

Tutorial 12: Graphs

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Introduction to Graphs

A **graph** $G = (V, E)$ consists of a set V of vertices (nodes) and a set E of edges connecting pairs of vertices. Graphs can be directed or undirected:

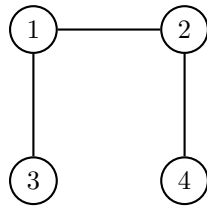
- **Directed Graph:** Edges have direction.
- **Undirected Graph:** Edges have no direction.

Graph Representation

Adjacency Matrix

A square matrix A , where $A[i, j] = 1$ if an edge exists from vertex i to vertex j ; otherwise, 0.

Example for an undirected graph:



Adjacency Matrix:

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

Adjacency List

An array of lists, where each index represents a vertex and stores a list of its adjacent vertices.

Adjacency List for the same graph:

$1 \rightarrow [2, 3]$
 $2 \rightarrow [1, 4]$
 $3 \rightarrow [1]$
 $4 \rightarrow [2]$

Breadth-First Search (BFS)

BFS explores vertices level by level, using a queue structure.

Algorithm 1 BFS(G, s)

```
1: for each vertex  $u \in V - \{s\}$  do
2:   color[ $u$ ] = WHITE
3:   distance[ $u$ ] =  $\infty$ , parent[ $u$ ] = NIL
4: end for
5: color[ $s$ ] = GRAY, distance[ $s$ ] = 0, parent[ $s$ ] = NIL
6:  $Q \leftarrow$  empty queue; enqueue( $Q, s$ )
7: while  $Q$  is not empty do
8:    $u \leftarrow$  dequeue( $Q$ )
9:   for each vertex  $v$  adjacent to  $u$  do
10:    if color[ $v$ ] = WHITE then
11:      color[ $v$ ] = GRAY, distance[ $v$ ] = distance[ $u$ ] + 1
12:      parent[ $v$ ] =  $u$ , enqueue( $Q, v$ )
13:    end if
14:  end for
15:  color[ $u$ ] = BLACK
16: end while
```

Example: BFS from vertex 1

Initial graph (above example):

Queue: $Q = [1]$; Visited: 1 (distance=0)

Dequeue 1: enqueue neighbors 2,3

Queue: $Q = [2, 3]$; Visited: 1,2,3

Dequeue 2: enqueue neighbor 4

Queue: $Q = [3, 4]$; Visited: 1,2,3,4

Dequeue 3: no unvisited neighbors

Queue: $Q = [4]$

Dequeue 4: no unvisited neighbors

Queue empty, BFS complete.

BFS traversal: $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$.

Depth-First Search (DFS)

DFS explores as far as possible along each branch before backtracking, using recursion or a stack.

Algorithm 2 DFS(G)

```
1: for each vertex  $u \in V$  do
2:   color[ $u$ ] = WHITE, parent[ $u$ ] = NIL
3: end for
4: time = 0
5: for each vertex  $u \in V$  do
6:   if color[ $u$ ] = WHITE then
7:     DFS-Visit( $u$ )
8:   end if
9: end for
```

Algorithm 3 DFS-Visit(u)

```
1: color[ $u$ ] = GRAY, time = time + 1, discovery[ $u$ ] = time
2: for each vertex  $v$  adjacent to  $u$  do
3:   if color[ $v$ ] = WHITE then
4:     parent[ $v$ ] =  $u$ , DFS-Visit( $v$ )
5:   end if
6: end for
7: color[ $u$ ] = BLACK, time = time + 1, finish[ $u$ ] = time
```

Example: DFS starting from vertex 1

- Visit 1 (time=1), explore neighbor 2
- Visit 2 (time=2), explore neighbor 4
- Visit 4 (time=3), no unvisited neighbors, finish 4 (time=4)
- Finish 2 (time=5), back to 1, explore neighbor 3
- Visit 3 (time=6), finish 3 (time=7)
- Finish 1 (time=8)

DFS traversal: $1 \rightarrow 2 \rightarrow 4 \rightarrow 3$.

Discovery/Finish times:

1 : (1, 8) 2 : (2, 5) 4 : (3, 4) 3 : (6, 7)