

# Certification of Termination Proofs for Term Rewriting

A short story of a long battle...

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# Who am I?



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- (1)  $\lambda(x) \circ y \rightarrow \lambda(x \circ (1 * (y \circ 1)))$
- (2)  $(x * y) \circ z \rightarrow (x \circ z) * (y \circ z)$
- (3)  $(x \circ y) \circ z \rightarrow x \circ (y \circ z)$
- (4)  $\text{id} \circ x \rightarrow x$
- (5)  $1 \circ \text{id} \rightarrow 1$
- (6)  $\uparrow \circ \text{id} \rightarrow \uparrow$
- (7)  $1 \circ (x * y) \rightarrow x$
- (8)  $\uparrow \circ (x * y) \rightarrow y$

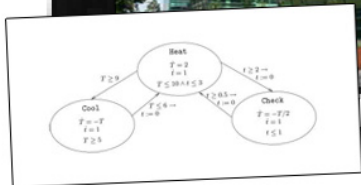
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- 1 Background: termination of term rewriting
- 2 CoLoR project: certification of termination proofs
  - Why?... motivation
  - How?... CoLoR's approach to certification
  - When?... history of the project
  - What?... overview of the content
  - Related work
  - Certified competition
- 3 Conclusions... sort of

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# Introduction to term rewriting

## Example (Quick sort)

$\text{qsort}(\text{nil}) \rightarrow \text{nil}$

$\text{qsort}(x :: xs) \rightarrow \text{append}(\text{qsort}(\text{filterLe}(x, xs)), x :: \text{qsort}(\text{filterGe}(x, xs)))$

$\text{append}(\text{nil}, l) \rightarrow l$

$\text{append}(x :: xs, l) \rightarrow x :: \text{append}(xs, l)$

$\text{filterLe}(n, \text{nil}) \rightarrow \text{nil}$

$\text{filterLe}(n, x :: xs) \rightarrow \text{filter}(\text{le}(x, n), x, \text{filterLe}(n, xs))$

$\text{filterGe}(n, \text{nil}) \rightarrow \text{nil}$

$\text{filterGe}(n, x :: xs) \rightarrow \text{filter}(\text{ge}(x, n), x, \text{filterGe}(n, xs))$

$\text{ge}(x, y) \rightarrow \text{le}(y, x)$

$\text{filter}(\text{false}, x, xs) \rightarrow xs$

$\text{le}(0, y) \rightarrow \text{true}$

$\text{filter}(\text{true}, x, xs) \rightarrow x :: xs$

$\text{le}(s(x), 0) \rightarrow \text{false}$

$\text{le}(s(x), s(y)) \rightarrow \text{le}(x, y)$

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$$\text{append}(\text{nil } l) \rightarrow l$$

## Example (Collatz conjecture)

$$\text{collatz}(s(s(x))) \rightarrow f(\text{even}(x), s(s(x)))$$

even(0)  $\rightarrow$  true

$$f(\text{true}, x) \rightarrow \text{collatz}(\text{half}(x))$$
$$\text{even}(s(0)) \rightarrow \text{false}$$
$$f(\text{false}, x) \rightarrow \text{collatz}(s(\text{triple}(x))) \quad \text{even}(s(s(x))) \rightarrow \text{even}(x)$$
$$\text{half}(0) \rightarrow 0$$
$$\text{triple}(0) \rightarrow 0$$
$$\text{half}(s(s(x))) \rightarrow s(\text{half}(x))$$
$$\text{triple}(s(x)) \rightarrow s(s(s(\text{triple}(x))))$$

80(1), 77(1)

$$\text{filter}(\text{false}, x, xs) \rightarrow xs$$
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$\text{f}(\text{true}, x) \rightarrow$

$\text{f}(\text{false}, x) \rightarrow$

### Definition

A TRS is **terminating** iff it does not admit infinite reductions.

$\text{half}(0) \rightarrow 0$

$\text{half}(s(s(x))) \rightarrow s(\text{half}(x))$

$\text{triple}(0) \rightarrow 0$

$\text{triple}(s(x)) \rightarrow s(s(s(\text{triple}(x))))$

$\text{le}(0, y) \rightarrow \text{true}$

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# Termination of rewriting

## Termination of rewriting:

- **Is undecidable.**
- Is an important topic in term rewriting.
- Many methods exist and new ones are constantly being developed.
- Recently the emphasis is on automation.
- There exists a number of tools for proving termination.
- Stimulated by the termination competition.
- Tools (and proofs that they produce) are getting more and more complex, so reliability is an issue (tools disqualifications in the competition).
- In 2007 a new category of certified termination introduced in the competition.

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`http://color.loria.fr`

CoLoR: Coq Library on Rewriting and Termination.

Goal: certification of termination proofs produced by various termination provers.

- Increasing reliability of termination provers.
- Common proof format for termination provers:
  - common tools (proof presentation, manipulation, ...),
  - control language for provers (integration of tools)
- Extension of proof assistance kernels.

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## How to certify termination results?

- Possibility: certification of tools source code.  
⇒ difficult, tool dependent, extra work with every change, ...
- CoLoR's approach:
  - TPG: common format for termination proofs.
  - Tools output proofs in TPG format.
  - CoLoR: a Coq library of results on termination.
  - Rainbow: a tool for translation from proofs in TPG format to Coq proofs, using results from CoLoR.

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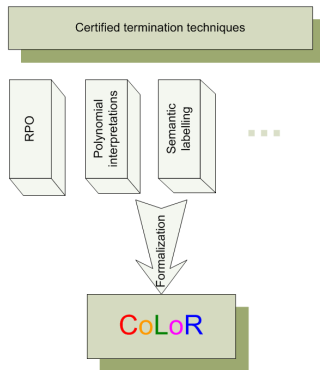
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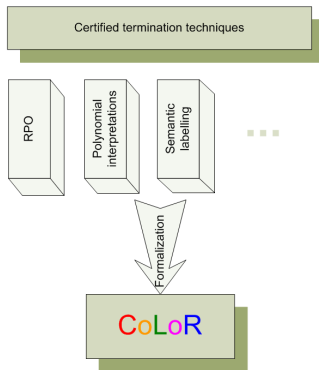
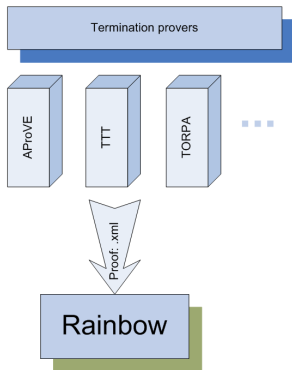
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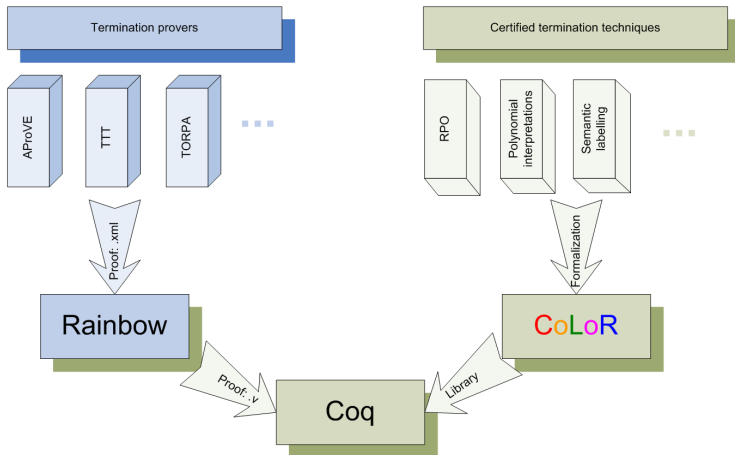
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- First release March 2005
- First certified proofs July 2006
- First certification workshop May 2007
- First certified competition June 2007

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- polynomial interpretations [Hinderer]
- multiset ordering [Koprowski]
- recursive path ordering [Coupet-Grimal, Delobel]
- higher-order recursive path ordering [Koprowski]
- dependency graph cycles [Blanqui]
- matrix interpretations [Koprowski, Zantema]
- arctic interpretations [Koprowski, Waldmann]

- Transformation techniques:

- dependency pairs [Blanqui]
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- semi-rings [Koprowski, Zantema]
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- lists, vectors, relations, etc.

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- simply typed lambda-terms
- varyadic terms
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In total:

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- $\approx 1.000$  definitions and  $\approx 3.000$  lemmas.
- Only 20% of that is the code for actual termination methods!

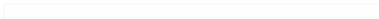

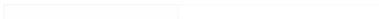
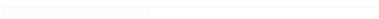
Size comparison with other libraries:

- |                        |                                   |
|------------------------|-----------------------------------|
| • Coq standard library | <div><div></div></div>            |
| • C-CoRN               | <div><div></div><div></div></div> |
| • COMPCERT             | <div><div></div><div></div></div> |
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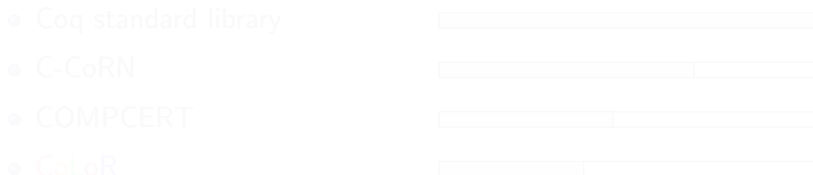
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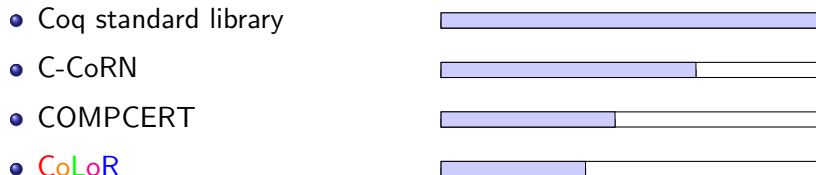
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Proof assistant: Coq
- A3PAT project  
Authors: Contejean, ...  
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# Certified competition

- In the 2007 termination competition a new “certified” category was introduced.
- Participants 2007 (975 problems):

- TPA + CoLoR
- C/ME + A3PAT
- $T_1T_2$  + CoLoR

- Participants 2008 (1391 problems):

- AProVE + C/ME + A3PAT
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- AProVE + C/ME



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• <i>AProVE</i> ( <i>non-certified</i> )		723
• TPA + CoLoR		354
• CiME + A3PAT		317
• TT <sub>2</sub> + CoLoR		289

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- AProVE + CoLoR
- AProVE + A3PAT
- CiME3 + A3PAT
- Matchbox + CoLoR


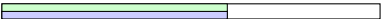

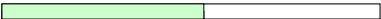
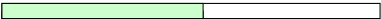

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• AProVE + CoLoR		580
• AProVE + A3PAT		532
• CiME3 + A3PAT		531
• Matchbox + CoLoR		458

- 1 Background: termination of term rewriting
- 2 CoLoR project: certification of termination proofs
  - Why?... motivation
  - How?... CoLoR's approach to certification
  - When?... history of the project
  - What?... overview of the content
  - Related work
  - Certified competition
- 3 Conclusions... sort of

- Lesson 1

If it is possible do (involved) computations/reasoning in an unsafe setting and verify the results in Coq a posteriori.

That requires some notion of a certificate.

Proof search is usually much more complex than proof verification.

We see that even in theorem provers — proof checking VS proof searching.

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It is not unusual for software projects to be behind schedule / run out of budget.

It is even more so for Coq projects.

Why?

- algorithm  $\mapsto$  program
- paper proof  $\mapsto$  formal proof in Coq
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- Lesson 3

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You want to make the right choices. You really do.

Because that will have a tremendous impact on the reasoning about those definitions that you are going to do for long hours afterwards.

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Thank you for your attention.