

1. Which Parser is used for the implementation of recursive descent parsing? Draw the model diagram for that parser. Construct the parsing table for the grammar

$$E \rightarrow E+T \mid T$$

$$T \rightarrow T*F \mid F$$

$$F \rightarrow (E) \mid id$$

Also construct the LR (0) Parsing table of the above grammar.

Answer

A grammar is said to be left recursive if it has a non-terminal A such that there is a derivation

$A \Rightarrow A\alpha$ for some string α . Top-down parsing methods cannot handle left-recursive grammars.

Hence, left recursion can be eliminated as follows:

If there is a production $A \rightarrow A\alpha \mid \beta$ it can be replaced with a sequence of two productions

$$A \rightarrow \beta A' \quad A' \rightarrow \alpha A' \mid \epsilon$$

Without changing the set of strings derivable from A.

Consider the following grammar for arithmetic expressions:

$$E \rightarrow E+T \mid T$$

$$T \rightarrow T*F \mid F$$

$$F \rightarrow (E) \mid id$$

First eliminate the left recursion for E as

$$E \rightarrow TE'$$

$$E' \rightarrow +TE' \mid \epsilon$$

Then eliminate for T as

$$T \rightarrow FT'$$

$$T' \rightarrow *FT' \mid \epsilon$$

Thus, the obtained grammar after eliminating left recursion is

$$E \rightarrow TE'$$

$$E' \rightarrow +TE' \mid \epsilon$$

$$T \rightarrow FT'$$

$$T' \rightarrow *FT' \mid \epsilon$$

$$F \rightarrow (E) \mid id$$

This is not LR(0)
Grammar.

	FIRST	FOLLOW
E	{(, id}	{ \$,) }
E'	{+, ε }	{ \$,) }
T	{(, id}	{ +, \$,) }
T'	{*, ε }	{ +, \$,) }
F	{(, id}	{ *, +, \$,) }

Answer CONT...

	INPUT SYMBOLS					
	+	*	()	id	\$
E			$E \rightarrow TE'$		$E \rightarrow TE'$	
E'	$E' \rightarrow +TE'$			$E' \rightarrow \varepsilon$		$E' \rightarrow \varepsilon$
T			$T \rightarrow FT'$		$T \rightarrow FT'$	
T'	$T' \rightarrow \varepsilon$	$T' \rightarrow *FT'$		$T' \rightarrow \varepsilon$		$T' \rightarrow \varepsilon$
F			$F \rightarrow (E)$		$F \rightarrow id$	

Stack implementation:

stack	Input	Output
\$E	id+id*id \$	
\$E'T	id+id*id \$	$E \rightarrow TE'$
\$E'T'F	id+id*id \$	$T \rightarrow FT'$
\$E'T'id	id+id*id \$	$F \rightarrow id$
\$E'T'	+id*id \$	
\$E'	+id*id \$	$T' \rightarrow \varepsilon$
\$E'T+	+id*id \$	$E' \rightarrow +TE'$
\$E'T	id*id \$	
\$E'T'F	id*id \$	$T \rightarrow FT'$
\$E'T'id	id*id \$	$F \rightarrow id$
\$E'T'	*id \$	
\$E'T'F*	*id \$	$T' \rightarrow *FT'$
\$E'T'F	id \$	
\$E'T'id	id \$	$F \rightarrow id$
\$E'T'	\$	
\$E'	\$	$T' \rightarrow \varepsilon$
\$	\$	$E' \rightarrow \varepsilon$

$$2. S \rightarrow iEtSS' | a$$

$$S' \rightarrow Es | \varepsilon$$

$$E \rightarrow b$$

Is this grammar LL (1). Describe it.

	FIRST	FOLLOW
S	{i ,a}	[\$,b]
S'	{b, ε }	[\$,b]
E	{b}	{t}

Parsing Table:

Non-terminal	i	t	a	b	s	\$
S	$S \rightarrow iEtSS'$		$S \rightarrow a$			
S'				$S' \rightarrow \varepsilon$ $S' \rightarrow Es$		$S' \rightarrow \varepsilon$
E				$E \rightarrow b$		

$M[S', E] = 2$

So this grammar is not an LL(1) grammar. Because $M[S', E]$ has two entries.