GTU Department of Computer Engineering
CSE344 - Spring 2021
Midterm Report

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1 Problem Definition

The problem is to implement **producer-consumer** problem by simulating covid19 vaccination flow by using shared memory and posix semaphore.

2 Solution

The homework was finished as expected in homework pdf file.

2.1 Some Problems and Solutions

2.1.1 Semaphore Usage

In order to provide synchronization between processes I used **7 posix named semaphore**. Followings are my semaphores;

```
1    /*create named semphores*/
2    sem_mutex = sem_open("mutex344", O_CREAT, 0666, 1);
3    if (sem_mutex == SEM_FAILED)
4        errExit("sem_open error!");
5
6    sem_vac = sem_open("wait_vaccinator", O_CREAT, 0666, 0);
7    if (sem_vac == SEM_FAILED)
8        errExit("sem_open error!");
9
10   sem_cit = sem_open("wait_citizen", O_CREAT, 0666, 0);
11   if (sem_cit == SEM_FAILED)
12       errExit("sem_open error!");
13
14   sem_full1 = sem_open("full344", O_CREAT, 0666, 0);
15   if (sem_full1 == SEM_FAILED)
16       errExit("sem_open error!");
17   sem_empty = sem_open("empty344", O_CREAT, 0666, bufferSize);
18   if (sem_empty == SEM_FAILED)
19       errExit("sem_open error!");
20   sem_full2 = sem_open("full3442", O_CREAT, 0666, 0);
21   if (sem_full2 == SEM_FAILED)
22       errExit("sem_open error!");
23   sem_run = sem_open("whorunfirst", O_CREAT, 0666, 0);
24   if (sem_run == SEM_FAILED)
25       errExit("sem_open error!");
```

Usage purpose of each semaphore will be explained in detail coming pages.

2.1.2 Shared Memory Usage

In order to use shared memory correct way between process I made following struct to keep in shared memory for each process. Note that this memory is generic between all process.

```
1  typedef struct GTU344
2  {
3      args givenParams;
4      int dose1;
5      int dose2;
6      int totalLeft;
7      int isRead;
8      int fd;
9      int leftCiti;
10 } clinic;
```

In this struct **dose1** and **dose2** denotes number of vaccine 1 and vaccine 2 in buffer(clinic), **isRead** denotes reading is done or not **totalLeft** denotes number of left vaccined to be made.

Usage:

```
static char memoryName[50];
static clinic *biontech;
mode_t mode = S_IRUSR | S_IWUSR | S_IRGRP | S_IROTH | S_IRWXU;
strcpy(memoryName, "clinic_sinovac344");
int memFd = shm_open(memoryName, O_CREAT | O_RDWR, mode);
if (memFd == -1)
    errExit("shm_open error!");
if (ftruncate(memFd, sizeof(*biontech)) == -1)
    errExit("ftruncate error");

clinic *gata = (clinic *)mmap(NULL, sizeof(*biontech), PROT_READ | PROT_WRITE, MAP_SHARED, memFd, 0);
if (gata == MAP_FAILED)
    errExit("mmap");

gata->givenParams = givenArgs;
gata->dose1 = 0;
gata->dose2 = 0;
gata->totalLeft = 2 * (givenArgs.tArg * givenArgs.cArg);
gata->isRead = 0;
gata->leftCiti = givenArgs.cArg;
gata->fd = safeOpen(givenArgs.iArg, O_RDONLY);
```

2.1.3 Communication between Nurses and Vaccinators

This problem is nothing but simply producer consumer problem. There are 4 semaphore used in this problem which are **empty, full1, full2, mutex**.

2.1.4 Solution

In classical solution of this problem we are using 3 semaphore which are **full, empty, mutex** but these are not enough to solve our problem because vaccinator needs to wait until there is at least 1 shot(1 and 2) in buffer. In order to solve this problem I used 1 extra **full** semaphore, in order to wait on both vaccine1 and vaccine2. This means that consumer in order to consume needs to get both full1 and full2 semaphore. Normally, producer(nurse) post full semaphore directly but here it will post full1 or full2 semaphore not both of them. Simple example is following;

Producer(Nurse)

```
if (sem_wait(sem_empty) == -1)
    errExit("sem_wait");
if (sem_wait(sem_mutex) == -1)
    errExit("sem_wait");
if (!biontech->isRead)
{
    char vaccine = readOneChar(biontech->fd);
    if (vaccine == '1')
    {
        biontech->dose1 = biontech->dose1 + 1;
        printNurseMsg(process->index, process->pid, '1', biontech);
        if (sem_post(sem_full1) == -1)
            errExit("sem_post");
    }
    else if (vaccine == '2')
    {
        biontech->dose2 = biontech->dose2 + 1;
        printNurseMsg(process->index, process->pid, '2', biontech);
        if (sem_post(sem_full2) == -1)
            errExit("sem_post");
    }
    else if (vaccine == 'x')
    {
        nurseLeaveMsg();
        biontech->isRead = 1;
        break;
    }
    else
        errExit("vaccine is wrong!!");
    if (sem_post(sem_mutex) == -1)
        errExit("sem_post");
```

Consumer(Vaccinator)

```
1           if (sem_wait(sem_full1) == -1)
2               errExit("sem_wait");
3           if (sem_wait(sem_full2) == -1)
4               errExit("sem_wait");
5           if (sem_wait(sem_mutex) == -1)
6               errExit("sem_wait");
7           if (biontech->totalLeft > 0)
8           {
9               if (sem_post(sem_vac) == -1)
10                  errExit("sem_post");
11              if (sem_wait(sem_cit) == -1) //wait for cit update its pid
12                  errExit("sem_wait");
13              printVaccinatorMsg(process->index, process->pid, *currentCitPid);
14              biontech->dose1 = biontech->dose1 - 1;
15              biontech->dose2 = biontech->dose2 - 1;
16              biontech->totalLeft = biontech->totalLeft - 2;
17              counter++;
18              if (sem_post(sem_run) == -1)
19                  errExit("sem_post");
20              if (sem_wait(sem_cit) == -1) //wait for cit give its message
21                  errExit("sem_wait");
22              if (sem_post(sem_mutex) == -1)
23                  errExit("sem_post");
24              if (sem_post(sem_empty) == -1)
25                  errExit("sem_post");
26              if (sem_post(sem_empty) == -1)
27                  errExit("sem_post");
28          }
29          else
30          {
31              break;
32          }
```

In this code piece most semaphores belongs to vaccinator and citizen communication. Also, We are making 2 empty semaphore post successive because we are deleting two element from buffer.

2.1.5 Communication between Vaccinators and Citizens

This problem again simple producer consumer problem.

2.1.6 Solution

There are 2 critic part in this problem;

▶ Getting waken citizen pid.

▶ Adjusment of message printing on screen

The solution of these problems is following;

Vaccinator

```
1       ...
2       //wake the citizen
3       if (sem_post(sem_vac) == -1)
4           errExit("sem_post");
5           //wait for citizen write its pid into shared memory
6       if (sem_wait(sem_cit) == -1)
7           errExit("sem_wait");
8
9       ...(remove vaccines from buffer)
10
11      //wake citizen again to write
12      //its message on the screen
13      if (sem_post(sem_run) == -1)
14          errExit("sem_post");
15      //wait for citizen finish its job
16      if (sem_wait(sem_cit) == -1)
17          errExit("sem_wait");
18      ...
```

Citizen

```
1    ...
2    //wait an invite from vaccinator
3    if (sem_wait(sem_vac) == -1)
4        errExit("sem_wait");
5
6    //update pid in shared memory
7    *currentCitPid = process->pid;
8    left --;
9
10   //post and wait again
11   if (sem_post(sem_cit) == -1)
12       errExit("sem_post");
13   if (sem_wait(sem_run) == -1)
14       errExit("sem_wait");
15   //write your mesage
16   printCitizenMsg(process->index, process->pid, total - left, biontech);
17   //post finished semaphore
18   if (sem_post(sem_cit) == -1)
19       errExit("sem_post");
20   ...
```

In this solution 3 **binary** semaphore was used. These are **vac, cit, run**.

Note: In this solution by using semaphores, I provided print message on screen is happening like a queue. In other words, each citizen will print its message directly after getting an invite from a vaccinator and vaccinator message. I had to do this otherwise messages on the screen will be messed up.

2.1.7   Removing Resources and Exiting

In order to run a cleaner method after finish, I used a cleaner method like this;

```
1    void removeAll()
2    {
3        if (sem_close(sem_full1) == -1)
4            errExit("sem_close");
5        if (sem_close(sem_empty) == -1)
6            errExit("sem_close");
7        if (sem_close(sem_full2) == -1)
8            errExit("sem_close");
9        if (sem_close(sem_run) == -1)
10           errExit("sem_close");
11       if (sem_close(sem_mutex) == -1)
12           errExit("sem_close");
13       if (sem_close(sem_vac) == -1)
14           errExit("sem_close");
15       if (sem_close(sem_cit) == -1)
16           errExit("sem_close");
17       sem_unlink("mutex344");
18       sem_unlink("full344");
19       sem_unlink("empty344");
20       sem_unlink("full3442");
21       sem_unlink("whorunfirst");
22       sem_unlink("wait_vaccinator");
23       sem_unlink("wait_citizen");
24       shm_unlink(memoryName); //generic data
25       shm_unlink(helperMemory); //current cit pid
26   }
```

Usage:

```
1    void cleanAndExit()
2    {
3
4        if (processInfo.type == PARENT)
5        {
6            reapDeadChildren();
7            if (close(biontech->fd) < 0)
8            {
9                errExit("close file error!");
10           }
11       }
12       else
13       {
14           removeAll();
15       }
16       exit(EXIT_SUCCESS);
17   }
```

2.1.8 Reaping Dead Children

In order not to let zombie process, I used waitpid to wait dead children. Example;

```
1    void reapDeadChildren()
2    {
3        pid_t childPid;
4        int status;
5        while ((childPid = waitpid(-1, &status, 0)) > 0);
6        if (childPid == -1 && errno != ECHILD)
7            errExit("waitpid");
8    }
```

2.1.9 CTRL-C Handling

In order to give a message on CTRL-C interrupt, I used **sigaction** function from **signal.h** library. Also, in midterm pdf there is no force to print a message on CTRL-C interrup, therefore I didn't give any message.

```
1    //handler function
2    void exitHandler(int signal)
3    {
4        if (signal == SIGINT)
5        {
6            int savedErrno = errno;
7            cleanAndExit();
8            errno = savedErrno;
9        }
10   }
```

There is no message, give resources back and exit elegantly.

3   References that was used

While doing this howemork following references was used;

▶ Course Textbook Listing 26-5: Reaping dead children via a handler for SIGCHLD

▶ Week-8 slides synchronization: Bounded buffer case in producer consumer problem.

4 Test Result

I tried to obey output example in midterm pdf. In following output all decisions belongs to CPU.

A simple test result is following;


```
akif@ubuntu:~/Desktop/system-programming/Midterm Project$ ./program -n 3 -v 2 -c 3 -b 11 -t 3 -i /home/akif/Desktop/source/mysource
Welcome to the GTU344 clinic. Number of citizens to vaccinate c=3 with t=3 doses.
Nurse 1 (pid=20123) has brought vaccine 1:the clinic has 1 vaccine1 and 0 vaccine2.
Nurse 1 (pid=20123) has brought vaccine 1:the clinic has 2 vaccine1 and 0 vaccine2.
Nurse 1 (pid=20123) has brought vaccine 1:the clinic has 3 vaccine1 and 0 vaccine2.
Nurse 1 (pid=20123) has brought vaccine 1:the clinic has 4 vaccine1 and 0 vaccine2.
Nurse 2 (pid=20124) has brought vaccine 1:the clinic has 5 vaccine1 and 0 vaccine2.
Nurse 2 (pid=20124) has brought vaccine 1:the clinic has 6 vaccine1 and 0 vaccine2.
Nurse 2 (pid=20124) has brought vaccine 1:the clinic has 7 vaccine1 and 0 vaccine2.
Nurse 2 (pid=20124) has brought vaccine 1:the clinic has 8 vaccine1 and 0 vaccine2.
Nurse 2 (pid=20124) has brought vaccine 1:the clinic has 9 vaccine1 and 0 vaccine2.
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 9 vaccine1 and 1 vaccine2.
Vaccinator 2 (pid=20122) is inviting citizen pid=20119 to the clinic
Citizen 2 (pid=20119) is vaccinated for the 1st time: the clinic has 8 vaccine1 and 0 vaccine2
Nurse 1 (pid=20123) has brought vaccine 2:the clinic has 8 vaccine1 and 1 vaccine2.
Nurse 1 (pid=20123) has brought vaccine 2:the clinic has 8 vaccine1 and 2 vaccine2.
Nurse 1 (pid=20123) has brought vaccine 2:the clinic has 8 vaccine1 and 3 vaccine2.
Vaccinator 2 (pid=20122) is inviting citizen pid=20118 to the clinic
Citizen 1 (pid=20118) is vaccinated for the 1st time: the clinic has 7 vaccine1 and 2 vaccine2
Vaccinator 2 (pid=20122) is inviting citizen pid=20120 to the clinic
Citizen 3 (pid=20120) is vaccinated for the 1st time: the clinic has 6 vaccine1 and 1 vaccine2
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 6 vaccine1 and 2 vaccine2.
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 6 vaccine1 and 3 vaccine2.
Vaccinator 2 (pid=20122) is inviting citizen pid=20119 to the clinic
Citizen 2 (pid=20119) is vaccinated for the 2nd time: the clinic has 5 vaccine1 and 2 vaccine2
Vaccinator 2 (pid=20122) is inviting citizen pid=20118 to the clinic
Citizen 1 (pid=20118) is vaccinated for the 2nd time: the clinic has 4 vaccine1 and 1 vaccine2
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 4 vaccine1 and 2 vaccine2.
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 4 vaccine1 and 3 vaccine2.
Nurse 3 (pid=20125) has brought vaccine 2:the clinic has 4 vaccine1 and 4 vaccine2.
Nurses have carried all vaccines to the buffer, terminating.
Vaccinator 1 (pid=20121) is inviting citizen pid=20120 to the clinic
Citizen 3 (pid=20120) is vaccinated for the 2nd time: the clinic has 3 vaccine1 and 3 vaccine2
Vaccinator 1 (pid=20121) is inviting citizen pid=20119 to the clinic
Citizen 2 (pid=20119) is vaccinated for the 3rd time: the clinic has 2 vaccine1 and 2 vaccine2
Citizen is leaving. Remaining citizens to vaccinate: 2
Vaccinator 1 (pid=20121) is inviting citizen pid=20118 to the clinic
Citizen 1 (pid=20118) is vaccinated for the 3rd time: the clinic has 1 vaccine1 and 1 vaccine2
Citizen is leaving. Remaining citizens to vaccinate: 1
Vaccinator 1 (pid=20121) is inviting citizen pid=20120 to the clinic
Citizen 3 (pid=20120) is vaccinated for the 3rd time: the clinic has 0 vaccine1 and 0 vaccine2
Citizen is leaving. Remaining citizens to vaccinate: 0
All citizens have been vaccinated.
Vaccinator 1 (pid=20121) vaccinated 4 doses.
Vaccinator 2 (pid=20122) vaccinated 5 doses.
The clinic is now closed. Stay healthy.
akif@ubuntu:~/Desktop/system-programming/Midterm Project$
```

Figure 1