

Propositional Logic p, q

$\underline{p \vee q}$, $\underline{p \wedge q}$, $\underline{p \oplus q}$, $\neg p$

$\underline{p \oplus q}$ p XOR q

$p \oplus q$ is true when either p or q is true but not both. otherwise false

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

Biconditional statement

$p \leftrightarrow q$
 p if and only if q
"p iff q"

$p \leftrightarrow q$ is true (when both p, q are true or both p, q are false.)

p	q	$p \leftrightarrow q$	$\neg(p \oplus q)$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	T	T

$p \oplus q$ is true when $p \leftrightarrow q$ is false

$p \oplus q$ is false when $p \leftrightarrow q$ is true

$\neg(p \oplus q) \equiv (p \leftrightarrow q)$
 equivalent to

$\neg(p \oplus q)$ and $(p \leftrightarrow q)$ give identical truth values for every possible combination of truth values of p and q .

That is why $\neg(p \oplus q)$ and $(p \leftrightarrow q)$ are logically equivalent.

Tautology: If a compound proposition evaluates to true for all possible combinations of propositional variables, then it is called a tautology.

Hence, $\neg(p \rightarrow q) \rightarrow \neg q$ is a tautology.

Example: $\neg(p \rightarrow q) \rightarrow \neg q$

p	q	$\neg q$	$p \rightarrow q$	$\neg(p \rightarrow q)$	$\neg(p \rightarrow q) \rightarrow \neg q$
T	T	F	T	F	T
T	F	T	F	T	T
F	T	F	T	F	T
F	F	T	T	F	T

every possible combination of p and q evaluates to true

$\neg(p \rightarrow q) \rightarrow \neg q$

$\neg(\neg(p \rightarrow q) \rightarrow \neg q)$ is a contradiction.

A compound proposition is a contradiction if it evaluates to false for every possible combination of truth values of propositional variables.

How to prove that a compound proposition is a tautology?

Construct the truth table for that compound proposition.

How to disprove that a compound proposition is a tautology?

Construct truth table.

Mention which particular values of propositional variables evaluate to false.

How to prove that two compound propositions are logically equivalent?

Construct truth tables for both the propositions.

In truth table, it must ensure that every possible truth values evaluates the same.

As $p \vee \neg(p \vee q)$ and $(\neg p \wedge \neg q)$ evaluate same for every combination of truth values.

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

Truth table for $\neg(p \vee q)$

p	q	$\neg p$	$\neg q$	$p \vee q$	$\neg p \wedge \neg q$	$\neg(p \vee q)$
T	T	F	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	T	F	F
F	F	T	T	F	T	T

$\neg(p \vee q)$ and $(\neg p \wedge \neg q)$ are logically equivalent.

$$\begin{aligned} \neg(p \vee q) &\equiv (\neg p \wedge \neg q) \\ \neg(p \wedge q) &\equiv (\neg p \vee \neg q) \end{aligned} \rightarrow \text{De-Morgan's Law}$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

Exercise? use the truth table to prove

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

Distributive Laws

Exercise: Prove that $(\neg p \vee q) \equiv (p \rightarrow q)$

$p \wedge \neg p$	p	$\neg p$	$p \wedge \neg p$
	T	F	F
	F	T	F

If p and q are logically equivalent then $(p \leftrightarrow q)$ is a tautology.

$$p \equiv q$$

Conditional disjunction equivalence.

What we know?

$$(i) (p \rightarrow q) \equiv \neg p \vee q$$

$$(ii) p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$(iii) p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

Proof: $(p \rightarrow q) \equiv (\neg p \vee q)$ [due to (i)]

$$(p \rightarrow r) \equiv (\neg p \vee r)$$
 [due to (i)]

$$\text{LHS} = (p \rightarrow q) \wedge (p \rightarrow r)$$

$$\equiv (\neg p \vee q) \wedge (\neg p \vee r)$$

$$\equiv \neg p \vee (q \wedge r)$$
 [due to (ii)]

$$\equiv p \rightarrow (q \wedge r)$$
 [due to (i)] = RHS

Every tautology is satisfiable.

But converse is not true.

An algorithm involves checking all possible combination of truth values of its propositional variables

n variables $\Rightarrow 2^n$ possible assignments \Rightarrow each have to be evaluated.

$$(i) ((p \vee \neg q) \wedge (q \vee \neg r) \wedge (r \vee \neg p))$$

AND OF OR CONJUNCTION (and) of disjunction (or)

CONJUNCTIVE NORMAL FORM (CNF)

Is it satisfiable?

all variables true $(q \rightarrow p) \wedge (r \rightarrow q) \wedge (p \rightarrow r)$
then it evaluates to true. all variables false

$$(iv) \neg(p \vee q) \equiv \neg p \wedge \neg q$$

$$(v) \neg(p \wedge q) \equiv \neg p \vee \neg q$$

Using (i), (ii), (iii), (iv), (v)

prove that

$$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$$

A compound proposition is satisfiable if there is an assignment of truth values of propositional variables that evaluates to true.

SATISFIABILITY:

Input: A compound proposition p

Question: Is there an assignment of propositional variables that evaluates p to true?

Fact: Every compound proposition can be represented as conjunctive normal form.

"CNF boolean formula"

"boolean formula in CNF"

CLAUSE

② $(p \vee q \vee r) \wedge (\neg p \vee \neg q \vee \neg r) \wedge (q \vee \neg r) \wedge (r \vee \neg p) \wedge (p \vee \neg q) \equiv \phi$

If all p, q, r are true then $(\neg p \vee \neg q \vee \neg r)$ is false
 If all p, q, r are false, then $(p \vee q \vee r)$ is false.

Is ϕ satisfiable? $\begin{cases} \text{all } p, q, r \text{ true} \\ \text{all } p, q, r \text{ false} \end{cases}$

Hence, no matter of the assignment of values, ϕ is false.

Hence, ϕ is unsatisfiable \equiv contradiction

Section 1.3 Rosen's book

6 Groups of Tutorial
25-30

Next Class: 23rd August 9:30 AM.
(Monday time table)

Tutorial-1: 23rd Aug 1:30 pm