

Computer Science & Engineering Department
I. I. T. Kharagpur

Compilers Laboratory: CS39003

3rd Year CSE: 5th Semester

Assignment – 6: Target Code Generator for **tinyC** Marks: 100
Assign Date: 19th October, 2016 Submit Date: 23:55, 8th November, 2016

1 Preamble – **tinyC**

The Lexical Grammar (Assignment 3) and the Phase Structure Grammar (Assignment 4) for **tinyC** have already been defined as subsets of the **C** language specification from the International Standard **ISO/IEC 9899:1999 (E)**. Finally, the 3-Address Code structure and a further subset of **tinyC** has been specified (Assignment 5) for translating the input **tinyC** program to an array of 3-Address quad's, a supporting symbol table, and other auxiliary data structures.

In this assignment you will write a target code translator from the array of 3-Address quad's (with the supporting symbol table, and other auxiliary data structures) to the assembly language of **x86-64**. The translation is now machine-specific and your generated assembly code would be translated with the **gcc** assembler to produce the final executable codes for the **tinyC** program.

2 Scope of Target Translation

- For simplicity restrict **tinyC** further:
 1. Skip **shift** and **bit** operators.
 2. Support only **void**, **int**, and **char** types. Skip **double** type.
 3. Support only one-dimensional arrays.
 4. Support only **void**, **int**, **char**, **void***, **int***, and **char*** types for returns types of functions.
 5. No type conversion to be supported.
- For I/O, provide a library (as created in Assignment 2) using in-line assembly language program of **x86-64** along with **syscall** for **gcc** assembler.:
 - **int prints(char *)** – prints a string of characters. The parameter is terminated by **'\0'**. The return value is the number of characters printed.
 - **int printi(int n)** – prints the integer value of **n** (no newline). It returns the number of characters printed.
 - **int readi(int *eP)** – reads an integer (signed) and returns it. The parameter is for error (**ERR = 1**, **OK = 0**).

The header file **my1.h** of the library will be as follows:

```
#ifndef _MYL_H
#define _MYL_H
#define ERR 1
#define OK 0
int prints(char *);
int printi(int);
int readi(int *eP); // *eP is for error, if the input is not an integer
#endif
```

3 Design of the Translator

The steps for target code generation were outlined in **Target Code Generation** lecture presentations. In this assignment, however, you do not need to deal with any machine-independent or machine-specific optimization. Hence the translation comprises the following major steps only:

Memory Binding This deals with the design of the allocation schema of variables (including parameters and constants) that associates each variable to the respective address expression or register. This needs to handle the following:

- *Handle local variables, parameters, and return value for a function.* These are automatic and reside in the **Activation Record (AR)** of the function. Various design schema for AR are possible based on the calling sequence protocol. A sample AR design could be as follows:

| Offset | Stack Item | Responsibility |
|--------|------------------------|---|
| -ve | Saved Registers | Callee Saves & Restores |
| -ve | Callee Local Data | Callee defines and uses |
| 0 | Base Pointer of Caller | Callee Saves & Restores |
| | Return Address | Saved by <code>call</code> , used by <code>ret</code> |
| +ve | Return Value | Callee writes, Caller reads |
| +ve | Parameters | Caller writes, Callee reads |

Activation Record Structure with Management Protocol

- Offset's in the AR are with respect to the Base Pointer of Callee.
- Return Value can alternatively be returned through a register (like accumulator).
- The AR will be populated from the Symbol Table of the function.
- Symbol Tables of nested blocks will be flattened and its variables allocated within the Symbol Table (and hence the AR) of the function where there occur in. Necessary name mangling will be performed to take care of same lexical name for different variables in different nested scopes.
- *Handle global variables* (note that local static variables are not allowed in `tinyC`) as static and generate allocations in static area. This will be populated from `ST.gbl`.
- *Generate Constants from Table of Constants* – handle string constants as assembler symbols in Data Segments and integer constants as parts of target code (Text Segment)
- *Register Allocations & Assignment:* Create memory binding for variables in registers:
 - After a load / store the variable on the activation record and the register have identical values
 - Registers can be used to store temporary computed values
 - Register allocations are often used to pass `int` or pointer parameters
 - Register allocations are often used to return `int` or pointer values

Note: Refer to **Run-Time Environment** lecture presentations for details and examples on memory binding.

Code Translation This deals with the translation of 3-Address quad's to x86-64 assembly code. This needs to handle:

- *Generation of Function Prologue* – few lines of code at the beginning of a function, which prepare the stack and registers for use within the function.
- *Generate Function Epilogue* – appears at the end of the function, and restores the stack and registers to the state they were in before the function was called.

- *Map 3-Address Code to Assembly* – to translate the function body do:
 - Choose optimized assembly instructions for every expression, assignment and control **quad**.
 - Use algebraic simplification & reduction of strength for choice of assembly instructions from a **quad**.
 - Use Machine Idioms (like **inc** for **i++** or **++i** in place of **add reg, 1**).

Note: Refer to Target Code Generation lecture presentations for details.

Target Code Integrate all the above code into an Assembly File for gcc assembler.

4 The Assignment

1. Write a target code (x86-64) translator from the 3-Address **quad**'s generated from the flex and yacc specifications of **tinyC** (with restrictions as mentioned in Section 2). Assume that the input **tinyC** file is lexically, syntactically, and semantically correct. Hence no error handling and / or recovery is expected.
2. Prepare a Makefile to compile and test the project.
3. Prepare test input files **ass6_roll_test<number>.c** to test the target code translation and generate the translation output in **ass6_roll <number>.asm**.
4. Name your files as follows:

| File | Naming |
|--|--|
| Flex Specification | ass6_roll.l |
| Yacc Specification | ass6_roll.y |
| Data Structures (Class Definitions) & Global Function Prototypes | ass6_roll_translator.h |
| Data Structures, Function Implementations & 3-Address Translator | ass6_roll_translator.cxx |
| Target Translator & x86-64 Translator main() | ass6_roll_target_translator.cxx |
| Test Inputs | ass6_roll_test<number>.c |
| 3-Address Test Outputs | ass6_roll_quads<number>.out |
| Test Outputs | ass6_roll <number>.asm |

5. Prepare a tar-archive with the name **ass6_roll.tar** containing all the files and upload to Moodle.

5 Credits

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|--|---------------------------------|
| Design of Memory Binding: | 15 + 5 + 5 + 5 + 10 = 40 |
| <i>Handling of Activation Records</i> | |
| <i>Handling of Nested Symbol Tables</i> | |
| <i>Handling of Static Memory & Binding</i> | |
| <i>Handling of Constants</i> | |
| <i>Handling of Register Allocation & Assignment</i> | |
| Design of Code Translation: | 5 + 5 + 10 = 20 |
| <i>Handling of Prologue</i> | |
| <i>Handling of Epilogue</i> | |
| <i>Handling of Function Body</i> | |
| Design of Target Code Management: | 10 |
| <i>Integration of translated codes into an assembly file</i> | |
| Design of Test files and correctness of outputs: | 10 + 10 = 20 |
| <i>Test at least 5 i/p files covering all rules</i> | |
| <i>Shortcoming and / or bugs, if any, should be high-lighted</i> | |
| Integrated interface of the tinyC Compiler: | 10 |