insn objects modeled closely after assembly-like instructions
- Static single assignment (SSA) form



### Optimizations—High-level IR

```
: optimize-tree ( nodes -- nodes' )
     analyze-recursive
     normalize
     propagate
      cleanup
     dup run-escape-analysis? [
          escape-analysis
          unbox-tuples
      ] when
     apply-identities
      compute-def-use
     remove-dead-code
     ?check
      compute-def-use
     optimize-modular-arithmetic
     finalize
   with-scope ;
```

### Optimizations—Low-level IR

```
: optimize-cfg ( cfg -- cfg' )
   optimize-tail-calls
   delete-useless-conditionals
   split-branches
   join-blocks
   normalize-height
   construct-ssa
   alias-analysis
   value-numbering
   copy-propagation
   eliminate-dead-code:
```

- - Structure
  - Optimizations
- Value Numbering
  - Local Value Numbering
  - Global Value Numbering

### Value Numbering

### Idea: assign each variable a value number

- Equal value numbers ⇒ equal at runtime
- Turn recomputations into ##copy instructions, saving time

### General problem is undecidable

- Seek conservative solution
- Discover Herbrand equivalences
- Consider two values congruent if
  - They're computed by the same operator
  - Their operands are congruent

- Thought to be invented by Balke in the 1960s
- Largely credited to Cocke & Schwartz in the 1970s
- Current implementation Factor uses

Pro: Easy to understand, implement, and extend

Con: Is local and pessimistic, discovering fewer congruences

### Implementation

Expressions are constructed from instructions

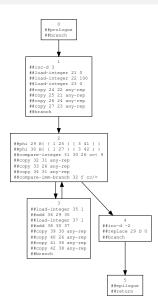
### Example

```
T{ ##add { dst 1 } { src1 2 } { src2 3 } } >expr { ##add 2 3 }
```

- Expression graph = 3 global hash tables
  - vregs>vns
  - exprs>vns
  - vns>insns
- If possible, instructions are simplified using data from expression graph

```
Example (In Factor)
0 100 [ 1 fixnum+fast ] times
Example (In Java)
```

```
int i = 0;
for (int j = 0; j < 100; j++) {
  i += 1;
```



### Basic Block 1

```
vregs>vns = H{ }
exprs>vns = H{ }
```

```
##inc-d 3
##load-integer 21 0
##load-integer 22 100
##load-integer 23 0
##copy 24 22 any-rep
##copy 25 21 any-rep
##copy 26 24 any-rep
##copy 27 23 any-rep
##branch
```

(no-op)

```
vregs>vns = H{ { 21 21 } }
exprs>vns = H\{ \{ 0 21 \} \}
```

```
##inc-d 3
##load-integer 21 0
##load-integer 22 100
##load-integer 23 0
##copy 24 22 any-rep
##copy 25 21 any-rep
##copy 26 24 any-rep
##copy 27 23 any-rep
##branch
```

```
(no-op)
> expr = 0
```

```
vregs>vns = H{ { 21 21 } { 22 22 } }
exprs>vns = H\{ \{ 0 21 \} \{ 100 22 \} \}
```

```
##inc-d 3
##load-integer 21 0
##load-integer 22 100
##load-integer 23 0
##copy 24 22 any-rep
##copy 25 21 any-rep
##copy 26 24 any-rep
##copy 27 23 any-rep
##branch
```

```
(no-op)
> expr = 0
> expr = 100
```

```
vregs>vns = H{ { 21 21 } { 22 22 } { 23 21 } }
exprs>vns = H\{ \{ 0 21 \} \{ 100 22 \} \}
```

```
##inc-d 3
##load-integer 21 0
##load-integer 22 100
##load-integer 23 0
##copy 24 22 any-rep
##copy 25 21 any-rep
##copy 26 24 any-rep
##copy 27 23 any-rep
##branch
```

```
(no-op)
> expr = 0
> expr = 100
> expr = 0
```

```
vregs>vns = H{ { 21 21 } { 22 22 } { 23 21 } { 24 22 } ... }
exprs>vns = H\{ \{ 0 21 \} \{ 100 22 \} \}
```

```
##inc-d 3
##load-integer 21 0
##load-integer 22 100
##copy 23 21 any-rep
##copy 24 22 any-rep
##copy 25 21 any-rep
##copy 26 24 any-rep
##copy 27 23 any-rep
##branch
```

```
(no-op)
> expr = 0
> expr = 100
> expr = 0
. . .
. . .
```

```
vregs>vns = H{ }
exprs>vns = H{ }
```

```
##phi 29 H{ { 1 25 } { 3 41 } }
##phi 30 H{ { 1 27 } { 3 42 } }
##compare-integer 31 30 26 cc< 9
##copy 32 31 any-rep
##copy 33 26 any-rep
##copy 34 31 any-rep
##compare-imm-branch 32 f cc/=
```

```
(no-op)
(no-op)
```

```
vregs>vns = H{ { 30 30 } { 26 26 } { 31 31 } }
exprs>vns = H{ { { ##compare-integer 30 26 cc< } 31 } }
```

```
##phi 29 H{ { 1 25 } { 3 41 } }
##phi 30 H{ { 1 27 } { 3 42 } }
##compare-integer 31 30 26 cc< 9
##copy 32 31 any-rep
##copy 33 26 any-rep
##copy 34 31 any-rep
##compare-imm-branch 32 f cc/=
```

```
(no-op)
(no-op)
>expr = ...
```

```
vregs>vns = H{ { 30 30 } { 26 26 } { 31 31 } ... }
exprs>vns = H{ { { ##compare-integer 30 26 cc< } 31 } }
```

```
##phi 29 H{ { 1 25 } { 3 41 } }
##phi 30 H{ { 1 27 } { 3 42 } }
##compare-integer 31 30 26 cc< 9
##copy 32 31 any-rep
##copy 33 26 any-rep
##copy 34 31 any-rep
##compare-imm-branch 32 f cc/=
```

```
(no-op)
(no-op)
>expr = ...
```

```
vregs>vns = H{ { 30 30 } { 26 26 } { 31 31 } ... }
exprs>vns = H{ { { ##compare-integer 30 26 cc< } 31 } }
```

```
##phi 29 H{ { 1 25 } { 3 41 } }
##phi 30 H{ { 1 27 } { 3 42 } }
##compare-integer 31 30 26 cc< 9
##copy 32 31 any-rep
##copy 33 26 any-rep
##copy 34 31 any-rep
##compare-integer-branch 30 26 cc<
```

```
(no-op)
(no-op)
>expr = ...
```

```
vregs>vns = H{ }
exprs>vns = H{ }
```

```
##load-integer 35 1
##add 36 29 35
##load-integer 37 1
##add 38 30 37
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

#### Basic Block 3

```
vregs>vns = H\{ \{ 35 35 \} \}
exprs>vns = H\{ \{ 1 35 \} \}
```

```
##load-integer 35 1
##add 36 29 35
##load-integer 37 1
##add 38 30 37
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

> expr = 1

```
vregs>vns = H{ { 35 35 } { 29 29 } { 36 36 } }
exprs>vns = H{ { 1 35 } { { ##add-imm 29 1 } 36 } }
```

```
##load-integer 35 1
##add 36 29 35
##load-integer 37 1
##add 38 30 37
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

```
> expr = 1
>expr = { ##add-imm 29 1 }
```

```
vregs>vns = H{ { 35 35 } { 29 29 } { 36 36 } { 37 35 } }
exprs>vns = H{ { 1 35 } { { ##add-imm 29 1 } 36 } }
```

```
##load-integer 35 1
##add-imm 36 29 1
##load-integer 37 1
##add 38 30 37
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

```
> expr = 1
>expr = { ##add-imm 29 1 }
> expr = 1
```

```
vregs>vns = H{ { 35 35 } { 29 29 } { 36 36 } { 37 35 } ... }
exprs>vns = H{ { 1 35 } { { ##add-imm 29 1 } 36 } ... }
```

```
##load-integer 35 1
##add-imm 36 29 1
##copy 37 35 any-rep
##add 38 30 37
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

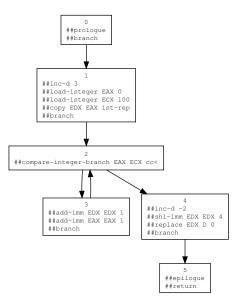
```
> expr = 1
>expr = { ##add-imm 29 1 }
> expr = 1
>expr = { ##add-imm 30 1 }
```

```
vregs>vns = H{ { 35 35 } { 29 29 } { 36 36 } { 37 35 } ... }
exprs>vns = H{ { 1 35 } { { ##add-imm 29 1 } 36 } ... }
```

```
##load-integer 35 1
##add-imm 36 29 1
##copy 37 35 any-rep
##add-imm 38 30 1
##copy 39 30 any-rep
##copy 40 26 any-rep
##copy 41 36 any-rep
##copy 42 38 any-rep
##branch
```

```
> expr = 1
>expr = { ##add-imm 29 1 }
> expr = 1
>expr = { ##add-imm 30 1 }
. . .
```

### Local Value Numbering Results



- - Structure
  - Optimizations
- Value Numbering
  - Local Value Numbering
  - Global Value Numbering
- Results