

Assignment 1

- 1) De Morgan's Law given
Base case:
 $n=1$ satisfy condition for base case

$$\overline{U_p = 1, 2, 3, \dots, S_p} = \overline{S_1 \cup (U_p = 2, 3, \dots, S_p)}$$

$$\overline{S_1 \cup (U_p = 2, 3, \dots, S_p)} = \overline{S_1} \cap \overline{(U_p = 2, 3, \dots, S_p)} \quad (\because \text{induction hypothesis})$$

$$\overline{(U_p = 2, 3, \dots, S_p)} = \overline{N_p = 2, 3, \dots, S_p}$$

$$\begin{aligned} \overline{U_p = 1, 2, 3, \dots, S_p} &= \overline{S_1} \cap \overline{(N_p = 2, 3, \dots, S_p)} \\ &= \overline{N_p = 1, 2, 3, \dots, S_p} \quad \checkmark \end{aligned}$$

$$2) (S_1 \cup S_2) \cap (S_1 \cap S_2)' = S_2$$

$$S \cap T = S \cap T'$$

$$(S_1 \cup S_2) \cap (S_1 \cap S_2)'$$

$$(S_1 \cup S_2) \cap (S_1' \cup S_2')$$

$$(S_1 \cap S_1') \cup (S_1 \cap S_2')$$

$$\cup (S_2 \cap S_1') \cup (S_2 \cap S_2')$$

$$((S_1 \cap S_1') \cup (S_1 \cap S_2')) \cup ((S_2 \cap S_1') \cup (S_2 \cap S_2'))$$

$$(\emptyset \cup (S_1 \cap S_2')) \cup ((S_2 \cap S_1') \cup S_2)$$

$$(S_1 \cap S_2) \cup ((S_2 \cap S_1') \cup S_2)$$

$$(S_1 \cap S_2) \cup S_2$$

$$= S_2$$

\Leftarrow using distributive law

3) $S \rightarrow aA$ This gives strings $\{adb, aabab, abababab, \dots\}$
 $A \rightarrow bS$
 $S \rightarrow \lambda$ This grammar gives combinations of "aabb"

$$L(G) = (a^2b)^n \text{ or } (ab)^n : n \geq 0$$

4) $L = \{a^n \mid n \geq 0\}$

Gives string $A \rightarrow Ba \rightarrow Aaa \rightarrow Ba \rightarrow Aaaa \rightarrow aaaaa$

$$L(G) = \{a^\infty\}$$

This grammar gives infinite amount of "a".

$$5) a) S \rightarrow aA | bS$$

$$A \rightarrow aB | bA$$

$$B \rightarrow aB | bB | \lambda$$

The only way for a string to end is if it gets to B.

It can only do so if it has 2 "a"s.

$$b) S \rightarrow aA | bS | \lambda$$

$$A \rightarrow aB | bA | \lambda$$

$$B \rightarrow aC | bB | \lambda$$

$$C \rightarrow bC | \lambda$$

Once a string gets to 3 "a"s

it can only have

b's. It can

end anytime before that happens.