CSE 431/531: Algorithm Analysis and Design

Fall 2021

## Homework 3

Instructor: Shi Li Deadline: 11/4/2021

Your Name: \_\_\_\_\_ Your Student ID: \_\_\_\_\_

Problems	1	2	3	Total
Max. Score	20	30	30	80
Your Score				

**Problem 1.** For each of the following recurrences, use the master theorem to give the tight asymptotic upper bound.

(a) 
$$T(n) = 4T(n/3) + O(n)$$
.  $T(n) = O(\underline{\hspace{1cm}})$ .

(b) 
$$T(n) = 3T(n/3) + O(n)$$
.  $T(n) = O(\underline{\hspace{1cm}})$ .

(c) 
$$T(n) = 4T(n/2) + O(n^2\sqrt{n}).$$
  $T(n) = O(\underline{\hspace{1cm}}).$ 

(d) 
$$T(n) = 8T(n/2) + O(n^3)$$
.  $T(n) = O(\underline{\hspace{1cm}})$ .

**Problem 2.** We consider the following problem of counting stronger inversions. Given an array A of n positive integers, a pair  $i, j \in \{1, 2, 3, \dots, n\}$  of indices is called a *strong inversion* if i < j and A[i] > 2A[j]. The goal of the problem is to count the number of strong inversions for a given array A.

Give a divide-and-conquer algorithm that runs in  $O(n \log n)$  time to solve the problem. Write down the recurrence for the running time, and use the master theorem to show that the running time is indeed  $O(n \log n)$ .

**Problem 3.** Given two sorted arrays A and B with total size n, and a positive integer  $k \leq n$ , you need to design an  $O(\log n)$ -time algorithm that outputs the k-th smallest number in the union of A and B. You need to prove that the running time of your algorithm is indeed  $O(\log n)$ .

For example, if A = [3, 5, 12, 18, 50], B = [2, 7, 11, 30], and k = 4 then you need to output 7 since the union of A and B is [2, 3, 5, 7, 11, 12, 18, 30, 50] after sorting.