

Monday October 4

Find and fix any and all mistakes with the following:

- (a) $(1)_2 = (1)_8$ no mistakes: multiple representation of one.
 (b) $(142)_{10} \neq (142)_{16}$ $(142)_{16} = 1 \cdot 16^2 + 4 \cdot 16 + 2 \cdot 16^0 = 256 + 64 + 2 = 322$
 (c) $(20)_{10} = (10100)_2$ no mistakes: $(10100)_2 = 16 + 4 = 20$
 (d) $(35)_8 = (1D)_{16}$ no mistakes: $(35)_8 = 3 \cdot 8^1 + 5 \cdot 8^0 = 24 + 5 = 29$
 $(1D)_{16} = 1 \cdot 16^1 + 13 \cdot 16^0 = 16 + 13 = 29$

Recall the definition of base expansion we discussed:

Definition For b an integer greater than 1 and n a positive integer, the **base b expansion of n** is

$$(a_{k-1} \cdots a_1 a_0)_b$$

where k is a positive integer, a_0, a_1, \dots, a_{k-1} are nonnegative integers less than b , $a_{k-1} \neq 0$, and

$$n = \sum_{i=0}^{k-1} a_i b^i$$

Notice: The base b expansion of a positive integer n is a string over the alphabet $\{x \in \mathbb{N} \mid x < b\}$ whose leftmost character is nonzero.

Base b	Collection of possible coefficients in base b expansion of a positive integer
Binary ($b = 2$)	$\{0, 1\}$
Ternary ($b = 3$)	$\{0, 1, 2\}$
Octal ($b = 8$)	$\{0, 1, 2, 3, 4, 5, 6, 7\}$
Decimal ($b = 10$)	$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
Hexadecimal ($b = 16$)	$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$ letter coefficient symbols represent numerical values $(A)_{16} = (10)_{10}$ $(B)_{16} = (11)_{10}$ $(C)_{16} = (12)_{10}$ $(D)_{16} = (13)_{10}$ $(E)_{16} = (14)_{10}$ $(F)_{16} = (15)_{10}$

We write an algorithm for converting from base b_1 expansion to base b_2 expansion:

- Input: base b_1 expansion of positive integer n
- ① Convert to base 10 expansion using definition
 - ② Convert value to base b_2 expansion using procedure base2 and output its result.

Least Significant First

(^{most significant} — — — — ^{least significant})_b

Calculating base b expansion, from right

```

1 procedure baseb2( $n, b$ : positive integers with  $b > 1$ )
2    $q := n$ 
3    $k := 0$ 
4   while  $q \neq 0$ 
5      $a_k := q \bmod b$ 
6      $q := q \div b$ 
7      $k := k + 1$ 
8   return  $(a_{k-1}, \dots, a_0)$  ( $(a_{k-1} \dots a_0)_b$  is the base  $b$  expansion of  $n$ )

```

n	b	q	k	a_k	$q \neq 0?$
17	3	17	0	$a_0 = 2$	T
		5	1	$a_1 = 2$	T
		1	2	$a_2 = 1$	T
		0			F

$$\text{Idea: } n = a_{k-1} \underbrace{b^{k-1} + \dots + a_1 b}_{n \div b} + \underbrace{a_0}_{n \bmod b}$$

$$(17)_{10} = (1 \ 2 \ 2)_3$$

$$\begin{aligned} 17 &= \boxed{5} \cdot 3 + \boxed{2} \\ 5 &= \boxed{1} \cdot 3 + \boxed{2} \\ 1 &= \boxed{0} \cdot 3 + \boxed{1} \end{aligned}$$

For every positive integer n and base (an integer > 1) b , we can find a base b expansion of n .

number of columns
r

Definition For b an integer greater than 1, w a positive integer, and n a nonnegative integer with $n < b^w$ the **base b fixed-width w expansion** of n is ≥ 0

$$(a_{w-1} \cdots a_1 a_0)_{b,w}$$

where a_0, a_1, \dots, a_{w-1} are nonnegative integers less than b and

$$n = \sum_{i=0}^{w-1} a_i b^i$$

allowing leading zeros in fixed-width expansions.

$$\left(\frac{1}{8} \frac{1}{4} \frac{1}{2} \frac{1}{1}\right)_2 = 15$$

Decimal $b = 10$	Binary $b = 2$	Binary fixed-width 10 $b = 2, w = 10$	Binary fixed-width 7 $b = 2, w = 7$	Binary fixed-width 4 $b = 2, w = 4$
$(20)_{10}$	$(10100)_2$ (a)	$(0000010100)_{2,10}$ (b)	$(0010100)_{2,7}$ (c)	$(10100)_{2,4}$ (d)

Definition For b an integer greater than 1, w a positive integer, w' a positive integer, and x a real number the **base b fixed-width expansion of x with integer part width w and fractional part width w'** is $(a_{w-1} \cdots a_1 a_0 . c_1 \cdots c_{w'})_{b,w,w'}$ where $a_0, a_1, \dots, a_{w-1}, c_1, \dots, c_{w'}$ are nonnegative integers less than b and

$$x \geq \sum_{i=0}^{w-1} a_i b^i + \sum_{j=1}^{w'} c_j b^{-j} \quad \text{and} \quad x < \sum_{i=0}^{w-1} a_i b^i + \sum_{j=1}^{w'} c_j b^{-j} + b^{-w'}$$

3.75 in fixed-width binary, integer part width 2, fractional part width 8	$(\frac{1}{2^2} \frac{1}{2^1} \frac{1}{2^0} \frac{1}{2^{-1}} \frac{1}{2^{-2}} \frac{0}{2^{-3}} \frac{0}{2^{-4}} \frac{0}{2^{-5}} \frac{0}{2^{-6}} \frac{0}{2^{-7}} \frac{0}{2^{-8}})_{2,2,8}$ radix point base / fractional int width
0.1 in fixed-width binary, integer part width 2, fractional part width 8	$(\frac{0}{2^2} \frac{0}{2^1} \frac{0}{2^0} \frac{0}{2^{-1}} \frac{0}{2^{-2}} \frac{1}{2^{-3}} \frac{1}{2^{-4}} \frac{0}{2^{-5}} \frac{0}{2^{-6}} \frac{1}{2^{-7}})_{2,2,8}$ approximate as well as we can from below.

```
welcome $jshell
| Welcome to JShell -- Version 10.0.1
| For an introduction type: /help intro

[jshell> 0.1
$1 ==>

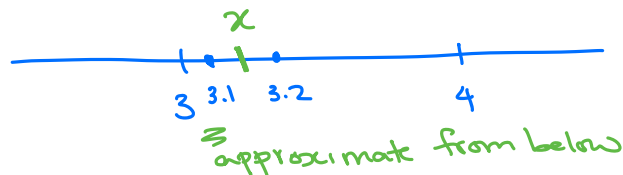
[jshell> 0.2
$2 ==>

[jshell> 0.1 + 0.2
$3 ==>

[jshell> Math.sqrt(2)
$4 ==>

[jshell> Math.sqrt(2)*Math.sqrt(2)
$5 ==>

[jshell>
```



Note: Java uses floating point, not fixed width representation, but similar rounding errors appear in both.

Review: Week 2 Monday

1. Recall the definitions from class for number representations for **base b expansion of n** , **base b fixed-width w expansion of n** , and **base b fixed-width expansion of x with integer part width w and fractional part width w'** .

For example, the base 2 (binary) expansion of 4 is $(100)_2$ and the base 2 (binary) fixed-width 8 expansion of 4 is $(00000100)_{2,8}$ and the base 2 (binary) fixed-width expansion of 4 with integer part width 3 and fractional part width 2 of 4 is $(100.00)_{2,3,2}$

Compute the listed expansions. Enter your number using the notation for base expansions with parentheses but without subscripts. For example, if your answer were $(100)_{2,3}$ you would type $(100)2,3$ into Gradescope.

- (a) Give the binary (base 2) expansion of the number whose octal (base 8) expansion is

$$(371)_8$$

- (b) Give the decimal (base 10) expansion of the number whose octal (base 8) expansion is

$$(371)_8$$

- (c) Give the octal (base 8) fixed-width 3 expansion of $(9)_{10}$?

- (d) Give the ternary (base 3) fixed-width 8 expansion of $(9)_{10}$?

- (e) Give the hexadecimal (base 16) fixed-width 6 expansion of $(16711935)_{10}$?¹

- (f) Give the hexadecimal (base 16) fixed-width 4 expansion of

$$(1011\ 1010\ 1001\ 0000)_2$$

Note: the spaces between each group of 4 bits above are for your convenience only. How might they help your calculations?

- (g) Give the binary fixed width expansion of 0.125 with integer part width 2 and fractional part width 4.

- (h) Give the binary fixed width expansion of 1 with integer part width 2 and fractional part width 3.

2. Select all and only the correct choices below.

- (a) Suppose you were told that the positive integer n_1 has the property that $n_1 \mathbf{div} 2 = 0$. Which of the following can you conclude?

- i. n_1 has a binary (base 2) expansion
- ii. n_1 has a ternary (base 3) expansion
- iii. n_1 has a hexadecimal (base 16) expansion
- iv. n_1 has a base 2 fixed-width 1 expansion
- v. n_1 has a base 2 fixed-width 20 expansion

- (b) Suppose you were told that the positive integer n_2 has the property that $n_2 \mathbf{mod} 4 = 0$. Which of the following can you conclude?

¹This matches a frequent debugging task – sometimes a program will show a number formatted as a base 10 integer that is much better understood with another representation.

- i. the leftmost symbol in the binary (base 2) expansion of n_2 is 1
- ii. the leftmost symbol in the base 4 expansion of n_2 is 1
- iii. the rightmost symbol in the base 4 expansion of n_2 is 0
- iv. the rightmost symbol in the octal (base 8) expansion of n_2 is 0

Wednesday October 6

- 10
sign magnitude

+10

base b expansion of n	base b fixed-width w expansion of n
For b an integer greater than 1 and n a positive integer, the base b expansion of n is $(a_{k-1} \cdots a_1 a_0)_b$ where k is a positive integer, a_0, a_1, \dots, a_{k-1} are nonnegative integers less than b , $a_{k-1} \neq 0$, and $n = a_{k-1}b^{k-1} + \cdots + a_1b + a_0$	For b an integer greater than 1, w a positive integer, and n a nonnegative integer with $n < b^w$, the base b fixed-width w expansion of n is $(a_{w-1} \cdots a_1 a_0)_{b,w}$ where a_0, a_1, \dots, a_{w-1} are nonnegative integers less than b and $n = a_{w-1}b^{w-1} + \cdots + a_1b + a_0$

Representing negative integers in binary: Fix a positive integer width for the representation w , $w > 1$.

	To represent a positive integer n	To represent a negative integer $-n$
Sign-magnitude	<p>$[0a_{w-2} \cdots a_0]_{s,w}$, where $n = (a_{w-2} \cdots a_0)_{2,w-1}$</p> <p>Example $n = 17$, $w = 7$:</p> <p>$[0 \ 01 \ 00 \ 00 \ 01]_{s,7}$</p> <p>$17 = (010001)_{2,6}$</p> <p><i>sign</i> <i>width</i> <i>binary fixed width expansion</i> <i>representation is sign magnitude</i></p>	<p>$[1a_{w-2} \cdots a_0]_{s,w}$, where $n = (a_{w-2} \cdots a_0)_{2,w-1}$</p> <p>Example $-n = -17$, $w = 7$:</p> <p>$-17 = [1 \ 010001]_{s,7}$</p> <p>$17 = (010001)_{2,6}$</p>
2s complement	<p>$[0a_{w-2} \cdots a_0]_{2c,w}$, where $n = (a_{w-2} \cdots a_0)_{2,w-1}$</p> <p>Example $n = 17$, $w = 7$:</p> <p>$17 = [0010001]_{2c,7}$</p>	<p>$[1a_{w-2} \cdots a_0]_{2c,w}$, where $2^{w-1} - n = (a_{w-2} \cdots a_0)_{2,w-1}$</p> <p>Example $-n = -17$, $w = 7$:</p> <p>$-17 = [1 \ 101111]_{2c,7}$</p> <p>$w=7$ $n=17$ $2^{w-1} - n = 2^6 - 17 = 64 - 17$ $= 47 = 32 + 8 + 4 + 2 + 1$ $= 2^5 + 0 \cdot 2^4 + 2^3 + 2^2 + 2^1 + 2^0$</p>
Extra example: 1s complement	<p>$[0a_{w-2} \cdots a_0]_{1c,w}$, where $n = (a_{w-2} \cdots a_0)_{2,w-1}$</p> <p>Example $n = 17$, $w = 7$:</p>	<p>$[1\bar{a}_{w-2} \cdots \bar{a}_0]_{1c,w}$, where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ and we define $\bar{0} = 1$ and $\bar{1} = 0$.</p> <p>Example $-n = -17$, $w = 7$:</p>

$$-17 = [1101111]_{2,7}$$

prespecifying number of columns

For positive integer n , to represent $-n$ in 2s complement with width w ,

Use definition

- Calculate $2^{w-1} - n$, convert result to binary fixed-width $w - 1$, pad with leading 1, or negative weight
- Express $-n$ as a sum of powers of 2, where the leftmost 2^{w-1} is negative weight, or (1) $\frac{1}{2} \frac{1}{2^2} \dots \frac{1}{2^{w-1}}$
- Convert n to binary fixed-width, flip bits, add 1 (ignore overflow)

Want -17 so $n=17 = (0010001)_{2,7}$

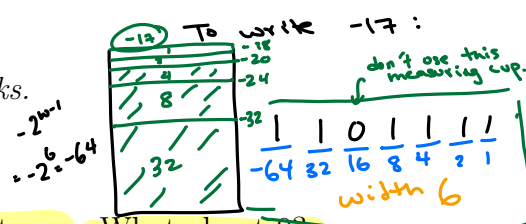
Flip bits: 1101110

Add 1: 1101111

Challenge: use definitions to explain why each of these approaches works.

Representing 0:

So far, the definitions of base expansions treat positive and negative integers. What about 0?



	To represent a non-negative integer n	To represent a non-positive integer $-n$
Sign-magnitude	$[0a_{w-2} \dots a_0]_{s,w}$, where $n = (a_{w-2} \dots a_0)_{2,w-1}$ Example $n = 0$, $w = 7$: $[0000000]_{s,7}$ (a)	$[1a_{w-2} \dots a_0]_{s,w}$, where $n = (a_{w-2} \dots a_0)_{2,w-1}$ Example $-n = 0$, $w = 7$: $[1000000]_{s,7}$ (b)
2s complement	$[0a_{w-2} \dots a_0]_{2c,w}$, where $n = (a_{w-2} \dots a_0)_{2,w-1}$ Example $n = 0$, $w = 7$: $[0000000]_{2c,7}$ (c)	 $[1a_{w-2} \dots a_0]_{2c,w}$, where $2^{w-1} - n = (a_{w-2} \dots a_0)_{2,w-1}$ Example $-n = 0$, $w = 7$: $[1000000]_{2c,7}$ (d)

wasteful

$$2^6 - 0 = 2^{w-1} - n = ()_{2,6}$$

width 4

Integers that can be represented in

... sign magnitude width 4

$$-7 = [1111]_{s,4}$$

$$[0000]_{s,4} = [1000]_{s,4}$$

$$[0111]_{s,4} = 7$$

... 2s complement width 4

$$-8 = [1000]_{2c,4}$$

$$[0000]_{2c,4} = 0$$

$$[0111]_{2c,4} = 7$$

$$n = 16 \quad w = 7$$

$$-n = [\quad \quad \quad]_{2^{c,7}}$$

$$n = (0010000)_{2,7}$$

$$\text{Flip bits: } 11'0'1'1'1$$

$$\text{Add 1: } 1110000_{\text{r}}$$

$$-16 = [\underset{-64}{1} \underset{32}{1} \underset{16}{1} \underset{8}{0} \underset{4}{0} \underset{2}{0} \underset{1}{0}]_{2^{c,7}}$$

Fixed-width addition: adding one bit at time, using the usual column-by-column and carry arithmetic, and dropping the carry from the leftmost column so the result is the same width as the summands. Does this give the right value for the sum?

$$\begin{array}{r}
 \text{Summand}_1 \quad (1\ 1\ 0\ 1\ 0\ 0)_{2,6} \\
 + \text{Summand}_2 \quad (0\ 0\ 0\ 1\ 0\ 1)_{2,6} \\
 \hline
 (1\ 1\ 1\ 0\ 0\ 1)_{2,6} \\
 \text{32 + 16 + 8 + 1} \\
 = 57
 \end{array}$$

$$\text{Summand}_1 \quad 32 + 16 + 4 = 52$$

$$\text{Summand}_2 \quad 4 + 1 = 5$$

$$\text{expected sum} \quad 52 + 5 = 57$$

fixed width
addition gave
expected sum

$$\begin{array}{r}
 \text{Summand}_1 \quad [1\ 1\ 0\ 1\ 0\ 0]_{s,6} \\
 + \text{Summand}_2 \quad [0\ 0\ 0\ 1\ 0\ 1]_{s,6} \\
 \hline
 [1\ 1\ 1\ 0\ 0\ 1]_{s,6} \\
 \text{sign} \quad [1\ 1\ 1\ 0\ 0\ 1]_{s,6} = - (16 + 8 + 1) \\
 = -25
 \end{array}$$

$$\text{Summand}_1 \quad - (16 + 4) = -20$$

$$\text{Summand}_2 \quad 4 + 1 = 5$$

$$\text{expected sum} \quad -15$$

Not what we expected!

$$\begin{array}{r}
 \text{Summand}_1 \quad [1\ 1\ 0\ 1\ 0\ 0]_{2c,6} \\
 + \text{Summand}_2 \quad [0\ 0\ 0\ 1\ 0\ 1]_{2c,6} \\
 \hline
 [1\ 1\ 1\ 0\ 0\ 1]_{2c,6}
 \end{array}$$

$$\begin{array}{r}
 -32 + 16 + 8 + 1 \\
 = -7
 \end{array}$$

$$\text{Summand}_1 \quad -32 + 16 + 4 = -12$$

$$\text{Summand}_2 \quad 4 + 1 = 5$$

$$\text{expected sum} \quad -7$$

Yes! 2s complement
fixed width addition
gave us expected
sum in this
example.

Review: Week 2 Wednesday

1. Recall the definitions of signed integer representations from class: sign-magnitude and 2s complement.
 - (a) Give the 2s complement width 6 representation of the number represented in binary fixed-width 5 representation as $(00101)_{2,5}$.
 - (b) Give the 2s complement width 6 representation of the number represented in binary fixed-width 5 representation as $(10101)_{2,5}$.
 - (c) Give the 2s complement width 4 representation of the number represented in sign-magnitude width 4 as $[1111]_{s,4}$.
 - (d) Give the sign magnitude width 4 representation of the number represented in 2s complement width 4 as $[1111]_{2c,4}$.
 - (e) Give the sign magnitude width 6 representation of the number represented in sign magnitude width 4 as $[1111]_{s,4}$.
 - (f) Give the 2s complement width 6 representation of the number represented in 2s complement width 4 as $[1111]_{2c,4}$.
2. Recall the definitions of signed integer representations from class: sign-magnitude and 2s complement.
 - (a) In binary fixed-width addition (adding one bit at time, using the usual column-by-column and carry arithmetic, and ignoring the carry from the leftmost column), we get:

$$\begin{array}{rcl} 1110 & \text{first summand} \\ +0100 & \text{second summand} \\ \hline 0010 & \text{result} \end{array}$$

Select all and only the true statements below:

- i. When interpreting each of the summands and the result in binary fixed-width 4, the result represents the actual value of the sum of the summands.
 - ii. When interpreting each of the summands and the sum in sign-magnitude width 4, the result represents the actual value of the sum of the summands.
 - iii. When interpreting each of the summands and the sum in 2s complement width 4, the result represents the actual value of the sum of the summands.
- (b) In binary fixed-width addition (adding one bit at time, using the usual column-by-column and carry arithmetic, and ignoring the carry from the leftmost column), we get:

$$\begin{array}{rcl} 0110 & \text{first summand} \\ +0111 & \text{second summand} \\ \hline 1101 & \text{result} \end{array}$$

Select all and only the true statements below:

- i. When interpreting each of the summands and the result in binary fixed-width 4, the result represents the actual value of the sum of the summands.
- ii. When interpreting each of the summands and the sum in sign-magnitude width 4, the result represents the actual value of the sum of the summands.
- iii. When interpreting each of the summands and the sum in 2s complement width 4, the result represents the actual value of the sum of the summands.

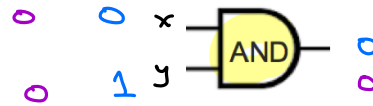
Friday October 8

In a **combinatorial circuit** (also known as a **logic circuit**), we have **logic gates** connected by **wires**. The inputs to the circuits are the values set on the input wires: possible values are 0 (low) or 1 (high). The values flow along the wires from left to right. A wire may be split into two or more wires, indicated with a filled-in circle (representing solder). Values stay the same along a wire. When one or more wires flow into a gate, the output value of that gate is computed from the input values based on the gate's definition table. Outputs of gates may become inputs to other gates.

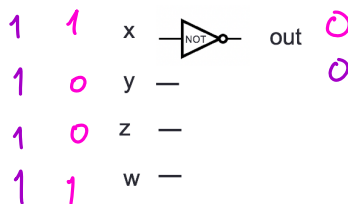
Inputs		Output	
x	y	x AND y	
1	1	1	
1	0	0	
0	1	0	
0	0	0	

Inputs		Output	
x	y	x XOR y	
1	1	0	
1	0	1	
0	1	1	
0	0	0	

Input	Output
x	NOT x
1	0
0	1

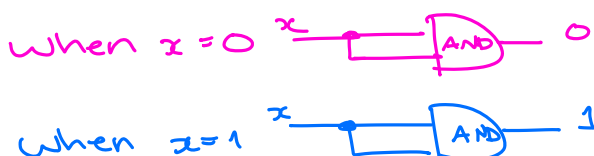
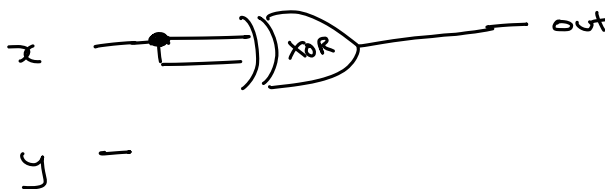


Example digital circuit:



Output when $x = 1, y = 0, z = 0, w = 1$ is 0
 Output when $x = 1, y = 1, z = 1, w = 1$ is 0
 Output when $x = 0, y = 0, z = 0, w = 1$ is 1

Draw a logic circuit with inputs x and y whose output is always 0. Can you use exactly 1 gate?



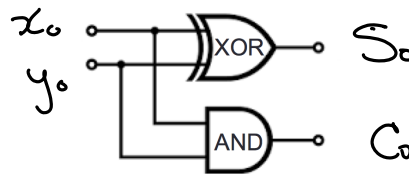
NOTE
~~is not a valid circuit~~

Fixed-width addition: adding one bit at time, using the usual column-by-column and carry arithmetic, and dropping the carry from the leftmost column so the result is the same width as the summands. In many cases, this gives representation of the correct value for the sum when we interpret the summands in fixed-width binary or in 2s complement.

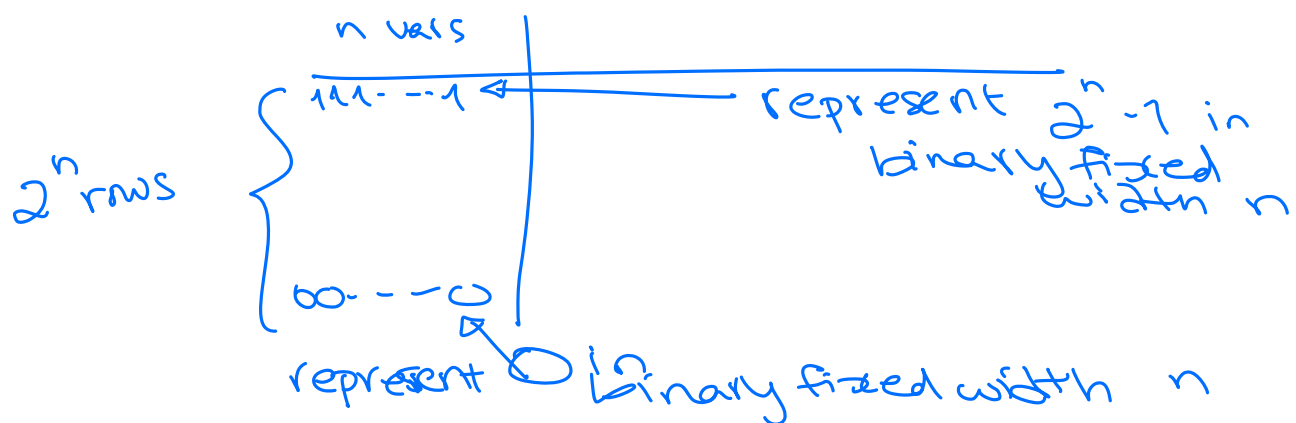
For single column:

$$\begin{array}{r}
 (x_0)_{2,1} \\
 + (y_0)_{2,1} \\
 \hline
 C_0 \quad S_0
 \end{array}$$

Input		Output	
x_0	y_0	c_0	s_0
1	1	1	0
1	0	0	1
0	1	0	1
0	0	0	0



Notice: when we write input-output table, typically organize the rows:



binary addition:
of two numbers expressed in
fixed-width binary.

$(x_1 y_0)_{2,2}$
 $(y_1 y_0)_{2,2}$ (may require up to width 3)
upto 3 bits.

So
 $(x_1 x_0)_{2,2}$
is between
0 and 3
and same
with
 $(y_1 y_0)_{2,2}$

- First approach:* half-adder for each column, then combine carry from right column with sum of left column

values:

$$\begin{array}{cc} & \cancel{G} \\ & (x_1 \quad x_0)_{2,2} \\ + & (y_1 \quad y_0)_{2,2} \\ \hline \text{ } & \underline{s_1} \quad s_0 \end{array}$$

The diagram illustrates a 2-bit ripple-carry adder. It takes four inputs: x_0, y_0, x_1, y_1 . The first stage (bit 0) uses XOR gates to produce the sum s_0 and AND gates to produce the carry c_0 . The second stage (bit 1) uses XOR gates to produce the sum s_1 and AND gates to produce the carry c_1 . The carry c_0 from the first stage is connected to the carry input of the second stage. The final outputs are $z_0 = s_0$, $z_1 = s_1$, and $z_2 = c_1$.

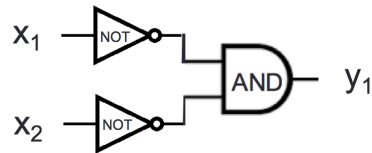
Write expressions for the circuit output values in terms of input values:

$$\begin{aligned} z_0 &= x_0 \oplus y_0 \\ z_1 &= ((x_0 \wedge y_0) \oplus x_1) \oplus y_1 \\ z_2 &= ((x_0 \wedge y_0) \wedge x_1) \oplus (((x_0 \wedge y_0) \oplus x_1) \wedge y_1) \end{aligned}$$

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Review: Week 2 Friday

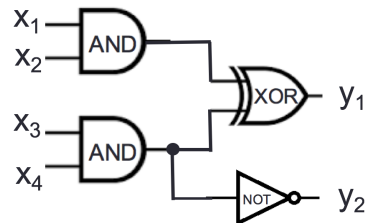
1. (a) Consider the logic circuit



Calculate the value of the output of this circuit (y_1) for each of the following settings(s) of input values.

- i. $x_1 = 1, x_2 = 1$
- ii. $x_1 = 1, x_2 = 0$
- iii. $x_1 = 0, x_2 = 1$
- iv. $x_1 = 0, x_2 = 0$

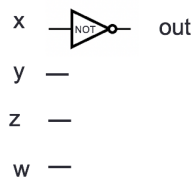
- (b) Consider the logic circuit



For which of the following settings(s) of input values is the output $y_1 = 0, y_2 = 1$? (Select all and only those that apply.)

- i. $x_1 = 0, x_2 = 0, x_3 = 0$, and $x_4 = 0$
- ii. $x_1 = 1, x_2 = 1, x_3 = 1$, and $x_4 = 1$
- iii. $x_1 = 1, x_2 = 0, x_3 = 0$, and $x_4 = 1$
- iv. $x_1 = 0, x_2 = 0, x_3 = 1$, and $x_4 = 1$

2. Recall this circuit from class:



Which of the following is true about all possible input values x, y, z, w ? (Select all and only choices that are true for all values.)

- (a) The output *out* is set to 1 exactly when x is 0, and it is set to 0 otherwise.
- (b) The output *out* is set to 1 exactly when $(xyzw)_{2,4} < 8$, and it is set to 0 otherwise.
- (c) The output *out* is set to 1 exactly when $(wzyx)_{2,4}$ is an even integer, and it is set to 0 otherwise.