CSE 431/531: Algorithm Analysis and Design (Spring 2022)

Divide-and-Conquer

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Outline

Greedy Algorithm

- mainly for combinatorial optimization problems
- trivial algorithm runs in exponential time
- greedy algorithm gives an efficient algorithm
- main focus of analysis: correctness of algorithm

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Divide-and-Conquer

- not necessarily for combinatorial optimization problems
- trivial algorithm already runs in polynomial time
- divide-and-conquer gives a more efficient algorithm
- main focus of analysis: running time

Divide-and-Conquer

- **Divide**: Divide instance into many smaller instances
- **Conquer**: Solve each of smaller instances recursively and separately
- **Combine**: Combine solutions to small instances to obtain a solution for the original big instance

merge-sort(A, n)

if $n = 1$ **then**

  return A

else

  $B \leftarrow \text{merge-sort}\Big(A\big[1..\lfloor n/2\rfloor\big], \lfloor n/2\rfloor\Big)$

  $C \leftarrow \text{merge-sort}\Big(A\big[\lfloor n/2\rfloor + 1..n\big], \lceil n/2\rceil\Big)$

  return merge$(B, C, \lfloor n/2\rfloor, \lceil n/2\rceil)$

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1: if n = 1 then
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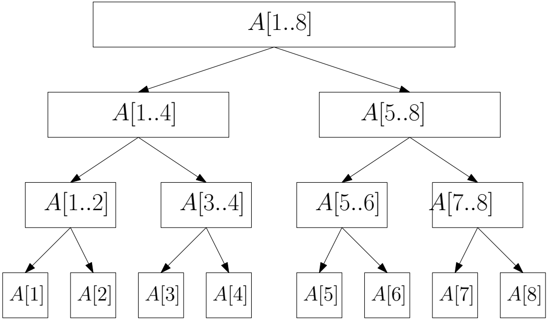
6:     return merge(B, C, ⌊n/2⌋, ⌈n/2⌉)
```

- Divide: trivial
- Conquer: 4, 5
- Combine: 6

Running Time for Merge-Sort



- Each level takes running time $O(n)$
- There are $O(\lg n)$ levels
- Running time $= O(n \lg n)$
- Better than insertion sort

Running Time for Merge-Sort Using Recurrence

- $T(n) =$ running time for sorting n numbers,then

$$T(n) = \begin{cases} O(1) & \text{if } n = 1 \\ T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + O(n) & \text{if } n \geq 2 \end{cases}$$

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- Even simpler: $T(n) = 2T(n/2) + O(n)$. (Implicit assumption: $T(n) = O(1)$ if n is at most some constant.)
- Solving this recurrence, we have $T(n) = O(n \lg n)$ (we shall show how later)

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Example:

$$10 \quad 8 \quad 15 \quad 9 \quad 12$$

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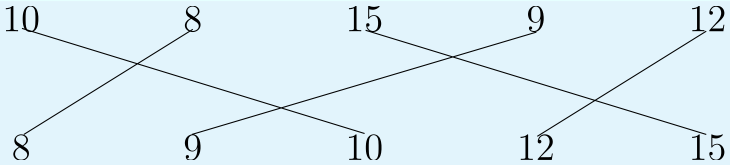

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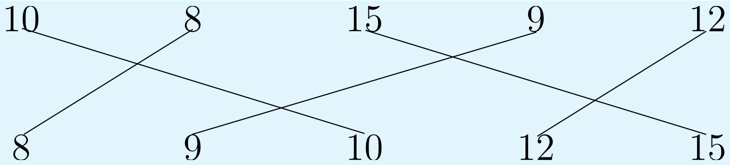


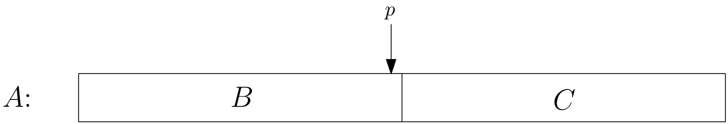
- 4 inversions (for convenience, using numbers, not indices): $(10, 8), (10, 9), (15, 9), (15, 12)$

Naive Algorithm for Counting Inversions

count-inversions(A, n)

```
c ← 0
for every i ← 1 to n − 1 do
    for every j ← i + 1 to n do
        if A[i] > A[j] then c ← c + 1
return c
```

Divide-and-Conquer



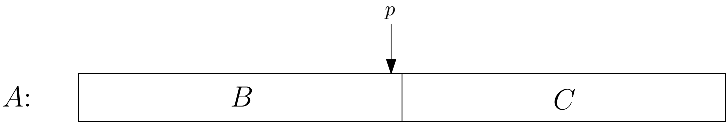
- $p = \lfloor n/2 \rfloor, B = A[1..p], C = A[p+1..n]$
- $$\#\mathsf{invs}(A) = \#\mathsf{invs}(B) + \#\mathsf{invs}(C) + m$$
  $$m = \Big|\big\{(i,j) : B[i] > C[j]\big\}\Big|$$

Q: How fast can we compute m, via trivial algorithm?

A: $O(n^2)$

- Can not improve the $O(n^2)$ time for counting inversions.

Divide-and-Conquer



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Lemma If both B and C are sorted, then we can compute m in $O(n)$ time!

Counting Inversions between B and C

Count pairs i, j such that $B[i] > C[j]$:

B:

3	8	12	20	32	48

total$=$ 0

C:

5	7	9	25	29

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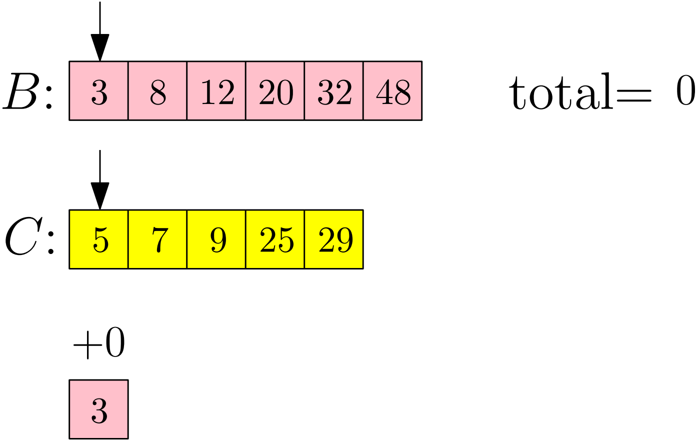

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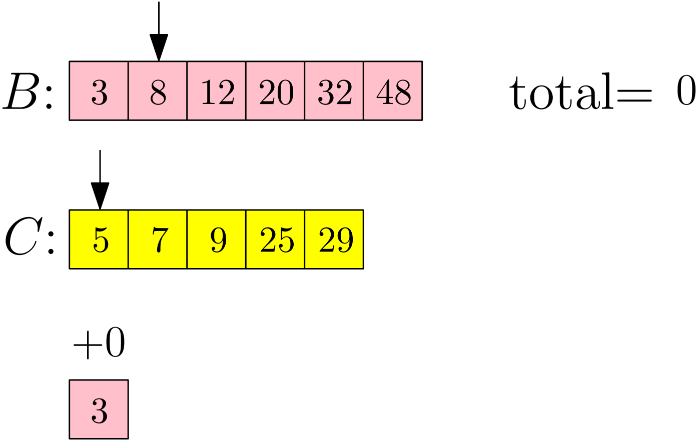

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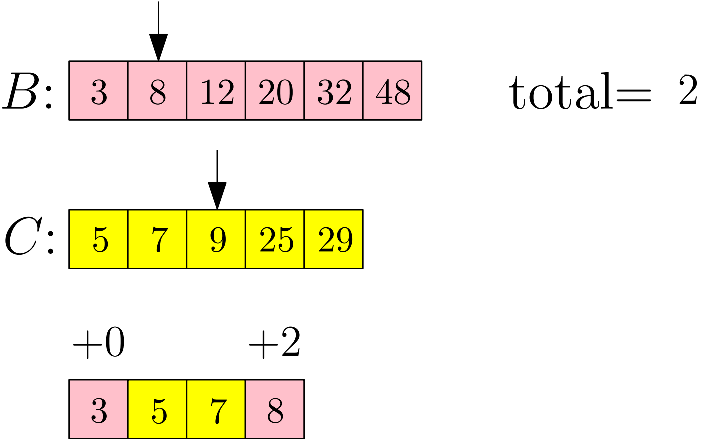

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Count pairs i, j such that $B[i] > C[j]$:

B: | 3 | 8 | 12 | 20 | 32 | 48 |

total$=$ 2

C: | 5 | 7 | 9 | 25 | 29 |

$+0$ $\quad$ $+2$

| 3 | 5 | 7 | 8 |

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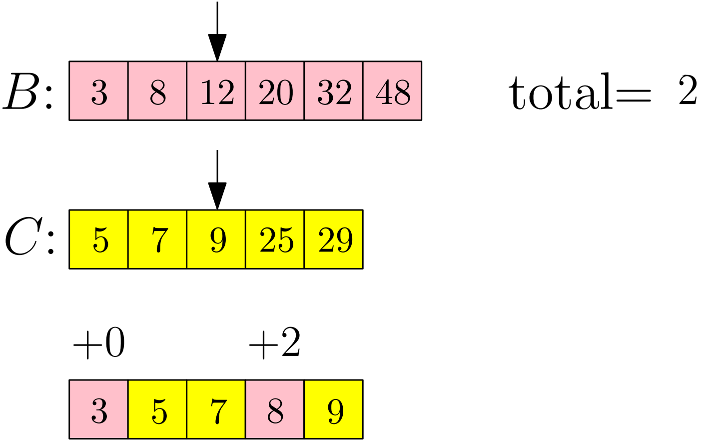

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$+0$			$+2$	
3	5	7	8	9

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Count pairs i, j such that $B[i] > C[j]$:

B:

3	8	12	20	32	48

total= 5

C:

5	7	9	25	29

$+0$			$+2$		$+3$
3	5	7	8	9	12

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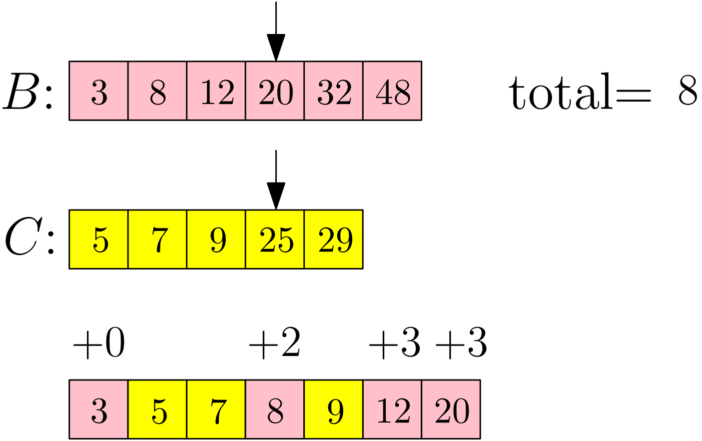

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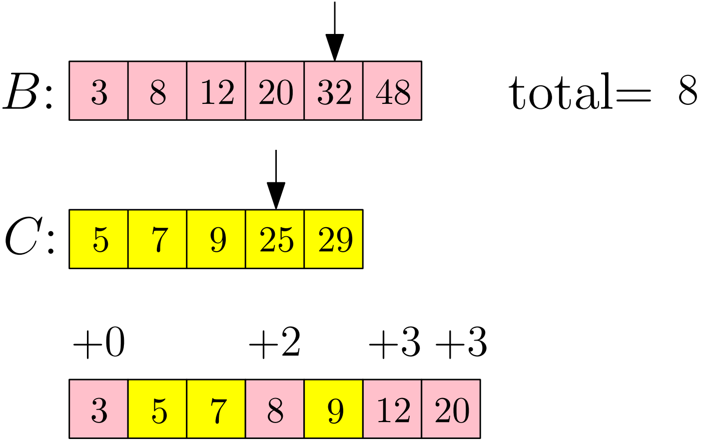

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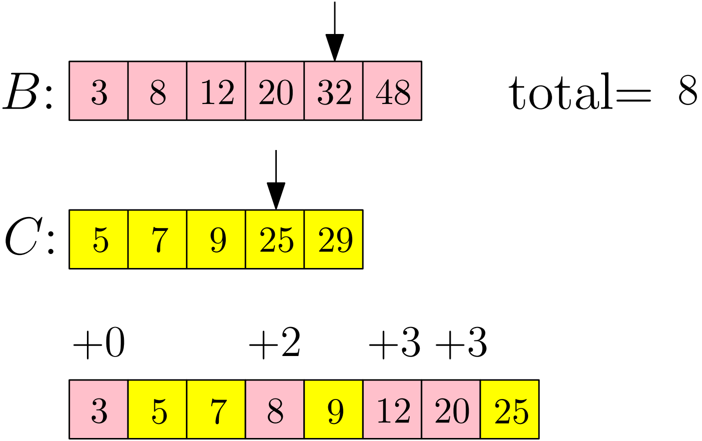

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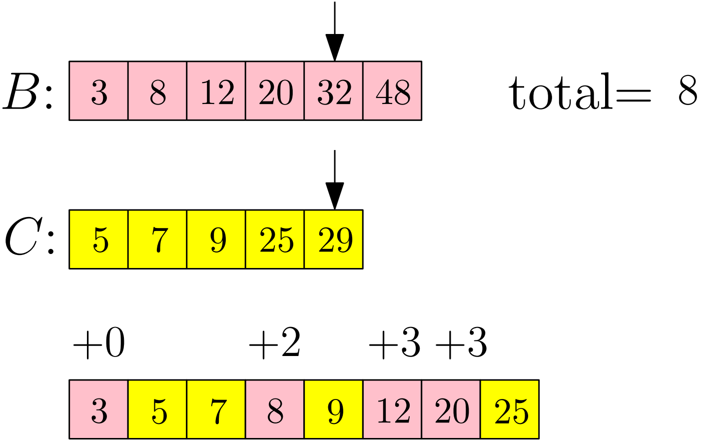

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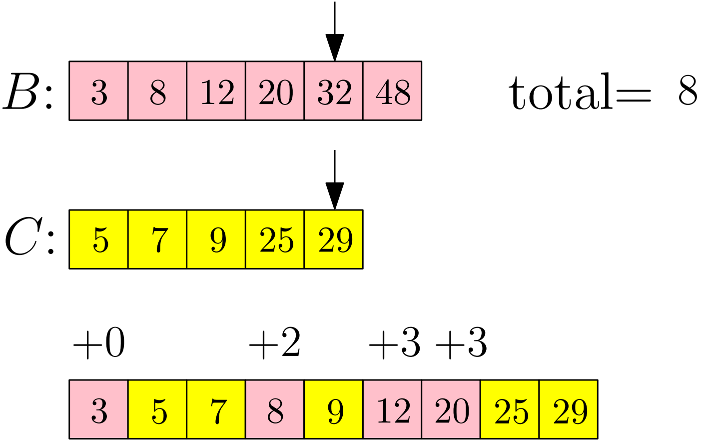

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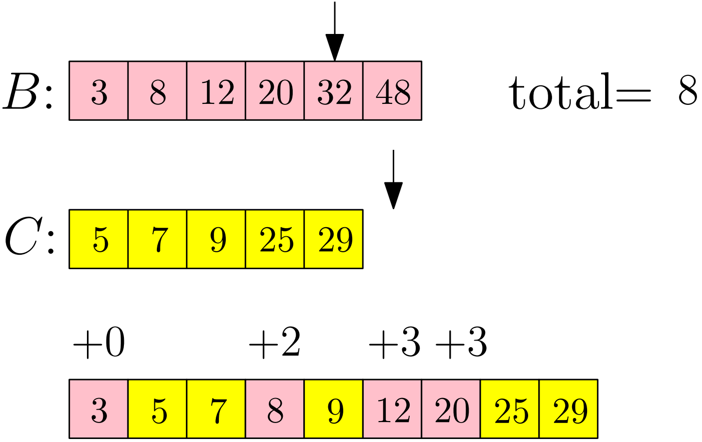

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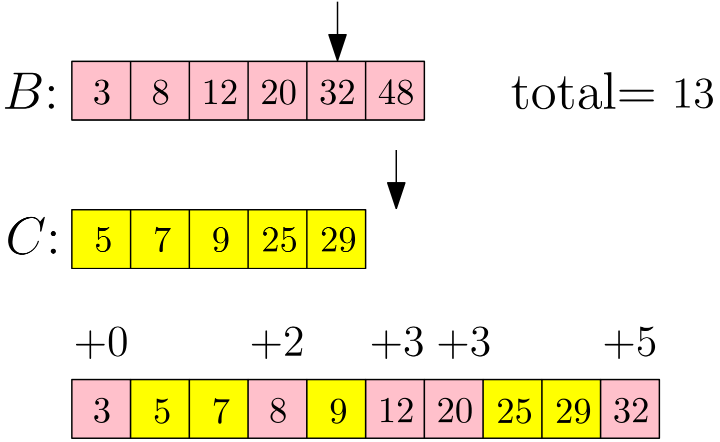

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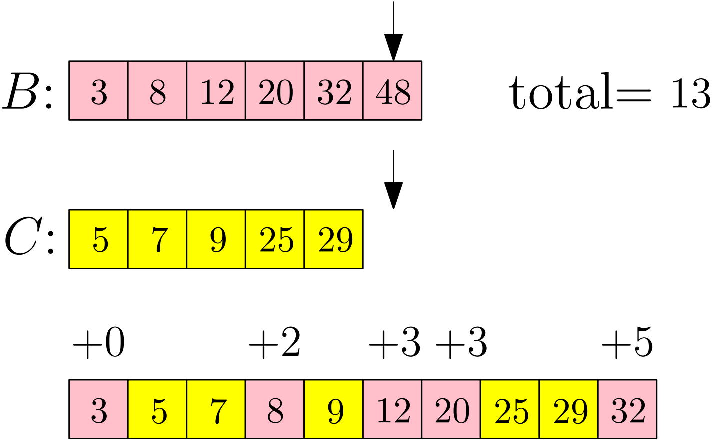

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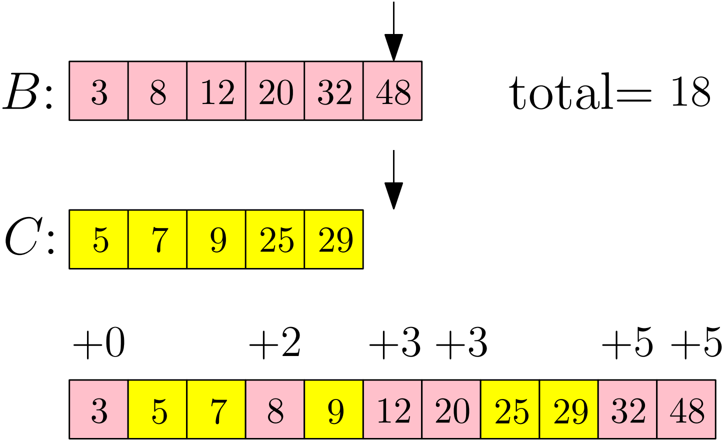

Counting Inversions between B and C

Count pairs i, j such that $B[i] > C[j]$:

B:	3	8	12	20	32	48

total $= 18$

C:	5	7	9	25	29

$+0$			$+2$		$+3$	$+3$			$+5$	$+5$
3	5	7	8	9	12	20	25	29	32	48

Count Inversions between B and C

- Procedure that merges B and C and counts inversions between B and C at the same time

merge-and-count(B, C, n_1, n_2)

> $count \leftarrow 0$;
> $A \leftarrow$ array of size $n_1 + n_2$; $i \leftarrow 1$; $j \leftarrow 1$
> **while** $i \leq n_1$ or $j \leq n_2$ **do**
>   **if** $j > n_2$ or $(i \leq n_1$ and $B[i] \leq C[j])$ **then**
> $A[i + j - 1] \leftarrow B[i]$; $i \leftarrow i + 1$
> $count \leftarrow count + (j - 1)$
>   **else**
> $A[i + j - 1] \leftarrow C[j]$; $j \leftarrow j + 1$
> **return** $(A, count)$

Sort and Count Inversions in A

- A procedure that returns the sorted array of A and counts the number of inversions in A:

sort-and-count(A, n)

if $n = 1$ **then**
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else
    $(B, m_1) \leftarrow$ sort-and-count$\Big(A\big[1..\lfloor n/2\rfloor\big], \lfloor n/2\rfloor\Big)$
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    $(A, m_3) \leftarrow$ merge-and-count$(B, C, \lfloor n/2\rfloor, \lceil n/2\rceil)$
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Quicksort vs Merge-Sort

	Merge Sort	**Quicksort**
Divide	Trivial	Separate small and big numbers
Conquer	Recurse	Recurse
Combine	Merge 2 sorted arrays	Trivial

Quicksort Example

Assumption We can choose median of an array of size n in $O(n)$ time.

29	82	75	64	38	45	94	69	25	76	15	92	37	17	85

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25	15	17	29	38	45	37	64	82	75	94	92	69	76	85

Quicksort

quicksort(A, n)

1: if $n \leq 1$ then return A
2: $x \leftarrow$ lower median of A
3: $A_L \leftarrow$ elements in A that are less than x \\ Divide
4: $A_R \leftarrow$ elements in A that are greater than x \\ Divide
5: $B_L \leftarrow$ quicksort$(A_L, A_L$.size$)$ \\ Conquer
6: $B_R \leftarrow$ quicksort$(A_R, A_R$.size$)$ \\ Conquer
7: $t \leftarrow$ number of times x appear A
8: return the array obtained by concatenating B_L, the array containing t copies of x, and B_R

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Q: How to remove this assumption?

A:

1. There is an algorithm to find median in $O(n)$ time, using divide-and-conquer (we shall not talk about it; it is complicated and not practical)
2. Choose a pivot randomly and pretend it is the median (it is practical)

Quicksort Using A Random Pivot

quicksort(A, n)

if $n \leq 1$ then return A
$x \leftarrow$ a random element of A (x is called a pivot)
$A_L \leftarrow$ elements in A that are less than x \\ Divide
$A_R \leftarrow$ elements in A that are greater than x \\ Divide
$B_L \leftarrow$ quicksort$(A_L, A_L.\text{size})$ \\ Conquer
$B_R \leftarrow$ quicksort$(A_R, A_R.\text{size})$ \\ Conquer
$t \leftarrow$ number of times x appear A
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Randomized Algorithm Model

Assumption There is a procedure to produce a random real number in $[0, 1]$.

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- In practice: use pseudo-random-generator, a deterministic algorithm returning numbers that "look like" random

Randomized Algorithm Model

Assumption There is a procedure to produce a random real number in $[0,1]$.

Q: Can computers really produce random numbers?

A: No! The execution of a computer programs is deterministic!

- In practice: use pseudo-random-generator, a deterministic algorithm returning numbers that "look like" random
- In theory: assume they can.

Quicksort Using A Random Pivot

quicksort(A, n)

if $n \leq 1$ then return A
$x \leftarrow$ a random element of A (x is called a pivot)
$A_L \leftarrow$ elements in A that are less than x \\ Divide
$A_R \leftarrow$ elements in A that are greater than x \\ Divide
$B_L \leftarrow$ quicksort$(A_L, A_L.\text{size})$ \\ Conquer
$B_R \leftarrow$ quicksort$(A_R, A_R.\text{size})$ \\ Conquer
$t \leftarrow$ number of times x appear A
return the array obtained by concatenating B_L, the array containing t copies of x, and B_R

Lemma The expected running time of the algorithm is $O(n \lg n)$.

Quicksort Can Be Implemented as an "In-Place" Sorting Algorithm

- In-Place Sorting Algorithm: an algorithm that only uses "small" extra space.

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- To partition the array into two parts, we only need $O(1)$ extra space.

partition(A, ℓ, r)

$p \leftarrow$ random integer between ℓ and r, swap $A[p]$ and $A[\ell]$
$i \leftarrow \ell, j \leftarrow r$
while true **do**
    while $i < j$ and $A[i] < A[j]$ do $j \leftarrow j - 1$
    if $i = j$ **then** break
    swap $A[i]$ and $A[j]$; $i \leftarrow i + 1$
    while $i < j$ and $A[i] < A[j]$ **do** $i \leftarrow i + 1$
    if $i = j$ **then** break
    swap $A[i]$ and $A[j]$; $j \leftarrow j - 1$
return i

In-Place Implementation of Quick-Sort

quicksort(A, ℓ, r)

```
if ℓ ≥ r then return
m ← patition(A, ℓ, r)
quicksort(A, ℓ, m − 1)
quicksort(A, m + 1, r)
```

- To sort an array A of size n, call quicksort$(A, 1, n)$.

Note: We pass the array A by reference, instead of by copying.

Merge-Sort is Not In-Place

- To merge two arrays, we need a third array with size equaling the total size of two arrays

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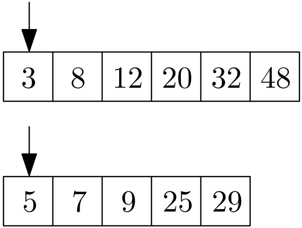

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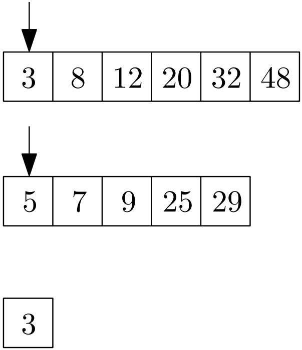

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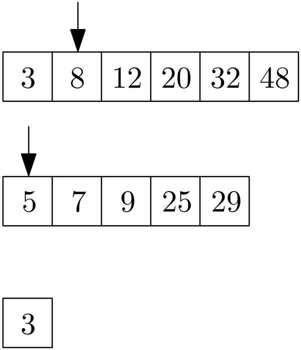

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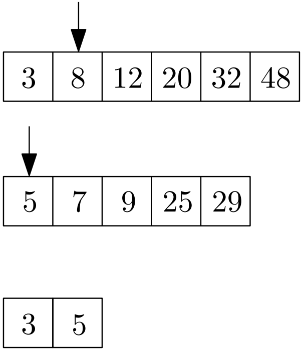

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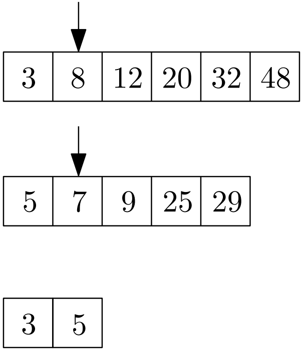

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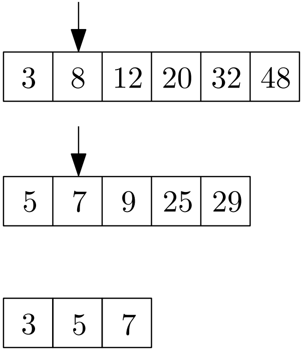

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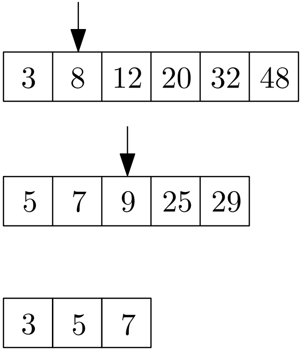

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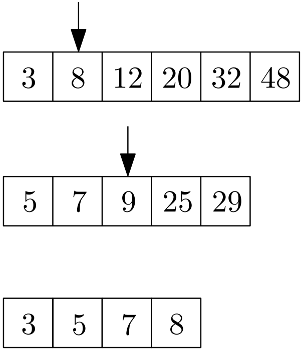

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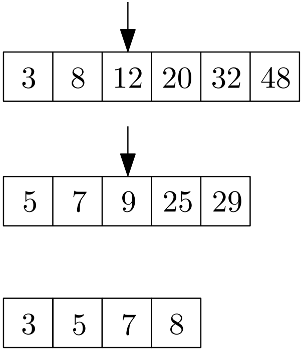

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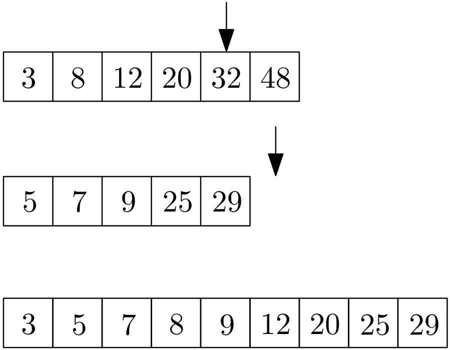

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Outline

Comparison-Based Sorting Algorithms

Q: Can we do better than $O(n \log n)$ for sorting?

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Comparison-Based Sorting Algorithms

- To sort, we are only allowed to compare two elements
- We can not use "internal structures" of the elements

Lemma The (worst-case) running time of any comparison-based sorting algorithm is $\Omega(n \lg n)$.

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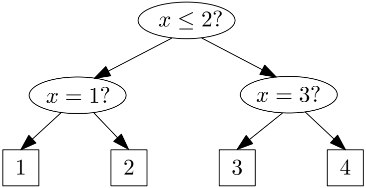

Comparison-Based Sorting Algorithms

Q: Can we do better than $O(n \log n)$ for sorting?

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Selection Problem

Input: a set A of n numbers, and $1 \le i \le n$

Output: the i-th smallest number in A

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- Our goal: $O(n)$ running time

Recall: Quicksort with Median Finder

quicksort(A, n)

if $n \leq 1$ **then return** A
$x \leftarrow$ lower median of A
$A_L \leftarrow$ elements in A that are less than x ▷ Divide
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Selection Algorithm with Median Finder

selection(A, n, i)

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if n = 1 then return A
x ← lower median of A
A_L ← elements in A that are less than x                        ▷ Divide
A_R ← elements in A that are greater than x                     ▷ Divide
if i ≤ A_L.size then
    return selection(A_L, A_L.size, i)                          ▷ Conquer
else if i > n − A_R.size then
    return selection(A_R, A_R.size, i − (n − A_R.size))         ▷ Conquer
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- Solving recurrence: $T(n) = O(n)$

Randomized Selection Algorithm

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- expected running time $= O(n)$

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Input: two polynomials of degree $n-1$

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&(3x^3 + 2x^2 - 5x + 4) \times (2x^3 - 3x^2 + 6x - 5) \\
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& + 4x^5 - 6x^4 + 12x^3 - 10x^2 \\
& - 10x^4 + 15x^3 - 30x^2 + 25x \\
& + 8x^3 - 12x^2 + 24x - 20 \\
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\end{aligned}
$$

- **Input**: $(4, -5, 2, 3), (-5, 6, -3, 2)$
- **Output**: $(-20, 49, -52, 20, 2, -5, 6)$

Naïve Algorithm

polynomial-multiplication(A, B, n)

let $C[k] \leftarrow 0$ for every $k = 0, 1, 2, \cdots, 2n-2$
for $i \leftarrow 0$ to $n-1$ **do**
    for $j \leftarrow 0$ to $n-1$ **do**
        $C[i+j] \leftarrow C[i+j] + A[i] \times B[j]$
return C

Naïve Algorithm

polynomial-multiplication(A, B, n)

```
let C[k] ← 0 for every k = 0, 1, 2, ··· , 2n − 2
for i ← 0 to n − 1 do
    for j ← 0 to n − 1 do
        C[i + j] ← C[i + j] + A[i] × B[j]
return C
```

Running time: $O(n^2)$

Divide-and-Conquer for Polynomial Multiplication

$$p(x) = 3x^3 + 2x^2 - 5x + 4 = (3x + 2)x^2 + (-5x + 4)$$
$$q(x) = 2x^3 - 3x^2 + 6x - 5 = (2x - 3)x^2 + (6x - 5)$$

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- $p(x)$: degree of $n-1$ (assume n is even)
- $p(x) = p_H(x)x^{n/2} + p_L(x)$,
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$$pq = \left(p_H x^{n/2} + p_L\right)\left(q_H x^{n/2} + q_L\right)$$

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- $p(x) = p_H(x)x^{n/2} + p_L(x)$,
- $p_H(x), p_L(x)$: polynomials of degree $n/2 - 1$.

$$\begin{aligned} pq &= \left(p_H x^{n/2} + p_L\right)\left(q_H x^{n/2} + q_L\right) \\ &= p_H q_H x^n + \left(p_H q_L + p_L q_H\right) x^{n/2} + p_L q_L \end{aligned}$$

Divide-and-Conquer for Polynomial Multiplication

$$\begin{aligned} pq &= \left(p_H x^{n/2} + p_L\right)\left(q_H x^{n/2} + q_L\right) \\ &= p_H q_H x^n + \left(p_H q_L + p_L q_H\right) x^{n/2} + p_L q_L \end{aligned}$$

Divide-and-Conquer for Polynomial Multiplication

$$
\begin{aligned}
pq &= \big(p_H x^{n/2} + p_L\big)\big(q_H x^{n/2} + q_L\big) \\
&= p_H q_H x^n + \big(p_H q_L + p_L q_H\big) x^{n/2} + p_L q_L
\end{aligned}
$$

$$
\begin{aligned}
\mathsf{multiply}(p,q) = \mathsf{multiply}(p_H, q_H) \times x^n \\
&+ \big(\mathsf{multiply}(p_H, q_L) + \mathsf{multiply}(p_L, q_H)\big) \times x^{n/2} \\
&+ \mathsf{multiply}(p_L, q_L)
\end{aligned}
$$

Divide-and-Conquer for Polynomial Multiplication

$$
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pq &= \left(p_H x^{n/2} + p_L\right)\left(q_H x^{n/2} + q_L\right) \\
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&+ \mathsf{multiply}(p_L, q_L)
\end{aligned}
$$

- Recurrence: $T(n) = 4T(n/2) + O(n)$

Divide-and-Conquer for Polynomial Multiplication

$$
\begin{aligned}
pq &= \big(p_H x^{n/2} + p_L\big)\big(q_H x^{n/2} + q_L\big) \\
&= p_H q_H x^n + \big(p_H q_L + p_L q_H\big) x^{n/2} + p_L q_L
\end{aligned}
$$

$$
\begin{aligned}
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& + \textsf{multiply}(p_L, q_L)
\end{aligned}
$$

- Recurrence: $T(n) = 4T(n/2) + O(n)$
- $T(n) = O(n^2)$

Reduce Number from 4 to 3

Reduce Number from 4 to 3

$$
\begin{aligned}
pq &= \big(p_H x^{n/2} + p_L\big)\big(q_H x^{n/2} + q_L\big) \\
&= p_H q_H x^n + \big(p_H q_L + p_L q_H\big) x^{n/2} + p_L q_L
\end{aligned}
$$

Reduce Number from 4 to 3

$$\begin{aligned} pq &= \left(p_H x^{n/2} + p_L\right)\left(q_H x^{n/2} + q_L\right) \\ &= p_H q_H x^n + \left(p_H q_L + p_L q_H\right) x^{n/2} + p_L q_L \end{aligned}$$

- $p_H q_L + p_L q_H = (p_H + p_L)(q_H + q_L) - p_H q_H - p_L q_L$

Divide-and-Conquer for Polynomial Multiplication

Divide-and-Conquer for Polynomial Multiplication

$$
\begin{aligned}
r_H &= \mathsf{multiply}(p_H, q_H) \\
r_L &= \mathsf{multiply}(p_L, q_L)
\end{aligned}
$$

Divide-and-Conquer for Polynomial Multiplication

$$r_H = \mathsf{multiply}(p_H, q_H)$$
$$r_L = \mathsf{multiply}(p_L, q_L)$$

$$\begin{aligned}\mathsf{multiply}(p, q) = {} & r_H \times x^n \\ & + \big(\mathsf{multiply}(p_H + p_L, q_H + q_L) - r_H - r_L\big) \times x^{n/2} \\ & + r_L\end{aligned}$$

Divide-and-Conquer for Polynomial Multiplication

$$r_H = \mathsf{multiply}(p_H, q_H)$$
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- Solving Recurrence: $T(n) = 3T(n/2) + O(n)$

Divide-and-Conquer for Polynomial Multiplication

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- Solving Recurrence: $T(n) = 3T(n/2) + O(n)$
- $T(n) = O(n^{\lg_2 3}) = O(n^{1.585})$

Assumption n is a power of 2. Arrays are 0-indexed.

multiply(A, B, n)

```
if n = 1 then return (A[0]B[0])
A_L ← A[0 .. n/2 − 1], A_H ← A[n/2 .. n − 1]
B_L ← B[0 .. n/2 − 1], B_H ← B[n/2 .. n − 1]
C_L ← multiply(A_L, B_L, n/2)
C_H ← multiply(A_H, B_H, n/2)
C_M ← multiply(A_L + A_H, B_L + B_H, n/2)
C ← array of (2n − 1) 0's
for i ← 0 to n − 2 do
    C[i] ← C[i] + C_L[i]
    C[i + n] ← C[i + n] + C_H[i]
    C[i + n/2] ← C[i + n/2] + C_M[i] − C_L[i] − C_H[i]
return C
```

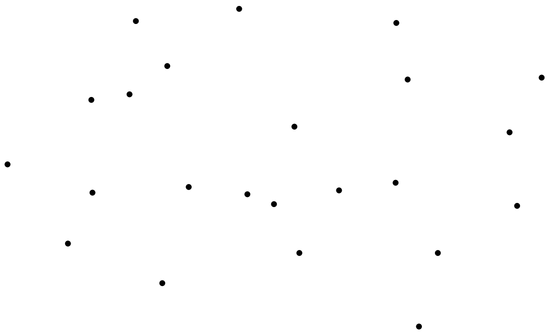
Outline

- Closest pair
- Convex hull
- Matrix multiplication
- FFT(Fast Fourier Transform): polynomial multiplication in $O(n \lg n)$ time

Closest Pair

Input: n points in plane: $(x_1, y_1), (x_2, y_2), \cdots, (x_n, y_n)$

Output: the pair of points that are closest



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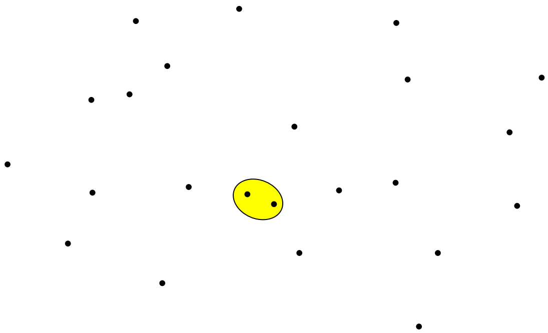

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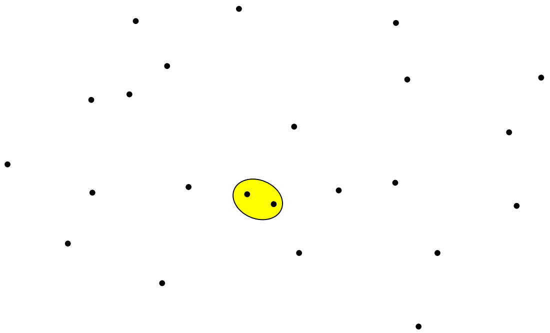
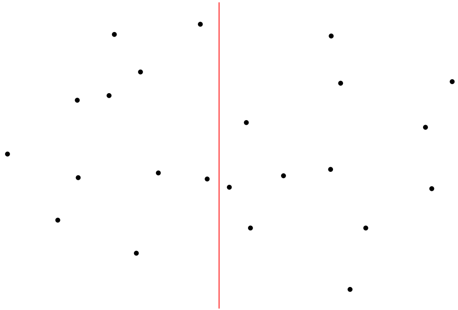


- Trivial algorithm: $O(n^2)$ running time

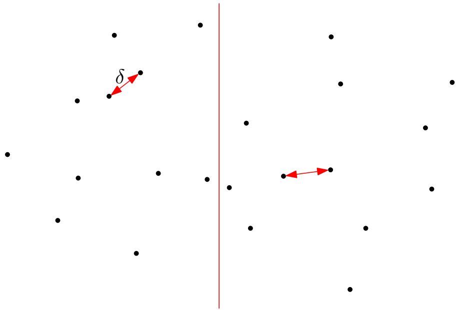
Divide-and-Conquer Algorithm for Closest Pair

- **Divide**: Divide the points into two halves via a vertical line



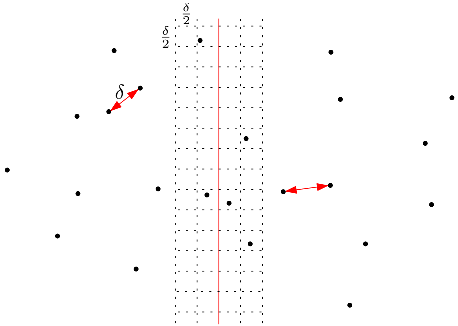
Divide-and-Conquer Algorithm for Closest Pair

- **Divide**: Divide the points into two halves via a vertical line
- **Conquer**: Solve two sub-instances recursively

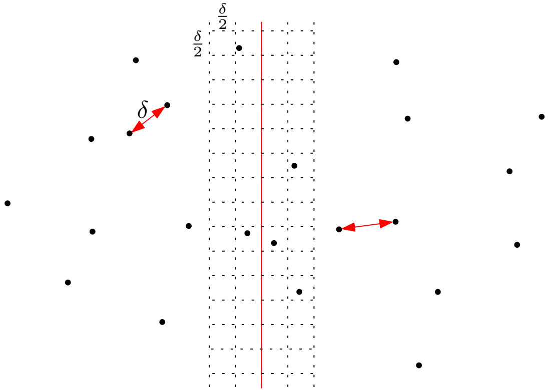


Divide-and-Conquer Algorithm for Closest Pair

- **Divide**: Divide the points into two halves via a vertical line
- **Conquer**: Solve two sub-instances recursively
- **Combine**: Check if there is a closer pair between left-half and right-half



Divide-and-Conquer Algorithm for Closest Pair



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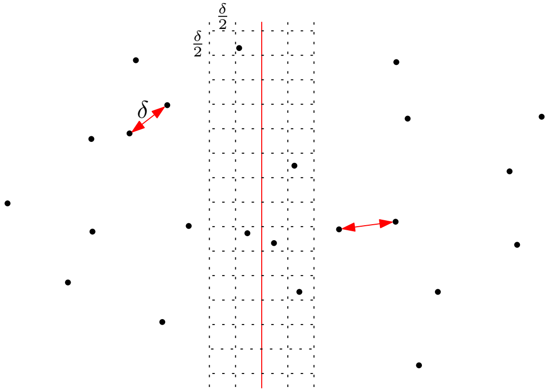
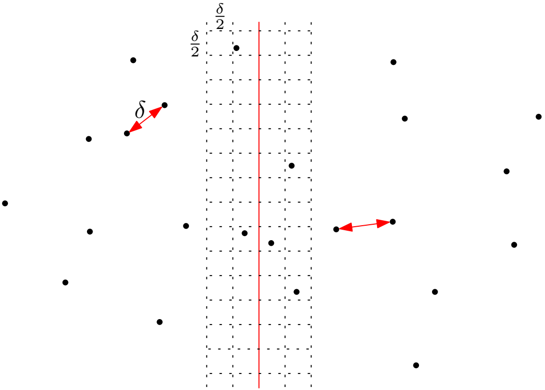


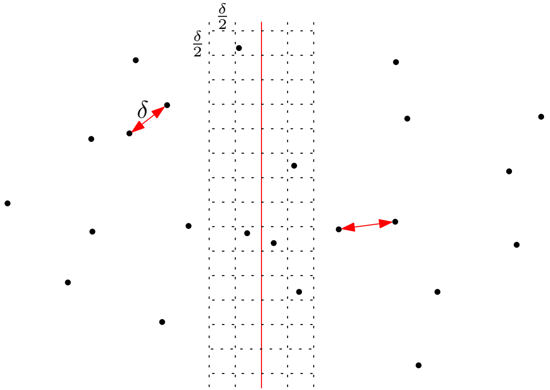
- Each box contains at most one pair

Divide-and-Conquer Algorithm for Closest Pair



- Each box contains at most one pair
- For each point, only need to consider $O(1)$ boxes nearby

Divide-and-Conquer Algorithm for Closest Pair



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- time for combine $= O(n)$ (many technicalities omitted)

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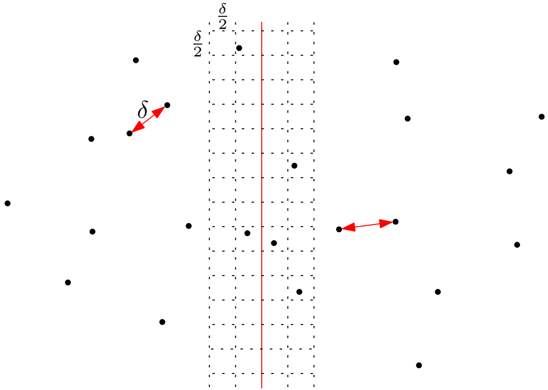

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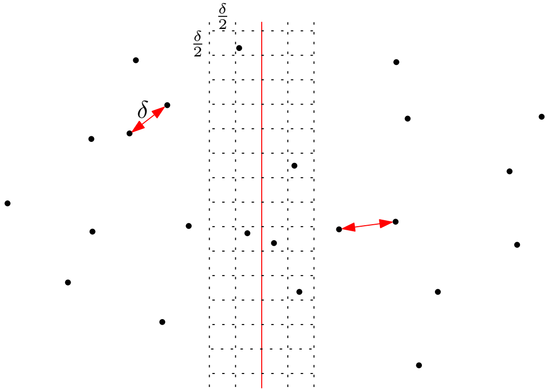
Divide-and-Conquer Algorithm for Closest Pair



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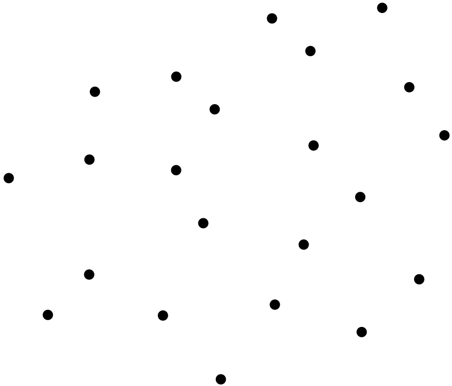
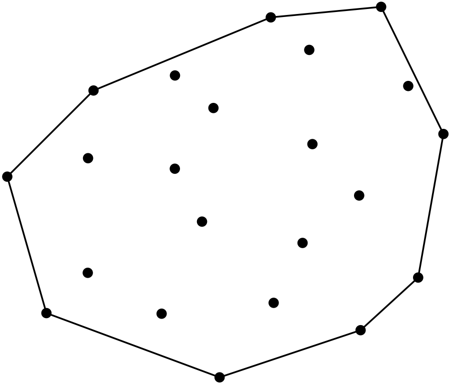
- Recurrence: $T(n) = 2T(n/2) + O(n)$
- Running time: $O(n \lg n)$

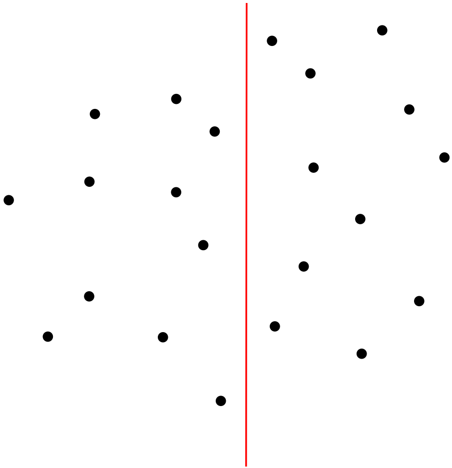
$O(n \lg n)$-Time Algorithm for Convex Hull



$O(n \lg n)$-Time Algorithm for Convex Hull



$O(n \lg n)$-Time Algorithm for Convex Hull



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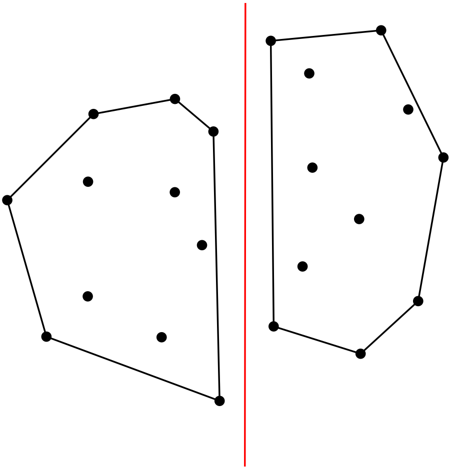
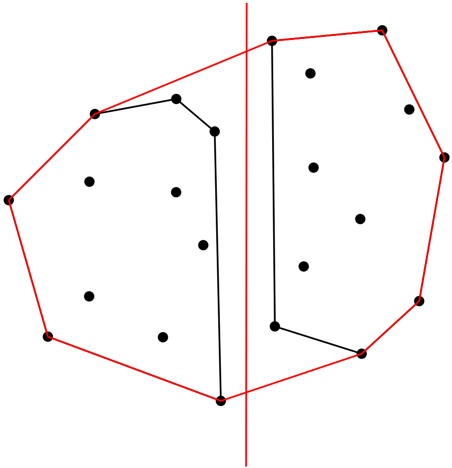


$O(n \lg n)$-Time Algorithm for Convex Hull



Strassen's Algorithm for Matrix Multiplication

Matrix Multiplication

Input: two $n \times n$ matrices A and B

Output: $C = AB$

Strassen's Algorithm for Matrix Multiplication

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Output: $C = AB$

Naive Algorithm: matrix-multiplication(A, B, n)

```
for i ← 1 to n do
    for j ← 1 to n do
        C[i, j] ← 0
        for k ← 1 to n do
            C[i, j] ← C[i, j] + A[i, k] × B[k, j]
return C
```

Strassen's Algorithm for Matrix Multiplication

Matrix Multiplication

Input: two $n \times n$ matrices A and B

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    for $j \leftarrow 1$ to n **do**
        $C[i, j] \leftarrow 0$
        for $k \leftarrow 1$ to n **do**
            $C[i, j] \leftarrow C[i, j] + A[i, k] \times B[k, j]$
return C

- running time $= O(n^3)$

Try to Use Divide-and-Conquer

$$A = \begin{array}{|c|c|l}\hline A_{11} & \overbrace{A_{12}}^{n/2} & \}\,n/2 \\ \hline A_{21} & A_{22} & \\ \hline \end{array} \qquad B = \begin{array}{|c|c|l}\hline B_{11} & \overbrace{B_{12}}^{n/2} & \}\,n/2 \\ \hline B_{21} & B_{22} & \\ \hline \end{array}$$

- $C = \begin{pmatrix} A_{11}B_{11} + A_{12}B_{21} & A_{11}B_{12} + A_{12}B_{22} \\ A_{21}B_{11} + A_{22}B_{21} & A_{21}B_{12} + A_{22}B_{22} \end{pmatrix}$
- matrix_multiplication(A, B) recursively calls matrix_multiplication(A_{11}, B_{11}), matrix_multiplication(A_{12}, B_{21}), $\cdots$

Try to Use Divide-and-Conquer

$$A = \begin{array}{|c|c|l}\multicolumn{1}{c}{} & \multicolumn{1}{c}{\overbrace{\hphantom{A_{12}}}^{n/2}} & \\ \hline A_{11} & A_{12} & \}\,n/2 \\ \hline A_{21} & A_{22} & \\ \hline \end{array} \qquad B = \begin{array}{|c|c|l}\multicolumn{1}{c}{} & \multicolumn{1}{c}{\overbrace{\hphantom{B_{12}}}^{n/2}} & \\ \hline B_{11} & B_{12} & \}\,n/2 \\ \hline B_{21} & B_{22} & \\ \hline \end{array}$$

- $C = \begin{pmatrix} A_{11}B_{11} + A_{12}B_{21} & A_{11}B_{12} + A_{12}B_{22} \\ A_{21}B_{11} + A_{22}B_{21} & A_{21}B_{12} + A_{22}B_{22} \end{pmatrix}$
- matrix_multiplication(A, B) recursively calls matrix_multiplication(A_{11}, B_{11}), matrix_multiplication(A_{12}, B_{21}), $\cdots$
- Recurrence for running time: $T(n) = 8T(n/2) + O(n^2)$
- $T(n) = O(n^3)$

Strassen's Algorithm

- $T(n) = 8T(n/2) + O(n^2)$
- Strassen's Algorithm: improve the number of multiplications from 8 to 7!
- New recurrence: $T(n) = 7T(n/2) + O(n^2)$

Strassen's Algorithm

- $T(n) = 8T(n/2) + O(n^2)$
- Strassen's Algorithm: improve the number of multiplications from 8 to 7!
- New recurrence: $T(n) = 7T(n/2) + O(n^2)$
- Solving Recurrence $T(n) = O(n^{\log_2 7}) = O(n^{2.808})$

Outline

Methods for Solving Recurrences

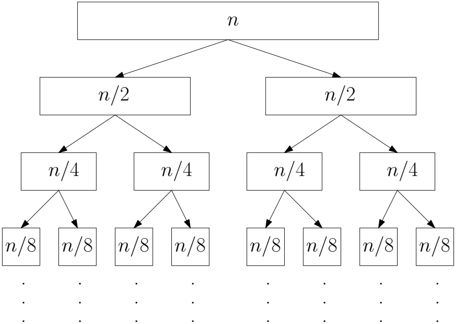
- The recursion-tree method
- The master theorem

Recursion-Tree Method

- $T(n) = 2T(n/2) + O(n)$

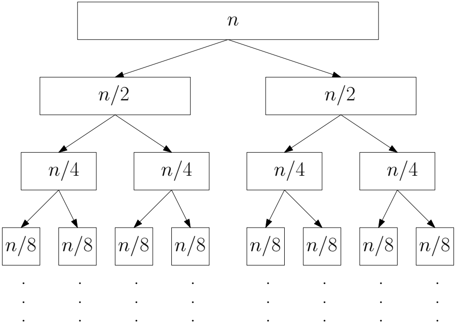
Recursion-Tree Method

- $T(n) = 2T(n/2) + O(n)$



Recursion-Tree Method

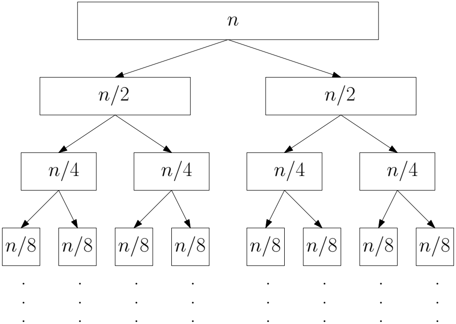
- $T(n) = 2T(n/2) + O(n)$



- Each level takes running time $O(n)$

Recursion-Tree Method

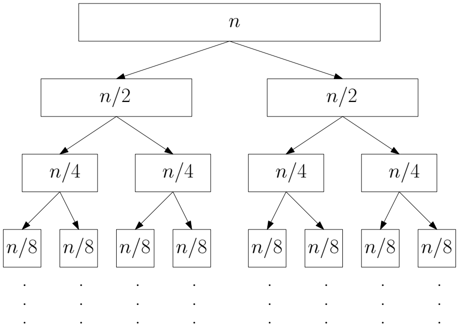
- $T(n) = 2T(n/2) + O(n)$



- Each level takes running time $O(n)$
- There are $O(\lg n)$ levels

Recursion-Tree Method

- $T(n) = 2T(n/2) + O(n)$



- Each level takes running time $O(n)$
- There are $O(\lg n)$ levels
- Running time $= O(n \lg n)$

Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$

Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$

$$\boxed{n}$$

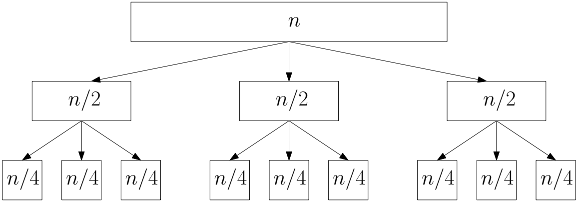
Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$

$$
\begin{array}{ccc}
 & n & \\
\swarrow & \downarrow & \searrow \\
n/2 & n/2 & n/2
\end{array}
$$

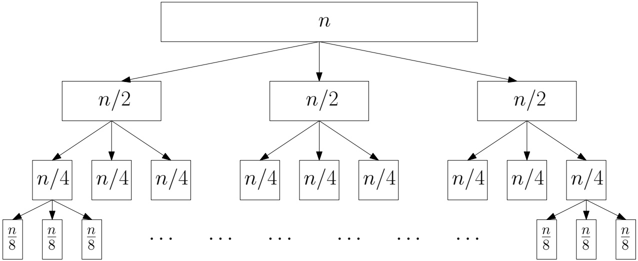
Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$



Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$



Recursion-Tree Method

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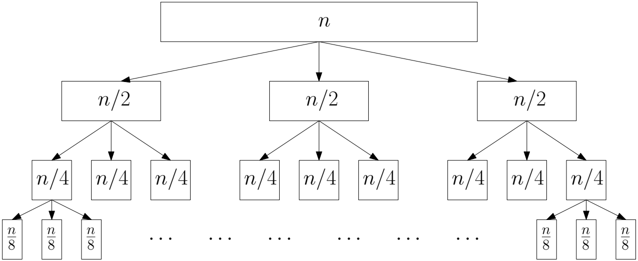
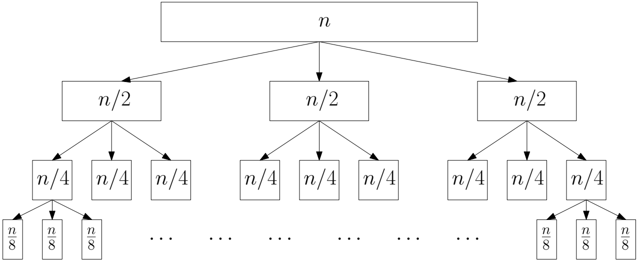


- Total running time at level i?

Recursion-Tree Method

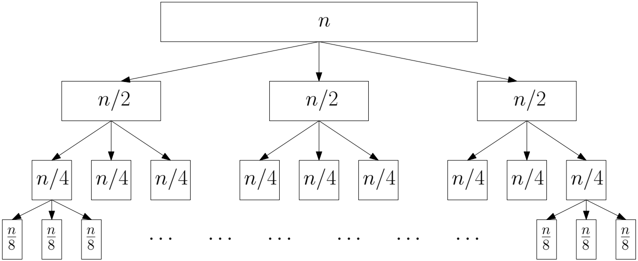
- $T(n) = 3T(n/2) + O(n)$



- Total running time at level i? $\frac{n}{2^i} \times 3^i = \left(\frac{3}{2}\right)^i n$

Recursion-Tree Method

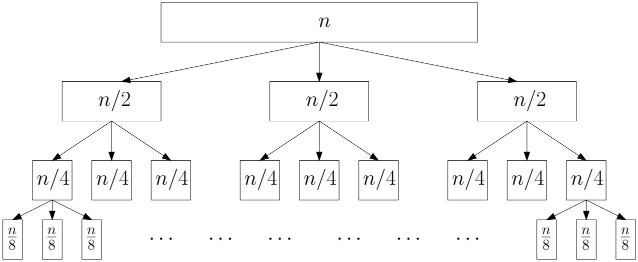
- $T(n) = 3T(n/2) + O(n)$



- Total running time at level i? $\frac{n}{2^i} \times 3^i = \left(\frac{3}{2}\right)^i n$
- Index of last level?

Recursion-Tree Method

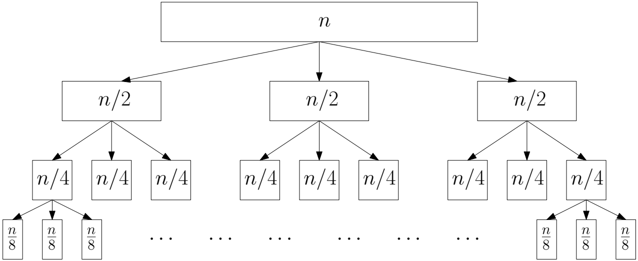
- $T(n) = 3T(n/2) + O(n)$



- Total running time at level i? $\frac{n}{2^i} \times 3^i = \left(\frac{3}{2}\right)^i n$
- Index of last level? $\lg_2 n$

Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n)$



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Recursion-Tree Method

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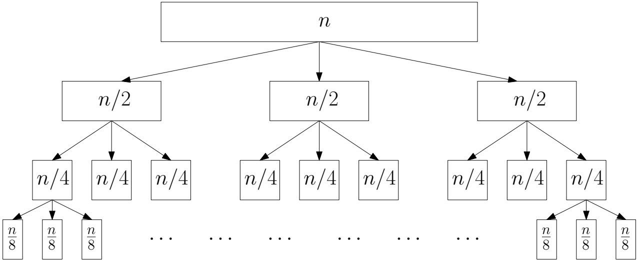


- Total running time at level i? $\frac{n}{2^i} \times 3^i = \left(\frac{3}{2}\right)^i n$
- Index of last level? $\lg_2 n$
- Total running time?

$$\sum_{i=0}^{\lg_2 n} \left(\frac{3}{2}\right)^i n = O\left(n\left(\frac{3}{2}\right)^{\lg_2 n}\right) = O(3^{\lg_2 n}) = O(n^{\lg_2 3}).$$

Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n^2)$

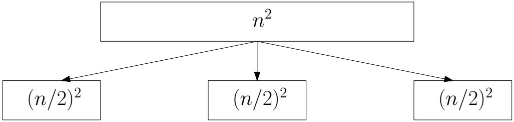
Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n^2)$

$$\boxed{n^2}$$

Recursion-Tree Method

- $T(n) = 3T(n/2) + O(n^2)$



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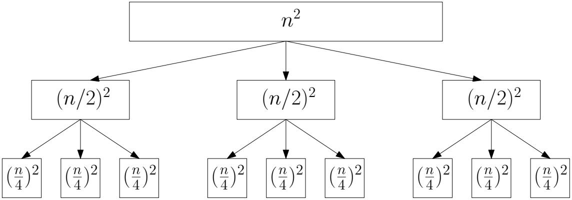

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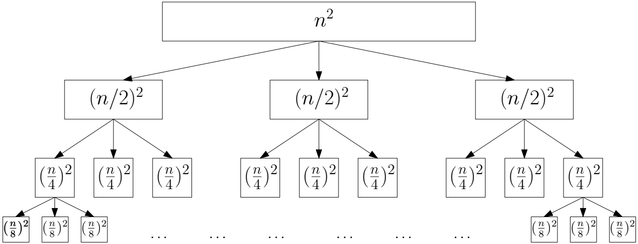

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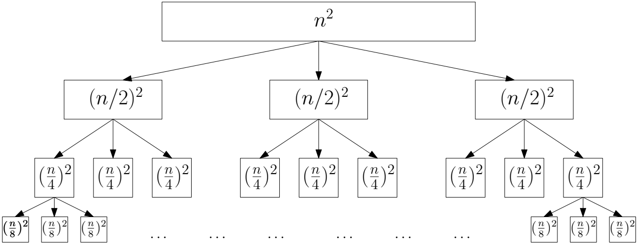

- Total running time at level i?

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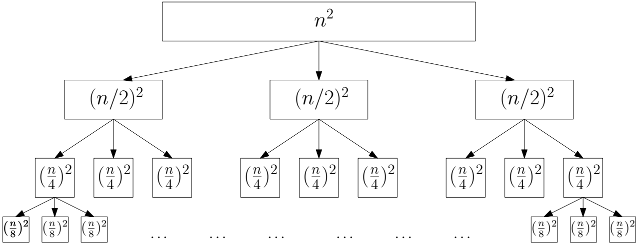

- Total running time at level i? $\left(\frac{n}{2^i}\right)^2 \times 3^i = \left(\frac{3}{4}\right)^i n^2$

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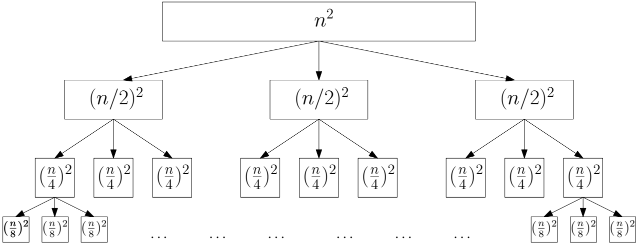

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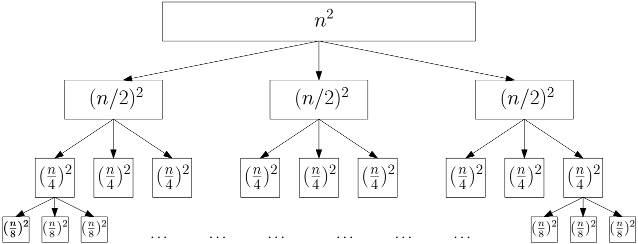

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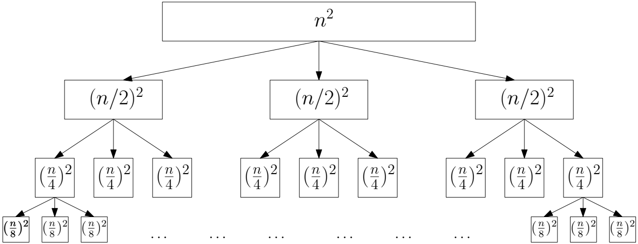

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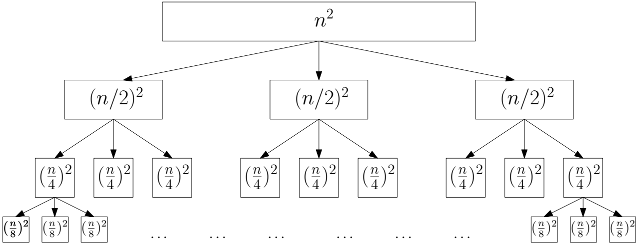

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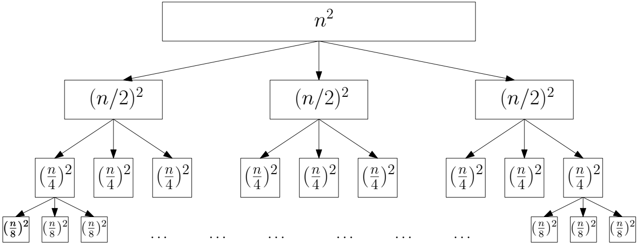

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Master Theorem

Recurrences	a	b	c	time
$T(n) = 2T(n/2) + O(n)$				$O(n \lg n)$
$T(n) = 3T(n/2) + O(n)$				$O(n^{\lg_2 3})$
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Theorem $T(n) = aT(n/b) + O(n^c)$, where $a \geq 1, b > 1, c \geq 0$ are constants. Then,

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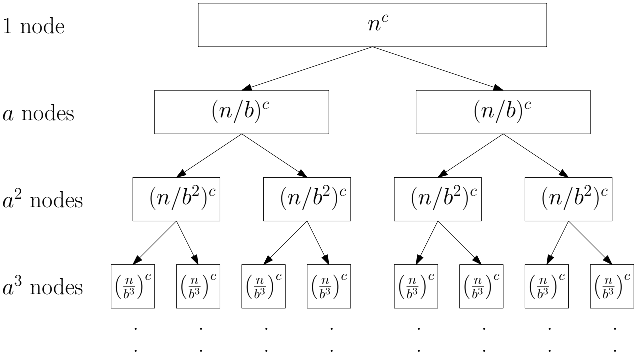
Proof of Master Theorem Using Recursion Tree

$$T(n) = aT(n/b) + O(n^c)$$



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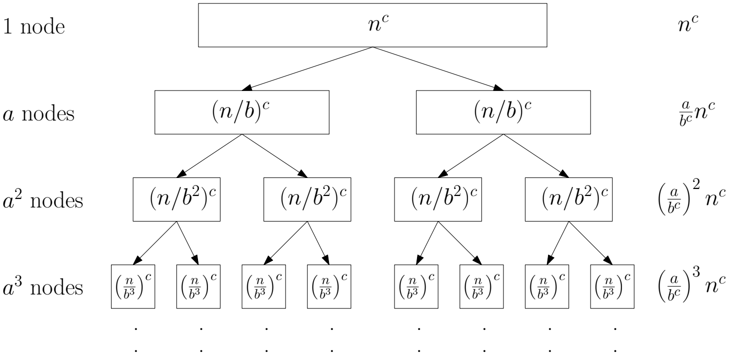

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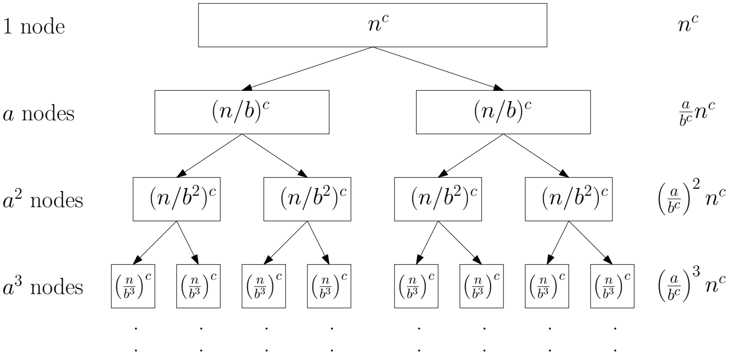

- $c < \lg_b a$: bottom-level dominates: $\left(\frac{a}{b^c}\right)^{\lg_b n} n^c = n^{\lg_b a}$

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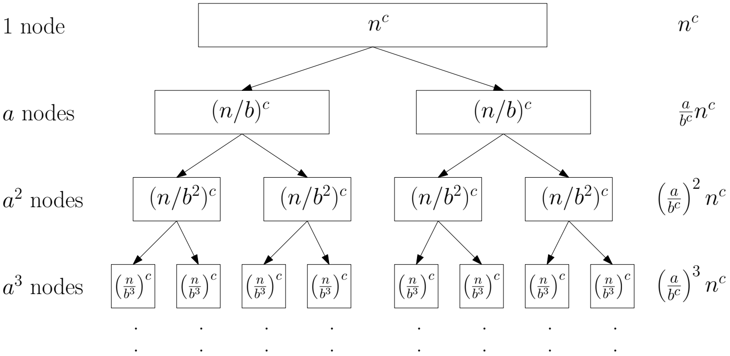

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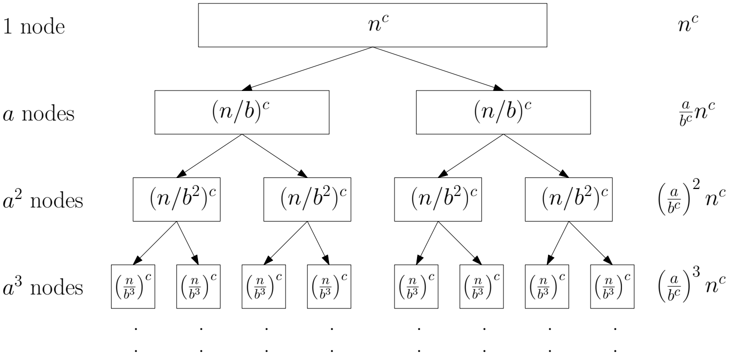

- $c < \lg_b a$: bottom-level dominates: $\left(\frac{a}{b^c}\right)^{\lg_b n} n^c = n^{\lg_b a}$
- $c = \lg_b a$: all levels have same time: $n^c \lg_b n = O(n^c \lg n)$
- $c > \lg_b a$: top-level dominates: $O(n^c)$

Outline

Fibonacci Numbers

- $F_0 = 0, F_1 = 1$
- $F_n = F_{n-1} + F_{n-2}, \forall n \geq 2$
- Fibonacci sequence: $0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, \cdots$

n-th Fibonacci Number

Input: integer $n > 0$

Output: F_n

Computing F_n : Stupid Divide-and-Conquer Algorithm

Fib(n)

if $n = 0$ return 0
if $n = 1$ return 1
return Fib$(n-1)$ + Fib$(n-2)$

Q: Is the running time of the algorithm polynomial or exponential in n?

Computing F_n : Stupid Divide-and-Conquer Algorithm

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- Running time is at least $\Omega(F_n)$

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Q: Is the running time of the algorithm polynomial or exponential in n?

A: Exponential

- Running time is at least $\Omega(F_n)$
- F_n is exponential in n

Computing F_n: Reasonable Algorithm

Fib(n)

$F[0] \leftarrow 0$
$F[1] \leftarrow 1$
for $i \leftarrow 2$ to n **do**
    $F[i] \leftarrow F[i-1] + F[i-2]$
return $F[n]$

- Dynamic Programming

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- Dynamic Programming
- Running time = ?

Computing F_n: Reasonable Algorithm

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for i ← 2 to n do
    F[i] ← F[i − 1] + F[i − 2]
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```

- Dynamic Programming
- Running time $= O(n)$

Computing F_n: Even Better Algorithm

$$\begin{pmatrix} F_n \\ F_{n-1} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F_{n-1} \\ F_{n-2} \end{pmatrix}$$

$$\begin{pmatrix} F_n \\ F_{n-1} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^2 \begin{pmatrix} F_{n-2} \\ F_{n-3} \end{pmatrix}$$

$$\cdots$$

$$\begin{pmatrix} F_n \\ F_{n-1} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^{n-1} \begin{pmatrix} F_1 \\ F_0 \end{pmatrix}$$

power(n)

if $n = 0$ then return $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$
$R \leftarrow$ power$(\lfloor n/2 \rfloor)$
$R \leftarrow R \times R$
if n is odd then $R \leftarrow R \times \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$
return R

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- Recurrence for running time?

power(n)

```
if n = 0 then return ( 1 0 ; 0 1 )
R ← power(⌊n/2⌋)
R ← R × R
if n is odd then R ← R × ( 1 1 ; 1 0 )
return R
```

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Fib(n)

if $n = 0$ then return 0
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- Recurrence for running time? $T(n) = T(n/2) + O(1)$
- $T(n) = O(\lg n)$

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- Even printing $F(n)$ requires time much larger than $O(\lg n)$

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Fixing the Problem

To compute F_n, we need $O(\lg n)$ basic arithmetic operations on integers

Summary: Divide-and-Conquer

- **Divide**: Divide instance into many smaller instances
- **Conquer**: Solve each of smaller instances recursively and separately
- **Combine**: Combine solutions to small instances to obtain a solution for the original big instance

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- **Divide**: Divide instance into many smaller instances
- **Conquer**: Solve each of smaller instances recursively and separately
- **Combine**: Combine solutions to small instances to obtain a solution for the original big instance
- Write down recurrence for running time
- Solve recurrence using master theorem

Summary: Divide-and-Conquer

- Merge sort, quicksort, count-inversions, closest pair, $\cdots$:
  $T(n) = 2T(n/2) + O(n) \Rightarrow T(n) = O(n \lg n)$

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- Usually, designing better algorithm for "combine" step is key to improve running time