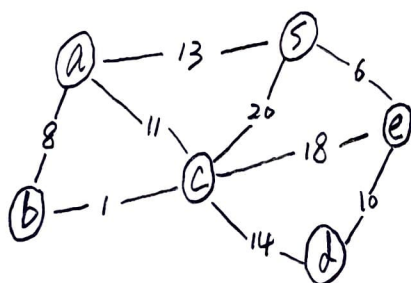


Feng Wei 50428109

Prim's Algorithm



lev_i = the weight of the lightest edge between v and s

$\pi[v]$ = the lightest edge between v and S ($\pi[v], v$).

[illegible]

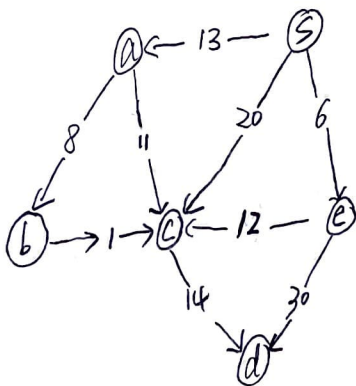
Problem 2.

Dijkstra's Algorithm

iteration	Vertex added to S in iteration i	a		b		c		d		e	
		d	π	d	π	d	π	d	π	d	π
		13	S	∞	L	20	S	∞	L	6	S
1	S	13	S	∞	L	18	e	36	e		
2	e	13	S	∞	L	18	e	36	e		
3	a			21	a	18	e	36	e		
4	c			21	a			32	c		
5	b							32	c		
6	d										

$d[v]$ = the shortest path from S to v

$\pi[v]$ = the parent of v in the shortest path tree



Problem 3.

Undirected Graph $G=(V, E)$, Non-negative edge weights $(W_e)_{e \in E}$

(3a.) Yes.

proof: Changing the weight of every edge e from W_e to W_e^2 doesn't change the ordering of the edge weights.

Let T be the original MST of G . And T^* be the after MST.

For every cut $(U, V \setminus U)$, T has one edge e and T^* has one edge e^* to connect the cut. Let $e=(u, v)$ and $e^*=(u^*, v^*)$. Assume $e \neq e^*$

According to the definition $W_{(u, v)} < W_{(u^*, v^*)}$ and $W_{(u, v)}^2 > W_{(u^*, v^*)}^2$

① ②

Since W_e are non-negative and different.

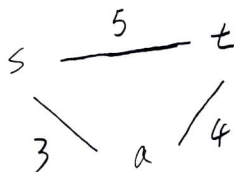
Then ① and ② cannot be correct at the same time.

As a result the assumption is wrong. We have $e = e^* \Rightarrow T = T^*$

T is still the minimum spanning tree of G .

(36.) No

Counter - example



original



After

Assume we only have three vertices s , t , and a .

In the original Graph the shortest path from s to t is $p(s-t) = 5$

In the After graph, p is no longer unique.

Since $\cancel{s-t} = 25 = \cancel{s-a} + \cancel{a-t}$

Problem 4.

$$d[u] + w(u,v) < d[v]$$

Assume for every edge $(u,v) \in E$, that $d[u] + w(u,v) \geq d[v]$

Let the negative cycle be $C = (a_1, a_2, \dots, a_n)$ $n > 1$. So we have

$$0 > w(a_1, a_2) + w(a_2, a_3) + \dots + w(a_{n-1}, a_n) + w(a_n, a_1)$$

According to the assumption:

$$d(a_1) + w(a_1, a_2) \geq d(a_2)$$

$$\vdots$$

$$d(a_{n-1}) + w(a_{n-1}, a_n) \geq d(a_n)$$

$$d(a_n) + w(a_n, a_1) \geq d(a_1)$$

$$\Rightarrow [d(a_1) + w(a_1, a_2)] + \dots [d(a_n) + w(a_n, a_1)] \geq d(a_2) + \dots + d(a_n) + d(a_1)$$

$$\Rightarrow \underbrace{[d(a_1) + d(a_2) + \dots + d(a_n)]}_{\textcircled{1}} + \underbrace{[w(a_1, a_2) + \dots + w(a_n, a_1)]}_{\textcircled{2}} \geq \underbrace{[d(a_1) + \dots + d(a_n)]}_{\textcircled{3}}$$

obviously $\textcircled{1} = \textcircled{3}$ and $\textcircled{2}$ is the negative cycle. So $\textcircled{2} < 0$

As a result the assumption is ~~wrong~~. So there exists edge such that $d[u] + w(u,v) < d[v]$