

solc-verify: source-level formal verification for Solidity

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NEMZETI KUTATÁSI, FEJLESZTÉSI
ÉS INNOVÁCIÓS HIVATAL



Új Nemzeti
Kiválóság Program

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Verification Landscape

SMTChecker

KSolidity

KEVM

VeriSolid

Slither

solc-verify

MythX

Certora

VeriSol

VerX

Securify

Truffle

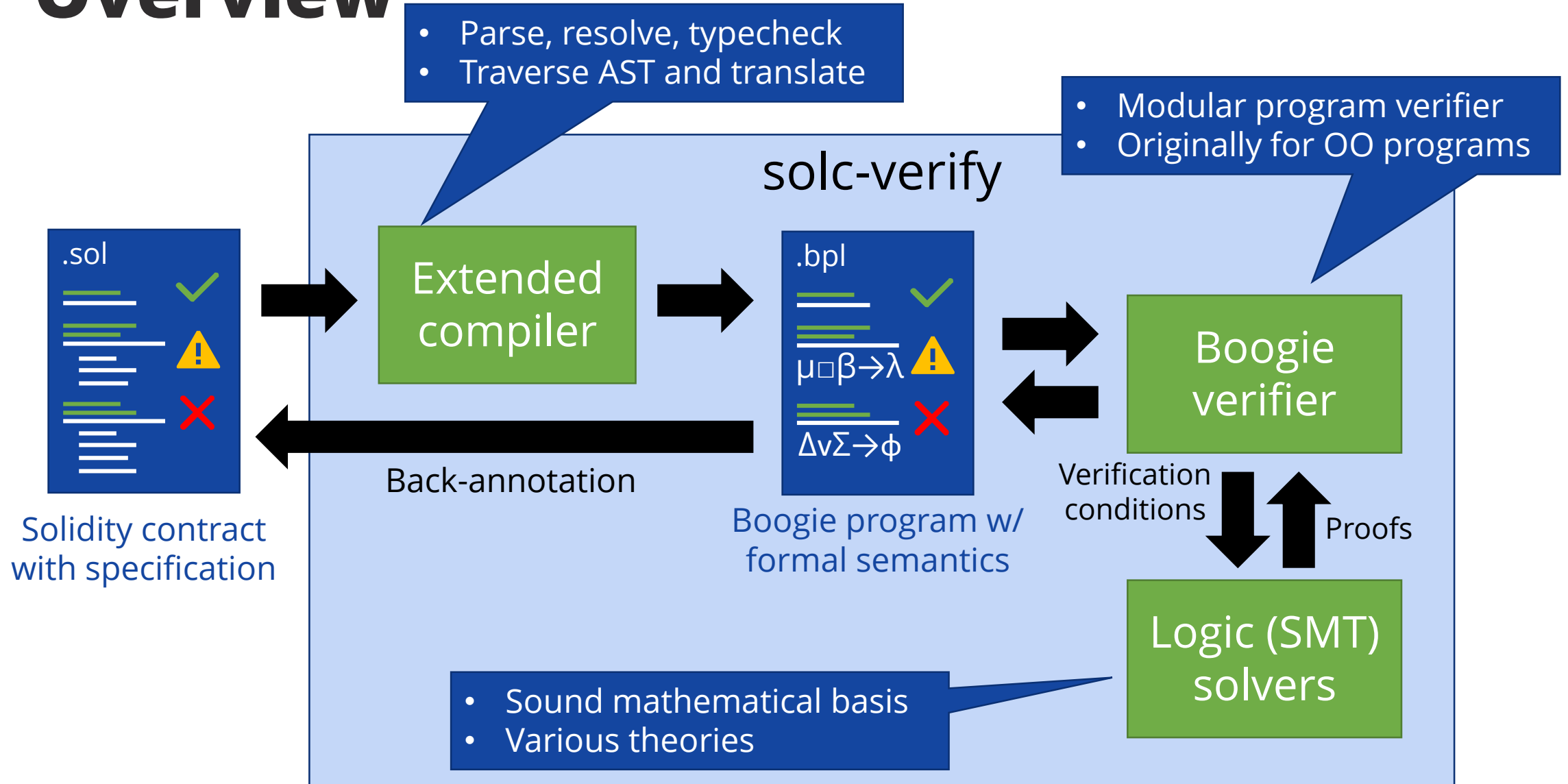
... and many more



DEMO

github.com/hajduakos/solidity-summit-demo

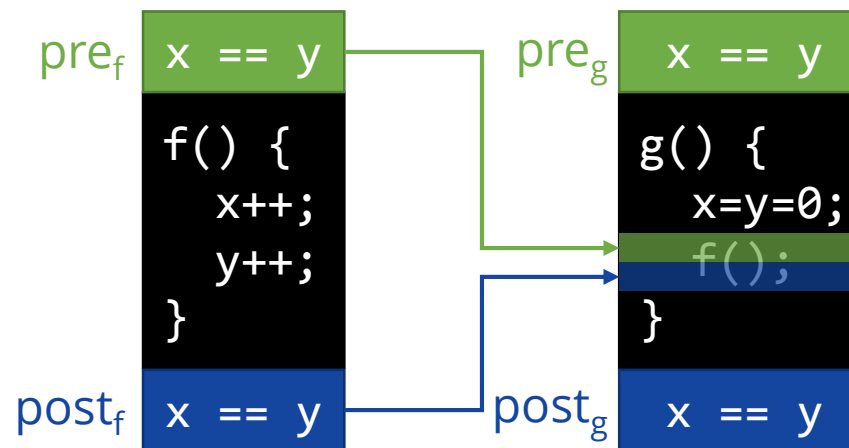
Overview



Formal Verification

- **Functional correctness** w.r.t specification
 - **Implicit**: assertion, overflow
 - **Explicit**: pre/postconditions, invariants, ...

- **Modular verification**
 - pre + body \rightarrow post
 - Discharge with SMT solvers
 - Replace calls with specification



If $x_0 = y_0$ and $x_1 = x_0 + 1$ and $y_1 = y_0 + 1$
then $x_1 = y_1$?

Specification Annotations

- Solidity provides (implicit)
 - require, assert
- Annotation language (explicit)
 - Features
 - Pre/postconditions
 - Contract level invariants
 - Loop invariants
 - Access control (modifies)
 - Events *in progress*
 - Solidity expressions (side effect free)
 - Extra: sum over collections, previous values
 - Quantifiers *in progress*

```
/// @notice invariant x == y
contract C {
    int x;
    int y;

    /// @notice precondition x == y
    /// @notice postcondition x == (y + n)
    /// @notice modifies x
    function add_to_x(int n) internal {
        x = x + n;
        require(x >= y);
    }

    /// @notice modifies x if n > 0
    /// @notice modifies y if n > 0
    function add(int n) public {
        require(n >= 0);
        add_to_x(n);
        /// @notice invariant y <= x
        while (y < x) {
            y = y + 1;
        }
    }
}
```



DEMO

github.com/hajduakos/solidity-summit-demo

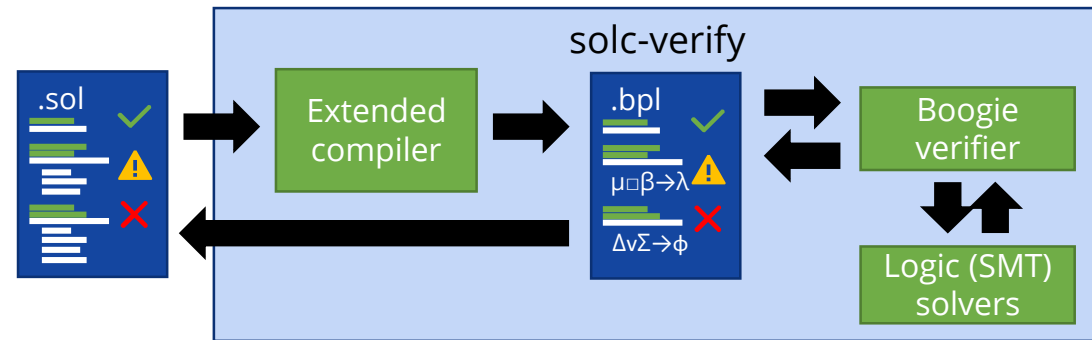
Summary

- **solc-verify**: source-level formal verification for Solidity

```
/// @notice invariant x == y
contract C {
  int x; int y;

  /// @notice precondition x == y
  /// @notice postcondition x == (y + n)
  /// @notice modifies x
  function add_to_x(int n) internal {
    x = x + n;
    require(x >= y);
  }

  /// @notice modifies x if n > 0
  /// @notice modifies y if n > 0
  function add(int n) public {
    require(n >= 0);
    add_to_x(n);
    /// @notice invariant y <= x
    while (y < x) { y = y + 1; }
  }
}
```



github.com/SRI-CSL/solidity



hajduakos.github.io



[@himynameisakos](https://twitter.com/himynameisakos)

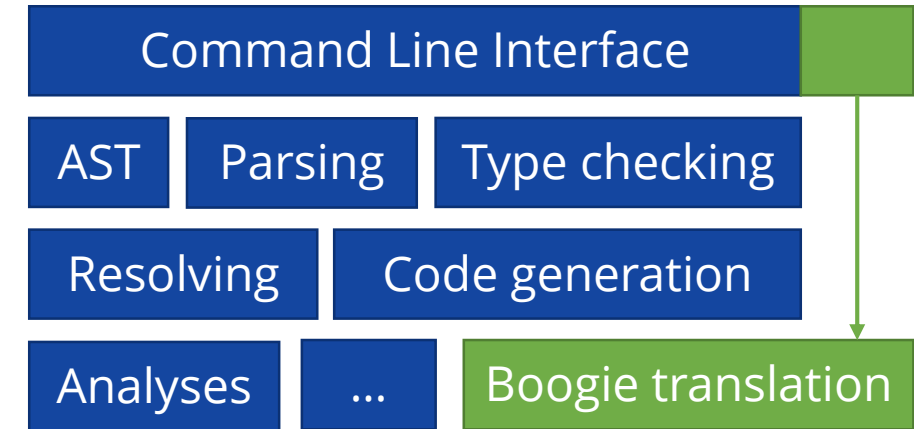
Examples: github.com/hajduakos/solidity-summit-demo

Tool paper: arxiv.org/abs/1907.04262

Formalizing memory model: arxiv.org/abs/2001.03256

Without Modifying the Compiler?

- More like **extending** than modifying
 - Compiler works same, with new options
- In principle could be done externally
 - Parse JSON AST
- **Benefits** of being inside compiler
 - More **robust** to changes, type safety
 - **Reuse** modules (e.g. parse specs)
- Extensible compiler infrastructure?
 - Like LLVM



Relationship with Act

- Act
 - Language **independent**
 - **Separate** specs
- solc-verify
 - Specs in the **same language**, Solidity
 - Code and specs **together**
- In principle, specs could come from (a subset) of Act

```
behaviour init of StateMachine
interface constructor()

creates
    uint256 x := 0
invariants
    x <= 1

behaviour f of StateMachine
interface f()

case x == 0:
    storage
        x => 1
case _:
    noop
ensures
    (x == 0) or (x == 1)
```

Relationship with SMTChecker

- SMTChecker (to the best of our knowledge)
 - Built-in analyzer in the compiler
 - [Direct translation](#) to SMT
 - Implicit specifications (overflow, assert)
 - Intra-function analysis
- We did some [experiments](#)
 - v0.5.10: arxiv.org/abs/1907.04262
 - v0.5.12: arxiv.org/abs/2001.03256
 - [Unsupported](#) features, [false alarms](#) for our use cases
 - E.g., overflows, memory model, external calls