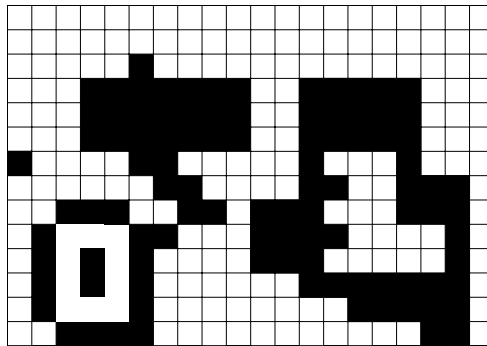


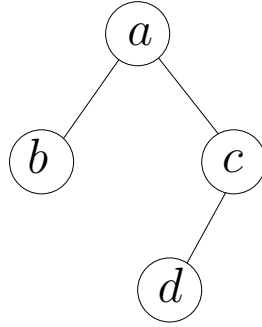
Assignment - 02
Advanced Data Structures and Algorithms
Indian Institute of Information Technology Kalyani

Due date: August 01, 2021

1. Write codes for insertion and deletion to a BST. Use your code to insert the following keys to an initially empty BST: 12, 5, 4, 7, 15, 20, 2, 24, and 5. Delete 15 and 20 from the BST now.
2. Modify your codes written for Prob-1 such that the search tree always grows as a height-balanced tree.
3. In the figure shown below, you have to count the number of connected components present. Consider the image as a 2D array, where the black squares are represented by 1 and the white squares by 0. Write a program, and your program should use both 4-N and 8-N connectivity separately for counting components.



4. In a representation of binary trees, a node **a** with left child and right child null is written as $(a())()$. The following tree is presented by the string: $(a((b())((c((d())())))))$. Given a string S in that form, you have to check if S represents a BST or not. Write a code for the same.



DFS and BFS Traversal:

You may use these traversal methods to solve Prob 3. In 4-N, square a is adjacent to square b if b is immediate *south* or *north* or *east* or *west* square of a . In 8-N, we include *north-east*, *south-east*, *south-west* and *north-west* with the definition of 4-N.

Algorithm 1: Depth-First-Search

Input: $G(V, E)$

```

1 for all  $v \in V$  do
2    $v_{status} \leftarrow \text{WHITE}$ 
3 for each  $v \in V$  do
4   if  $v_{status} == \text{WHITE}$  then
5      $DFS(v)$ 

```

Procedure $DFS(v)$

```

1  $v_{status} \leftarrow \text{GRAY}$ 
2 for each  $u \in Adj(v)$  do
3   if  $u_{status} == \text{WHITE}$  then
4      $DFS(u)$ 
5  $v_{status} \leftarrow \text{BLACK}$ 

```

Algorithm 2: Breadth-First-Search

Input: $G(V, E)$

```
1 for all  $v \in V$  do
2    $v_{status} \leftarrow \text{WHITE}$ 
3 for each  $v \in V$  do
4   if  $v_{status} == \text{WHITE}$  then
5      $BFS(v)$ 
```

Procedure $BFS(w)$

```
1  $w_{status} \leftarrow \text{GRAY}$ 
2  $ENQUEUE(w, Q)$ 
3 while  $Q \neq \phi$  do
4    $v \leftarrow DEQUEUE(Q)$ 
5   for each  $u \in Adj(v)$  do
6     if  $u_{status} == \text{WHITE}$  then
7        $u_{status} \leftarrow \text{GRAY}$ 
8        $ENQUEUE(u, Q)$ 
9    $v_{status} \leftarrow \text{BLACK}$ 
```
