

### Iteration Method

It means to expand the recurrence and express it as a summation of terms of  $n$  and initial condition.

1.

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + 1 & n > 0 \end{cases}$$

$$T(n-2) = T(n-3) + 1$$

$$\begin{aligned} T(n) &= T(n-1) + 1 \dots\dots\dots 1^{\text{st}} \text{ time} \\ &= T(n-2) + 1 + 1 \dots\dots\dots 2^{\text{nd}} \text{ time} \\ &= T(n-3) + 1 + 1 + 1 \dots\dots\dots 3^{\text{rd}} \text{ time} \end{aligned}$$

$$\begin{aligned} &\downarrow \text{ k times} \\ &= T(n-k) + k \end{aligned}$$

Assume  $n-k = 0$

$$\begin{aligned} k &= n \\ &= T(n-n) + n \\ &= T(0) + n \end{aligned}$$

$$T(n) = 1 + n \dots \text{exact answer}$$

$$T(n) = O(n)$$

2.

$$T(n) = \begin{cases} 1, n=0 \\ T(n-1) + n, n>1 \end{cases}$$

$$T(n) = T(n-2) + n-1 + \dots + n^2$$

$$\begin{aligned} T(n) &= T(n-1) + n \dots\dots\dots 1^{\text{ST}} \text{ TIME} \\ &= T(n-2) + n-1 + n \dots\dots\dots 2^{\text{ND}} \text{ TIME} \\ &= T(n-3) + n-2 + n-1 + n \dots\dots\dots 3^{\text{RD}} \text{ TIME} \end{aligned}$$



$$\begin{aligned} &= T(n-k) + n-(k-1) + n-(k-2) + \dots + n-1 + n \\ &\text{assume } n-k = 0 \text{ so } k = n \text{ and } T(0) = 1 \\ &= 1 + 1 + 2 + 3 + \dots + n \\ &= 1 + n(n+1)/2 \\ &= 1 + n^2/2 + n/2 \dots\dots \text{exact answer} \end{aligned}$$

we take the dominating term which is  $n^2 \cdot 1/2$  therefor  $1/2 = \text{big O}$   
 $T(n) = O(n^2)$

$$3. \quad T(n) = \begin{cases} 1, n=0 \\ T(n-1) + \log n, n>0 \end{cases}$$

$$\begin{aligned} T(n) &= T(n-1) + \log n \dots\dots\dots 1^{\text{st}} \text{ time} \\ &= T(n-2) + \log(n-1) + \log n \dots\dots\dots 2^{\text{nd}} \text{ time} \\ &= T(n-3) + \log(n-2) + \log(n-1) + \log n \dots\dots\dots 3^{\text{rd}} \text{ time} \end{aligned}$$



$$\begin{aligned}
&= T(n-k) + \log(n-(k-1)) + \log(n-(k-2)) + \dots + \log(n-1) + \log n \\
&\text{Assume } n-k = 0 \text{ so } k = n \text{ and } T(0) = 1 \\
&= 1 + \log 1 + \log 2 + \dots + \log n \\
&= 1 + \log(1 * 2 * 3 * \dots * n) \\
&= 1 + \log n! \\
&= 1 + \log n^n \\
&= 1 + n \log n \dots \text{exact answer} \\
&= O(n \log n)
\end{aligned}$$

4.

$$T(n) = \begin{cases} 1, & n=0 \\ T(n-2) + n^2, & n>0 \end{cases}$$

$$T(n-2) = T(n-4) + (n-2)^2$$

$$T(n) = T(n-2 \cdot 1) + n^2 \dots \dots \dots 1^{\text{st}} \text{ time}$$

$$= T(n-4 \cdot 2) + (n-2 \cdot 2)^2 + n^2 \dots \dots \dots 2^{\text{nd}} \text{ time}$$

$$= T(n-6 \cdot 3) + (n-4 \cdot 2)^2 + (n-2 \cdot 1)^2 + n^2 \dots \dots \dots 3^{\text{rd}} \text{ time}$$

time

$$= T(n-8 \cdot 4) + (n-6 \cdot 3)^2 + (n-4 \cdot 2)^2 + (n-2 \cdot 1)^2 + n^2$$



$$= T(n-2k) + (n-2(k-1))^2 + (n-2(k-2))^2 + \dots + (n-0)^2$$

Assume  $n-2k=0$

$$n=2k$$

$$k=n/2$$

$$= T(0) + 2^2 + 4^2 + \dots + n^2$$

$$= 1 + (2 \cdot 1)^2 + (2 \cdot 2)^2 + (2 \cdot 3)^2 + \dots + (2 \cdot n/2)^2$$

$$\begin{aligned}
&= 1 + 2^2 (1^2 + 2^2 + 3^2 + \dots + (n/2)^2) \\
&= 1 + 2^2 [n/2 (n/2 + 1) (2n/2 + 1)]/6 \\
&= O(n^3)
\end{aligned}$$