



Today's agenda

↳ $\text{Pow}(a, n)$

↳ TC & SC of Recursion



AlgoPrep



Q) Given a and N , calculate a^N .

Ex:

a n

2 3

→

8

3 5

→

243

```
int Pow (int a, int n) {  
    if (n == 1) { return a; }  
}
```

Faith: Given a and n , calculate and return a^n .

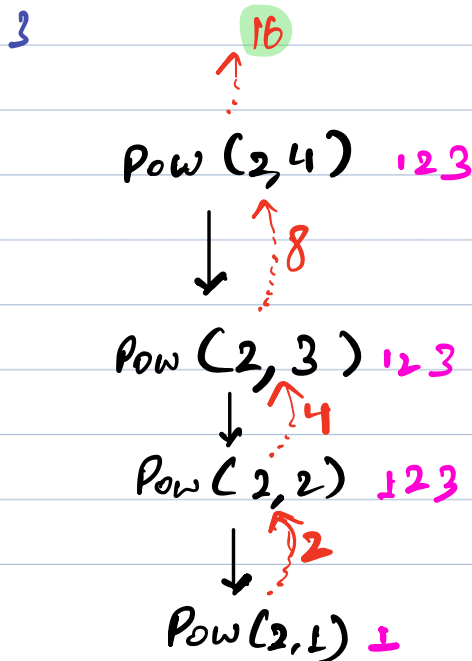
```
int temp = Pow(a, n-1);  
return temp * a;
```

Main logic:

$a^n \rightarrow temp * a$
 $a^{n-1} \rightarrow temp$

Base case:

if (n == 1) { return a; }



T.C. of 1 function: $O(1) \rightarrow n$
No. of function: $N \rightarrow y$

T.C. of recursion: $n * y = O(1) * N = O(N)$

S.C. of recursion: $O(1) * N = O(N)$



// do something better than $O(n)$.

$$\rightarrow a^n = a^{n-1} * a$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} \quad \text{if } n \text{ is even}$$

$$\text{ex: } 2^8 = 2^4 * 2^4$$

$$2^9 = 2^4 * 2^4 * 2$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} * a \quad \text{if } n \text{ is odd}$$

$$2^9 = 2^4 * 2^4 * 2$$

```
int Pow (int a, int n) {
    if (n == 1) { return a; }
```

Faith: Given a and n , calculate and return a^n .

```
    int temp = Pow(a, n/2);
    if (n % 2 == 0) { return temp * temp; }
```

```
    else { return temp * temp * a; }
```

Main logic:

a^n

\rightarrow n is even $= \text{temp} * \text{temp}$

\rightarrow n is odd $= \text{temp} * \text{temp} * a$

$a^{n/2} \rightarrow \text{temp}$

3

Base case:

```
if (n == 1) { return a; }
```



```
int Pow (int a, int n) {
```

```
1 if (n == 1) { return a; }
```

```
2 int temp = Pow (a, n/2);
```

```
3 { if (n%2 == 0) { return temp * temp; }  
  else { return temp * temp * a; } }
```

```
}
```

$a = 2$

$n = 37$

$Pow(2, 37)$ 1 2 3

↓ 262144

$Pow(2, 18)$ 1 2 3

↓ 512

$Pow(2, 9)$ 1 2 3

↓ 16

$Pow(2, 4)$ 1 2 3

↓ 4

$Pow(2, 2)$ 1 2 3

↓ 2

$Pow(2, 1)$ 1

overall T.C:

T.C of 1 function: $O(1)$

No. of function: $O(\log N)$

$\approx O(\log N)$

overall s.c: $O(1) * \log N \approx O(\log N)$

if n is negative

$2^{-37} \Rightarrow 2^{37} \rightarrow$ recursive soln \rightarrow ans

↓

return $\frac{1}{ans}$

$\rightarrow a^n = a^{n/2} * a^{n/2} \rightarrow$ most optimal here.

$a^n = a^{n/3} * a^{n/3} * a^{n/3}$

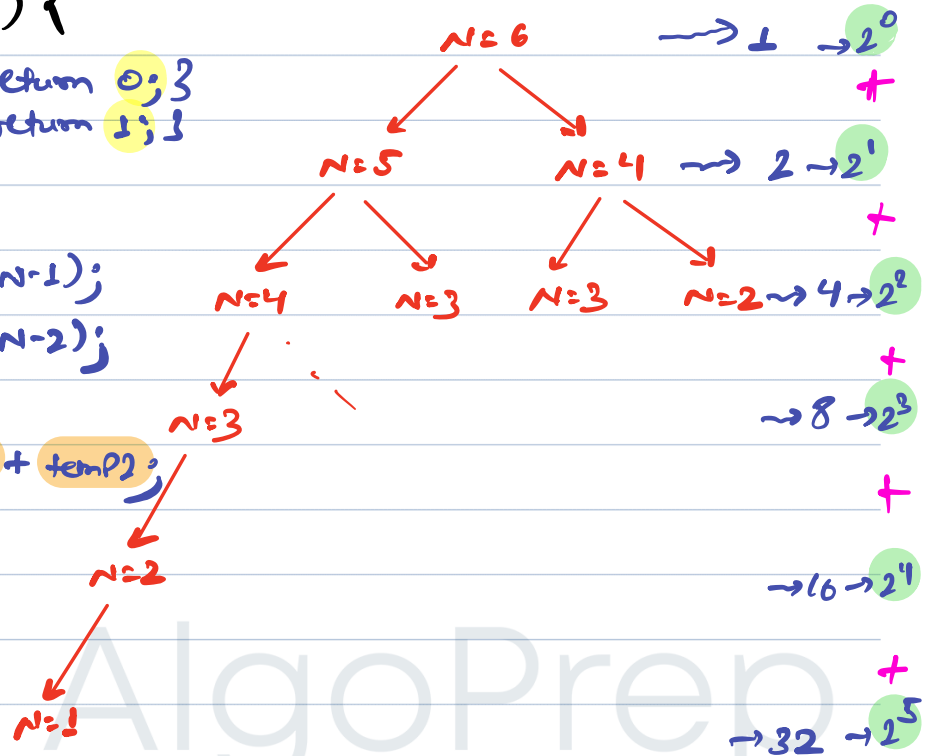
$a^n = a^{n/4} * a^{n/4} * a^{n/4} * a^{n/4}$

$a^n = a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5}$

$a^n = a^{n/n} * a^{n/n} * a^{n/n} * a^{n/n} \dots a^{n/2} \rightarrow$ ideal $\rightarrow O(n)$



```
int fib(int n) {  
1 if (n == 0) { return 0; }  
  if (n == 1) { return 1; }  
2 int temp1 = fib(n-1);  
3 int temp2 = fib(n-2);  
4 return temp1 + temp2;  
}
```



$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1}$$

$$\text{Sum of G.P.} = a * \frac{(2^n - 1)}{2 - 1}$$

$$= 1 * \frac{2^n - 1}{2 - 1} = 2^n - 1$$

$$\text{Overall T.C: } O(1) * 2^n \approx O(2^n)$$



*

No. of Calls = n

Count of levels = n

no. of functions = n^n

Break till 9:49 PM.



AlgoPrep



Q) Given an array, check if it is Palindrome or not?
↳ recursion

Ex: MALAYALAM → true

Ex: a a b b a → false

//idea 1

→ MALAYALAM

→ MALAYALAM

→ true

a a b b a

a b b a a

→ false

//idea 2

0 n-1
M A L A Y A L A M

T.C: $O(1) * N/2$

$\approx O(N)$

S.C: $O(1) * N/2 \approx O(N)$

boolean isPalindrome(char[] ch, int s, int e) { Faith: Check and

if (s == e) { return true; }

return whether

if (s > e) { return true; }

ch array is Palindrome

betⁿ 0^s to n-1^e.

if (ch[s] == ch[e]) {

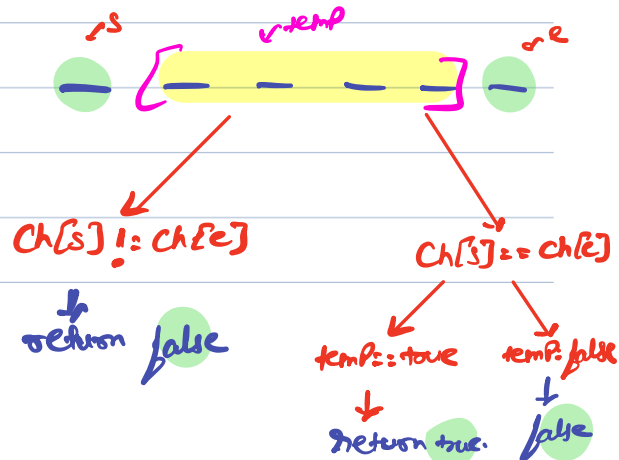
boolean temp = isPalindrome(ch, s+1, e-1); main logic:

return temp;

}

else { return false; }

}



base case:



boolean isPalindrome(char[] ch, int s, int e)

ch: M A L A Y A L A M

```
1 if (s == e) { return true; }
  if (s > e) { return true; }
```

```
2 if (ch[s] == ch[e]) {
  boolean temp = isPalindrome(ch, s+1, e-1);
  return temp;
}
```

```
3 else { return false; }
3
```

isPalindrome(0, 8) → true

isPalindrome(1, 7) → true

isPalindrome(2, 6) → true

isPalindrome(3, 5) → true

isPalindrome(4, 4) → true

ch: M A L A A L A M

isPalin(0, 7) → true

isPalin(1, 6) → true

isPalin(2, 5) → true

isPalin(3, 4) → true

isPalin(4, 3) → true



ch: ⁰M ¹A ²L ³A ⁴A ⁵L ⁶B ⁷M

isPalin(0, 7) ^{false}

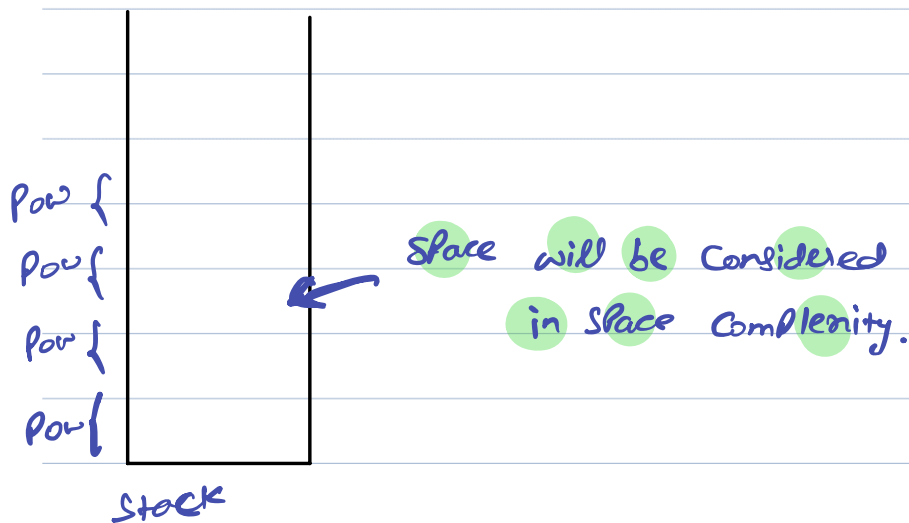
↓ ^{false}
isPalin(1, 6)

write iterative code → for (int i=0; i<n; i++) {

}

write recursive code: → within 1 function space complexity is like iteration.

↳ you create multiple instance of same function.





S.C: (space used in 1 function) & (max no. of function
in stack at any point
of time.)



AlgoPrep