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EE-360 Algorithms Assignment 1  
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Gale-Shapley Pseudocode:  
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This implementation of the Gale-Shapley algorithm puts students in a queue and goes through their preferred internship list until they find an internship that will accept them. They are temporarily paired until all students are matched or have reached the end of their list, in which case they are left unemployed. An internship accepts a student if it has room for a student or the student is better than their current worst student (tracked using a min heap).

```
// let M be the matching
// let M.m and M.n be the number of internships and students respectively
// let M.internshipSlots be the number of slots in each internship
// let M.studentPref be the student preference list
// let M.internPref be the internship preference list
```

Gale-Shapley(M):

```
    O = Optimization(M)

    studentQueue = [0, 1, 2, ... , M.n] // add all students
    studentMatching = [-1, -1, -1, ... , -1] // -1 for all students

    // the current index of internship preference the student is at
    studentCIPI = [0, 0, 0, ... , 0] // 0 for all students

    internshipMatchings = []
    for internship in (0, M.m):
        internshipMatchings.add(new InternshipMatching(M, O,
internship))

    while studentQueue is not empty:
        student = studentQueue.removeMin()
        internshipIndex = studentCIPI[student]

        matched = false

        while not matched:
            if internshipIndex is out of range of
internships:
                matched = true
                studentMatching[student] = -1
            else:
                internship =
M.studentPref[student][internshipIndex]
```

```

        if
internshipMatchings[internship].canAdd(student):
        rejectedStudent =
internshipMatchings[internship].add(student)
        if rejectedStudent is not -
1:
        studentQueue.add(rejectedStudent)
        studentMatching[student] =
internship
        return studentMatching

```

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class InternshipMatching:
    studentHeap = []

    init(M, O, internship):
        initilize studentHeap as min heap to prioritize least
preferred students

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        canAdd(student):
            if studentHeap.size < number of slots in internship:
                return true
            else if studentHeap.peek() is less preferred than
student:
                return true
            else:
                return false

```

```

        addStudent(student):
            rejectedStudent = -1

            studentHeap.add(student)

            if studentHeap.size >= number of slots in internship
                rejectedStudent = studentHeap.removeMin()

            return rejectedStudent

```

```

//optimizes access times for finding the preference index in a preference
list

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Optimization(M):
    O.invStudentPref = Invert all lists (M.studentPref)
    O.invInternPref = Invert all lists (M.internPref)
    return O

```

```

// inverts the given list and maps the value at an index to the index

```

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Invert(list):
    invertedList = list
    for i in (0, list.length):
        invertedList[list[i]] = i

```

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return invertedList
```

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***Overall Time Complexity***  
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The overall time complexity of this program is:  
 $O(m * n * \log k)$  where  $m$ ,  $n$ ,  $k$  are the number of internships, students,  
and the max number of slots an internship has respectively.

This complexity simplifies to  $O(m * n)$  if  $k \ll m$ ,  $n$ .

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Time Complexity Explanation:  
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- Invert takes  $O(n)$ , given list size  $n$
- Optimization takes  $n*m + m*n = O(m*n)$ , given number of students and internships,  $n$ ,  $m$
- Creating a InternshipMatching takes  $O(1)$
- GaleShapley takes  $O(\text{Optimization}) + O(3*n) + O(m) + O(\text{main matching loop})$
- main matching loop takes  $O(m * n * \log k)$ , where  $m$  is the number of internships,  $n$  is the number of students, and  $k$  the largest number of slots an internship has.
  - this is because, at the worst case, the loops could iterate through each student and their list of internship preferences ( $m * n$ )
  - within each of these, we are finding the the least preferred student for a given internship. Using a min heap to keep track of the least preferred students, this process will take  $\log k$  time, where  $k$  is the size of the heap, and worst case is equal to the number of slots the internship has available.
  - multiplying it out, we have that the main loop takes  $O(m * n * \log k)$  in time