# Homework 8 Colorado CSCI 5454

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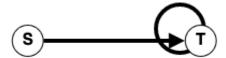
People I studied with for this homework: Gautam G, Praddunya

### Problem 1

#### **Solution**:

 $\pi$  for S and T, where the graph contains a directed edge from S to T and A self-edge to itself to T.

let's take an example of the below graph:



The Normalized Adjacency Matrix WG is -  $\begin{bmatrix} 0 & 1 \\ 0 & 1 \end{bmatrix}$ 

The above matrix  $\pi(t) \approx 1$ . Since there are no edges out of t, the probability of a vertex landing on t remains the same once it arrives t.

As we know that  $\pi$  is a stationary distribution if  $\pi W = \pi$  Let's assume the initial distribution to be  $(1/2 \ 1/2)$ 

$$\pi * W = \begin{pmatrix} 1/2 & 1/2 \end{pmatrix} * \begin{bmatrix} 0 & 1 \\ 0 & 1 \end{bmatrix}$$

$$= \begin{pmatrix} 0 & 1 \end{pmatrix}$$

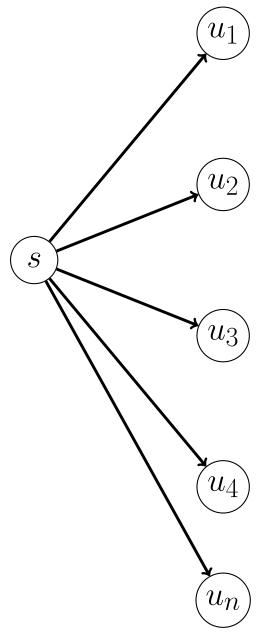
It will carry forward the same for the other examples, as we have the self-loop for the sink t. So, for normalized Adjacency Matrix has the value 1 in that position, when we multiply with W, it will have that value carry to that position. So  $\pi$  is the stationary distribution as it won't go to a different vertex, as it has no other vertex than itself.

From this, we can say that the probability is 1. and  $\pi$  is the stationary distribution.

## Problem 2

#### Solution:

• Let's consider the start graph which we have already used previously



• Now the normalized adjacency matrix for the above graph is written below

- consider initial stationary distribution as  $(1/2 \ 1/2n \ 1/2n \ 1/2n \ . \ . \ . \ 1/2n)$
- if we perform  $\pi W_G =$

$$= (1/2 \ 1/2n \ 1/2n \ 1/2n \ \dots \ 1/2n)$$

- As after matrix multiplication we get  $\pi$ . So we can say that  $\pi(W)$  is a stationary distribution.
- This can be claimed with certainty since the only other option, after reaching a leaf node, is to return to the central node and then leap to the leaf node. The central node won't have a 1/2 probability of choosing any of the leaf nodes in this scenario.
- So, from above our consideration for  $\pi(w) = 0.5$ .

### Problem 3

lets say adjacency matrix:

$$A_G[i,j] = \begin{cases} 1 & \text{if } i,j \in E \\ 0 & \text{otherwise} \end{cases}$$

• As given  $A_G[i,j] = A_G[j,i]$  as it's an undirected and an unweighted graph, it means  $A_G[i,j] = A_G[j,i] = 1$ .

- So, as  $A_G[i,j] = A_G[j,i]$  for all the vertices, we can say that  $A_G$  will be equal to it's transpose. It means that  $A_G = A_G^T$ .
- So, as  $A_G[i,j] = A_G[j,i]$  for all the vertices, we can say that  $A_G$  will be equal to it's transpose. So,  $A_G = A_G^T$ .

we can say that:

$$B = A_G * A_G$$
$$=> B = A_G * A_G^T$$

- Now, we need the sum of all the elements at [i, i] position. That means that we are finding the dot product of the  $i^{th}$  row and the  $i^{th}$  column.
- Now as we know that the adjacency matrix is equal to its transpose, we can say that the  $i^{th}$  row in the matrix and the  $i^{th}$  column in the transpose will be equal to each other.
- So basically when we consider the dot product of  $i^{th}$  row and the  $i^{th}$  column, we will be multiplying 1 with 1 and 0's with 0's (as we multiply with the same vector).
- So, we can say that:

 $B[i,i] = \text{sum of 1's in the } i^{th} \text{ row in the original adjacency matrix}$ 

$$\sum_{i=1}^{n} B[i, i] = \text{sum of 1's in all the rows}$$

- Now we know that we have m edges. For suppose for an edge between i and j, we can say that it will have a 1 in the matrix in 2 places.
- They are  $A_G[i, j]$  and  $A_G[j, i]$ . As we know that  $i \neq j$ , we can say that for m edges, there will be 2m 1's. So, the sum of all the 1's in the matrix will be,

$$\sum_{i=1}^{n} B[i, i] = \text{sum of 1's in all the rows}$$

$$= 2m * 1$$

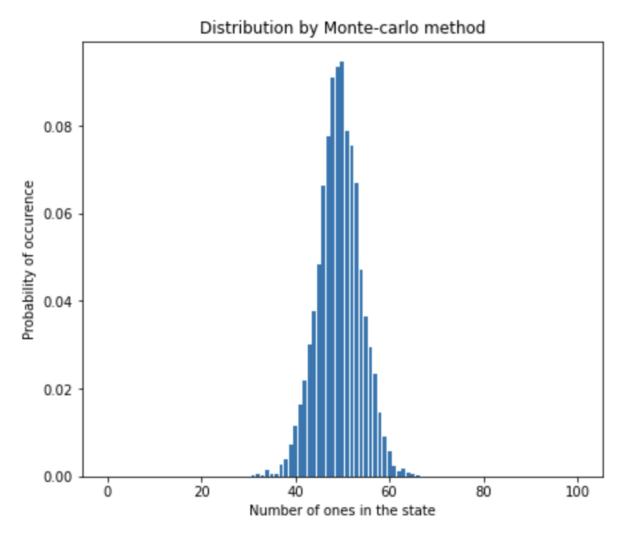
$$\sum_{i=1}^{n} B[i, i] = 2m$$

### Problem 4

#### Part a

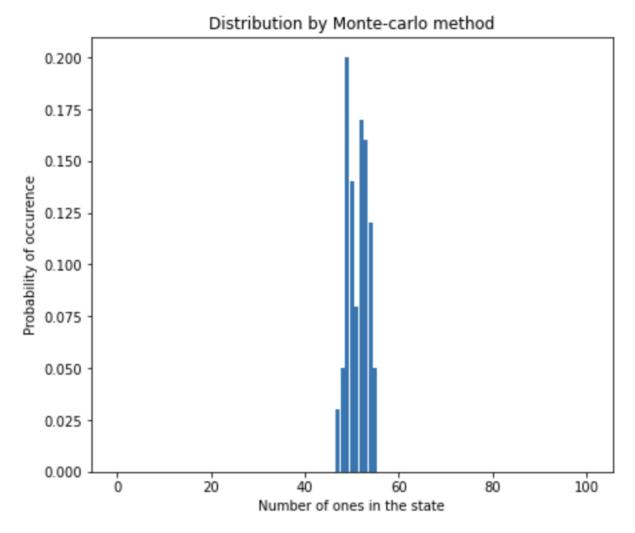
- The state of the system can be described by a vector  $x \in \{0,1\}^n$ , i.e. a list of n bits. So the number of states of the system is  $2^n$ .
- the graph in the Metropolis-Hastings algorithm. For this, we represent a state as a vertex and connect a state to all the states where they differ in their state at exactly one position. This is how we find neighbors of a state.
- So to find the degree of a vertex, we see that for a state, its neighbor can differ at one of the positions in its bits and there are n positions. So maximum degree r is n.

Part c



the histogram for distribution of Monte-carlo method is attached for T = 10000.

Part d



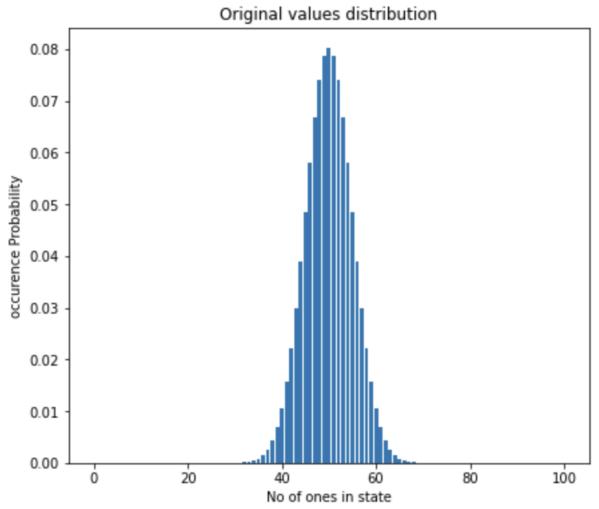
The histogram attached for T = 100 is right skewed. we know that original distribution is around 50. From above T=10000 is closer to the original distribution but for T=100 it's on the right skew.

#### part e

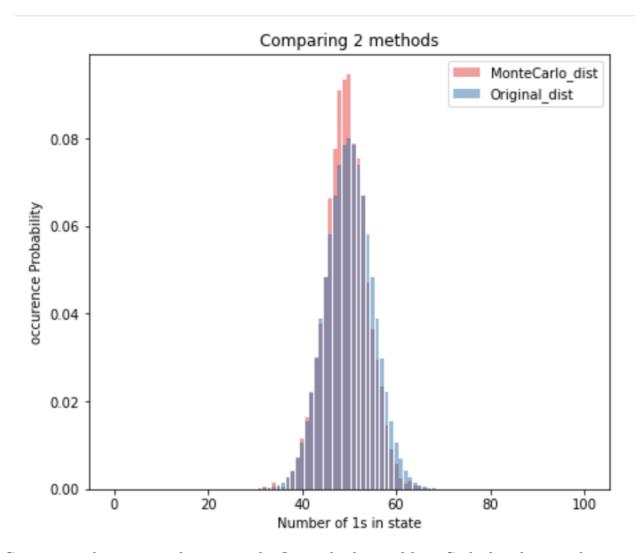
When we substitute w(x) with the (x/n) then we get the below equation. we need to calculate the original probability for the number of ones in the state.

So we have to choose x ones from n, these many states are possible which have the number of 1's as x. So the probability of this occurring is proportional to (nx)H2(x/n) To get the original probability we need to scale this.

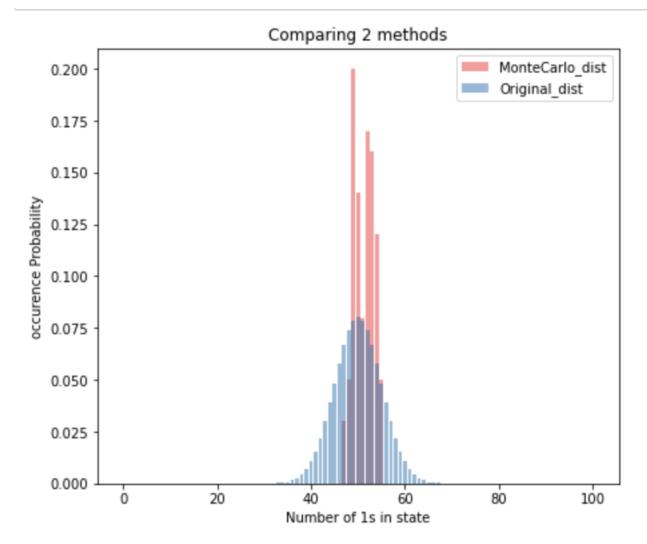
$$Pr(X = x) = \binom{n}{x} * H_2(x/n) / \sum_{x=0}^{n} \binom{n}{x} * H_2(x/n)$$
 (1)



Below is the comparison with the original graph can be found below.



Comparision between two histograms for Original values vs MonteCarlo distribution when



T=100

Part b

The complete code will be there at the end of the pdf with all the distribution and graph codes