Data Science & Analytics

Hiren Deliwala & Dr. Jian ZHANG

Agenda

- Continue from last class
 - Review Table Structures
 - Review SQL Syntax
 - Review MySQL commands
- Run Queries from World Database & Employees Database
- Review Advanced Concepts
 - Joins,
 - Group Bys
 - Nested Queries
 - Nulls

Table Structures

- World Database
 - Country Country identified by Code
 - City Joins with Country using City.CountryCode = Country.Code
 - CountryLanguage Joins with Country using CountryLanguage.CountryCode = Country.Code
- Employees Database
 - employees Employees identified by emp_no
 - departments department identified by dept_no
 - dept_manager
 - Joins with department using dept_manager.dept_no = departments.dept_no
 - ▶ Joins with employees using dept_manager.emp_no = employees.emp_no
 - dept_emp
 - Joins with department using dept_emp.dept_no = departments.dept_no
 - Joins with employees using dept_emp.emp_no = employees.emp_no
 - titles
 - ▶ Joins with employees using titles.emp_no = employees.emp_no

SQL Syntax

- SELECT <attributes>
- FROM <one or more relations>
- WHERE <conditions>

```
SELECT PName, Price, Manufacturer
```

```
FROM Product
```

WHERE Price > 100

```
SELECT PName, Price
```

FROM Product, Company

WHERE Manufacturer=CName AND Country='Japan'

AND Price <= 200

MySQL Basics

show databases;	show which databases are available
select database();	show the name of the current database
select version();	show mysql version
show tables;	show the tables in your current database
help	shows help options and available commands
help select	(for example) shows help for the select command
\P	sets the Pager so that results are displayed one page at a time; if this is on, you need to type 'q' to exit the result display
quit	to exit mysql
describe [table name];	to describe the table with that name
use [database name];	to change to the database with that name
source [filename.sql];	execute the commands in the file with that name

SQL JOIN OPERATION

SQL: Join operation

A join can be specified in the FROM clause which list the two input relations and the WHERE clause which lists the join condition. Example:

	CITIP	
ID	State	Dept
1000	CA	
1001	MA	1001
1002	TN	1002

Emn

• F .	
ID	Name
1001	IT
1002	Sales
1003	Biotech

Dept

SQL: Join operation (cont.)

inner join = join
 SELECT *
 FROM emp join dept (or FROM emp, dept)
 on emp.id = dept.id;

Emp		Dept		
ID	State	Dept	ID	Name
1000	CA		1001	IT
1001	MA	1001	1002	Sales
1002	TN	1002	1003	Biotech

Emp.ID	Emp.State	Dept.ID	Dept.Name
1001	MA	1001	IT
1002	TN	1002	Sales

SQL: Join operation (cont.)

left outer join = left join
SELECT *
FROM emp left join dept
on emp.id = dept.id;

ID	State	Dept
1000	CA	
1001	MA	1001
1002	TN	1002

Emp

9 7 P .	
ID	Name
1001	IT
1002	Sales

Biotech

Dept

1003

Emp.ID	Emp.State	Dept.ID	Dept.Name
1000	CA	null	null
1001	MA	1001	IT
1002	TN	1002	Sales

Show all employees even if they have no department assigned

SQL: Join operation (cont.)

right outer join = right join **SELECT** * FROM emp right join dept on emp.id = dept.id;

Emp			Dep	ot
ID	State	Dept	ID	Name
1000	CA		1001	IT
1001	MA	1001	1002	Sales
1002	TN	1002	1003	Biote

ch

Emp.ID	Emp.State	Dept.ID	Dept.Name	Э 📗
1001	MA	1001	IT	
1002	TN	1002	Sales	
null	null	1003	Biotech	

Show all departments even if they have no employees

GROUPING AND AGGREGATION

Purchase

ProdName	Store	Date	Price	Quantity
----------	-------	------	-------	----------

Grouping and Aggregation

Purchase(product, date, price, quantity)

Find total sales after 01/1/2014 per product.

SELECT product, Sum(price*quantity) AS TotalSales

FROM Purchase

WHERE date > '01/1/2014'

GROUP BY product

Let's see what this means...

Grouping and Aggregation

- 1. Compute the FROM and WHERE clauses.
- 2. Group by the attributes in the GROUPBY
- 3. Compute the SELECT clause: grouped attributes and aggregates.

HAVING Clause

Same query, except that we consider only products that had at least 30 buyers.

SELECT product, Sum(price * quantity)

FROM Purchase

WHERE date > '01/1/2014'

GROUP BY product

HAVING Sum(quantity) > 3

HAVING clause contains conditions on aggregates.

General form of Grouping and Aggregation

SELECT S

FROM $R_1,...,R_n$

WHERE C1

GROUP BY $a_1,...,a_k$

HAVING C2

 $S = may contain attributes a_1,...,a_k and/or any aggregates but NO OTHER ATTRIBUTES$

C1 = is any condition on the attributes in $R_1,...,R_n$

C2 = is any condition on aggregate expressions

Why?

General form of Grouping and Aggregation

```
\begin{array}{ccc} \textbf{SELECT} & \textbf{S} \\ \textbf{FROM} & \textbf{R}_1, \dots, \textbf{R}_n \\ \textbf{WHERE} & \textbf{C1} \\ \textbf{GROUP BY } \textbf{a}_1, \dots, \textbf{a}_k \\ \textbf{HAVING} & \textbf{C2} \\ \end{array}
```

Evaluation steps:

- 1. Evaluate FROM-WHERE, apply condition C1
- 2. Group by the attributes $a_1,...,a_k$
- 3. Apply condition C2 to each group (may have aggregates)
- 4. Compute aggregates in S and return the result

NESTED QUERIES

Reserves

Details

 sid
 bid
 day

 22
 101
 10/10/96

 95
 103
 11/12/96

Sailors

We will use these instances of relations in our examples.

Boats

sid	sname	rating	age
22	Dustin	7	45.0
31	Lubber	8	55.5
95	Bob	3	63.5

<u>bid</u>	bname	color
101	Interlake	blue
102	Interlake	red
103	Clipper	green
104	Marine	red

Nested Queries

- Powerful feature of SQL: WHERE clause can itself contain an SQL query!
 - Actually, so can FROM and HAVING clauses.

Names of sailors who've reserved boat #103:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN (SELECT R.sid
FROM Reserves R
WHERE R.bid=103)
```

- ▶ To find sailors who've **not** reserved #103, use NOT IN.
- To understand semantics of nested queries:
 - think of a nested loops evaluation: For each Sailors tuple, check the qualification by computing the subquery.

Nested Queries with Correlation

Find names of sailors who've reserved boat #103:

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS (SELECT *
FROM Reserves R)
WHERE R.bid=103 AND
S.sid=R.sid)
```

- EXISTS is another set comparison operator, like IN.
- Can also specify NOT EXISTS
- ► If UNIQUE is used, and * is replaced by R.bid, finds sailors with at most one reservation for boat #103.
 - ▶ UNIQUE checks for duplicate tuples in a subquery;
- Subquery must be recomputed for each Sailors tuple.
 - Think of subquery as a function call that runs a query!

Subqueries producing One Value

Usually subqueries produce a relation as an answer. However, sometimes we expect them to produce single values

```
SELECT Purchase.product
FROM Purchase
WHERE buyer =
   (SELECT buyer
FROM Person
WHERE Home_Zip = '70810');
```

In this case, the subquery returns one value. If it returns more, it's a run-time error.

GROUP BY v.s. Nested Quereis

SELECT product, Sum(price*quantity) **AS** TotalSales

FROM Purchase

WHERE date > '10/1/2014'

GROUP BY product

SELECT DISTINCT x.product, (**SELECT Sum**(y.price*y.quantity)

FROM Purchase y

WHERE x.product = y.product

AND y.date > 10/1/2014

AS TotalSales

FROM Purchase x

WHERE x.date > '10/1/2010'

Group-by v.s. Nested Query

Author(<u>login</u>,name) Wrote(login,url)

- Find authors who wrote ≥ 10 documents:
- Attempt 1: with nested queries

```
FROM Author

WHERE count(SELECT Wrote.url

FROM Wrote

WHERE Author.login=Wrote.login)

> 10
```

Group-by v.s. Nested Query

- Find all authors who wrote at least 10 documents:
- Attempt 2: SQL style (with GROUP BY)

SELECT Author.name

FROM Author, Wrote

WHERE Author.login=Wrote.login

GROUP BY Author.name

HAVING count(wrote.url) > 10

No need for DISTINCT: automatically from GROUP BY

Group-by v.s. Nested Query

Author(<u>login</u>,name)

Wrote(login,url)

Mentions(url,word)

Find authors with vocabulary ≥ 10000 words:

SELECT Author.name

FROM Author, Wrote, Mentions

WHERE Author.login=Wrote.login AND Wrote.url=Mentions.url

GROUP BY Author.name

HAVING count(distinct Mentions.word) > 10000

Store(sid, sname)
Product(pid, pname, price, sid)

Find all stores that sell *only* products with price > 100

same as:

Find all stores s.t. all their products have price > 100)

FROM Store.name
FROM Store, Product
WHERE Store.sid = Product.sid
GROUP BY Store.sid, Store.name
HAVING 100 < min(Product.price)</pre>

```
SELECT Store.name
FROM Store
WHERE Store.sid NOT IN

(SELECT Product.sid
FROM Product
WHERE Product.price <= 100)
```

Store(<u>sid</u>, sname)
Product(<u>pid</u>, pname, price, sid)

For each store, find its most expensive product

This is easy but doesn't do what we want:

```
SELECT Store.sname, max(Product.price)
FROM Store, Product
WHERE Store.sid = Product.sid
GROUP BY Store.sid, Store.sname
```

Better:

But may return multiple product names per store

Finally, choose some pid arbitrarily, if there are many with highest price:

NULLS in SQL

NULLS in SQL

- Whenever we don't have a value, we can put a NULL
- Can mean many things:
 - Value does not exists
 - Value exists but is unknown
 - Value not applicable
 - Etc.
- The schema specifies for each attribute if can be null (nullable attribute) or not
- How does SQL cope with tables that have NULLs?

- If x = NULL then 4*(3-x)/7 is still NULL
- ► If x= NULL then x="Joe" is UNKNOWN
- In SQL there are three boolean values:

```
FALSE = 0
UNKNOWN = 0.5
TRUE = 1
```

- \triangleright C1 AND C2 = min(C1, C2)
- \triangleright C1 OR C2 = max(C1, C2)
- NOT C1 = 1 C1

```
SELECT *
FROM Person
WHERE (age < 25) AND
(height > 6 OR weight > 190)
```

E.g.
age=20
heigth=NULL
weight=200

Rule in SQL: include only tuples that yield TRUE

Unexpected behavior:

```
SELECT *
FROM Person
WHERE age < 25 OR age >= 25
```

Can test for NULL explicitly:

- > x IS NULL
- > x IS NOT NULL

```
SELECT *
FROM Person
WHERE age < 25 OR age >= 25 OR age IS NULL
```

Now it includes all Persons