

Logic: Passing Criteria:

- (i) Score at least 30% aggregate (p)
and
(ii) Score at least 25% in midsem (q)
or

Score at least 25% in endsem (r)

$$b \wedge (q \vee r)$$

Every computer is connected to the university network.

CS3 is connected to the university network.

How to represent using logical expression] Domain \rightarrow the set of all computers from the university network

Example: Domain = set of all integers. " $x > 3$ ". variable x

There exists x $\exists x$

For all x $\forall x$

There is a computer of the university network that is under malware attack.

the set of computers from the university network

Computer x is under malware attack

↓ SUBJECT

↓ PREDICATE

a property that the subject of the state can have

$P(x) \rightarrow$ Computer x is under malware attack x

There is x s.t. $P(x)$

For all x $P(x)$ For all x : $x > 3$

$\forall x P(x) \rightarrow$ false

$P(5) : 5 > 3 \rightarrow$ true

$P(1) : 1 > 3 \rightarrow$ false

$P(x) : x^2 \geq 0$ where the domain is the set of all real numbers

$P(0) : 0^2 \geq 0 \rightarrow$ true
 $P(1) : 1^2 \geq 0 \rightarrow$ true

$P(-2) :$
 $4 \geq 0$
true

square of every real number is non-negative (≥ 0)

$\mathbb{N} = \{0, 1, 2, \dots\} =$ natural numbers
= non-negative integers

$\forall x P(x)$

$P(x) : x^2 \geq 0$
true

Quantifiers:

$Q(x) : x < 0$

$Q(-3) : -3 < 0$

$Q(2) : 2 < 0$

→ false

$\exists x Q(x) \rightarrow$ true because

$\forall x$

→ universal quantifier

$\exists x$

→ existential quantifier.

$P(x, y) \equiv x = y + 3$, domain

is the set of all real numbers.

$$\begin{cases} x = 7 \\ y = 4 \end{cases}$$

$$P(7, 4) : 7 = 4 + 3$$

$P(7, 4)$ is true

$P(2, 3)$

$$\begin{cases} x = 2 \\ y = 3 \end{cases}$$

$$\int 2 = 3 + 3$$

is false

false.

The truth value of $P(x, y)$ depends on the values of x and y .

$P(x)$ is a predicate where x is a variable from some domain.

Then, $\forall x P(x)$ is true when $P(x)$ is true for all values of x in the domain

Example: $P(x) : x > 2$ domain

is the set of all real numbers.

Why called first order logic -

Predicate logic quantification

$\forall x \exists x$ are over the variables from some domain

$\forall x P(x)$ is true if and only if $\exists x \neg P(x)$ is false

$\exists x \neg P(x)$ is false. No matter of which x is considered $\neg P(x)$ is false

For all x $\neg P(x)$ is false

if $x = 3$, then $P(3)$ is true

$P(2, 3)$

$$\begin{cases} x = 2 \\ y = 3 \end{cases}$$

$$\int 2 = 3 + 3$$

is false

false.

$\forall x P(x)$ is false when there is

an element x in the domain such that $P(x)$ is false

equivalently, $\exists x \neg P(x)$ is true

$$P(x) : x > 2 \rightarrow \neg P(x) : x \leq 2$$

$\forall x P(x)$ is false

$$\exists x \neg P(x)$$

But $\exists x P(x)$ is true

is true

\downarrow
 $x = 4$ is a value

such that $P(4)$ is true

\downarrow
 $x = 0$ is a value

$\neg P(x)$ is true.

$\forall x P(x)$ is false if and only if

$\exists x \neg P(x)$ is true

$$\forall x P(x) \equiv \neg (\exists x \neg P(x))$$

$\exists x \neg P(x)$ is true when there

are some elements a_1, a_2, \dots, a_n of the domain such that $\neg P(x)$ is true.

$\exists x \neg P(x)$ is false then even

\downarrow
 $P(x)$ is true

For the elements other than a_1, a_2, \dots, a_n $\neg P(x)$ is false

for the elements a_1, a_2, \dots, a_n from the domain $\neg P(x)$ is false. \checkmark

For the elements other than a_1, a_2, \dots, a_n from the domain $\neg P(x)$ is false. \checkmark

For all elements in the domain $\neg P(x)$ is false. \rightarrow for all elements in the domain $P(x)$ is true.

$\forall x P(x)$ is false when there is x in the domain s.t. $P(x)$ is false.

$\exists x P(x)$ is true when there is some element x in the domain for which $P(x)$ is true.

$\exists x P(x)$ is false when $P(x)$ is false for all elements in the domain.

$\forall x$ = universal quantifier

$\exists x$ = existential quantifier.

$A(x)$: x is divisible by 2

the domain is set of all integers

$\forall x (A(x) \rightarrow B(x))$ \rightarrow True

For every integer x , if x is divisible by 2, then x is even

$\forall x (B(x) \rightarrow A(x))$ \rightarrow is True

For every integer x , if x is even then x is divisible by 2

$\forall x P(x)$ The domain is the

$\forall x Q(x)$ same.

$B(x)$: x is an even number.

a is divisible by b if there is an integer k such that $a = kb$

$b \mid a$ \rightarrow b divides a
 $\rightarrow a$ is divisible by b

$\forall x ((A(x) \rightarrow B(x)) \wedge (B(x) \rightarrow A(x)))$

$\equiv \forall x (A(x) \leftrightarrow B(x))$

If $\forall x P(x)$ is true and $\forall x Q(x)$ is true then can we say $\forall x (P(x) \wedge Q(x))$

Yes, because for every possible element x in the domain, $P(x)$ is true and $Q(x)$ is true.

Therefore $P(x) \wedge Q(x)$ is true.

Hence $\forall x (P(x) \wedge Q(x))$ is true.

The particular element for which $P(x)$ is true is a

That does not imply that $P(x)$ and $Q(x)$ both are true

That is why $\exists x (P(x) \wedge Q(x))$ cannot be true.

$P(x)$: x is an even integer

$Q(x)$: x is an odd integer

Domain: set of all integers.

$\exists x P(x)$ is true

$\exists x Q(x)$ is true

$\exists x (P(x) \wedge Q(x))$

is false

is true?
domain is the same
 $\exists x P(x)$ is true
 $\exists x Q(x)$ is true
can we say that $\exists x (P(x) \wedge Q(x))$ is true? NO

The particular element for which $P(x)$ is true is a .

for the same element in the domain.

Suppose $\exists x (P(x) \wedge Q(x))$ is true.

Then, $(\exists x P(x)) \wedge (\exists x Q(x))$ is true

Because there ^{element} b such that $P(b) \wedge Q(b)$ is true. $\exists x P(x)$ is true
 \downarrow
 $P(b)$ is true and $Q(b)$ is true.

Section 1.4 \rightarrow Rosen's Book

Predicate Logic