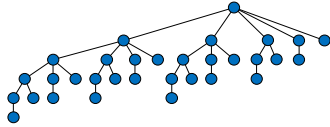


Lecture 1. Course Overview, Motivation, Fundamental Concepts

CpSc 8400: Algorithms and Data Structures
Brian C. Dean



School of Computing
Clemson University
Spring, 2021

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Introductions

- **Instructor:** Dr. Brian C. Dean
 - Ph.D. from MIT in 2005.
 - **Email:** bcdean@clemson.edu
 - **Office:** McAdams 221 (but not in-person this semester)
 - **Office Hours:** Tuesday after class.
 - **Research Interests:** Algorithms (optimization, data mining and analytics, randomization, data structures, heuristics, matching, geometry) and their applications (medical informatics, networking, scheduling).
 - **Educational Interests** in algorithmic CS education and problem-solving at the high-school level; director of USA Computing Olympiad (usaco.org).

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Course Overview

- This course provides a fun, fast-paced, modern theory-oriented study of algorithms.
- We will learn:
 - Algorithm design techniques
Divide and conquer, greedy algorithms, dynamic programming, iterative refinement, randomized incremental construction, ...
 - Common algorithms for common algorithmic problems
Sorting and selection, graph problems, FFTs, optimization, ...
 - Mathematical tools for algorithm analysis
Probability theory, amortized analysis, recurrences, ...
 - Interesting algorithmic subfields in which you may want to pursue further research/study.
Computational geometry, cryptography, approximation and randomized algorithms, data structures, bioinformatics, ...

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Why Learn Algorithms?...

- Algorithms are the heart and soul of computing!
 - Algorithmic computing now plays a key role in nearly everything – proficiency opens many doors.
- 
- ```
graph LR; S[Science] <--> ACS["(Algorithmic) CS"]; E[Engineering] <--> ACS; A[Arts] <--> ACS; ACS <--> C[Commerce]; ACS <--> H[Health]; ACS <--> S2[Society];
```
- The diagram illustrates the central role of (Algorithmic) CS in various fields. A central box labeled "(Algorithmic) CS" is connected by double-headed arrows to six surrounding boxes: Science, Engineering, and Arts on the left; and Commerce, Health, and Society on the right.
- Algorithmic proficiency differentiates a true “computer scientist” from a run-of-the-mill “programmer”, and provides the foundation for a long-term career in computing that can thrive as technology changes.
  - Theory meets practice. Algorithms have motivated the development of some truly elegant theoretical results in mathematics, for those who appreciate mathematics.
  - Algorithms are fun!

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## Programming ≠ Algorithmic Problem Solving

- “Computer Science is no more about computers than astronomy is about telescopes” – folklore, sometimes attributed to E. Dijkstra.
- Programming and software engineering are certainly an important part of most computing projects and careers.
- Problem-solving skills and programming skills are different, but re-inforce each-other.
- Both are crucial for success as a computing expert.

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## Course Details

- **Prerequisites**
  - A reasonable amount of mathematical maturity.
  - Enthusiasm, willingness to challenge yourself and ask questions, and a good work ethic!
- **Course Materials**
  - *Algorithms Explained*, currently being written by the instructor – portions may be made available electronically on Canvas (don't redistribute!)
  - Lecture slides and videos to appear on Canvas.
  - I can suggest supplemental reading, if interested.
  - I may also post prominent research papers relevant to the course content.

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## Assignments and Grades

- **Homework (35%)**
  - Typically focused on mathematical analysis of algorithms, not implementation.
  - Solutions must be typeset and submitted as a PDF file using [handin.cs.clemson.edu](mailto:handin.cs.clemson.edu) by email *before the start of class on the day they are due*.
- **2 Quizzes (2 x 20%) and a Final (25%)**
  - All quizzes/exams are cumulative.
- Appropriate letter grade cutoffs set by instructor at end of semester.

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## Course Conduct

- **Academic Integrity:**
  - Do not cheat.
  - Do not plagiarize (pass off the work of others as your own without appropriate attribution).
- **Collaboration:**
  - Highly encouraged, but each student should write up final solutions independently. Please list your collaborators.
  - Do not consult homework solutions from previous semesters, and do not use the web for anything but general reference.
- **Feedback:**
  - Please feel welcome to ask for feedback at any time. The instructor always appreciates constructive feedback.

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## A Good Algorithm (or Data Structure)...

- Always terminates and produces **correct** output.
  - A “close enough” answer is sometimes fine.
  - Some types of randomized algorithms can fail, but only with miniscule probability.
- Makes **efficient** use of computational resources.
  - Minimizes running time, memory usage, processors, bandwidth, power consumed, heat produced.
- Is **simple** to describe, understand, analyze, implement, and debug. Elegance matters!

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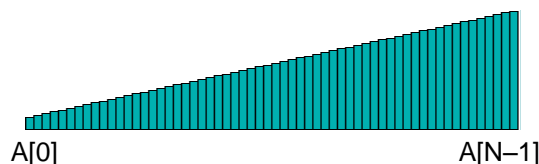
## Example: Searching an Array

- **Linear** search: runs in  $N$  “steps” in the worst case.

```
for i = 0..N-1:
 if target = A[i], found it!
```

- **Binary** search:  $\leq \log_2 N$  “steps” in worst case.  
(requires our array to be sorted).

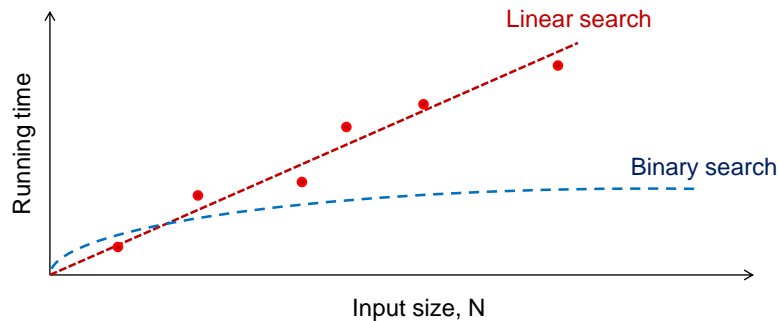
```
if target = middle element, found it!
else recursively search first or second half
array, as appropriate.
```



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## Empirical Performance Testing



- Choose inputs carefully, since often some inputs are much easier than others.
- Do you want to measure “average case” or “worst-case” performance...?

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## Asymptotic Analysis

- **Linear** search:  $O(N)$  time.
- **Binary** search:  $O(\log N)$  time.
- $O(f(N))$  means “upper-bounded by a constant times  $f(N)$  as  $N$  grows large”.
- Provides an asymptotic upper bound on running time, where constant factors and lower-order terms don’t matter.
- Captures what usually matters most about algorithm performance: how worst-case running time scales with input size.
- However, this can lose important information...  
(e.g., sequential vs. non-sequential linear search)

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## Asymptotic Notation

- $O()$  provides an asymptotic upper bound.
- $\Omega()$  provides an asymptotic lower bound.
- $\Theta()$  means both a lower and upper bound.  
(think of these as “ $\leq$ ”, “ $\geq$ ”, and “ $=$ ”)

Usage examples:

- “The running time of our algorithm is  $O(n^2)$ .”
- “The worst-case running time of our algorithm is  $\Theta(n^2)$ .”
- “This algorithm uses  $\Omega(n^2)$  memory”.
- “ $17n^2 - 5n + 200 = \Theta(n^2)$ ”
- “Consider the polynomial  $5x^{10} - 3x^9 + o(x^9)$ .”
- Can you solve the problem in  $o(n^2)$  time?

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## Running Times

- We almost always focus on **worst case** running times. Why?
- Common running times:
  - Constant:  $O(1)$
  - Logarithmic:  $O(\log n)$
  - Linear:  $O(n)$
  - Polynomial:  $O(n \log n)$ ,  $O(n^2)$ ,  $O(n^3)$ ,  $O(n^{100})$ , ...
  - Exponential:  $O(2^n)$ ,  $O(3^n)$ , ...
  - Worse than exponential:  $O(n!)$ ,  $O(n^n)$ .
- Logs: base doesn't matter in  $O()$ , as long as not in an exponent. By  $\log n$  we usually mean  $\log_2 n$ .

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## A Few More Symbols...

- $O()$  provides an asymptotic upper bound.
- $\Omega()$  provides an asymptotic lower bound.
- $\Theta()$  means both a lower and upper bound.  
(think of these as “ $\leq$ ”, “ $\geq$ ”, and “ $=$ ”)
- **New additions:**  
 $o()$  and  $\omega()$  are like “ $<$ ” and “ $>$ ”

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## Asymptotic Analysis Practice

- Consider this code:

```
For i = 1 to n
 For j = 1 to n
 Increment counter
```

- Which of these expressions describe the running time correctly?

|          |               |               |          |               |
|----------|---------------|---------------|----------|---------------|
| $O(n)$   | $\Omega(n)$   | $\Theta(n)$   | $o(n)$   | $\omega(n)$   |
| $O(n^2)$ | $\Omega(n^2)$ | $\Theta(n^2)$ | $o(n^2)$ | $\omega(n^2)$ |
| $O(n^3)$ | $\Omega(n^3)$ | $\Theta(n^3)$ | $o(n^3)$ | $\omega(n^3)$ |

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