

BOOLEAN ALGEBRA

➤ Introduction:

- An algebra that deals with binary number system is called “Boolean Algebra”.
- It is very power in designing logic circuits used by the processor of computer system.
- The logic gates are the building blocks of all the circuit in a computer.
- Boolean algebra derives its name from the mathematician **George Boole** (1815-1864) who is considered the “**Father of symbolic logic**”.
- Boolean algebra deals with truth table TRUE and FALSE.
- It is also called as “**Switching Algebra**”.

➤ Binary Valued Quantities – Variable and Constants:

- A variable used in Boolean algebra or Boolean equation can have only one of two variables. The two values are FALSE (0) and TRUE (1)
- A Sentence which can be determined to be TRUE or FALSE are called logical statements or truth functions and the results TRUE or FALSE is called Truth values.
- The variables which can store the truth values are called logical variables or binary valued variables. These can store one of the two values 1 or 0.
- The decision which results into either YES (TRUE or 1) or NO (FALSE or 0) is called Binary decision.

➤ Truth Table:

- A truth table is a mathematical table used in logic to computer functional values of logical expressions.
- A truth table is a table whose columns are statements and whose rows are possible scenarios.
- Example: Consider the logical expression

Logical Statement: Meals = “Ram prefer rice and roti for the meal”

Y = A AND B (Logical Variables: Y, A, B, Logical Operator AND)

Ram Prefer Rice	Ram Prefer Roti	Meals
FALSE	FALSE	FALSE
FALSE	TRUE	FALSE
TRUE	FALSE	FALSE
TRUE	TRUE	TRUE

A	B	Y
0	0	0
0	1	0
1	0	0
1	1	0

- If result of any logical statement or expression is always TRUE or 1, it is called **Tautology** and if the result is always FALSE or 0, it is called **Fallacy**.

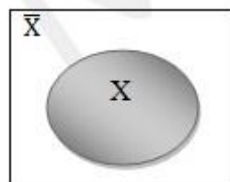
➤ Logical Operators:

- There are three logical operator, NOT, OR and AND.
- These operators are now used in computer construction known as switching circuits.

➤ NOT Operator:

- The Not operator is a unary operator. This operator operates on single variable.
- The operation performed by Not operator is called **complementation**.
- The symbol we use for it is bar.
- \bar{X} means complementation of X
- If $X=1, \bar{X}=0$ If $X=0, \bar{X}=1$
- The Truth table and the Venn diagram for the NOT operator is:

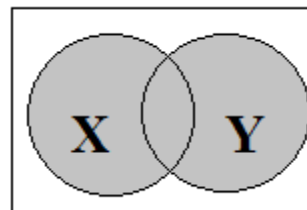
X	\bar{X}
1	0
0	1



➤ OR Operator:

- The OR operator is a binary operator. This operator operates on two variables.
- The operation performed by OR operator is called **logical addition**.
- The symbol we use for it is '+'. Example: $X + Y$ can be read as **X OR Y**
- The Truth table and the Venn diagram for the NOT operator is:

X	Y	$X + Y$
0	0	0
0	1	1
1	0	1
1	1	1

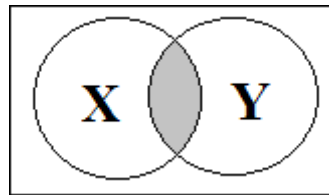


➤ AND Operator:

- The AND operator is a binary operator. This operator operates on two variables.
- The operation performed by AND operator is called **logical multiplication**.
- The symbol we use for it is '·'.
- Example: $X \cdot Y$ can be read as **X AND Y**

- The Truth table and the Venn diagram for the NOT operator is:

X	Y	$X \cdot Y$
0	0	0
0	1	0
1	0	0
1	1	1



➤ Evaluation of Boolean Expression using Truth Table:

- To create a truth table, follow the steps given below.*
- Step 1: Determine the number of variables, for n variables create a table with 2^n rows.
 - For two variables i.e. X, Y then truth table will need 2^2 or 4 rows.
 - For three variables i.e. X, Y, Z, then truth table will need 2^3 or 8 rows.
- Step 2: List the variables and every combination of 1 (TRUE) and 0 (FALSE) for the given variables
- Step 3: Create a new column for each term of the statement or argument.
- Step 4: If two statements have the same truth values, then they are equivalent.

➤ Example: Consider the following Boolean Expression $F = X + \bar{Y}$

- Step 1: This expression as two variables X and Y, then 2^2 or 4 rows.
- Step 2: List the variables and every combination of X and Y.
- Step 3: Create a new column \bar{Y} of the statement, and then fill the truth values of Y in that column.
- Step 4: The final column contain the values of $X + \bar{Y}$

X	Y	\bar{Y}	$X + \bar{Y}$
0	0	1	1
0	1	0	0
1	0	1	1
1	1	0	1

➤ Exercise Problems:

- Prepare a table of combination for the following Boolean algebra expressions.
 - $\bar{X}\bar{Y} + \bar{X}Y$
 - $XY\bar{Z} + \bar{X}\bar{Y}Z$
- Verify using truth table for the following Boolean algebra.
 - $X + XY = X$
 - $\overline{X + Y} = \bar{X} \cdot \bar{Y}$

➤ Boolean Postulates:

- The fundamental laws of Boolean algebra are called as the postulates of Boolean algebra.
- These postulates for Boolean algebra originate from the three basic logic functions AND, OR and NOT.
- **Properties of 0 and 1:**
 - I. If $X \neq 0$ then $X = 1$, and If $X \neq 1$ then $X = 0$
 - II. OR relation (Logical Addition)
 - a. $0 + 0 = 0$
 - b. $0 + 1 = 1$
 - c. $1 + 0 = 1$
 - d. $1 + 1 = 1$
 - III. AND relation (Logical Multiplication)
 - a. $0 \cdot 0 = 0$
 - b. $0 \cdot 1 = 0$
 - c. $1 \cdot 0 = 0$
 - d. $1 \cdot 1 = 1$
 - IV. Complement Rules
 - a. $\bar{0} = 1$
 - b. $\bar{1} = 0$

➤ Principle of Duality Theorem:

- This is very important principle used in Boolean algebra.
- Principle of Duality states that;
 - Changing each OR sign (+) to an AND sign (.)
 - Changing each AND sign (.) to an OR sign (+)
 - Replacing each 0 by 1 and each 1 by 0.
- The derived relation using duality principle is called dual of original expression.
- Example: Take postulate II, related to logical addition:
 - 1) $0 + 0 = 0$ 2) $0 + 1 = 1$ 3) $1 + 0 = 1$ 4) $1 + 1 = 1$
- 2. Now working according to above relations, + is changed to . and 0's replaced by 1's
 - a) $1 \cdot 1 = 1$ b) $1 \cdot 0 = 0$ c) $0 \cdot 1 = 0$ d) $0 \cdot 0 = 0$
- which are nothing but same as that of postulate III related to logical multiplication.
- So 1, 2, 3, 4, are the duals of a, b, c, d.
- Example: Find the duals for the following Boolean Expression

SI No	Boolean Expression	Duals
1	$X + 0 = X$	$X \cdot 1 = X$
2	$X + 1 = 1$	$X \cdot 0 = 0$
3	$X \cdot \bar{X} = 0$	$X + \bar{X} = 1$
4	$X \cdot (Y + Z)$	$X + (Y \cdot Z)$
5	$X + X \cdot Y = X + Y$	$X \cdot (X + Y) = X \cdot Y$

➤ **Boolean Theorems:**

- Boolean Theorem can be proved by substituting all possible values of the variable that are 0 and 1.
- This technique of proving theorem is called **Proof by perfect induction.**

Sl No	Theorem	Sl No	Theorem
Properties of 0 and 1		Associative Law	
1	$0 + X = X$	12	$X \cdot (Y \cdot Z) = (X \cdot Y) \cdot Z$
2	$1 + X = 1$	13	$(X + Y) \cdot Z = X + (Y \cdot Z)$
3	$0 \cdot X = 0$	Distributive Law	
4	$1 \cdot X = X$	14	$X \cdot (Y + Z) = X \cdot Y + X \cdot Z$
Idempotence Law		15	$X + Y \cdot Z = (X + Y) \cdot (X + Z)$
5	$X + X = X$	Absorption Law	
6	$X \cdot X = X$	16	$X + XY = X$
Complementary Law		17	$X(X + Y) = X$
7	$X + \bar{X} = 1$	18	$XY + \bar{X}Y = Y$
8	$X \cdot \bar{X} = 0$	19	$(X + Y)(X + \bar{Y}) = X$
Involution Law		20	$X + \bar{X}Y = X + Y$
9	$\bar{\bar{X}} = X$	21	$X(\bar{X} + Y) = XY$
Commutative Law			
10	$X + Y = Y + X$		
11	$X \cdot Y = Y \cdot X$		

➤ **Theorem 1: $0 + X = X$**

Proof: If $X = 0$ then LHS $= 0 + X$ $= 0 + 0$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= 0 + X$ $= 0 + 1$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>0</th><th>X</th><th>0+X</th></tr> <tr> <td>0</td><td>0</td><td>0</td></tr> <tr> <td>0</td><td>1</td><td>1</td></tr> </table>	0	X	0+X	0	0	0	0	1	1
0	X	0+X									
0	0	0									
0	1	1									

➤ **Theorem 2: $1 + X = 1$**

Proof: If $X = 0$ then LHS $= 1 + X$ $= 1 + 0$ $= 1$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= 1 + X$ $= 1 + 1$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>1</th><th>X</th><th>1+X</th></tr> <tr> <td>1</td><td>0</td><td>1</td></tr> <tr> <td>1</td><td>1</td><td>1</td></tr> </table>	1	X	1+X	1	0	1	1	1	1
1	X	1+X									
1	0	1									
1	1	1									

➤ **Theorem 3:** $0 \cdot X = 0$

Proof: If $X = 0$ then LHS $= 0 \cdot X$ $= 0 \cdot 0$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= 0 \cdot X$ $= 0 \cdot 1$ $= 0$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>0</th><th>X</th><th>0.X</th></tr> <tr> <td>0</td><td>0</td><td>0</td></tr> <tr> <td>0</td><td>1</td><td>0</td></tr> </table>	0	X	0.X	0	0	0	0	1	0
0	X	0.X									
0	0	0									
0	1	0									

➤ **Theorem 4:** $1 \cdot X = X$

Proof: If $X = 0$ then LHS $= 1 \cdot X$ $= 1 \cdot 0$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= 1 \cdot X$ $= 1 \cdot 1$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>1</th><th>X</th><th>0.X</th></tr> <tr> <td>1</td><td>0</td><td>0</td></tr> <tr> <td>1</td><td>1</td><td>1</td></tr> </table>	1	X	0.X	1	0	0	1	1	1
1	X	0.X									
1	0	0									
1	1	1									

➤ **Idempotence Law:** “This law states that when a variable is combines with itself using OR or AND operator, the output is the same variable”.

➤ **Theorem 5:** $X + X = X$

Proof: If $X = 0$ then LHS $= X + X$ $= 0 + 0$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= X + X$ $= 1 + 1$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>X</th><th>X</th><th>X+X</th></tr> <tr> <td>0</td><td>0</td><td>0</td></tr> <tr> <td>1</td><td>1</td><td>1</td></tr> </table>	X	X	X+X	0	0	0	1	1	1
X	X	X+X									
0	0	0									
1	1	1									

➤ **Theorem 6:** $X \cdot X = X$

Proof: If $X = 0$ then LHS $= X \cdot X$ $= 0 \cdot 0$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= X \cdot X$ $= 1 \cdot 1$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>X</th><th>X</th><th>X.X</th></tr> <tr> <td>0</td><td>0</td><td>0</td></tr> <tr> <td>1</td><td>1</td><td>1</td></tr> </table>	X	X	X.X	0	0	0	1	1	1
X	X	X.X									
0	0	0									
1	1	1									

➤ **Complementary Law:** “This law states that when a variable is And ed with its complement is equal to 0 and a variable is OR ed with its complement is equal to 1”.

➤ **Theorem 7:** $X + \bar{X} = 1$

Proof: If $X = 0$ then LHS $= X + \bar{X}$ $= 0 + 1$ $= 1$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= X + \bar{X}$ $= 1 + 0$ $= 1$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>X</th><th>\bar{X}</th><th>$X + \bar{X}$</th></tr> <tr> <td>0</td><td>1</td><td>1</td></tr> <tr> <td>1</td><td>0</td><td>1</td></tr> </table>	X	\bar{X}	$X + \bar{X}$	0	1	1	1	0	1
X	\bar{X}	$X + \bar{X}$									
0	1	1									
1	0	1									

➤ **Theorem 8:** $X \cdot \bar{X} = 0$

Proof: If $X = 0$ then LHS $= X \cdot \bar{X}$ $= 0 \cdot 1$ $= 0$ $= \text{RHS}$	Proof: If $X = 1$ then LHS $= X \cdot \bar{X}$ $= 1 \cdot 0$ $= 0$ $= \text{RHS}$	Using Truth Table <table border="1"> <tr> <th>X</th><th>\bar{X}</th><th>$X \cdot \bar{X}$</th></tr> <tr> <td>0</td><td>1</td><td>0</td></tr> <tr> <td>1</td><td>0</td><td>0</td></tr> </table>	X	\bar{X}	$X \cdot \bar{X}$	0	1	0	1	0	0
X	\bar{X}	$X \cdot \bar{X}$									
0	1	0									
1	0	0									

➤ **Involution Law:** “This law states that when a variable is inverted twice is equal to the original variable”.

➤ **Theorem 9:** $\bar{\bar{X}} = X$

Proof: If $X = 0$, then $\bar{X} = 1$ Take complement again, then $\bar{\bar{X}} = 0$ i.e. X If $X = 1$, then $\bar{X} = 0$ Take complement again, then $\bar{\bar{X}} = 1$ i.e. X	Using Truth Table <table><tr><th>X</th><th>\bar{X}</th><th>$\bar{\bar{X}}$</th></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>1</td></tr></table>	X	\bar{X}	$\bar{\bar{X}}$	0	1	0	1	0	1
X	\bar{X}	$\bar{\bar{X}}$								
0	1	0								
1	0	1								

➤ **Commutative Law:** “This law states that the order in which two variable are Or ed or AND ed make no difference”.

➤ **Theorem 10:** $X + Y = Y + X$

Proof: If $Y = 0$	Proof: If $Y = 1$	Using Truth Table																				
<div>then LHS</div> <div>$= X + Y$ $= X + 0$ $= X$</div> <div>RHS</div> <div>$= Y + X$ $= 0 + X$ $= X$</div> <div>Therefore LHS = RHS</div>	<div>then LHS</div> <div>$= X + Y$ $= X + 1$ $= 1$</div> <div>RHS</div> <div>$= Y + X$ $= 1 + X$ $= 1$</div> <div>Therefore LHS = RHS</div>	<table><tr><th>X</th><th>Y</th><th>X+Y</th><th>Y+X</th></tr><tr><td>0</td><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>1</td><td>1</td><td>1</td></tr></table>	X	Y	X+Y	Y+X	0	0	0	0	0	1	1	1	1	0	1	1	1	1	1	1
X	Y	X+Y	Y+X																			
0	0	0	0																			
0	1	1	1																			
1	0	1	1																			
1	1	1	1																			

➤ **Theorem 11:** $X \cdot Y = Y \cdot X$

Proof: If $Y = 0$ then LHS $= X \cdot Y$ $= X \cdot 0$ $= 0$ RHS $= Y \cdot X$ $= 0 \cdot X$ $= 0$ Therefore LHS = RHS	Proof: If $Y = 1$ then LHS $= X \cdot Y$ $= X \cdot 1$ $= X$ RHS $= Y \cdot X$ $= 1 \cdot X$ $= X$ Therefore LHS = RHS	Using Truth Table <table><tr><th>X</th><th>Y</th><th>X.Y</th><th>Y.X</th></tr><tr><td>0</td><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>1</td><td>1</td></tr></table>	X	Y	X.Y	Y.X	0	0	0	0	0	1	0	0	1	0	0	0	1	1	1	1
X	Y	X.Y	Y.X																			
0	0	0	0																			
0	1	0	0																			
1	0	0	0																			
1	1	1	1																			

➤ **Associative Law:** “This law allows the removal of brackets from an expression and regrouping of the variables”.

➤ **Theorem 12:** $X \cdot (Y \cdot Z) = (X \cdot Y) \cdot Z$

Proof: If $Y = 0$	Proof: If $Y = 1$	Using Truth Table						
LHS = $X.(Y.Z)$	LHS = $X.(Y.Z)$	X	Y	Z	XY	YZ	X.(Y.Z)	(X.Y).Z
$= X.(0.Z)$	$= X.(1.Z)$	0	0	0	0	0	0	0
$= X.0$	$= X.Z$	0	0	1	0	0	0	0
$= 0$	$= XZ$	0	1	0	0	0	0	0
RHS = $(X.Y).Z$	RHS = $(X.Y).Z$	0	1	1	0	1	0	0
$= (X.0).Z$	$= (X.1).Z$	1	0	0	0	0	0	0
$= 0.Z$	$= X.Z$	1	0	1	0	0	0	0
$= 0$	$= XZ$	1	1	0	1	0	0	0
Therefore LHS = RHS	Therefore LHS = RHS	1	1	1	1	1	1	1

➤ **Theorem 13:** $X + (Y + Z) = (X + Y) + Z$

<p>Proof: If $Y = 0$</p> <p>LHS = $X+(Y+Z)$</p> <p>$= X+(0+Z)$</p> <p>$= X+Z$</p> <p>RHS = $(X+Y)+Z$</p> <p>$= (X+0)+Z$</p> <p>$= X+Z$</p> <p>Therefore LHS = RHS</p>	<p>Proof: If $Y = 1$</p> <p>LHS = $X+(Y+Z)$</p> <p>$= X+(1+Z)$</p> <p>$= X+1$</p> <p>$= 1$</p> <p>RHS = $(X+Y)+Z$</p> <p>$= (X+1).Z$</p> <p>$= 1+Z$</p> <p>$= 1$</p> <p>Therefore LHS = RHS</p>	<p>Using Truth Table</p> <table><tr><th>X</th><th>Y</th><th>Z</th><th>X+Y</th><th>Y+Z</th><th>X+(Y+Z)</th><th>(X+Y)+Z</th></tr><tr><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>0</td><td>1</td><td>0</td><td>1</td><td>1</td><td>1</td></tr><tr><td>0</td><td>1</td><td>0</td><td>1</td><td>1</td><td>1</td><td>1</td></tr><tr><td>0</td><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>0</td><td>1</td><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>1</td><td>0</td><td>1</td><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td><td>1</td></tr></table>	X	Y	Z	X+Y	Y+Z	X+(Y+Z)	(X+Y)+Z	0	0	0	0	0	0	0	0	0	1	0	1	1	1	0	1	0	1	1	1	1	0	1	1	1	1	1	1	1	0	0	1	0	1	1	1	0	1	1	1	1	1	1	1	0	1	1	1	1	1	1	1	1	1	1	1
X	Y	Z	X+Y	Y+Z	X+(Y+Z)	(X+Y)+Z																																																											
0	0	0	0	0	0	0																																																											
0	0	1	0	1	1	1																																																											
0	1	0	1	1	1	1																																																											
0	1	1	1	1	1	1																																																											
1	0	0	1	0	1	1																																																											
1	0	1	1	1	1	1																																																											
1	1	0	1	1	1	1																																																											
1	1	1	1	1	1	1																																																											

➤ **Distributive Law:** “This law allows the multiplying or factoring out an expression”.

➤ **Theorem 14:** $X.(Y+Z) = XY + XZ$

Proof: If $X = 0$ $LHS = X.(Y+Z)$ $= 0.(Y+Z)$ $= 0$ $RHS = XY + XZ$ $= 0.Y + 0.Z$ $= 0$ Therefore $LHS = RHS$	Proof: If $X = 1$ $LHS = X.(Y+Z)$ $= 1.(Y+Z)$ $= Y+Z$ $RHS = XY + XZ$ $= 1.Y + 1.Z$ $= Y+Z$ Therefore $LHS = RHS$
--	--

➤ **Theorem 15:** $(X + Y) (X + Z) = X + YZ$

$$\begin{aligned}
 LHS: (X + Y) (X + Z) &= XX + XZ + XY + YZ \\
 &= X + XZ + XY + YZ \\
 &= X(1 + Z) + XY + YZ \\
 &= X + XY + YZ \\
 &= X(1 + Y) + YZ \\
 &= X + YZ \\
 &= RHS
 \end{aligned}$$

➤ **Absorption Law:** “This law enables a reduction of complicated expression to a simpler one by absorbing common terms”.

16) $X+XY = X$ $LHS = X + XY$ $= X (1 + Y)$ $= X$ $= RHS$	17) $X (X+Y) = X$ $LHS = X (X+Y)$ $= XX + XY$ $= X + XY$ $= X (1+Y)$ $= X$	18) $XY + X\bar{Y} = X$ $LHS = XY + X\bar{Y}$ $= X(Y + \bar{Y})$ $= X.1$ $= X$ $= RHS$
19) $(X+Y) (X + \bar{Y}) = X$ $LHS = (X+Y) (X + \bar{Y})$ $= XX + X\bar{Y} + XY + Y\bar{Y}$ $= X + X\bar{Y} + XY + 0$ $= X (1 + \bar{Y} + Y)$ $= X.1$ $= X$	20) $X + \bar{X}Y = X+Y$ $LHS = X + \bar{X}Y$ $= (X + \bar{X}) (X+Y)$ $= 1. (X+Y)$ $= X+Y$ $= RHS$	21) $X (\bar{X}+Y) = XY$ $LHS = X (\bar{X}+Y)$ $= X.\bar{X} + X.Y$ $= 0 + XY$ $= XY$ $= RHS$

➤ **DeMorgan's Theorem:**

• **DeMorgan's First Theorem:**

- **Statement:** “When the OR sum of two variables is inverted, this is same as inverting each variable individually and then AND ing these inverted variables”
- This can be written as $\overline{X + Y} = \bar{X} \cdot \bar{Y}$
- We can prove the DeMorgan's First theorem by using Truth Table is

X	Y	\bar{X}	\bar{Y}	X+Y	\overline{XY}	$\bar{X}\bar{Y}$
0	0	1	1	0	1	1
0	1	1	0	1	0	0
1	0	0	1	1	0	0
1	1	0	0	1	0	0

- Compare the column $\overline{X + Y}$ and $\bar{X} \cdot \bar{Y}$. Both of these are identical. Hence the DeMorgan's first theorem is proved.

• **DeMorgan's Second Theorem:**

- **Statement:** “When the AND product of two variables is inverted, this is same as inverting each variable individually and then OR ing these inverted variables”
- This can be written as $\overline{XY} = \bar{X} + \bar{Y}$
- We can prove the DeMorgan's Second theorem by using Truth Table is:

X	Y	\bar{X}	\bar{Y}	X.Y	\overline{XY}	$\bar{X} + \bar{Y}$
0	0	1	1	0	1	1
0	1	1	0	0	1	1
1	0	0	1	0	1	1
1	1	0	0	1	0	0

- Compare the column \overline{XY} and $\bar{X} + \bar{Y}$. Both of these are identical. Hence the DeMorgan's Second theorem is proved.

• **Application of DeMorgan's Theorem:**

- It is used in simplification of Boolean expression.
- DeMorgan's law commonly apply to text searching using Boolean operators AND, OR and NOT.
- It is useful in the implementation of the basic gates operations with alternative gates.

