

Binary Search Tree - II

Querying a BST : finding min, max, predecessor, successor, or search

Search : returns a pointer to a node with key k
if one exists in the subtree, else returns NIL.

Tree-Search (x, k)

root of a subtree in T
the value to search for

if $x == \text{NIL}$ or $k == x.\text{key}$
return x

if $k < x.\text{key}$
return Tree-Search ($x.\text{left}, k$)

else
return Tree-Search ($x.\text{right}, k$)

Running time of Tree-Search is $O(h)$
where h is the height of the tree

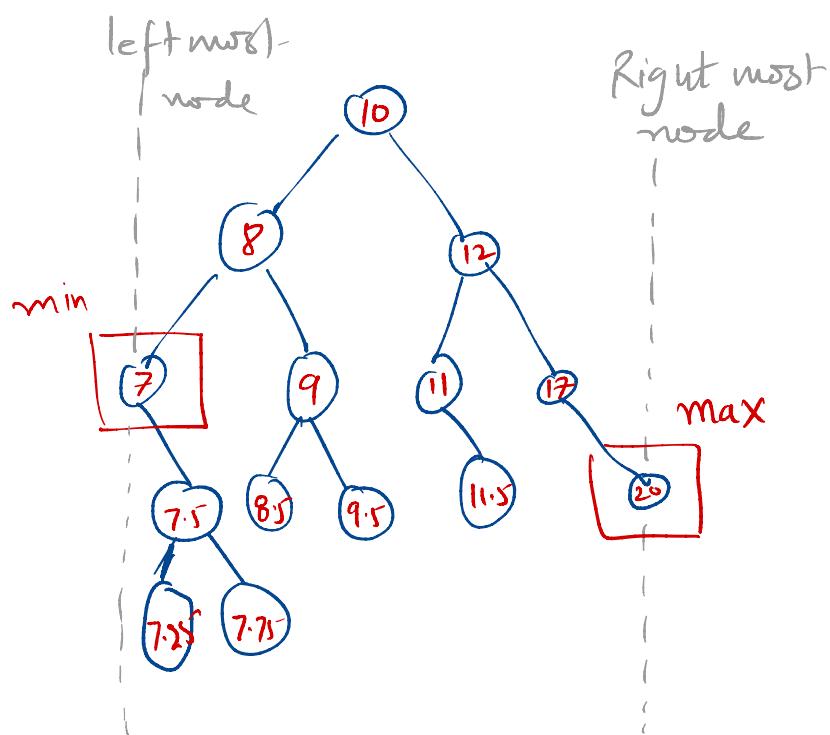
Iterative -Tree -Search (x, k)

while $x \neq \text{NIL}$ and $k \neq x.\text{key}$

|
if $k < x.\text{key}$
 $x = x.\text{left}$ $\rightarrow O(h)$
else
 $x = x.\text{right}$
return x

— Minimum / Maximum

→ Returns a pointer to min/max element in a subtree rooted at node x (assumed to be not NIL)



Tree - Minimum (x)

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while  $x.\text{left} \neq \text{NIL}$ 
|  $x = x.\text{left}$ 
return  $x$ 

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Tree - Maximum (x)

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while  $x.\text{right} \neq \text{NIL}$ 
|  $x = x.\text{right}$ 
return  $x$ 

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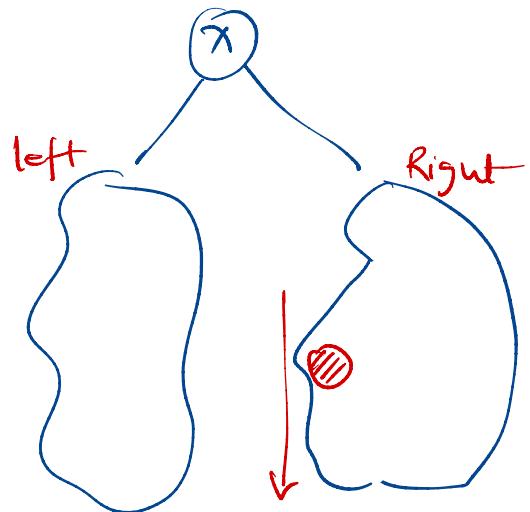
- Successor and Predecessor

→ Find successor of a given node in the sorted order determined by the inorder walk.



Define successor as the next node visited in an inorder tree walk.

Two cases : (I) If the right subtree of x is non empty



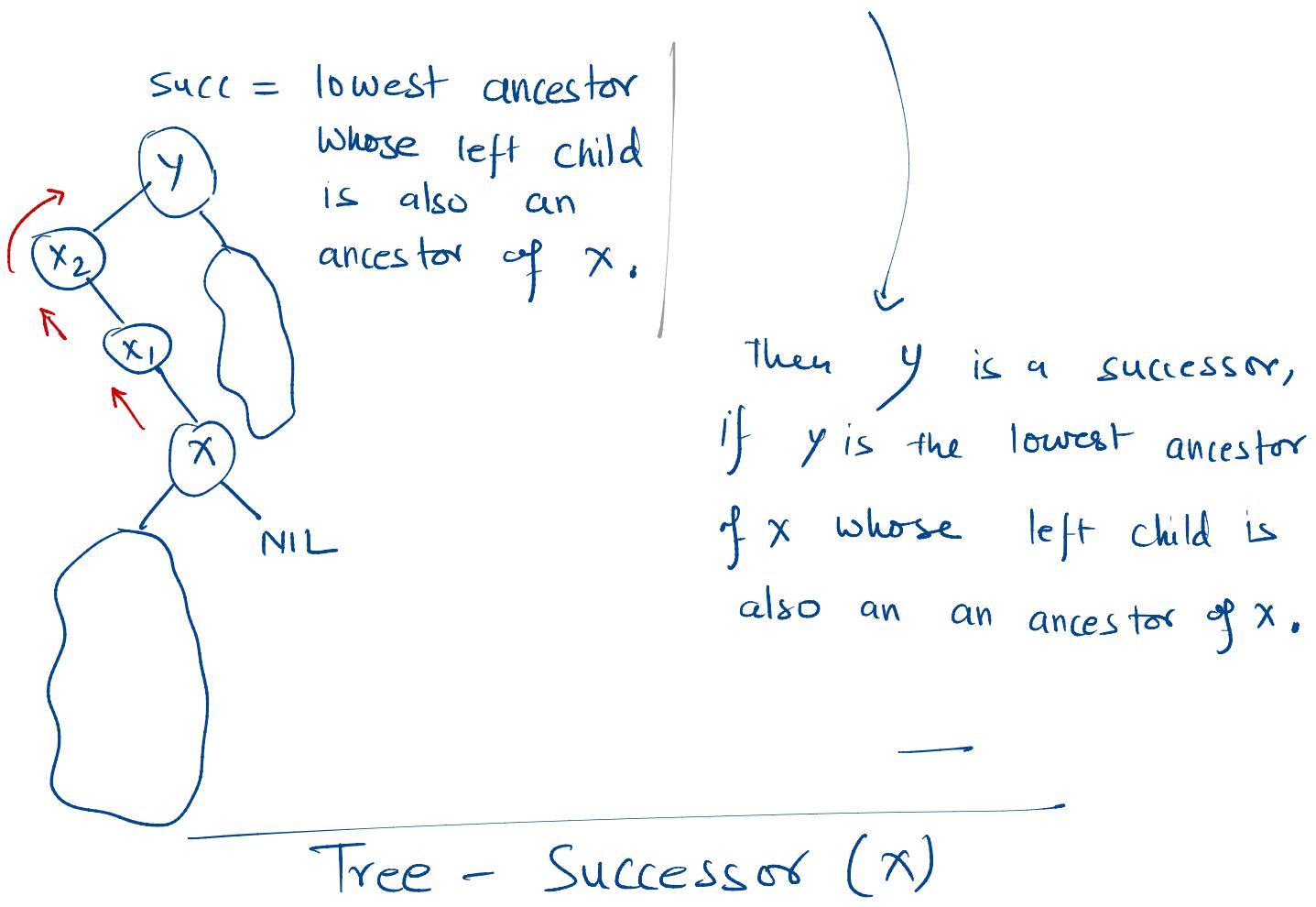
→ successor = leftmost node in x 's right subtree

↓
call

Tree-minimum ($x.\text{right}$)

Succ = leftmost node

(II) If the right subtree of x is empty



if $x.\text{right} \neq \text{NIL}$

return Tree-Minimum($x.\text{right}$)

else

$$y = x.P$$

While $y \neq \text{NIL}$ and $x == y.\text{right}$

$$x = y$$

$$y = y.P$$

return y
