

Syntax Direction Definitions: (SDD) \Rightarrow implemented in conjunction with LR Parser.

\rightarrow CFG with attributes and rules

\rightarrow Attributes - assoc. with grammar symbols

\rightarrow Rules - assoc. with productions

\rightarrow 2 types of attributes:

* Inherited attribute - for a non-terminal B at a parse-tree node N is defined by a semantic rule associated with the production at the parent of N.

* Synthesized attribute - for a non-terminal A at a parse-tree node N is defined by a semantic rule associated with the production at N.

\rightarrow Synthesized attributes at node N can be defined in terms of inherited attribute values at node N.

\rightarrow Terminals can have ~~inherited~~ ^{synthesized} attributes and ~~terminals~~ ~~cannot~~ have inherited attributes

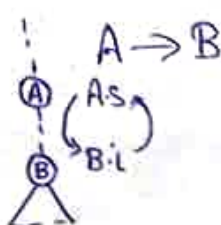
\rightarrow SDD with only synthesized attributes - S-attributed

\rightarrow SDD without side effects - attribute grammar.

Annotated PT:

\rightarrow A parse Tree showing values of its attributes is called annotated Parse Tree.

\rightarrow Circular dependency:



$$A.s = B.i$$

$$B.i = A.s + 1$$

Ex:

Production

Semantic Rules

1. $T \rightarrow FT'$

$$T.inh = F.val$$

$$T.val = T'.syn$$

2. $T' \rightarrow *FT'$

$$T'.inh = T.inh \times F.val$$

$$T'.syn = T'.syn$$

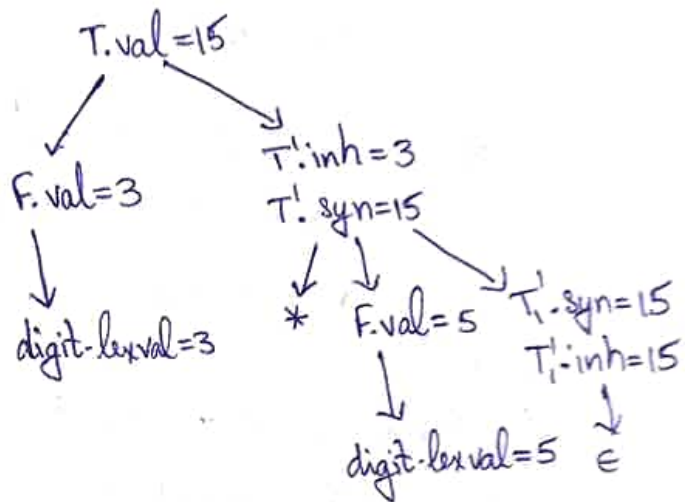
3. $T' \rightarrow \epsilon$

$$T'.syn = T.inh$$

4. $F \rightarrow \text{digit}$

$$F.val = \text{digit.lexval}$$

Annotated PT for $3*5$ is:



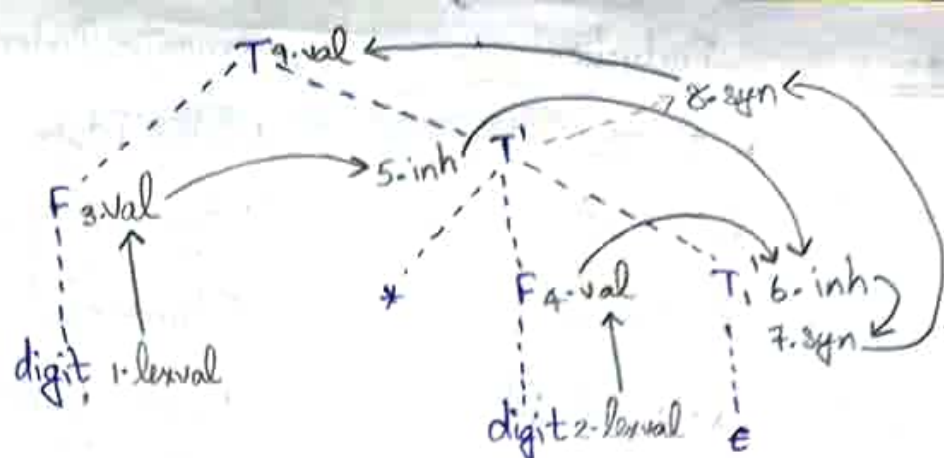
Evaluation order for SDD:

(i) Dependency Graph:

→ Depicts flow of information among attribute instances in a particular parse tree.

→ Edge from one attr. inst. to another means that the value of first needs to be computed before second.

→ Each attribute assoc. with a node X in Parse Tree is a node in Dep. Graph.



(ii) **Linear ordering:** Topological sort (without cycles)
 {Here, 135246789}

(iii) **S-Attributed Definitions:**

- No cycles in dep. graph.
- Evaluate in postorder sequence.

(iv) **L-Attributed Definitions:**

→ Each attribute must be either

1. Synthesized, or
2. Inherited with rules:

$A \rightarrow X_1 X_2 \dots X_n$, inherited attr. X_i - a uses

(a) inh. attr. assoc. with A

(b) inh./syn. attr. assoc. with $X_1 X_2 \dots X_{i-1}$

(c) inh./syn. attr. assoc. with X_i itself

such that there are no cycles in dep. graph.

Semantic Rules with controlled side-effects:

→ Attr. grammar - no side effects, allow any eval. order consistent with dep. graph.

→ Translation scheme - left-to-right eval., allow semantic actions to contain any pgm-fragment

Ex:

Production

Semantic Rules

1. $D \rightarrow TL$

$L.inh = T.type$

2. $T \rightarrow int$

$T.type = integer$

3. $T \rightarrow float$

$T.type = float$

4. $L \rightarrow L_1.id$

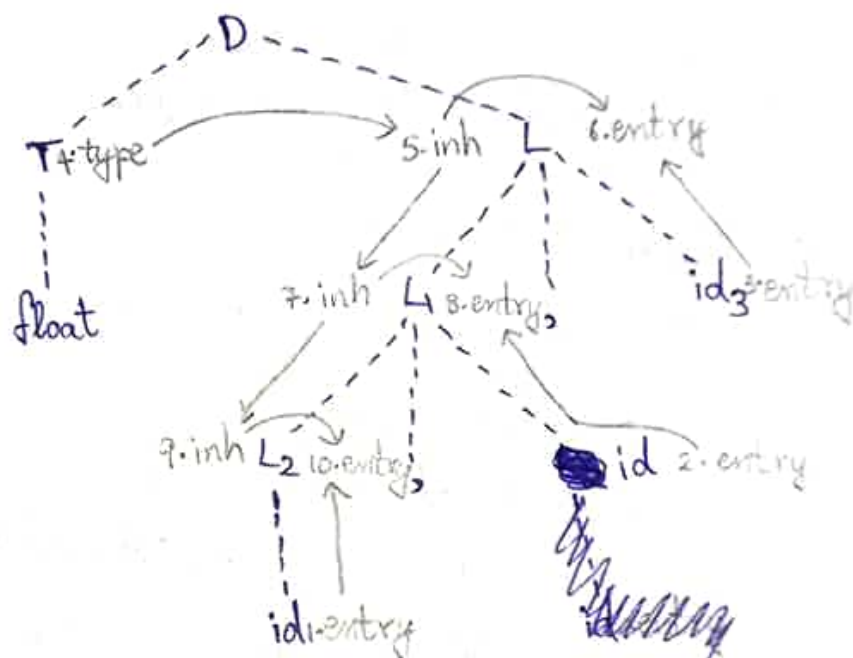
$L_1.inh = L.inh$

$addType(id.entry, L.inh)$

5. $L \rightarrow id$

$addType(id.entry, L.inh)$

(Dep-graph) APT for float id₁, id₂, id₃



Applications of Syntax Directed Translations:

→ Type checking, intermediate code generation

→ Const. of syntax tree. (Refer book Pg: 318 to 320).

Structure of a type:

→ Grammar for basic type/array type:

Production

Semantic Rules

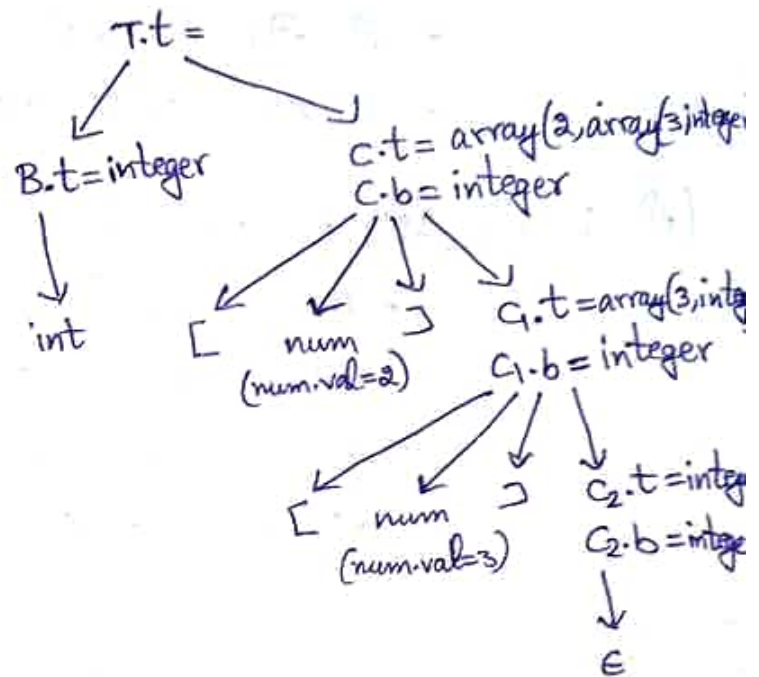
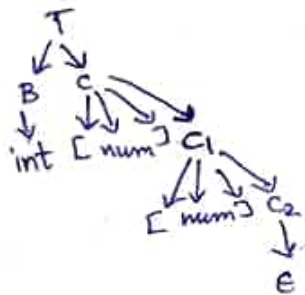
1. $T \rightarrow BC$

$T.t = C.t$

$C.b = B.t$

2. $B \rightarrow \text{int}$ $B.t = \text{integer}$
3. $B \rightarrow \text{float}$ $B.t = \text{float}$
4. $C \rightarrow [\text{num}] C_1$ $C.t = \text{array}(\text{num.val}, C_1.t)$
 $C.b = C.b$
5. $C \rightarrow \epsilon$ $C.t = C.b$

→ SDT for $\text{int}[2][3]$:



Syntax Directed Translation Schemes (SDT):

- CFG with pgm fragments embedded within production bodies.
- SDT's with all actions at right end of production bodies are called postfix SDT's.
- SDT with action inside productions:

$B \rightarrow X \{a\} Y$ a is performed

→ Bottom up Parser - as soon as X appears on top of parsing stack.

→ Top down Parser - ~~as soon as~~ before expanding this occ. of Y .

Implementation:

(i) L-attr. SDD by Recursive Descent Parsing:

- inh. attr. - arguments of functions for nonterminals.
- syn. attr. - returned by functions

(ii) L-attr. SDD on LL Grammar:

- Records to hold syn. attr. for non-terminal are placed below NT on stack.
- inh. attr. for a NT - stored with NT on stack.
- Action records are also placed on stack.

(iii) L-attr. SDD on an LL Grammar, Bottom Up:

- Marker non-terminals appears on Bottom Up
Parser's stack
 - inh. attr. of NT - above it on stack
 - syn. attr. - are kept with NT on stack.
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