

## Netflix intro

What data should we encode about each Netflix account holder to help us make effective recommendations?

In machine learning, clustering can be used to group similar data for prediction and recommendation. For example, each Netflix user's viewing history can be represented as a  $n$ -tuple indicating their preferences about movies in the database, where  $n$  is the number of movies in the database. People with similar tastes in movies can then be clustered to provide recommendations of movies for one another. Mathematically, clustering is based on a notion of distance between pairs of  $n$ -tuples.

# Set operations

To define a set we can use the roster method, set builder notation, a recursive definition, and also we can apply a set operation to other sets.

## New! Cartesian product of sets and set-wise concatenation of sets of strings

**Definition:** Let  $X$  and  $Y$  be sets. The **Cartesian product** of  $X$  and  $Y$ , denoted  $X \times Y$ , is the set of all ordered pairs  $(x, y)$  where  $x \in X$  and  $y \in Y$

$$X \times Y = \{(x, y) \mid x \in X \text{ and } y \in Y\}$$

**Definition:** Let  $X$  and  $Y$  be sets of strings over the same alphabet. The **set-wise concatenation** of  $X$  and  $Y$ , denoted  $X \circ Y$ , is the set of all results of string concatenation  $xy$  where  $x \in X$  and  $y \in Y$

$$X \circ Y = \{xy \mid x \in X \text{ and } y \in Y\}$$

**Pro-tip:** the meaning of writing one element next to another like  $xy$  depends on the data-types of  $x$  and  $y$ . When  $x$  and  $y$  are strings, the convention is that  $xy$  is the result of string concatenation. When  $x$  and  $y$  are numbers, the convention is that  $xy$  is the result of multiplication. This is (one of the many reasons) why is it very important to declare the data-type of variables before we use them.

Fill in the missing entries in the table:

Set	Example elements in this set:			
$B$	A	C	G	U
	(A, C)		(U, U)	
$B \times \{-1, 0, 1\}$				
$\{-1, 0, 1\} \times B$				
	(0, 0, 0)			
$\{A, C, G, U\} \circ \{A, C, G, U\}$				
	GGGG			

# Defining functions

**New! Defining functions** A function is defined by its (1) domain, (2) codomain, and (3) rule assigning each element in the domain exactly one element in the codomain.

The domain and codomain are nonempty sets.

The rule can be depicted as a table, formula, or English description.

The notation is

“Let the function  $\text{FUNCTION-NAME}: \text{DOMAIN} \rightarrow \text{CODOMAIN}$  be given by  
 $\text{FUNCTION-NAME}(x) = \dots$  for every  $x \in \text{DOMAIN}$ ”.

or

“Consider the function  $\text{FUNCTION-NAME}: \text{DOMAIN} \rightarrow \text{CODOMAIN}$  given by  
 $\text{FUNCTION-NAME}(x) = \dots$  for every  $x \in \text{DOMAIN}$ ”.

Example: The absolute value function

**Domain**

**Codomain**

**Rule**

# Defining functions recursively

When the domain of a function is a *recursively defined set*, the rule assigning images to domain elements (outputs) can also be defined recursively.

Recall: The set of RNA strands  $S$  is defined (recursively) by:

$$\begin{array}{ll} \text{Basis Step:} & \mathbf{A} \in S, \mathbf{C} \in S, \mathbf{U} \in S, \mathbf{G} \in S \\ \text{Recursive Step:} & \text{If } s \in S \text{ and } b \in B, \text{ then } sb \in S \end{array}$$

where  $sb$  is string concatenation.

**Definition** (Of a function, recursively) A function  $rnalen$  that computes the length of RNA strands in  $S$  is defined by:

$$\begin{array}{lll} & & rnalen : S \rightarrow \mathbb{Z}^+ \\ \text{Basis Step:} & \text{If } b \in B \text{ then} & rnalen(b) = 1 \\ \text{Recursive Step:} & \text{If } s \in S \text{ and } b \in B, \text{ then} & rnalen(sb) = 1 + rnalen(s) \end{array}$$

The domain of  $rnalen$  is

The codomain of  $rnalen$  is

Example function application:

$$rnalen(\mathbf{ACU}) =$$

*Extra example:* A function  $basecount$  that computes the number of a given base  $b$  appearing in a RNA strand  $s$  is defined recursively:

$$\begin{array}{lll} & & basecount : S \times B \rightarrow \mathbb{N} \\ \text{Basis Step:} & \text{If } b_1 \in B, b_2 \in B & basecount((b_1, b_2)) = \begin{cases} 1 & \text{when } b_1 = b_2 \\ 0 & \text{when } b_1 \neq b_2 \end{cases} \\ \text{Recursive Step:} & \text{If } s \in S, b_1 \in B, b_2 \in B & basecount((sb_1, b_2)) = \begin{cases} 1 + basecount((s, b_2)) & \text{when } b_1 = b_2 \\ basecount((s, b_2)) & \text{when } b_1 \neq b_2 \end{cases} \end{array}$$

$$basecount((\mathbf{ACU}, \mathbf{A})) = basecount((\mathbf{AC}, \mathbf{A})) = basecount((\mathbf{A}, \mathbf{A})) = 1$$

$$basecount((\mathbf{ACU}, \mathbf{G})) = basecount((\mathbf{AC}, \mathbf{G})) = basecount((\mathbf{A}, \mathbf{G})) = 0$$

*Extra example:* The function which outputs  $2^n$  when given a nonnegative integer  $n$  can be defined recursively, because its domain is the set of nonnegative integers.

# Why represent numbers

Modeling uses data-types that are encoded in a computer.

The details of the encoding impact the efficiency of algorithms we use to understand the systems we are modeling and the impacts of these algorithms on the people using the systems.

Case study: how to encode numbers?

## Base expansion definition

**Definition** For  $b$  an integer greater than 1 and  $n$  a positive integer, the **base  $b$  expansion of  $n$**  is

$$(a_{k-1} \cdots a_1 a_0)_b$$

where  $k$  is a positive integer,  $a_0, a_1, \dots, a_{k-1}$  are nonnegative integers less than  $b$ ,  $a_{k-1} \neq 0$ , and

$$n = \sum_{i=0}^{k-1} a_i b^i$$

Notice: *The base  $b$  expansion of a positive integer  $n$  is a string over the alphabet  $\{x \in \mathbb{N} \mid x < b\}$  whose leftmost character is nonzero.*

Base $b$	Collection of possible coefficients in base $b$ expansion of a positive integer
Binary ( $b = 2$ )	$\{0, 1\}$
Ternary ( $b = 3$ )	$\{0, 1, 2\}$
Octal ( $b = 8$ )	$\{0, 1, 2, 3, 4, 5, 6, 7\}$
Decimal ( $b = 10$ )	$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
Hexadecimal ( $b = 16$ )	$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$ letter coefficient symbols represent numerical values $(A)_{16} = (10)_{10}$ $(B)_{16} = (11)_{10}$ $(C)_{16} = (12)_{10}$ $(D)_{16} = (13)_{10}$ $(E)_{16} = (14)_{10}$ $(F)_{16} = (15)_{10}$

# Base expansion examples

Common bases:	Binary $b = 2$	Octal $b = 8$	Decimal $b = 10$	Hexadecimal $b = 16$
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*Examples:*

$(1401)_2$

$(1401)_{10}$

$(1401)_{16}$

## Algorithm definition

<b>New!</b> An algorithm is a finite sequence of precise instructions for solving a problem.
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## Algorithm half

Algorithm for calculating integer part of half the input

```
1  procedure half( $n$ : a positive integer)
2     $r := 0$ 
3    while  $n > 1$ 
4       $r := r + 1$ 
5       $n := n - 2$ 
6    return  $r$  { $r$  holds the result of the operation}
```

$n$	$r$	$n > 1?$
6		

$n$	$r$	$n > 1?$
5		

## Algorithm log

Algorithm for calculating integer part of log

```
1  procedure log( $n$ : a positive integer)
2     $r := 0$ 
3    while  $n > 1$ 
4       $r := r + 1$ 
5       $n := \text{half}(n)$ 
6    return  $r$  { $r$  holds the result of the log operation}
```

$n$	$r$	$n > 1?$
8		

$n$	$r$	$n > 1?$
6		

$2^0 = 1$	$2^1 = 2$	$2^2 = 4$	$2^3 = 8$	$2^4 = 16$	$2^5 = 32$	$2^6 = 64$	$2^7 = 128$	$2^8 = 256$	$2^9 = 512$	$2^{10} = 1024$
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## Division algorithm

**Integer division and remainders** (aka The Division Algorithm) Let  $n$  be an integer and  $d$  a positive integer. There are unique integers  $q$  and  $r$ , with  $0 \leq r < d$ , such that  $n = dq + r$ . In this case,  $d$  is called the divisor,  $n$  is called the dividend,  $q$  is called the quotient, and  $r$  is called the remainder. We write  $q = n \text{ div } d$  and  $r = n \text{ mod } d$ .

*Extra example:* How do **div** and **mod** compare to  $/$  and  $\%$  in Java and python?

## Base expansion algorithms

**Two algorithms for constructing base  $b$  expansion from decimal representation**

**Most significant first:** Start with left-most coefficient of expansion

Calculating integer part of  $\log_b$

```

1 procedure logb( $n, b$ : positive integers with  $b > 1$ )
2    $r := 0$ 
3   while  $n > b - 1$ 
4      $r := r + 1$ 
5      $n := n \text{ div } b$ 
6   return  $r$  { $r$  holds the result of the  $\log_b$  operation}
```

Calculating base  $b$  expansion, from left

```

1 procedure baseb1( $n, b$ : positive integers with  $b > 1$ )
2    $v := n$ 
3    $k := \log_b(n, b) + 1$ 
4   for  $i := 1$  to  $k$ 
5      $a_{k-i} := 0$ 
6     while  $v \geq b^{k-i}$ 
7        $a_{k-i} := a_{k-i} + 1$ 
8        $v := v - b^{k-i}$ 
9   return  $(a_{k-1}, \dots, a_0)\{(a_{k-1} \dots a_0)_b \text{ is the base } b \text{ expansion of } n\}$ 
```

**Least significant first:** Start with right-most coefficient of expansion

Idea: (when  $k > 1$ )

$$n = a_{k-1}b^{k-1} + \dots + a_1b + a_0$$

$$= b(a_{k-1}b^{k-2} + \dots + a_1) + a_0$$

so  $a_0 = n \text{ mod } b$  and  $a_{k-1}b^{k-2} + \dots + a_1 =$   
 $n \text{ div } b$ .

Calculating base  $b$  expansion, from right

```

1 procedure baseb2( $n, b$ : positive integers with  $b > 1$ )
2    $q := n$ 
3    $k := 0$ 
4   while  $q \neq 0$ 
5      $a_k := q \text{ mod } b$ 
6      $q := q \text{ div } b$ 
7      $k := k + 1$ 
8   return  $(a_{k-1}, \dots, a_0)\{(a_{k-1} \dots a_0)_b \text{ is the base } b \text{ expansion of } n\}$ 
```

## Base expansion review

Find and fix any and all mistakes with the following:

(a)  $(1)_2 = (1)_8$

(b)  $(142)_{10} = (142)_{16}$

(c)  $(20)_{10} = (10100)_2$

(d)  $(35)_8 = (1D)_{16}$



# Base conversion algorithm

Recall the definition of base expansion we discussed:

**Definition** For  $b$  an integer greater than 1 and  $n$  a positive integer, the **base  $b$  expansion of  $n$**  is

$$(a_{k-1} \cdots a_1 a_0)_b$$

where  $k$  is a positive integer,  $a_0, a_1, \dots, a_{k-1}$  are nonnegative integers less than  $b$ ,  $a_{k-1} \neq 0$ , and

$$n = \sum_{i=0}^{k-1} a_i b^i$$

Notice: The base  $b$  expansion of a positive integer  $n$  is a string over the alphabet  $\{x \in \mathbb{N} \mid x < b\}$  whose leftmost character is nonzero.

Base $b$	Collection of possible coefficients in base $b$ expansion of a positive integer
Binary ( $b = 2$ )	$\{0, 1\}$
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Hexadecimal ( $b = 16$ )	$\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$ letter coefficient symbols represent numerical values $(A)_{16} = (10)_{10}$ $(B)_{16} = (11)_{10}$ $(C)_{16} = (12)_{10}$ $(D)_{16} = (13)_{10}$ $(E)_{16} = (14)_{10}$ $(F)_{16} = (15)_{10}$

We write an algorithm for converting from base  $b_1$  expansion to base  $b_2$  expansion:

# Fixed width definition

**Definition** For  $b$  an integer greater than 1,  $w$  a positive integer, and  $n$  a nonnegative integer \_\_\_\_\_, the **base  $b$  fixed-width  $w$  expansion of  $n$**  is

$$(a_{w-1} \cdots a_1 a_0)_{b,w}$$

where  $a_0, a_1, \dots, a_{w-1}$  are nonnegative integers less than  $b$  and

$$n = \sum_{i=0}^{w-1} a_i b^i$$

# Fixed width example

Decimal $b = 10$	Binary $b = 2$	Binary fixed-width 10 $b = 2, w = 10$	Binary fixed-width 7 $b = 2, w = 7$	Binary fixed-width 4 $b = 2, w = 4$
$(20)_{10}$	(a)	(b)	(c)	(d)

# Fixed width fractional definition

**Definition** For  $b$  an integer greater than 1,  $w$  a positive integer,  $w'$  a positive integer, and  $x$  a real number the **base  $b$  fixed-width expansion of  $x$  with integer part width  $w$  and fractional part width  $w'$**  is  $(a_{w-1} \cdots a_1 a_0 . c_1 \cdots c_{w'})_{b,w,w'}$  where  $a_0, a_1, \dots, a_{w-1}, c_1, \dots, c_{w'}$  are nonnegative integers less than  $b$  and

$$x \geq \sum_{i=0}^{w-1} a_i b^i + \sum_{j=1}^{w'} c_j b^{-j} \quad \text{and} \quad x < \sum_{i=0}^{w-1} a_i b^i + \sum_{j=1}^{w'} c_j b^{-j} + b^{-w'}$$

3.75 in fixed-width binary, integer part width 2, fractional part width 8	
0.1 in fixed-width binary, integer part width 2, fractional part width 8	

```

|welcome $jshell
| Welcome to JShell -- Version 10.0.1
| For an introduction type: /help intro

[jshell> 0.1
$1 ==>

[jshell> 0.2
$2 ==>

[jshell> 0.1 + 0.2
$3 ==>

[jshell> Math.sqrt(2)
$4 ==>

[jshell> Math.sqrt(2)*Math.sqrt(2)
$5 ==>

[jshell> █

```

Note: Java uses floating point, not fixed width representation, but similar rounding errors appear in both.

## Expansion summary

base $b$ expansion of $n$	base $b$ fixed-width $w$ expansion of $n$
For $b$ an integer greater than 1 and $n$ a positive integer, the <b>base <math>b</math> expansion of <math>n</math></b> is $(a_{k-1} \cdots a_1 a_0)_b$ where $k$ is a positive integer, $a_0, a_1, \dots, a_{k-1}$ are nonnegative integers less than $b$ , $a_{k-1} \neq 0$ , and $n = a_{k-1}b^{k-1} + \cdots + a_1b + a_0$	For $b$ an integer greater than 1, $w$ a positive integer, and $n$ a nonnegative integer with $n < b^w$ , the <b>base <math>b</math> fixed-width <math>w</math> expansion of <math>n</math></b> is $(a_{w-1} \cdots a_1 a_0)_{b,w}$ where $a_0, a_1, \dots, a_{w-1}$ are nonnegative integers less than $b$ and $n = a_{w-1}b^{w-1} + \cdots + a_1b + a_0$

# Negative int expansions

**Representing negative integers in binary:** Fix a positive integer width for the representation  $w$ ,  $w > 1$ .

	To represent a positive integer $n$	To represent a negative integer $-n$
Sign-magnitude	$[0a_{w-2} \cdots a_0]_{s,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $n = 17$ , $w = 7$ :	$[1a_{w-2} \cdots a_0]_{s,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $-n = -17$ , $w = 7$ :
2s complement	$[0a_{w-2} \cdots a_0]_{2c,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $n = 17$ , $w = 7$ :	$[1a_{w-2} \cdots a_0]_{2c,w}$ , where $2^{w-1} - n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $-n = -17$ , $w = 7$ :
Extra example: 1s complement	$[0a_{w-2} \cdots a_0]_{1c,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $n = 17$ , $w = 7$ :	$[1\bar{a}_{w-2} \cdots \bar{a}_0]_{1c,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ and we define $\bar{0} = 1$ and $\bar{1} = 0$ . Example $-n = -17$ , $w = 7$ :

# Calculating 2s complement

For positive integer  $n$ , to represent  $-n$  in 2s complement with width  $w$ ,

- Calculate  $2^{w-1} - n$ , convert result to binary fixed-width  $w - 1$ , pad with leading 1, or
- Express  $-n$  as a sum of powers of 2, where the leftmost  $2^{w-1}$  is negative weight, or
- Convert  $n$  to binary fixed-width  $w$ , flip bits, add 1 (ignore overflow)

*Challenge: use definitions to explain why each of these approaches works.*

## Representing zero

**Representing 0:**

So far, we have representations for positive and negative integers. What about 0?

	To represent a <b>non-negative</b> integer $n$	To represent a <b>non-positive</b> integer $-n$
Sign-magnitude	$[0a_{w-2} \cdots a_0]_{s,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $n = 0$ , $w = 7$ :  (a)	$[1a_{w-2} \cdots a_0]_{s,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $-n = 0$ , $w = 7$ :  (b)
2s complement	$[0a_{w-2} \cdots a_0]_{2c,w}$ , where $n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $n = 0$ , $w = 7$ :  (c)	$[1a_{w-2} \cdots a_0]_{2c,w}$ , where $2^{w-1} - n = (a_{w-2} \cdots a_0)_{2,w-1}$ Example $-n = 0$ , $w = 7$ :  (d)

# Fixed width addition

**Fixed-width addition:** adding one bit at time, using the usual column-by-column and carry arithmetic, and dropping the carry from the leftmost column so the result is the same width as the summands. *Does this give the right value for the sum?*

$$\begin{array}{r} (1\ 1\ 0\ 1\ 0\ 0)_{2,6} \\ + (0\ 0\ 0\ 1\ 0\ 1)_{2,6} \\ \hline \end{array}$$

$$\begin{array}{r} [1\ 1\ 0\ 1\ 0\ 0]_{s,6} \\ + [0\ 0\ 0\ 1\ 0\ 1]_{s,6} \\ \hline \end{array}$$

$$\begin{array}{r} [1\ 1\ 0\ 1\ 0\ 0]_{2c,6} \\ + [0\ 0\ 0\ 1\ 0\ 1]_{2c,6} \\ \hline \end{array}$$

## Circuits basics

In a **combinatorial circuit** (also known as a **logic circuit**), we have **logic gates** connected by **wires**. The inputs to the circuits are the values set on the input wires: possible values are 0 (low) or 1 (high). The values flow along the wires from left to right. A wire may be split into two or more wires, indicated with a filled-in circle (representing solder). Values stay the same along a wire. When one or more wires flow into a gate, the output value of that gate is computed from the input values based on the gate's definition table. Outputs of gates may become inputs to other gates.

## Logic gates definitions

Inputs		Output
$x$	$y$	$x$ AND $y$
1	1	1
1	0	0
0	1	0
0	0	0



Inputs		Output
$x$	$y$	$x$ XOR $y$
1	1	0
1	0	1
0	1	1
0	0	0



Input	Output
$x$	NOT $x$
1	0
0	1



# Digital circuits basic examples

Example digital circuit:



Output when  $x = 1, y = 0, z = 0, w = 1$  is \_\_\_\_\_

Output when  $x = 1, y = 1, z = 1, w = 1$  is \_\_\_\_\_

Output when  $x = 0, y = 0, z = 0, w = 1$  is \_\_\_\_\_

Draw a logic circuit with inputs  $x$  and  $y$  whose output is always 0. *Can you use exactly 1 gate?*

## Half adder circuit

**Fixed-width addition:** adding one bit at time, using the usual column-by-column and carry arithmetic, and dropping the carry from the leftmost column so the result is the same width as the summands. In many cases, this gives representation of the correct value for the sum when we interpret the summands in fixed-width binary or in 2s complement.

For single column:

Input		Output	
$x_0$	$y_0$	$c_0$	$s_0$
1	1		
1	0		
0	1		
0	0		



## Two bit adder circuit

Draw a logic circuit that implements binary addition of two numbers that are each represented in fixed-width binary:

- Inputs  $x_0, y_0, x_1, y_1$  represent  $(x_1x_0)_{2,2}$  and  $(y_1y_0)_{2,2}$
- Outputs  $z_0, z_1, z_2$  represent  $(z_2z_1z_0)_{2,3} = (x_1x_0)_{2,2} + (y_1y_0)_{2,2}$  (may require up to width 3)

*First approach:* half-adder for each column, then combine carry from right column with sum of left column

Write expressions for the circuit output values in terms of input values:

$z_0 =$  \_\_\_\_\_

$z_1 =$  \_\_\_\_\_

$z_2 =$  \_\_\_\_\_



*Second approach:* for middle column, first add carry from right column to  $x_1$ , then add result to  $y_1$

Write expressions for the circuit output values in terms of input values:

$z_0 =$  \_\_\_\_\_

$z_1 =$  \_\_\_\_\_

$z_2 =$  \_\_\_\_\_

*Extra example* Describe how to generalize this addition circuit for larger width inputs.



# Logical operators

Logical operators aka propositional connectives

<b>Conjunction</b>	AND	$\wedge$	<code>\land</code>	2 inputs	Evaluates to $T$ exactly when <b>both</b> inputs are $T$
<b>Exclusive or</b>	XOR	$\oplus$	<code>\oplus</code>	2 inputs	Evaluates to $T$ exactly when <b>exactly one</b> of inputs is $T$
<b>Disjunction</b>	OR	$\vee$	<code>\lor</code>	2 inputs	Evaluates to $T$ exactly when <b>at least one</b> of inputs is $T$
<b>Negation</b>	NOT	$\neg$	<code>\lnot</code>	1 input	Evaluates to $T$ exactly when its input is $F$

## Logical operators truth tables

Truth tables: Input-output tables where we use  $T$  for 1 and  $F$  for 0.

Input		Output		
		<b>Conjunction</b>	<b>Exclusive or</b>	<b>Disjunction</b>
$p$	$q$	$p \wedge q$	$p \oplus q$	$p \vee q$
$T$	$T$	$T$	$F$	$T$
$T$	$F$	$F$	$T$	$T$
$F$	$T$	$F$	$T$	$T$
$F$	$F$	$F$	$F$	$F$





Input	Output
$p$	<b>Negation</b> $\neg p$
$T$	$F$
$F$	$T$



## Logical operators example truth table

Input			Output	
$p$	$q$	$r$	$(p \wedge q) \oplus ((p \oplus q) \wedge r)$	$(p \wedge q) \vee ((p \oplus q) \wedge r)$
$T$	$T$	$T$		
$T$	$T$	$F$		
$T$	$F$	$T$		
$T$	$F$	$F$		
$F$	$T$	$T$		
$F$	$T$	$F$		
$F$	$F$	$T$		
$F$	$F$	$F$		

# Truth table to compound proposition

Given a truth table, how do we find an expression using the input variables and logical operators that has the output values specified in this table?

*Application:* design a circuit given a desired input-output relationship.

Input		Output	
$p$	$q$	$mystery_1$	$mystery_2$
$T$	$T$	$T$	$F$
$T$	$F$	$T$	$F$
$F$	$T$	$F$	$F$
$F$	$F$	$T$	$T$

Expressions that have output  $mystery_1$  are

Expressions that have output  $mystery_2$  are

## Dnf cnf definition

**Definition** An expression built of variables and logical operators is in **disjunctive normal form** (DNF) means that it is an OR of ANDs of variables and their negations.

**Definition** An expression built of variables and logical operators is in **conjunctive normal form** (CNF) means that it is an AND of ORs of variables and their negations.

## Dnf cnf example

*Extra example:* An expression that has output ? is:

Input			Output
$p$	$q$	$r$	?
$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$
$T$	$F$	$T$	$F$
$T$	$F$	$F$	$T$
$F$	$T$	$T$	$F$
$F$	$T$	$F$	$F$
$F$	$F$	$T$	$T$
$F$	$F$	$F$	$F$

## Compound proposition definitions

**Proposition:** Declarative sentence that is true or false (not both).

**Propositional variable:** Variable that represents a proposition.

**Compound proposition:** New proposition formed from existing propositions (potentially) using logical operators. *Note:* A propositional variable is one example of a compound proposition.

**Truth table:** Table with one row for each of the possible combinations of truth values of the input and an additional column that shows the truth value of the result of the operation corresponding to a particular row.

## Logical equivalence

**Logical equivalence :** Two compound propositions are **logically equivalent** means that they have the same truth values for all settings of truth values to their propositional variables.

**Tautology:** A compound proposition that evaluates to true for all settings of truth values to its propositional variables; it is abbreviated  $T$ .

**Contradiction:** A compound proposition that evaluates to false for all settings of truth values to its propositional variables; it is abbreviated  $F$ .

**Contingency:** A compound proposition that is neither a tautology nor a contradiction.

# Tautology contradiction contingency examples

Label each of the following as a tautology, contradiction, or contingency.

$$p \wedge p$$

$$p \oplus p$$

$$p \vee p$$

$$p \vee \neg p$$

$$p \wedge \neg p$$

## Logical equivalence extra example

*Extra Example:* Which of the compound propositions in the table below are logically equivalent?

Input		Output				
$p$	$q$	$\neg(p \wedge \neg q)$	$\neg(\neg p \vee \neg q)$	$(\neg p \vee q)$	$(\neg q \vee \neg p)$	$(p \wedge q)$
$T$	$T$					
$T$	$F$					
$F$	$T$					
$F$	$F$					

## Logical operators full truth table

Input		Output				
$p$	$q$	Conjunction $p \wedge q$	Exclusive or $p \oplus q$	Disjunction $p \vee q$	Conditional $p \rightarrow q$	Biconditional $p \leftrightarrow q$
$T$	$T$	$T$	$F$	$T$	$T$	$T$
$T$	$F$	$F$	$T$	$T$	$F$	$F$
$F$	$T$	$F$	$T$	$T$	$T$	$F$
$F$	$F$	$F$	$F$	$F$	$T$	$T$
		“ $p$ and $q$ ”	“ $p$ xor $q$ ”	“ $p$ or $q$ ”	“if $p$ then $q$ ”	“ $p$ if and only if $q$ ”

## Hypothesis conclusion

The only way to make the conditional statement  $p \rightarrow q$  false is to \_\_\_\_\_

The **hypothesis** of  $p \rightarrow q$  is \_\_\_\_\_ The **antecedent** of  $p \rightarrow q$  is \_\_\_\_\_

The **conclusion** of  $p \rightarrow q$  is \_\_\_\_\_ The **consequent** of  $p \rightarrow q$  is \_\_\_\_\_

## Converse inverse contrapositive

The **converse** of  $p \rightarrow q$  is \_\_\_\_\_

The **inverse** of  $p \rightarrow q$  is \_\_\_\_\_

The **contrapositive** of  $p \rightarrow q$  is \_\_\_\_\_

## Compound propositions recursive definition

We can use a recursive definition to describe all compound propositions that use propositional variables from a specified collection. Here's the definition for all compound propositions whose propositional variables are in  $\{p, q\}$ .

Basis Step:	$p$ and $q$ are each a compound proposition
Recursive Step:	If $x$ is a compound proposition then so is $(\neg x)$ and if $x$ and $y$ are both compound propositions then so is each of $(x \wedge y)$ , $(x \oplus y)$ , $(x \vee y)$ , $(x \rightarrow y)$ , $(x \leftrightarrow y)$

## Compound propositions precedence

Order of operations (Precedence) for logical operators:

Negation, then conjunction / disjunction, then conditional / biconditionals.

Example:  $\neg p \vee \neg q$  means  $(\neg p) \vee (\neg q)$ .

# Logical equivalence identities

## (Some) logical equivalences

*Can replace  $p$  and  $q$  with any compound proposition*

$$\neg(\neg p) \equiv p$$

**Double negation**

$$p \vee q \equiv q \vee p$$

$$p \wedge q \equiv q \wedge p$$

**Commutativity** Ordering of terms

$$(p \vee q) \vee r \equiv p \vee (q \vee r)$$

$$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

**Associativity** Grouping of terms

$$p \wedge F \equiv F$$

$$p \vee T \equiv T$$

$$p \wedge T \equiv p$$

$$p \vee F \equiv p$$

**Domination** aka short circuit evaluation

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

**DeMorgan's Laws**

$$p \rightarrow q \equiv \neg p \vee q$$

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

**Contrapositive**

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

$$\neg(p \leftrightarrow q) \equiv p \oplus q$$

$$p \leftrightarrow q \equiv q \leftrightarrow p$$

*Extra examples:*

$p \leftrightarrow q$  is not logically equivalent to  $p \wedge q$  because \_\_\_\_\_

$p \rightarrow q$  is not logically equivalent to  $q \rightarrow p$  because \_\_\_\_\_

# Logical operators english synonyms

## Common ways to express logical operators in English:

**Negation**  $\neg p$  can be said in English as

- Not  $p$ .
- It's not the case that  $p$ .
- $p$  is false.

**Conjunction**  $p \wedge q$  can be said in English as

- $p$  and  $q$ .
- Both  $p$  and  $q$  are true.
- $p$  but  $q$ .

**Exclusive or**  $p \oplus q$  can be said in English as

- $p$  or  $q$ , but not both.
- Exactly one of  $p$  and  $q$  is true.

**Disjunction**  $p \vee q$  can be said in English as

- $p$  or  $q$ , or both.
- $p$  or  $q$  (inclusive).
- At least one of  $p$  and  $q$  is true.

**Conditional**  $p \rightarrow q$  can be said in English as

- |                               |                               |
|-------------------------------|-------------------------------|
| • if $p$ , then $q$ .         | • $q$ follows from $p$ .      |
| • $p$ is sufficient for $q$ . | • $p$ is sufficient for $q$ . |
| • $q$ when $p$ .              | • $q$ is necessary for $p$ .  |
| • $q$ whenever $p$ .          | • $p$ only if $q$ .           |
| • $p$ implies $q$ .           |                               |

**Biconditional**

- $p$  if and only if  $q$ .
- $p$  iff  $q$ .
- If  $p$  then  $q$ , and conversely.
- $p$  is necessary and sufficient for  $q$ .





# Algorithm redundancy

Real-life representations are often prone to corruption. Biological codes, like RNA, may mutate naturally<sup>1</sup> and during measurement; cosmic radiation and other ambient noise can flip bits in computer storage<sup>2</sup>. One way to recover from corrupted data is to introduce or exploit redundancy.

Consider the following algorithm to introduce redundancy in a string of 0s and 1s.

Create redundancy by repeating each bit three times

---

```
1 procedure redun3( $a_{k-1} \cdots a_0$ : a nonempty bitstring)
2 for  $i := 0$  to  $k-1$ 
3    $c_{3i} := a_i$ 
4    $c_{3i+1} := a_i$ 
5    $c_{3i+2} := a_i$ 
6 return  $c_{3k-1} \cdots c_0$ 
```

---

Decode sequence of bits using majority rule on consecutive three bit sequences

---

```
1 procedure decode3( $c_{3k-1} \cdots c_0$ : a nonempty bitstring whose length is an integer multiple of 3)
2 for  $i := 0$  to  $k-1$ 
3   if exactly two or three of  $c_{3i}, c_{3i+1}, c_{3i+2}$  are set to 1
4      $a_i := 1$ 
5   else
6      $a_i := 0$ 
7 return  $a_{k-1} \cdots a_0$ 
```

---

Give a recursive definition of the set of outputs of the *redun3* procedure, *Out*,

Consider the message  $m = 0001$  so that the sender calculates  $\text{redun3}(m) = \text{redun3}(0001) = 000000000111$ .

Introduce \_\_\_\_ errors into the message so that the signal received by the receiver is \_\_\_\_\_ but the receiver is still able to decode the original message.

*Challenge: what is the biggest number of errors you can introduce?*

Building a circuit for lines 3-6 in *decode* procedure: given three input bits, we need to determine whether the majority is a 0 or a 1.

$c_{3i}$	$c_{3i+1}$	$c_{3i+2}$	$a_i$	Circuit
1	1	1		
1	1	0		
1	0	1		
1	0	0		
0	1	1		
0	1	0		
0	0	1		
0	0	0		

---

<sup>1</sup>Mutations of specific RNA codons have been linked to many disorders and cancers.

<sup>2</sup>This RadioLab podcast episode goes into more detail on bit flips: <https://www.wnycstudios.org/story/bit-flip>

# Cartesian product definition

**Definition:** The **Cartesian product** of the sets  $A$  and  $B$ ,  $A \times B$ , is the set of all ordered pairs  $(a, b)$ , where  $a \in A$  and  $b \in B$ . That is:  $A \times B = \{(a, b) \mid (a \in A) \wedge (b \in B)\}$ . The Cartesian product of the sets  $A_1, A_2, \dots, A_n$ , denoted by  $A_1 \times A_2 \times \dots \times A_n$ , is the set of ordered n-tuples  $(a_1, a_2, \dots, a_n)$ , where  $a_i$  belongs to  $A_i$  for  $i = 1, 2, \dots, n$ . That is,

$$A_1 \times A_2 \times \dots \times A_n = \{(a_1, a_2, \dots, a_n) \mid a_i \in A_i \text{ for } i = 1, 2, \dots, n\}$$

## Algorithm rna mutation insertion deletion

Recall that  $S$  is defined as the set of all RNA strands, nonempty strings made of the bases in  $B = \{\text{A, U, G, C}\}$ . We define the functions

$$\begin{aligned} \text{mutation} : S \times \mathbb{Z}^+ \times B &\rightarrow S & \text{insertion} : S \times \mathbb{Z}^+ \times B &\rightarrow S \\ \text{deletion} : \{s \in S \mid \text{rnalength}(s) > 1\} \times \mathbb{Z}^+ &\rightarrow S & & \text{with rules} \end{aligned}$$

---

```
1 procedure mutation( $b_1 \dots b_n$ : a RNA strand,  $k$ : a positive integer,  $b$ : an element of  $B$ )
2 for  $i := 1$  to  $n$ 
3   if  $i = k$ 
4      $c_i := b$ 
5   else
6      $c_i := b_i$ 
7 return  $c_1 \dots c_n$  {The return value is a RNA strand made of the  $c_i$  values}
```

---

---

```
1 procedure insertion( $b_1 \dots b_n$ : a RNA strand,  $k$ : a positive integer,  $b$ : an element of  $B$ )
2 if  $k > n$ 
3   for  $i := 1$  to  $n$ 
4      $c_i := b_i$ 
5    $c_{n+1} := b$ 
6 else
7   for  $i := 1$  to  $k-1$ 
8      $c_i := b_i$ 
9    $c_k := b$ 
10  for  $i := k+1$  to  $n+1$ 
11     $c_i := b_{i-1}$ 
12 return  $c_1 \dots c_{n+1}$  {The return value is a RNA strand made of the  $c_i$  values}
```

---

---

```
1 procedure deletion( $b_1 \dots b_n$ : a RNA strand with  $n > 1$ ,  $k$ : a positive integer)
2 if  $k > n$ 
3    $m := n$ 
4   for  $i := 1$  to  $n$ 
5      $c_i := b_i$ 
6 else
7    $m := n-1$ 
8   for  $i := 1$  to  $k-1$ 
9      $c_i := b_i$ 
10  for  $i := k$  to  $n-1$ 
11     $c_i := b_{i+1}$ 
12 return  $c_1 \dots c_m$  {The return value is a RNA strand made of the  $c_i$  values}
```

---

---

# Rna mutation insertion deletion example

Trace the pseudocode to find the output of  $mutation( (AUC, 3, G) )$

Fill in the blanks so that  $insertion( (AUC, \_, \_) ) = AUCG$

Fill in the blanks so that  $deletion( (\_, \_) ) = G$

## Predicate definition

**Definition:** A **predicate** is a function from a given set (domain) to  $\{T, F\}$ .

A predicate can be applied, or **evaluated** at, an element of the domain.

Usually, a predicate *describes a property* that domain elements may or may not have.

Two predicates over the same domain are **equivalent** means they evaluate to the same truth values for all possible assignments of domain elements to the input. In other words, they are equivalent means that they are equal as functions.

To define a predicate, we must specify its domain and its value at each domain element. The rule assigning truth values to domain elements can be specified using a formula, English description, in a table (if the domain is finite), or recursively (if the domain is recursively defined).

## Predicate examples finite domain

Input $x$	Output		
	$V(x)$ $[x]_{2c,3} > 0$	$N(x)$ $[x]_{2c,3} < 0$	$Mystery(x)$
000	$F$		$T$
001	$T$		$T$
010	$T$		$T$
011	$T$		$F$
100	$F$		$F$
101	$F$		$T$
110	$F$		$F$
111	$F$		$T$

The domain for each of the predicates  $V(x)$ ,  $N(x)$ ,  $Mystery(x)$  is \_\_\_\_\_.

Fill in the table of values for the predicate  $N(x)$  based on the formula given.

## Predicate truth set definition

**Definition:** The **truth set** of a predicate is the collection of all elements in its domain where the predicate evaluates to  $T$ .

Notice that specifying the domain and the truth set is sufficient for defining a predicate.

## Predicate truth set example

The truth set for the predicate  $V(x)$  is \_\_\_\_\_.

The truth set for the predicate  $N(x)$  is \_\_\_\_\_.

The truth set for the predicate  $Mystery(x)$  is \_\_\_\_\_.

# Quantification definition

The **universal quantification** of predicate  $P(x)$  over domain  $U$  is the statement “ $P(x)$  for all values of  $x$  in the domain  $U$ ” and is written  $\forall x P(x)$  or  $\forall x \in U P(x)$ . When the domain is finite, universal quantification over the domain is equivalent to iterated *conjunction* (ands).

The **existential quantification** of predicate  $P(x)$  over domain  $U$  is the statement “There exists an element  $x$  in the domain  $U$  such that  $P(x)$ ” and is written  $\exists x P(x)$  for  $\exists x \in U P(x)$ . When the domain is finite, existential quantification over the domain is equivalent to iterated *disjunction* (ors).

An element for which  $P(x) = F$  is called a **counterexample** of  $\forall x P(x)$ .

An element for which  $P(x) = T$  is called a **witness** of  $\exists x P(x)$ .

## Quantification logical equivalence

Statements involving predicates and quantifiers are **logically equivalent** means they have the same truth value no matter which predicates (domains and functions) are substituted in.

**Quantifier version of De Morgan’s laws:**  $\boxed{\neg \forall x P(x) \equiv \exists x (\neg P(x))}$   $\boxed{\neg \exists x Q(x) \equiv \forall x (\neg Q(x))}$

## Quantification examples finite domain

Examples of quantifications using  $V(x), N(x), Mystery(x)$ :

**True or False:**  $\exists x ( V(x) \wedge N(x) )$

**True or False:**  $\forall x ( V(x) \rightarrow N(x) )$

**True or False:**  $\exists x ( N(x) \leftrightarrow Mystery(x) )$

Rewrite  $\neg \forall x ( V(x) \oplus Mystery(x) )$  into a logical equivalent statement.

Notice that these are examples where the predicates have *finite* domain. How would we evaluate quantifications where the domain may be infinite?

# Predicate rna example

**Example predicates on  $S$ , the set of RNA strands (an infinite set)**

$H : S \rightarrow \{T, F\}$  where  $H(s) = T$  for all  $s$ .

Truth set of  $H$  is \_\_\_\_\_

$F_A : S \rightarrow \{T, F\}$  defined recursively by:

Basis step:  $F_A(A) = T$ ,  $F_A(C) = F_A(G) = F_A(U) = F$

Recursive step: If  $s \in S$  and  $b \in B$ , then  $F_A(sb) = F_A(s)$ .

Example where  $F_A$  evaluates to  $T$  is \_\_\_\_\_

Example where  $F_A$  evaluates to  $F$  is \_\_\_\_\_

## Rna rna len basecount definitions

*Recall the definitions:* The set of RNA strands  $S$  is defined (recursively) by:

Basis Step:  $A \in S, C \in S, U \in S, G \in S$

Recursive Step: If  $s \in S$  and  $b \in B$ , then  $sb \in S$

where  $sb$  is string concatenation.

The function *rna len* that computes the length of RNA strands in  $S$  is defined recursively by:

		$rna len : S \rightarrow \mathbb{Z}^+$
Basis Step:	If $b \in B$ then	$rna len(b) = 1$
Recursive Step:	If $s \in S$ and $b \in B$ , then	$rna len(sb) = 1 + rna len(s)$

The function *basecount* that computes the number of a given base  $b$  appearing in a RNA strand  $s$  is defined recursively by:

		$basecount : S \times B \rightarrow \mathbb{N}$
Basis Step:	If $b_1 \in B, b_2 \in B$	$basecount((b_1, b_2)) = \begin{cases} 1 & \text{when } b_1 = b_2 \\ 0 & \text{when } b_1 \neq b_2 \end{cases}$
Recursive Step:	If $s \in S, b_1 \in B, b_2 \in B$	$basecount((sb_1, b_2)) = \begin{cases} 1 + basecount((s, b_2)) & \text{when } b_1 = b_2 \\ basecount((s, b_2)) & \text{when } b_1 \neq b_2 \end{cases}$

# Predicates example $rnalen$ basecount

Using functions to define predicates:

$L$  with domain  $S \times \mathbb{Z}^+$  is defined by, for  $s \in S$  and  $n \in \mathbb{Z}^+$ ,

$$L( (s, n) ) = \begin{cases} T & \text{if } rnalen(s) = n \\ F & \text{otherwise} \end{cases}$$

In other words,  $L( (s, n) )$  means  $rnalen(s) = n$

$BC$  with domain  $S \times B \times \mathbb{N}$  is defined by, for  $s \in S$  and  $b \in B$  and  $n \in \mathbb{N}$ ,

$$BC( (s, b, n) ) = \begin{cases} T & \text{if } basecount( (s, b) ) = n \\ F & \text{otherwise} \end{cases}$$

In other words,  $BC( (s, b, n) )$  means  $basecount( (s, b) ) = n$

Example where  $L$  evaluates to  $T$ : \_\_\_\_\_ Why?

Example where  $BC$  evaluates to  $T$ : \_\_\_\_\_ Why?

Example where  $L$  evaluates to  $F$ : \_\_\_\_\_ Why?

Example where  $BC$  evaluates to  $F$ : \_\_\_\_\_ Why?

$$\exists t \ BC(t) \qquad \exists (s, b, n) \in S \times B \times \mathbb{N} \ (basecount( (s, b) ) = n)$$

In English:

Witness that proves this existential quantification is true:

$$\forall t \ BC(t) \qquad \forall (s, b, n) \in S \times B \times \mathbb{N} \ (basecount( (s, b) ) = n)$$

In English:

Counterexample that proves this universal quantification is false:

# Predicates projecting example rna basecount

## New predicates from old

1. Define the **new** predicate with domain  $S \times B$  and rule

$$\text{basecount}( (s, b) ) = 3$$

Example domain element where predicate is  $T$ :

2. Define the **new** predicate with domain  $S \times \mathbb{N}$  and rule

$$\text{basecount}( (s, \mathbf{A}) ) = n$$

Example domain element where predicate is  $T$ :

3. Define the **new** predicate with domain  $S \times B$  and rule

$$\exists n \in \mathbb{N} (\text{basecount}( (s, b) ) = n)$$

Example domain element where predicate is  $T$ :

4. Define the **new** predicate with domain  $S$  and rule

$$\forall b \in B (\text{basecount}( (s, b) ) = 1)$$

Example domain element where predicate is  $T$ :

## Predicate notation

**Notation:** for a predicate  $P$  with domain  $X_1 \times \cdots \times X_n$  and a  $n$ -tuple  $(x_1, \dots, x_n)$  with each  $x_i \in X$ , we can write  $P(x_1, \dots, x_n)$  to mean  $P( (x_1, \dots, x_n) )$ .



# Nested quantifiers

## Nested quantifiers

$$\forall s \in S \forall b \in B \forall n \in \mathbb{N} (\text{basecount}(s, b) = n)$$

In English:

Counterexample that proves this universal quantification is false:

$$\forall n \in \mathbb{N} \forall s \in S \forall b \in B (\text{basecount}(s, b) = n)$$

In English:

Counterexample that proves this universal quantification is false:

# Alternating quantifiers strategies rna examples

## Alternating nested quantifiers

$$\forall s \in S \exists b \in B ( \text{basecount}(s, b) = 3 )$$

In English: For each RNA strand there is a base that occurs 3 times in this strand.

Write the negation and use De Morgan's law to find a logically equivalent version where the negation is applied only to the  $BC$  predicate (not next to a quantifier).

Is the original statement **True** or **False**?

$$\exists s \in S \forall b \in B \exists n \in \mathbb{N} ( \text{basecount}(s, b) = n )$$

In English: There is an RNA strand so that for each base there is some nonnegative integer that counts the number of occurrences of that base in this strand.

Write the negation and use De Morgan's law to find a logically equivalent version where the negation is applied only to the  $BC$  predicate (not next to a quantifier).

Is the original statement **True** or **False**?

## Proof strategies road map

We now have propositional and predicate logic that can help us express statements about any domain. We will develop proof strategies to craft valid argument for proving that such statements are true or disproving them (by showing they are false). We will practice these strategies with statements about sets and numbers, both because they are familiar and because they can be used to build cryptographic systems. Then we will apply proof strategies more broadly to prove statements about data structures and machine learning applications.

# Proof strategies quantification finite domain

When a predicate  $P(x)$  is over a **finite** domain:

- To show that  $\forall x P(x)$  is true: check that  $P(x)$  evaluates to  $T$  at each domain element by evaluating over and over.
- To show that  $\forall x P(x)$  is false: find one counterexample, a domain element where  $P(x)$  evaluates to  $F$ .
- To show that  $\exists x P(x)$  is true: find one witness, a domain element where  $P(x)$  evaluates to  $T$ .
- To show that  $\exists x P(x)$  is false: check that  $P(x)$  evaluates to  $F$  at each domain element by evaluating over and over.

## Proof strategy universal exhaustion

New! **Proof of universal by exhaustion:** To prove that  $\forall x P(x)$  is true when  $P$  has a finite domain, evaluate the predicate at **each** domain element to confirm that it is always  $T$ .

## Proof strategy universal generalization

New! **Proof by universal generalization:** To prove that  $\forall x P(x)$  is true, we can take an arbitrary element  $e$  from the domain of quantification and show that  $P(e)$  is true, without making any assumptions about  $e$  other than that it comes from the domain.

An **arbitrary** element of a set or domain is a fixed but unknown element from that set.

# Sets equality subset definition

## Definitions:

A **set** is an unordered collection of elements. When  $A$  and  $B$  are sets,  $A = B$  (set equality) means

$$\forall x(x \in A \leftrightarrow x \in B)$$

When  $A$  and  $B$  are sets,  $A \subseteq B$  (“ $A$  is a **subset** of  $B$ ”) means

$$\forall x(x \in A \rightarrow x \in B)$$

When  $A$  and  $B$  are sets,  $A \subsetneq B$  (“ $A$  is a **proper subset** of  $B$ ”) means

$$(A \subseteq B) \wedge (A \neq B)$$

# Proof strategies conditionals

**New! Proof of conditional by direct proof:** To prove that the conditional statement  $p \rightarrow q$  is true, we can assume  $p$  is true and use that assumption to show  $q$  is true.

**New! Proof of conditional by contrapositive proof:** To prove that the implication  $p \rightarrow q$  is true, we can assume  $q$  is false and use that assumption to show  $p$  is also false.

**New! Proof of disjunction using equivalent conditional:** To prove that the disjunction  $p \vee q$  is true, we can rewrite it equivalently as  $\neg p \rightarrow q$  and then use direct proof or contrapositive proof.

# Proof strategies proof by cases

**New! Proof by Cases:** To prove  $q$ , we can work by cases by first describing all possible cases we might be in and then showing that each one guarantees  $q$ . Formally, if we know that  $p_1 \vee p_2$  is true, and we can show that  $(p_1 \rightarrow q)$  is true and we can show that  $(p_2 \rightarrow q)$ , then we can conclude  $q$  is true.

## Proof strategies and

**New! Proof of conjunctions with subgoals:** To show that  $p \wedge q$  is true, we have two subgoals: subgoal (1) prove  $p$  is true; and, subgoal (2) prove  $q$  is true.

To show that  $p \wedge q$  is false, it's enough to prove that  $\neg p$ .

To show that  $p \wedge q$  is false, it's enough to prove that  $\neg q$ .

## Sets proof strategies

To prove that one set is a subset of another, e.g. to show  $A \subseteq B$ :

To prove that two sets are equal, e.g. to show  $A = B$ :

## Sets equality example

Example:  $\{43, 7, 9\} = \{7, 43, 9, 7\}$

## Sets basic proofs

**Prove or disprove:**  $\{A, C, U, G\} \subseteq \{AA, AC, AU, AG\}$

**Prove or disprove:** For some set  $B$ ,  $\emptyset \in B$ .

**Prove or disprove:** For every set  $B$ ,  $\emptyset \in B$ .

**Prove or disprove:** The empty set is a subset of every set.

**Prove or disprove:** The empty set is a proper subset of every set.

**Prove or disprove:**  $\{4, 6\} \subseteq \{n \mid \exists c \in \mathbb{Z}(n = 4c)\}$

**Prove or disprove:**  $\{4, 6\} \subseteq \{n \bmod 10 \mid \exists c \in \mathbb{Z}(n = 4c)\}$

## Proofs signposting

Consider ..., an **arbitrary** .... **Assume** ..., we **want to show** that .... Which is what was needed, so the proof is complete  $\square$ .

*or, in other words:*

Let ... be an **arbitrary** .... **Assume** ..., **WTS** that ... **QED**.



# Set operations union intersection powerset

**Cartesian product:** When  $A$  and  $B$  are sets,

$$A \times B = \{(a, b) \mid a \in A \wedge b \in B\}$$

Example:  $\{43, 9\} \times \{9, \mathbb{Z}\} =$

Example:  $\mathbb{Z} \times \emptyset =$

**Union:** When  $A$  and  $B$  are sets,

$$A \cup B = \{x \mid x \in A \vee x \in B\}$$

Example:  $\{43, 9\} \cup \{9, \mathbb{Z}\} =$

Example:  $\mathbb{Z} \cup \emptyset =$

**Intersection:** When  $A$  and  $B$  are sets,

$$A \cap B = \{x \mid x \in A \wedge x \in B\}$$

Example:  $\{43, 9\} \cap \{9, \mathbb{Z}\} =$

Example:  $\mathbb{Z} \cap \emptyset =$

**Set difference:** When  $A$  and  $B$  are sets,

$$A - B = \{x \mid x \in A \wedge x \notin B\}$$

Example:  $\{43, 9\} - \{9, \mathbb{Z}\} =$

Example:  $\mathbb{Z} - \emptyset =$

**Disjoint sets:** sets  $A$  and  $B$  are disjoint means  $A \cap B = \emptyset$

Example:  $\{43, 9\}, \{9, \mathbb{Z}\}$  are not disjoint

Example: The sets  $\mathbb{Z}$  and  $\emptyset$  are disjoint

**Power set:** When  $U$  is a set,  $\mathcal{P}(U) = \{X \mid X \subseteq U\}$

Example:  $\mathcal{P}(\{43, 9\}) =$

Example:  $\mathcal{P}(\emptyset) =$

## Sets basic proofs operations

Let  $W = \mathcal{P}(\{1, 2, 3, 4, 5\})$

Example elements in  $W$  are:

**Prove or disprove:**  $\forall A \in W \forall B \in W (A \subseteq B \rightarrow \mathcal{P}(A) \subseteq \mathcal{P}(B))$

*Extra example:* **Prove or disprove:**  $\forall A \in W \forall B \in W (\mathcal{P}(A) = \mathcal{P}(B) \rightarrow A = B)$

*Extra example:* **Prove or disprove:**  $\forall A \in W \forall B \in W \forall C \in W (A \cup B = A \cup C \rightarrow B = C)$

## Numbers facts

1. Addition and multiplication of real numbers are each commutative and associative.
2. The product of two positive numbers is positive, of two negative numbers is positive, and of a positive and a negative number is negative.
3. The sum of two integers, the product of two integers, and the difference between two integers are each integers.
4. For every integer  $x$  there is no integer strictly between  $x$  and  $x + 1$ ,
5. When  $x, y$  are positive integers,  $xy \geq x$  and  $xy \geq y$ .

# Factoring definition

**Definition:** When  $a$  and  $b$  are integers and  $a$  is nonzero,  $a$  **divides**  $b$  means there is an integer  $c$  such that  $b = ac$ .

Symbolically,  $F( (a, b) ) =$  \_\_\_\_\_ and is a predicate over the domain \_\_\_\_\_

Other (synonymous) ways to say that  $F( (a, b) )$  is true:

$a$  is a **factor** of  $b$        $a$  is a **divisor** of  $b$        $b$  is a **multiple** of  $a$        $a|b$

When  $a$  is a positive integer and  $b$  is any integer,  $a|b$  exactly when  $b \bmod a = 0$

When  $a$  is a positive integer and  $b$  is any integer,  $a|b$  exactly  $b = a \cdot (b \text{ div } a)$

## Factoring translation examples

*Translate these quantified statements by matching to English statement on right.*

$\exists a \in \mathbb{Z}^{\neq 0} ( F( (a, a) ) )$

Every nonzero integer is a factor of itself.

$\exists a \in \mathbb{Z}^{\neq 0} ( \neg F( (a, a) ) )$

No nonzero integer is a factor of itself.

$\forall a \in \mathbb{Z}^{\neq 0} ( F( (a, a) ) )$

At least one nonzero integer is a factor of itself.

$\forall a \in \mathbb{Z}^{\neq 0} ( \neg F( (a, a) ) )$

Some nonzero integer is not a factor of itself.

## Factoring basic claims

**Claim:** Every nonzero integer is a factor of itself.

**Proof:**

**Prove or Disprove:** There is a nonzero integer that does not divide its square.

**Prove or Disprove:** Every positive factor of a positive integer is less than or equal to it.

## Factoring basic claims continued

**Claim:** Every nonzero integer is a factor of itself and every nonzero integer divides its square.

## Factoring even odd

**Definition:** an integer  $n$  is **even** means that there is an integer  $a$  such that  $n = 2a$ ; an integer  $n$  is **odd** means that there is an integer  $a$  such that  $n = 2a + 1$ . Equivalently, an integer  $n$  is **even** means  $n \bmod 2 = 0$ ; an integer  $n$  is **odd** means  $n \bmod 2 = 1$ . Also, an integer is even if and only if it is not odd.

## Prime number definition

**Definition:** An integer  $p$  greater than 1 is called **prime** means the only positive factors of  $p$  are 1 and  $p$ . A positive integer that is greater than 1 and is not prime is called composite.

## Primes basic claims

*Extra examples:* Use the definition to prove that 1 is not prime, 2 is prime, 3 is prime, 4 is not prime, 5 is prime, 6 is not prime, and 7 is prime.

**True or False:** The statement “There are three consecutive positive integers that are prime.”

*Hint:* These numbers would be of the form  $p, p + 1, p + 2$  (where  $p$  is a positive integer).

**Proof:** We need to show \_\_\_\_\_

**True or False:** The statement “There are three consecutive odd positive integers that are prime.”

*Hint:* These numbers would be of the form  $p, p + 2, p + 4$  (where  $p$  is an odd positive integer).

**Proof:** We need to show \_\_\_\_\_

# Rna rna len basecount definitions

*Recall the definitions:* The set of RNA strands  $S$  is defined (recursively) by:

$$\begin{array}{ll} \text{Basis Step:} & \mathbf{A} \in S, \mathbf{C} \in S, \mathbf{U} \in S, \mathbf{G} \in S \\ \text{Recursive Step:} & \text{If } s \in S \text{ and } b \in B, \text{ then } sb \in S \end{array}$$

where  $sb$  is string concatenation.

The function  $rnalen$  that computes the length of RNA strands in  $S$  is defined recursively by:

$$\begin{array}{ll} & rnalen : S \rightarrow \mathbb{Z}^+ \\ \text{Basis Step:} & \text{If } b \in B \text{ then } rnalen(b) = 1 \\ \text{Recursive Step:} & \text{If } s \in S \text{ and } b \in B, \text{ then } rnalen(sb) = 1 + rnalen(s) \end{array}$$

The function  $basecount$  that computes the number of a given base  $b$  appearing in a RNA strand  $s$  is defined recursively by:

$$\begin{array}{ll} & basecount : S \times B \rightarrow \mathbb{N} \\ \text{Basis Step:} & \text{If } b_1 \in B, b_2 \in B \quad basecount( (b_1, b_2) ) = \begin{cases} 1 & \text{when } b_1 = b_2 \\ 0 & \text{when } b_1 \neq b_2 \end{cases} \\ \text{Recursive Step:} & \text{If } s \in S, b_1 \in B, b_2 \in B \quad basecount( (sb_1, b_2) ) = \begin{cases} 1 + basecount( (s, b_2) ) & \text{when } b_1 = b_2 \\ basecount( (s, b_2) ) & \text{when } b_1 \neq b_2 \end{cases} \end{array}$$

## Alternating quantifiers proofs rna examples

Which proof strategies could be used to prove each of the following statements?

*Hint: first translate the statements to English and identify the main logical structure.*

$$\forall s \in S \ ( \text{rnalen}(s) > 0 )$$

$$\forall b \in B \ \exists s \in S \ ( \text{basecount}( \ (s, b) \ ) > 0 )$$

$$\forall s \in S \ \exists b \in B \ ( \text{basecount}( \ (s, b) \ ) > 0 )$$

$$\exists s \in S \ ( \text{rnalen}(s) = \text{basecount}( \ (s, \mathbf{A}) \ ) )$$

$$\forall s \in S \ ( \text{rnalen}(s) \geq \text{basecount}( \ (s, \mathbf{A}) \ ) )$$



## Structural induction motivating example rna

**Claim**  $\forall s \in S ( \text{rnalen}(s) > 0 )$

**Proof:** Let  $s$  be an arbitrary RNA strand. By the recursive definition of  $S$ , either  $s \in B$  or there is some strand  $s_0$  and some base  $b$  such that  $s = s_0b$ . We will show that the inequality holds for both cases.

**Case:** Assume  $s \in B$ . We need to show  $\text{rnalen}(s) > 0$ . By the basis step in the definition of  $\text{rnalen}$ ,

$$\text{rnalen}(s) = 1$$

which is greater than 0, as required.

**Case:** Assume there is some strand  $s_0$  and some base  $b$  such that  $s = s_0b$ . We will show (*the stronger claim*) that

$$\forall u \in S \forall b \in B ( \text{rnalen}(u) > 0 \rightarrow \text{rnalen}(ub) > 0 )$$

Consider an arbitrary RNA strand  $u$  and an arbitrary base  $b$ , and assume towards a direct proof, that

$$\text{rnalen}(u) > 0$$

We need to show that  $\text{rnalen}(ub) > 0$ .

$$\text{rnalen}(ub) = 1 + \text{rnalen}(u) > 1 + 0 = 1 > 0$$

as required.

## Proof strategies structural induction

**Proof by Structural Induction** To prove a universal quantification over a recursively defined set:

**Basis Step:** Show the statement holds for elements specified in the basis step of the definition.

**Recursive Step:** Show that if the statement is true for each of the elements used to construct new elements in the recursive step of the definition, the result holds for these new elements.

# Structural induction example *rnalen* basecount

**Claim**  $\forall s \in S (rnalen(s) \geq basecount( (s, \mathbf{A}) ))$ :

**Proof:** We proceed by structural induction on the recursively defined set  $S$ .

**Basis Case:** We need to prove that the inequality holds for each element in the basis step of the recursive definition of  $S$ . Need to show

$$\begin{aligned} & ( rnalen(\mathbf{A}) \geq basecount( (\mathbf{A}, \mathbf{A}) ) ) \wedge ( rnalen(\mathbf{C}) \geq basecount( (\mathbf{C}, \mathbf{A}) ) ) \\ & \wedge ( rnalen(\mathbf{U}) \geq basecount( (\mathbf{U}, \mathbf{A}) ) ) \wedge ( rnalen(\mathbf{G}) \geq basecount( (\mathbf{G}, \mathbf{A}) ) ) \end{aligned}$$

We calculate, using the definitions of *rnalen* and *basecount*:

**Recursive Case:** We will prove that

$$\forall u \in S \forall b \in B ( rnalen(u) \geq basecount( (u, \mathbf{A}) ) \rightarrow rnalen(ub) \geq basecount( (ub, \mathbf{A}) ) )$$

Consider arbitrary RNA strand  $u$  and arbitrary base  $b$ . Assume, as the **induction hypothesis**, that  $rnalen(u) \geq basecount( (u, \mathbf{A}) )$ . We need to show that  $rnalen(ub) \geq basecount( (ub, \mathbf{A}) )$ .

Using the recursive step in the definition of the function *rnalen*:

$$rnalen(ub) = 1 + rnalen(u)$$

The recursive step in the definition of the function *basecount* has two cases. We notice that  $b = \mathbf{A} \vee b \neq \mathbf{A}$  and we proceed by cases.

*Case i.* Assume  $b = \mathbf{A}$ .

Using the first case in the recursive step in the definition of the function *basecount*:

$$basecount( (ub, \mathbf{A}) ) = 1 + basecount( (u, \mathbf{A}) )$$

By the **induction hypothesis**, we know that  $basecount( (u, \mathbf{A}) ) \leq rnalen(u)$  so:

$$basecount( (ub, \mathbf{A}) ) = 1 + basecount( (u, \mathbf{A}) ) \leq 1 + rnalen(u) = rnalen(ub)$$

and, thus,  $basecount( (ub, \mathbf{A}) ) \leq rnalen(ub)$ , as required.

*Case ii.* Assume  $b \neq \mathbf{A}$ .

Using the second case in the recursive step in the definition of the function *basecount*:

$$basecount( (ub, \mathbf{A}) ) = basecount( (u, \mathbf{A}) )$$

By the **induction hypothesis**, we know that  $basecount( (u, \mathbf{A}) ) \leq rnalen(u)$  so:

$$basecount( (ub, \mathbf{A}) ) = basecount( (u, \mathbf{A}) ) \leq rnalen(u) < 1 + rnalen(u) = rnalen(ub)$$

and, thus,  $basecount( (ub, \mathbf{A}) ) \leq rnalen(ub)$ , as required.

## Quiz basecount rnalen induction

Recall the definitions of the functions  $rnalen$  and  $basecount$  from class.

1. Select all and only options that give a witness for the existential quantification

$$\exists s \in S ( rnalen(s) = basecount( (s, \mathbf{U}) ) )$$

- (a)  $\mathbf{A}$
- (b)  $\mathbf{UU}$
- (c)  $\mathbf{CU}$
- (d)  $(\mathbf{U}, 1)$
- (e) None of the above.

2. Select all and only options that give a counterexample for the universal quantification

$$\forall s \in S ( rnalen(s) > basecount( (s, \mathbf{G}) ) )$$

- (a)  $\mathbf{U}$
- (b)  $\mathbf{GG}$
- (c)  $\mathbf{AG}$
- (d)  $\mathbf{CUG}$
- (e) None of the above.

3. Select all and only the true statements

- (a)  $\forall s \in S \exists b \in B ( rnalen(s) = basecount( (s, b) ) )$
- (b)  $\exists s \in S \forall b \in B ( rnalen(s) = basecount( (s, b) ) )$
- (c)

$$\begin{aligned} \forall s_1 \in S \forall s_2 \in S \forall b \in B ( ( rnalen(s_1) = basecount( (s_1, b) ) \\ \wedge rnalen(s_2) = basecount( (s_2, b) ) \wedge rnalen(s_1) = rnalen(s_2) ) \rightarrow s_1 = s_2 ) \end{aligned}$$

- (d) None of the above.

# Proofs signposting kinds of claims

To organize our proofs, it's useful to highlight which claims are most important for our overall goals. We use some terminology to describe different roles statements can have.

**Theorem:** Statement that can be shown to be true, usually an important one.

Less important theorems can be called **proposition**, **fact**, **result**, **claim**.

**Lemma:** A less important theorem that is useful in proving a theorem.

**Corollary:** A theorem that can be proved directly after another one has been proved, without needing a lot of extra work.

**Invariant:** A theorem that describes a property that is true about an algorithm or system no matter what inputs are used.

# Structural induction example robot grid



**Theorem:** A robot on an infinite 2-dimensional integer grid starts at  $(0,0)$  and at each step moves to diagonally adjacent grid point. This robot can / cannot (*circle one*) reach  $(1,0)$ .

**Definition** The set of positions the robot can visit  $P$  is defined by:

Basis Step:  $(0,0) \in P$

Recursive Step: If  $(x,y) \in P$ , then  $(x+1,y), (x-1,y), (x,y+1), (x,y-1)$  are also in  $P$

*Example elements of  $P$  are:*

**Lemma:**  $\forall (x,y) \in P (x+y \text{ is an even integer})$

*Why are we calling this a lemma?*

Proof of theorem using lemma: To show is  $(1,0) \notin P$ . Rewriting the lemma to explicitly restrict the domain of the universal, we have  $\forall (x,y) ((x,y) \in P \rightarrow (x+y \text{ is an even integer}))$ . Since the universal is true,  $((1,0) \in P \rightarrow (1+0 \text{ is an even integer}))$  is a true statement. Evaluating the conclusion of this conditional statement: By definition of long division, since  $1 = 0 \cdot 2 + 1$  (where  $0 \in \mathbb{Z}$  and  $1 \in \mathbb{Z}$  and  $0 \leq 1 < 2$  mean that 0 is the quotient and 1 is the remainder),  $1 \bmod 2 = 1$  which is not 0 so the conclusion is false. A true conditional with a false conclusion must have a false hypothesis. Thus,  $(1,0) \notin P$ , QED.  $\square$

Proof of lemma by structural induction:

**Basis Step:**

**Recursive Step:** Consider arbitrary  $(x,y) \in P$ . To show is:

$(x+y \text{ is an even integer}) \rightarrow (\text{sum of coordinates of next position is even integer})$

Assume as the induction hypothesis, **IH** that:

# Structural induction example sum of powers

The set  $\mathbb{N}$  is recursively defined. Therefore, the function  $sumPow : \mathbb{N} \rightarrow \mathbb{N}$  which computes, for input  $i$ , the sum of the nonnegative powers of 2 up to and including exponent  $i$  is defined recursively by

Basis step:  $sumPow(0) = 1$

Recursive step: If  $x \in \mathbb{N}$ , then  $sumPow(x + 1) = sumPow(x) + 2^{x+1}$

$sumPow(0) =$

$sumPow(1) =$

$sumPow(2) =$

Fill in the blanks in the following proof of

$$\forall n \in \mathbb{N} (sumPow(n) = 2^{n+1} - 1)$$

**Proof:** Since  $\mathbb{N}$  is recursively defined, we proceed by \_\_\_\_\_.

**Basis case:** We need to show that \_\_\_\_\_. Evaluating each side:  $LHS = sumPow(0) = 1$  by the basis case in the recursive definition of  $sumPow$ ;  $RHS = 2^{0+1} - 1 = 2^1 - 1 = 2 - 1 = 1$ . Since  $1 = 1$ , the equality holds.

**Recursive case:** Consider arbitrary natural number  $n$  and assume, as the \_\_\_\_\_ that  $sumPow(n) = 2^{n+1} - 1$ . We need to show that \_\_\_\_\_. Evaluating each side:

$$LHS = sumPow(n + 1) \stackrel{\text{rec def}}{=} sumPow(n) + 2^{n+1} \stackrel{\text{IH}}{=} (2^{n+1} - 1) + 2^{n+1}.$$

$$RHS = 2^{(n+1)+1} - 1 \stackrel{\text{exponent rules}}{=} 2 \cdot 2^{n+1} - 1 = (2^{n+1} + 2^{n+1}) - 1 \stackrel{\text{regrouping}}{=} (2^{n+1} - 1) + 2^{n+1}$$

Thus,  $LHS = RHS$ . The structural induction is complete and we have proved the universal generalization.  $\square$

# Proof strategy mathematical induction

## Proof by Mathematical Induction

To prove a universal quantification over the set of all integers greater than or equal to some base integer  $b$ ,

**Basis Step:** Show the property holds for  $b$ .

**Recursive Step:** Consider an arbitrary integer  $n$  greater than or equal to  $b$ , assume (as the **induction hypothesis**) that the property holds for  $n$ , and use this and other facts to prove that the property holds for  $n + 1$ .

## Linked lists definition

**Definition** The set of linked lists of natural numbers  $L$  is defined recursively by

Basis Step:  $[] \in L$

Recursive Step: If  $l \in L$  and  $n \in \mathbb{N}$ , then  $(n, l) \in L$

## Linked lists examples

Visually:

Example: the list with two nodes whose first node has 20 and whose second node has 42

## Linked list length definition

**Definition:** The length of a linked list of natural numbers  $L$ ,  $length : L \rightarrow \mathbb{N}$  is defined by

Basis Step:  $length([]) = 0$

Recursive Step: If  $l \in L$  and  $n \in \mathbb{N}$ , then  $length((n, l)) = 1 + length(l)$



## Linked lists prepend definition

**Definition:** The function  $prepend : L \times \mathbb{N} \rightarrow L$  that adds an element at the front of a linked list is defined by

## Linked list append definition

**Definition** The function  $append : L \times \mathbb{N} \rightarrow L$  that adds an element at the end of a linked list is defined by

Basis Step: If  $m \in \mathbb{N}$  then

Recursive Step: If  $l \in L$  and  $n \in \mathbb{N}$  and  $m \in \mathbb{N}$ , then

# Linked list append length claim proof

**Claim:**  $\forall l \in L \ ( \ length( \ append( \ (l, 100) \ ) \ ) > length(l) \ )$

**Proof:** By structural induction on  $L$ , we have two cases:

## Basis Step

1. **To Show**  $length( \ append( \ ( [], 100) \ ) \ ) > length( \ [] \ )$  Because  $[]$  is the only element defined in the basis step of  $L$ , we only need to prove that the property holds for  $[]$ .
2. **To Show**  $length( \ (100, []) \ ) > length( \ [] \ )$  By basis step in definition of *append*.
3. **To Show**  $(1 + length( \ [] \ )) > length( \ [] \ )$  By recursive step in definition of *length*.
4. **To Show**  $1 + 0 > 0$  By basis step in definition of *length*.
5.  $T$  By properties of integers

QED

Because we got to  $T$  only by rewriting **To Show** to equivalent statements, using well-defined proof techniques, and applying definitions.

## Recursive Step

Consider an arbitrary:  $l' \in L$ ,  $n \in \mathbb{N}$ , and we assume as the **induction hypothesis** that:

$$length( \ append( \ (l', 100) \ ) \ ) > length(l')$$

Our goal is to show that  $length( \ append( \ ( (n, l'), 100) \ ) \ ) > length( \ (n, l') \ )$  is also true. We start by working with one side of the candidate inequality:

$$\begin{aligned} LHS &= length( \ append( \ ( (n, l'), 100) \ ) \ ) \\ &= length( \ (n, append( \ (l', 100) \ ) \ ) \ ) && \text{by the recursive definition of } append \\ &= 1 + length( \ append( \ (l', 100) \ ) \ ) && \text{by the recursive definition of } length \\ &> 1 + length(l') && \text{by the induction hypothesis} \\ &= length((n, l')) && \text{by the recursive definition of } length \\ &= RHS \end{aligned}$$

## Linked list example each length

Prove or disprove:  $\forall n \in \mathbb{N} \exists l \in L ( \text{length}(l) = n )$

## Quiz linked list definitions

Recall the definition of linked lists from class.

Consider this (incomplete) definition:

**Definition** The function *increment* : \_\_\_\_\_ that adds 1 to the data in each node of a linked list is defined by:

$$\begin{array}{ll} \text{Basis Step:} & \text{increment} : \text{_____} \rightarrow \text{_____} \\ & \text{increment}(\[]) = \text{[]} \\ \text{Recursive Step: If } l \in L, n \in \mathbb{N} & \text{increment}((n, l)) = (1 + n, \text{increment}(l)) \end{array}$$

Consider this (incomplete) definition:

**Definition** The function *sum* :  $L \rightarrow \mathbb{N}$  that adds together all the data in nodes of the list is defined by:

$$\begin{array}{ll} \text{Basis Step:} & \text{sum} : L \rightarrow \mathbb{N} \\ & \text{sum}(\[]) = 0 \\ \text{Recursive Step: If } l \in L, n \in \mathbb{N} & \text{sum}((n, l)) = \text{_____} \end{array}$$

You will compute a sample function application and then fill in the blanks for the domain and codomain of each of these functions.

- Based on the definition, what is the result of  $\text{increment}((4, (2, (7, []))))$ ? Write your answer directly with no spaces.
- Which of the following describes the domain and codomain of *increment*?
  - The domain is  $L$  and the codomain is  $\mathbb{N}$
  - The domain is  $L$  and the codomain is  $\mathbb{N} \times L$
  - The domain is  $L \times \mathbb{N}$  and the codomain is  $L$
  - The domain is  $L \times \mathbb{N}$  and the codomain is  $\mathbb{N}$
  - The domain is  $L$  and the codomain is  $L$
  - None of the above
- Assuming we would like  $\text{sum}((5, (6, [])))$  to evaluate to 11 and  $\text{sum}((3, (1, (8, []))))$  to evaluate to 12, which of the following could be used to fill in the definition of the recursive case of *sum*?
  - $\begin{cases} 1 + \text{sum}(l) & \text{when } n \neq 0 \\ \text{sum}(l) & \text{when } n = 0 \end{cases}$
  - $1 + \text{sum}(l)$
  - $n + \text{increment}(l)$
  - $n + \text{sum}(l)$
  - None of the above

4. Choose only and all of the following statements that are **well-defined**; that is, they correctly reflect the domains and codomains of the functions and quantifiers, and respect the notational conventions we use in this class. Note that a well-defined statement may be true or false.

- |   |   |
|---|---|
| (a) $\forall l \in L (sum(l))$                  | (e) $\forall l \in L \forall n \in \mathbb{N} ((n \times l) \subseteq L)$ |
| (b) $\exists l \in L (sum(l) \wedge length(l))$ | (f) $\forall l_1 \in L \exists l_2 \in L (increment(sum(l_1)) = l_2)$     |
| (c) $\forall l \in L (sum(increment(l)) = 10)$  | (g) $\forall l \in L (length(increment(l)) = length(l))$                  |
| (d) $\exists l \in L (sum(increment(l)) = 10)$  |   |

5. Choose only and all of the statements in the previous part that are both well-defined and true.

## Proof strategy mathematical induction

### Proof by Mathematical Induction

To prove a universal quantification over the set of all integers greater than or equal to some base integer  $b$ ,

**Basis Step:** Show the property holds for  $b$ .

**Recursive Step:** Consider an arbitrary integer  $n$  greater than or equal to  $b$ , assume (as the **induction hypothesis**) that the property holds for  $n$ , and use this and other facts to prove that the property holds for  $n + 1$ .

## Proof strategy strong induction

### Proof by Strong Induction

To prove that a universal quantification over the set of all integers greater than or equal to some base integer  $b$  holds, pick a fixed nonnegative integer  $j$  and then:

**Basis Step:** Show the statement holds for  $b, b + 1, \dots, b + j$ .

**Recursive Step:** Consider an arbitrary integer  $n$  greater than or equal to  $b + j$ , assume (as the **strong induction hypothesis**) that the property holds for **each of**  $b, b + 1, \dots, n$ , and use this and other facts to prove that the property holds for  $n + 1$ .

# Binary expansions exist proof

**Theorem:** Every positive integer is a sum of (one or more) distinct powers of 2. *Binary expansions exist!*

Recall the definition for binary expansion:

**Definition** For  $n$  a positive integer, the **binary expansion of  $n$**  is

$$(a_{k-1} \cdots a_1 a_0)_b$$

where  $k$  is a positive integer,  $a_0, a_1, \dots, a_{k-1}$  are each 0 or 1,  $a_{k-1} \neq 0$ , and

$$n = \sum_{i=0}^{k-1} a_i b^i$$

The idea in the “Least significant first” algorithm for computing binary expansions is that the binary expansion of *half* a number becomes *part* of the binary expansion of the number of itself. We can use this idea in a proof by strong induction that binary expansions exist for all positive integers  $n$ .

**Proof by strong induction**, with  $b = 1$  and  $j = 0$ .

**Basis step:** WTS property is true about 1.

**Recursive step:** Consider an arbitrary integer  $n \geq 1$ .

Assume (as the strong induction hypothesis, IH) that the property is true about each of  $1, \dots, n$ .

WTS that the property is true about  $n + 1$ .

*Idea:* We will apply the IH to  $(n + 1) \text{ div } 2$ .

*Why is this ok?*

*Why is this helpful?*

By the IH, we can write  $(n+1) \text{ div } 2$  as a sum of powers of 2. In other words, there are values  $a_{k-1}, \dots, a_0$  such that each  $a_i$  is 0 or 1,  $a_{k-1} = 1$ , and

$$\sum_{i=0}^{k-1} a_i 2^i = (n+1) \text{ div } 2$$

Define the collection of coefficients

$$c_j = \begin{cases} a_{j-1} & \text{if } 1 \leq j \leq k \\ (n+1) \bmod 2 & \text{if } j = 0 \end{cases}$$

Calculating:

$$\begin{aligned} \sum_{j=0}^k c_j 2^j &= c_0 + \sum_{j=1}^k c_j 2^j = c_0 + \sum_{i=0}^{k-1} c_{i+1} 2^{i+1} && \text{re-indexing the summation} \\ &= c_0 + 2 \cdot \sum_{i=0}^{k-1} c_{i+1} 2^i && \text{factoring out a 2 from each term in the sum} \\ &= c_0 + 2 \cdot \sum_{i=0}^{k-1} a_i 2^i && \text{by definition of } c_{i+1} \\ &= c_0 + 2 \cdot ((n+1) \text{ div } 2) && \text{by IH} \\ &= ((n+1) \bmod 2) + 2 \cdot ((n+1) \text{ div } 2) && \text{by definition of } c_0 \\ &= n+1 && \text{by definition of long division} \end{aligned}$$

Thus,  $n+1$  can be expressed as a sum of powers of 2, as required.

# Strong induction making change proof idea

Suppose we had postage stamps worth 5 cents and 3 cents. Which number of cents can we form using these stamps? In other words, which postage can we pay?

11?

15?

4?

$$\begin{aligned} &CanPay(0) \wedge \neg CanPay(1) \wedge \neg CanPay(2) \wedge \\ &CanPay(3) \wedge \neg CanPay(4) \wedge CanPay(5) \wedge CanPay(6) \\ &\neg CanPay(7) \wedge \forall n \in \mathbb{Z}^{\geq 8} CanPay(n) \end{aligned}$$

where the predicate  $CanPay$  with domain  $\mathbb{N}$  is

$$CanPay(n) = \exists x \in \mathbb{N} \exists y \in \mathbb{N} (5x + 3y = n)$$

**Proof** (idea): First, explicitly give witnesses or general arguments for postages between 0 and 7. To prove the universal claim, we can use mathematical induction or strong induction.

*Approach 1, mathematical induction:* if we have stamps that add up to  $n$  cents, need to use them (and others) to give  $n + 1$  cents. How do we get 1 cent with just 3-cent and 5-cent stamps?

Either take away a 5-cent stamps and add two 3-cent stamps,

or take away three 3-cent stamps and add two 5-cent stamps.

The details of this proof by mathematical induction are making sure we have enough stamps to use one of these approaches.

*Approach 2, strong induction:* assuming we know how to make postage for **all** smaller values (greater than or equal to 8), when we need to make  $n + 1$  cents, add one 3 cent stamp to however we make  $(n + 1) - 3$  cents.

The details of this proof by strong induction are making sure we stay in the domain of the universal when applying the induction hypothesis.



## Strong induction nim

Consider the following game: two players start with two (equal) piles of jellybeans in front of them. They take turns removing any positive integer number of jellybeans at a time from one of two piles in front of them in turns.

The player who removes the last jellybean wins the game.

Which player (if any) has a strategy to guarantee to win the game?

Try out some games, starting with 1 jellybean in each pile, then 2 jellybeans in each pile, then 3 jellybeans in each pile. Who wins in each game?

Notice that reasoning about the strategy for the 1 jellybean game is easier than about the strategy for the 2 jellybean game.

*Formulate a winning strategy by working to transform the game to a simpler one we know we can win.*

*Player 2's Strategy:* Take the same number of jellybeans that Player 1 did, but from the opposite pile.

*Why is this a good idea:* If Player 2 plays this strategy, at the next turn Player 1 faces a game with the same setup as the original, just with fewer jellybeans in the two piles. Then Player 2 can keep playing this strategy to win.

**Claim:** Player 2's strategy guarantees they will win the game.

**Proof:** By strong induction, we will prove that for all positive integers  $n$ , Player 2's strategy guarantees a win in the game that starts with  $n$  jellybeans in each pile.

**Basis step:** WTS Player 2's strategy guarantees a win when each pile starts with 1 jellybean.

In this case, Player 1 has to take the jellybean from one of the piles (because they can't take from both piles at once). Following the strategy, Player 2 takes the jellybean from the other pile, and wins because this is the last jellybean.

**Recursive step:** Let  $n$  be a positive integer. As the strong induction hypothesis, assume that Player 2's strategy guarantees a win in the games where there are  $1, 2, \dots, n$  many jellybeans in each pile at the start of the game.

WTS that Player 2's strategy guarantees a win in the game where there are  $n + 1$  in the jellybeans in each pile at the start of the game.

In this game, the first move has Player 1 take some number, call it  $c$  (where  $1 \leq c \leq n + 1$ ), of jellybeans from one of the piles. Playing according to their strategy, Player 2 then takes the same number of jellybeans from the other pile.

Notice that  $(c = n + 1) \vee (c \leq n)$ .

*Case 1:* Assume  $c = n + 1$ , then in their first move, Player 2 wins because they take all of the second pile, which includes the last jellybean.

*Case 2:* Assume  $c \leq n$ . Then after Player 2's first move, the two piles have an equal number of jellybeans. The number of jellybeans in each pile is

$$(n + 1) - c$$

and, since  $1 \leq c \leq n$ , this number is between 1 and  $n$ . Thus, at this stage of the game, the game appears identical to a new game where the two piles have an equal number of jellybeans between 1 and  $n$ . Thus, the strong induction hypothesis applies, and Player 2's strategy guarantees they win.

## Proof strategy proof by contradiction

### New! Proof by Contradiction

To prove that a statement  $p$  is true, pick another statement  $r$  and once we show that  $\neg p \rightarrow (r \wedge \neg r)$  then we can conclude that  $p$  is true.

*Informally* The statement we care about can't possibly be false, so it must be true.

# Rational numbers definition

The **set of rational numbers**,  $\mathbb{Q}$  is defined as

$$\left\{ \frac{p}{q} \mid p \in \mathbb{Z} \text{ and } q \in \mathbb{Z} \text{ and } q \neq 0 \right\} \quad \text{or, equivalently,} \quad \{x \in \mathbb{R} \mid \exists p \in \mathbb{Z} \exists q \in \mathbb{Z}^+ (p = x \cdot q)\}$$

*Extra practice:* Use the definition of set equality to prove that the definitions above give the same set.

## Sets numbers subsets

We have the following subset relationships between sets of numbers:

$$\mathbb{Z}^+ \subsetneq \mathbb{N} \subsetneq \mathbb{Z} \subsetneq \mathbb{Q} \subsetneq \mathbb{R}$$

Which of the proper subset inclusions above can you prove?

# Proof by contradiction irrational

**Goal:** The square root of 2 is not a rational number. In other words:  $\neg \exists x \in \mathbb{Q}(x^2 - 2 = 0)$

**Attempted proof:** The definition of the set of rational numbers is the collection of fractions  $p/q$  where  $p$  is an integer and  $q$  is a nonzero integer. Looking for a **witness**  $p$  and  $q$ , we can write the square root of 2 as the fraction  $\sqrt{2}/1$ , where 1 is a nonzero integer. Since the numerator is not in the domain, this witness is not allowed, and we have shown that the square root of 2 is not a fraction of integers (with nonzero denominator). Thus, the square root of 2 is not rational.

*The problem in the above attempted proof is that* \_\_\_\_\_

**Lemma 1:** For every two integers  $a$  and  $b$ , not both zero, with  $\gcd(a, b) = 1$ , it is not the case that both  $a$  is even and  $b$  is even.

**Lemma 2:** For every integer  $x$ ,  $x$  is even if and only if  $x^2$  is even.

**Proof:** Towards a proof by contradiction, we will define a statement  $r$  such that  $\sqrt{2} \in \mathbb{Q} \rightarrow (r \wedge \neg r)$ .

Assume that  $\sqrt{2} \in \mathbb{Q}$ . Namely, there are positive integers  $p, q$  such that

$$\sqrt{2} = \frac{p}{q}$$

Let  $a = \frac{p}{\gcd(p, q)}$ ,  $b = \frac{q}{\gcd(p, q)}$ , then

$$\sqrt{2} = \frac{a}{b} \quad \text{and} \quad \gcd(a, b) = 1$$

By Lemma 1,  $a$  and  $b$  are not both even. We define  $r$  to be the statement “ $a$  is even and  $b$  is even”, and we have proved  $\neg r$ .

Squaring both sides and clearing denominator:  $2b^2 = a^2$ .

By definition of even, since  $b^2$  is an integer,  $a^2$  is even.

By Lemma 2, this guarantees that  $a$  is even too. So, by definition of even, there is some integer (call it  $c$ ), such that  $a = 2c$ .

Plugging into the equation:

$$2b^2 = a^2 = (2c)^2 = 4c^2$$

and dividing both sides by 2

$$b^2 = 2c^2$$

and since  $c^2$  is an integer,  $b^2$  is even. By Lemma 2,  $b$  is even too. Thus,  $a$  is even and  $b$  is even and we have proved  $r$ .

In other words, assuming that  $\sqrt{2} \in \mathbb{Q}$  guarantees  $r \wedge \neg r$ , which is impossible, so  $\sqrt{2} \notin \mathbb{Q}$ . QED