

Midterm

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1) Base Case:  $n=1 \quad l=2^0$

Induction step: Assume theorem to be true  
for all positive integers up to  $n$ .

Now for  $n+1$

$K = \text{largest power of } 2$

$K \leq n+1$

$K = 2^{a_1} + 2^{a_2} + 2^{a_3} + 2^{a_4} + 2^{a_x}$  where

$a_1 > a_2 > a_3 > a_4 > a_x$

$K = n+1 \checkmark$

If  $K < n+1$  we add  $2^b$  to  $K$

$b < a_1$

$b = \max(a_x - 1 : a_x - 1 > a_1 - 1) = b < a_1$

$n+1 = K + 2^b + (n+1-K-2^b)$

hence  $b < a_1$ , power of  $2$ ,  $2^b$  is distinct  
from it being used in representation of  $K$

$n+1 - K - 2^b$

2)  $S \rightarrow A|B|AB$   
 $A \rightarrow Sa|aa$   
 $B \rightarrow Sb|bb$

$S$  = starting symbol

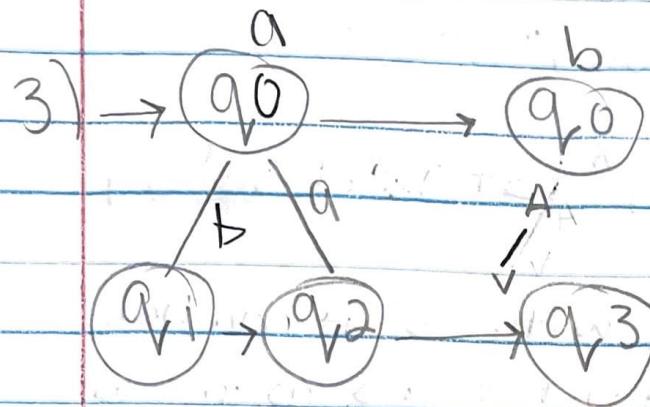
$A = "aa"$

$B = "bb"$

$S$  can generate "aa" for substring (A)  
or "bb" for substring (B)  
or both AB.

A can generate "a" followed by A (Sa)  
or produce "aa" (aa).

B can generate "b" followed by B (Sb)  
or produce "bb" (bb).



initial state  $q_0$

input a  $q_0 \rightarrow q_1$     input b  $q_0 \rightarrow q_0$   
 input a  $q_1 \rightarrow q_2$     input b  $q_1 \rightarrow q_0$   
 input a  $q_2 \rightarrow q_1$     input b  $q_2 \rightarrow q_3$   
 input a  $q_3 \rightarrow q_0$     input b  $q_3 \rightarrow q_0$

$$4) \cancel{(a^+ bb + b^+ a^* bb)} * a^* b^*$$

$$(b + aa^* bb)^* (\epsilon + aa^* + aa^* b)$$

If string starts with "b" it can have  
any amount of a's as long as  
it ends with "b".

If string starts with "a" it can have  
any number of a's but not end with "a"

$\epsilon$  matches the string b empty which  
also does not have aba in it.