

An Introduction to Cryptography

Part 3: Quantum & Post-Quantum Crypto

Anthony Kosednar

My Background

- Education
 - Number Theory, Field Theory, Cryptography and Statistics @ ASU
 - Aerospace Engineering - Astronautics @ ASU
 - DHS Cybersecurity Training for Industrial Control Systems – SCADA, DCS, PLC, Embedded Control
 - GIAC Exploit Researcher and Advanced Penetration Tester - GXPN
- Career Highlights
 - Built a natural air-cooled datacenter in under 6 months
 - Lead encryption for a large fortune 50 financial institution
 - Prevented cyber attacks for large events (Arizona Cardinals, Arizona Board of Tourism, NFL, Super Bowl XLIX, Fiesta Bowl)

Quantum Computers – Common Fears

Is the security of the internet at risk?

Will all financial institutions be at risk?

Will I never win the lottery?



I'm here to help!

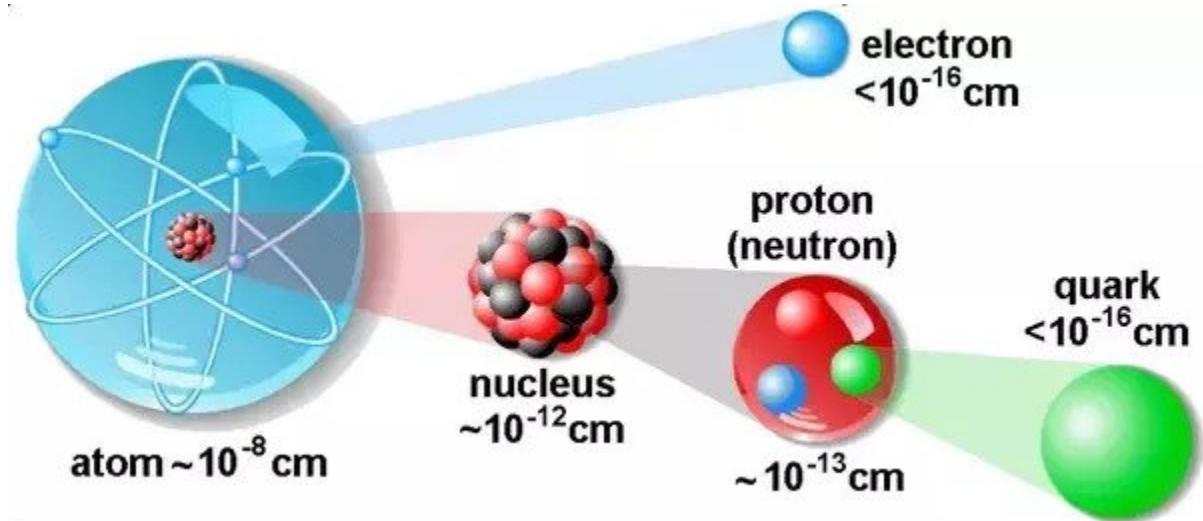


First some background

A Simple Overview Quantum Mechanics

What is Quantum Mechanics?

- Describes nature at the smallest scales of energy (atoms & subatomic particles)



Quantum Mechanics - Popular View

- Quantum Mechanics is weird!
- Quantum objects can both be waves and particles: ***wave-particle duality***
- Quantum objects can be in more than one state at once: ***superposition***
- You can't simultaneously know exactly two properties of a quantum object: ***Heisenberg's uncertainty principle***
- Quantum objects can affect one another instantly over huge distances ("spooky action at a distance"): ***entanglement***
- You can't measure anything without disturbing it, and so quantum mechanics is subjective

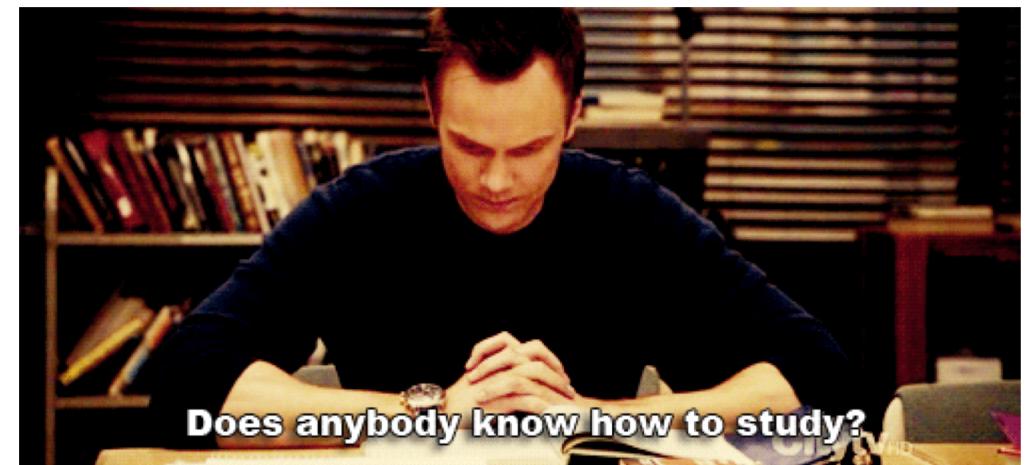
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Not Exactly Right

Why is Quantum Mechanics Hard?

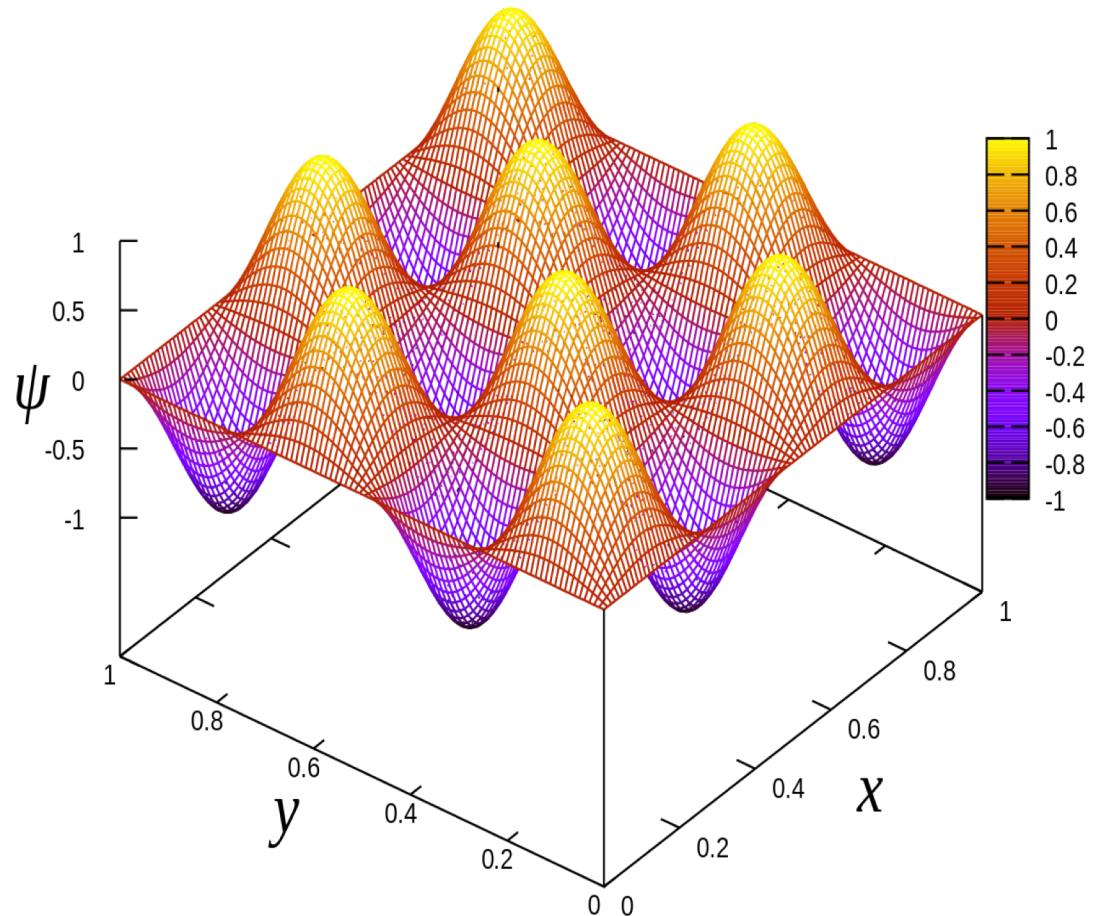
- Traditional Newtonian Mechanics – (what you are used to):
 - It's obvious
 - We don't need to ask questions around our definitions
 - We have an idea with a position, path and force are
- Quantum Mechanics:
 - Non-definitive
 - Influenced by us measuring
 - We don't know all the answers



Wave-Particle Duality

Schrodinger's Wave Function

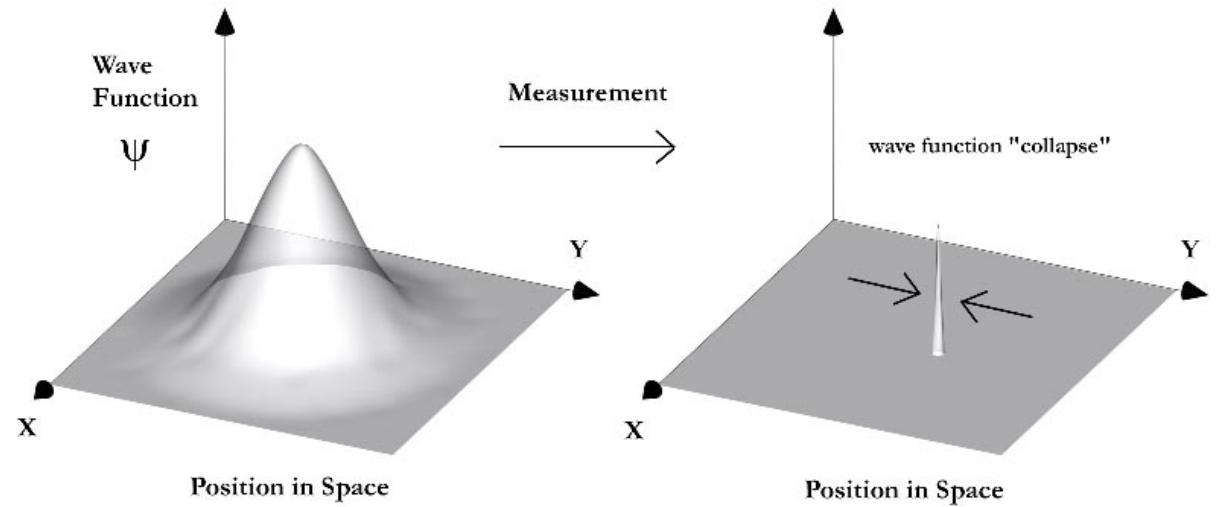
- Schrodinger created function to show particles might act as a wave
- Doesn't tell us the protectory of a particle but rather where we might find a particle; What properties it might have.
- Often said the particle is spread out all over the place
- However is a math functions to represent the possible measurement outcomes (and probability to make that result).
- Tells us the chance of where we find it if we look



Collapse of the Wave Function

- When we make a measurement suddenly it collapses to a single position because now we know where it is
- No way of describing this in Quantum Mechanics
- Quantum Mechanics says:
 - There were different probabilities that a measurement would reveal a particle at a particular place
 - After measuring you know it's "THERE"

The Copenhagen Interpretation:



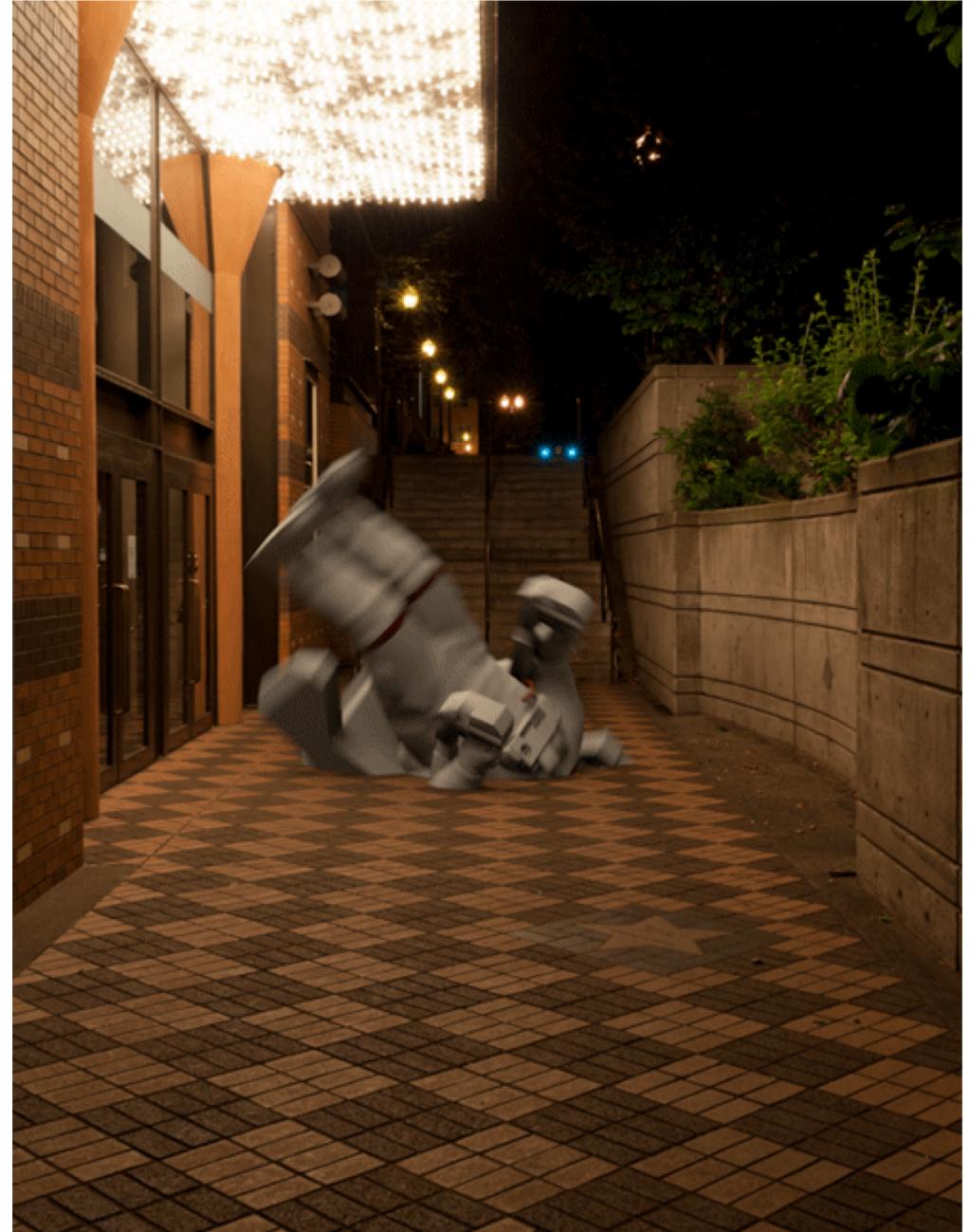
Spin!

Spin Up Spin Down

(1)



(0)



Spin!

Spin Up Spin Down

(1)

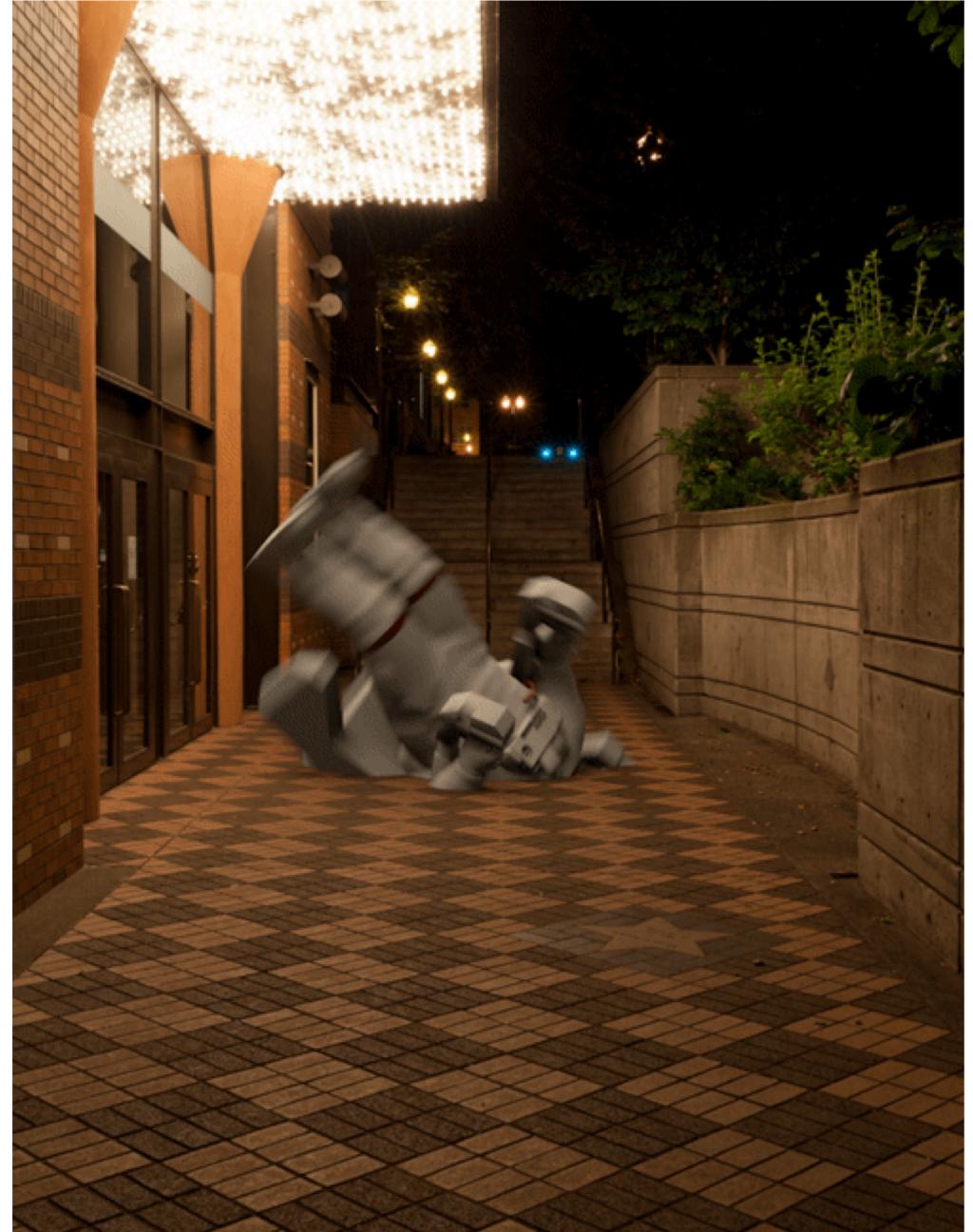


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Basis of quantum computers because
we can record binary data!

Act as quantum bits or QBits!



Spin!

Spin Up Spin Down

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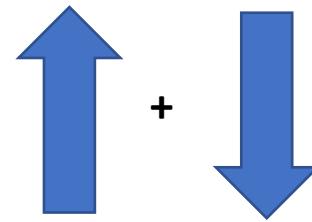


Basis of quantum computers because
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However it is said it can be often in both states at once

Superposition



Spin!

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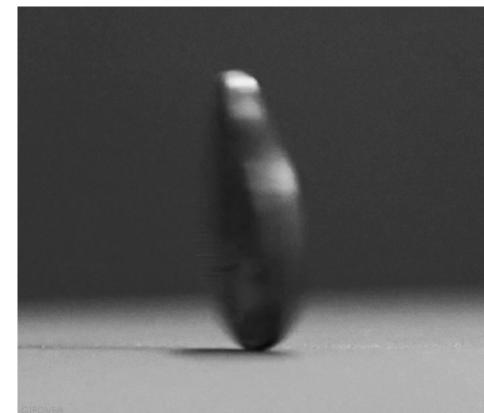
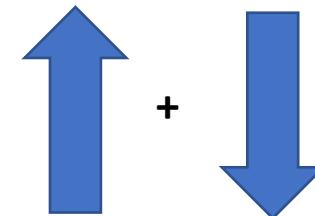


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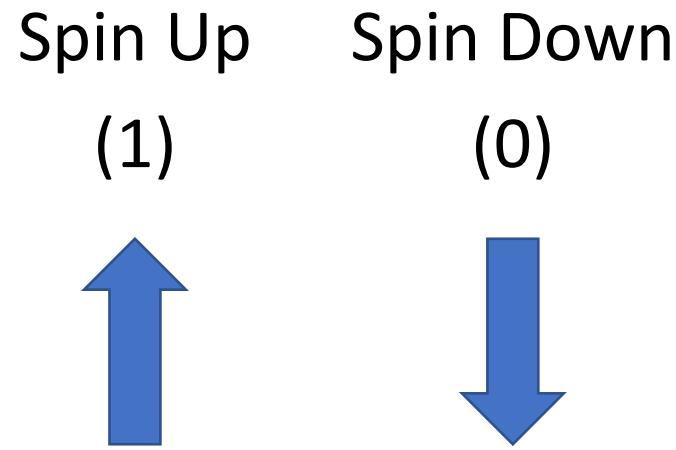
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Spin!



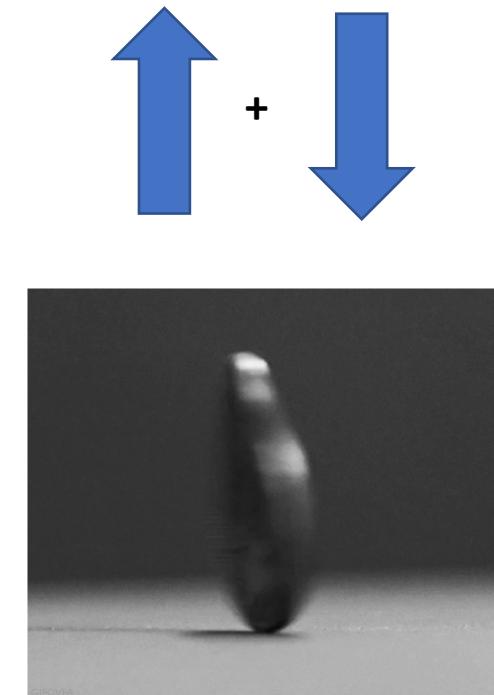
Basis of quantum computers because
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Act as quantum bits or QBits!

Wave function only tells us what to expect when
we make a measurement

~~WRONG It is said it can be often in both states at once~~

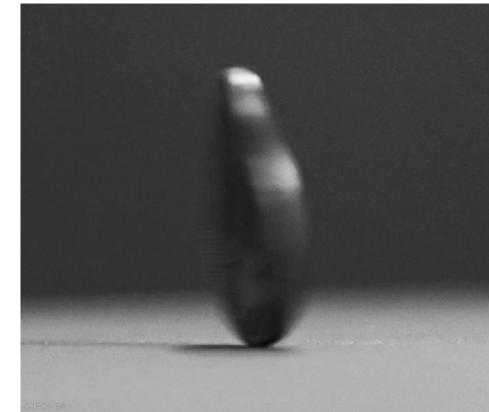
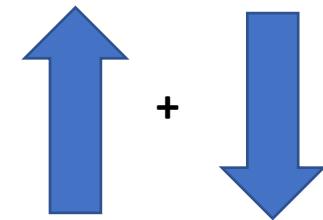
Superposition



Spin!

- What the wave function says:
 - A measurement might give us an up or a down spin
- Before measurement?
 - Quantum mechanics doesn't really tell us

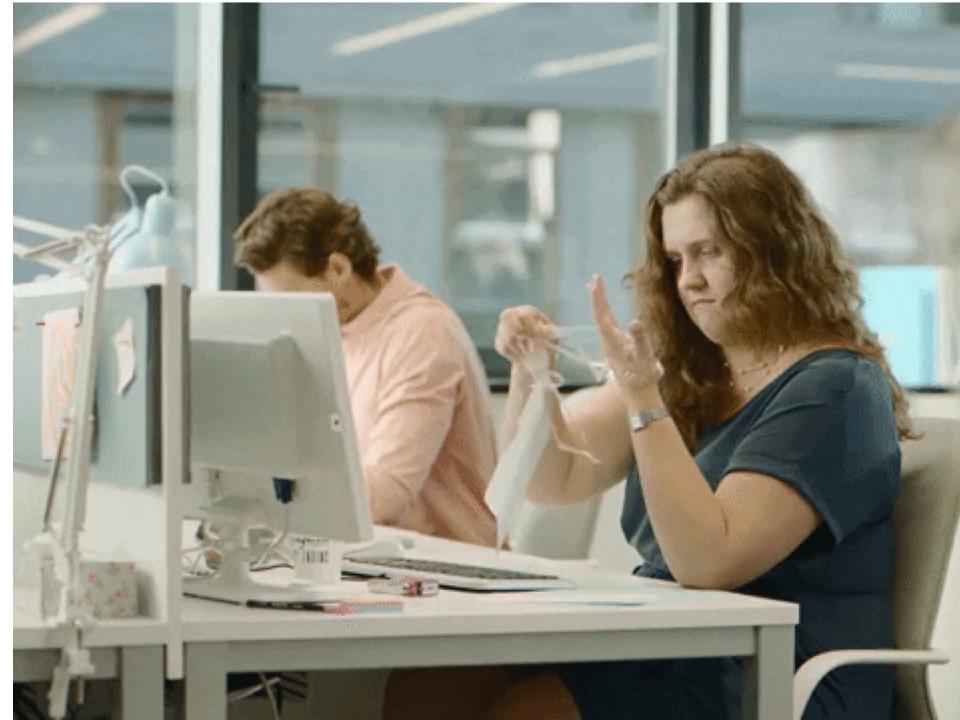
Superposition





Entanglement

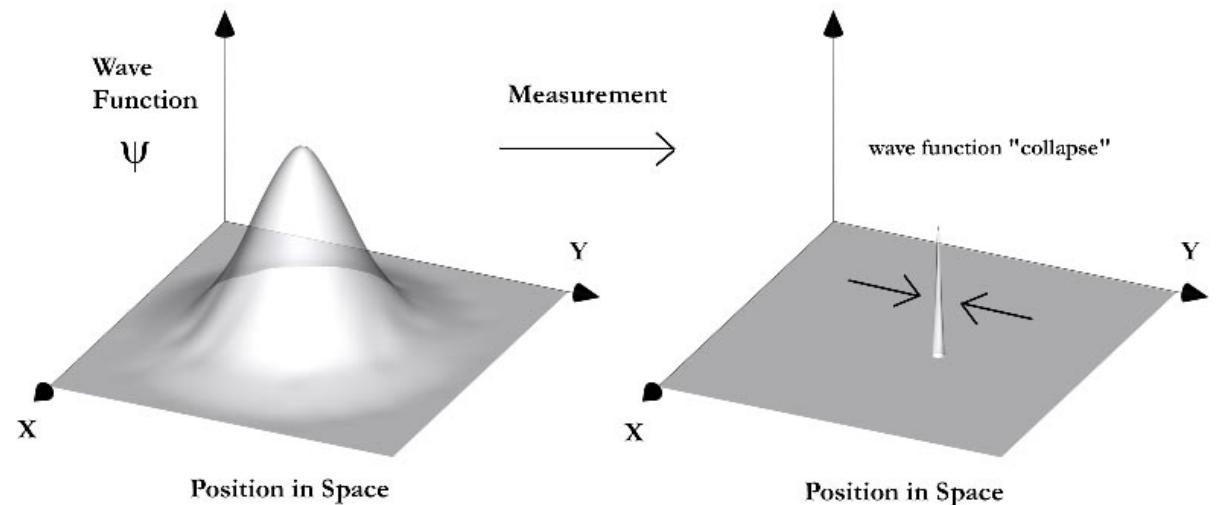
- Two quantum objects that collide become entangled or linked
- When measured you'd be able to predict the measurement of the other
- Quantum objects that are entangled do not have properties that are located in the object
 - Example: The blackness of two quantum entangled objects are not local to their locality but quantumly existing



More Quantum Properties

- As a quantum object interacts with its environment, it becomes harder to see its super position in the environment. The quantum-ness gets washed away – decoherence.
- After this, they start to become more like classical objects

The Copenhagen Interpretation:



Quantum Mechanics – Reality

- Not – here it is a particle, there it is a wave
- But Rather - if measured, a quantum object behaves in a manner we associate with particles; but if we measure it, it behaves as if it's a wave
- Not - the particle is in two states at once
- But – if measured, we will detect this state with probability X, and that state with probability Y

Classical vs Quantum Computing



A Note on Mindset

- A light bulb is not a more powerful version of the candle. They execute similar functions but are based on different technologies. They execute their task in different ways.

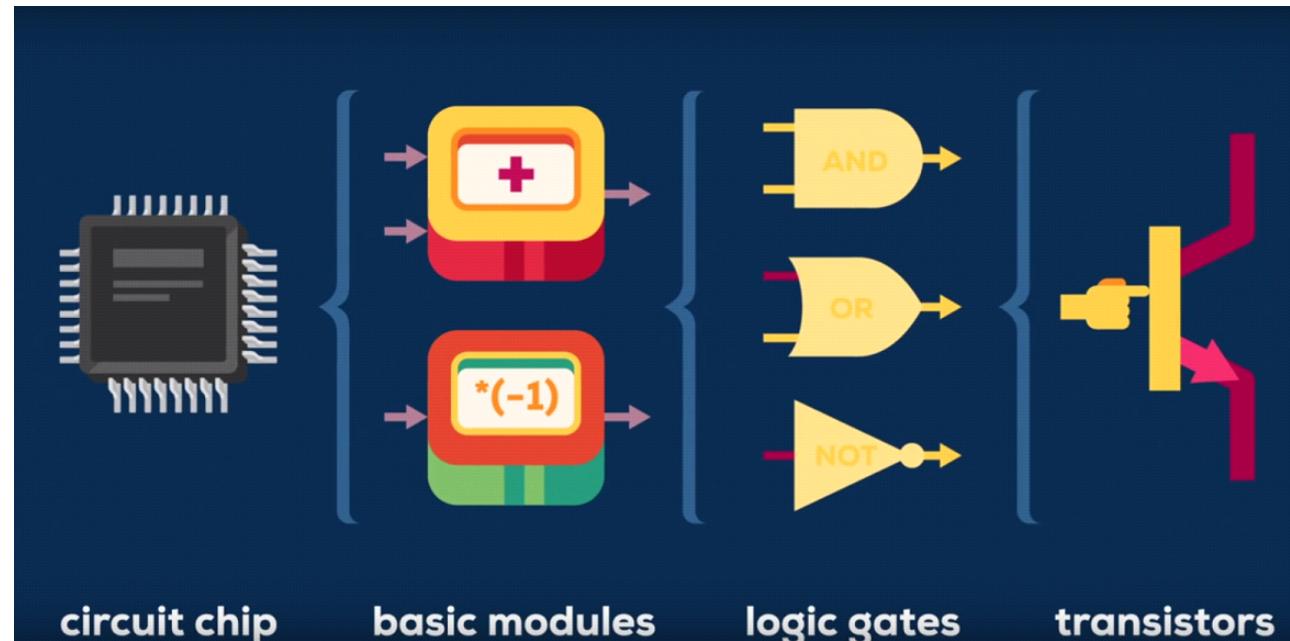
A Note on Mindset

- A light bulb is not a more powerful version of the candle. They execute similar functions but are based on different technologies. They execute their task in different ways.
- Similarly....a quantum computer is not a more powerful version of a classical computer!



How do classical computers work?

- Bits (16 combinations)
 - 1 or 0 - on or off
- Transistors
- Logic Gates
- Modules



Bits vs Qubits

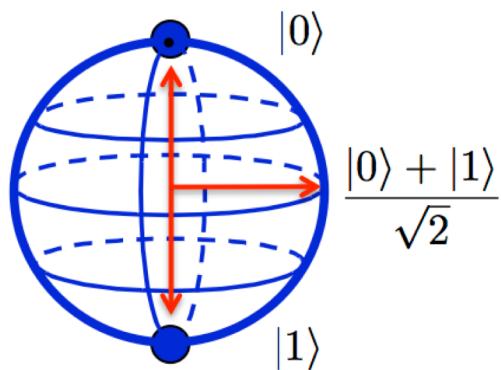
Bits:

- Computes by manipulating bits with the help of logical gates
- Has a memory made up of bits where each bit holds either a one, or a zero

0

1

Classical Bit



Qubit

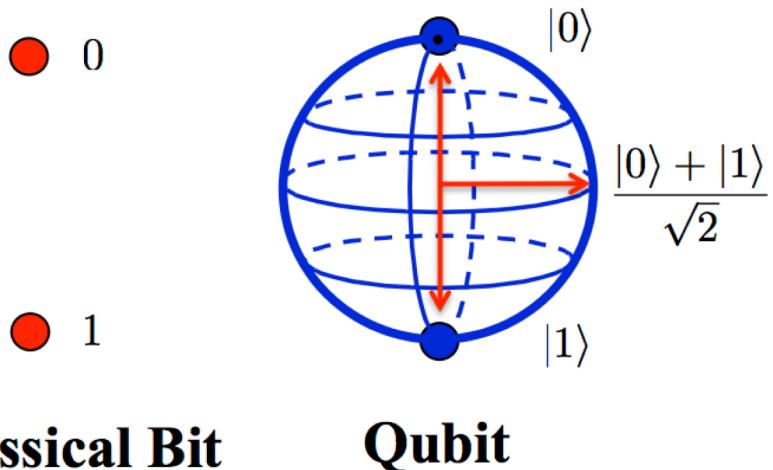
Qubits:

- Device computes by manipulating bits with the help of quantum logic gates.
- A qubit when can hold a one, zero, or a super position of those until measured
- Once measured it's non-quantum

Bits vs Qubits

Bits:

- Computes by manipulating bits with the help of logical gates
- Has a memory made up of bits where each bit holds either a one or a zero



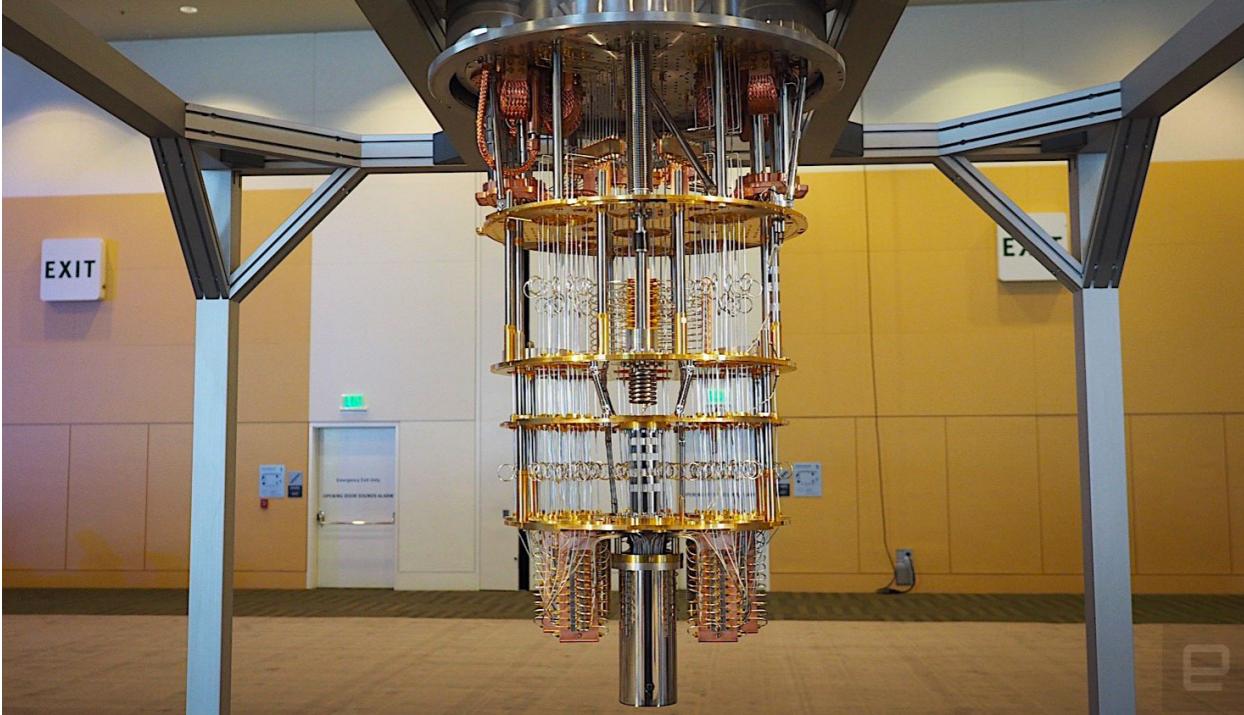
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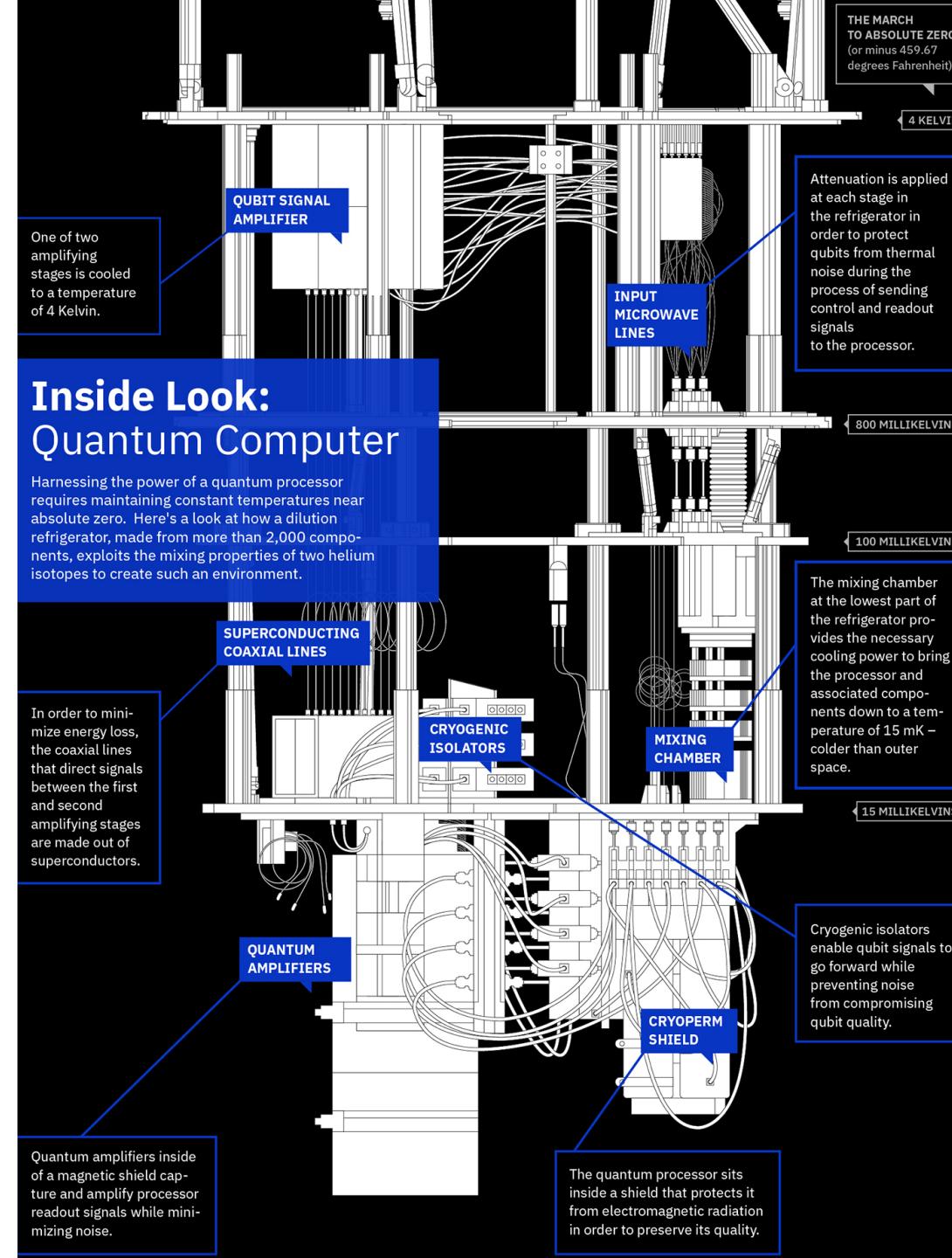
- Device computes by manipulating bits with the help of quantum logic gates.
- A qubit can hold a one, zero, or a superposition of those until measured
- Once measured it's non-quantum
- **We can only keep it quantum for so long!**

Quantum Computer!

- Has a physical qubit that can be put into the spin states
 - Typically atoms, ions, or other super conducting materials
 - Controlled and measured via microwave pulses
- Very sensitive to noise and environmental affects
- Can only keep information quantum for so long

Quantum Computer!





Breaking Cryptography with a Quantum Computer

Shor's Algorithm and Breaking RSA Encryption
(Asymmetric Cryptography)

Review!: RSA

- Widely used:
 - Digital signatures
 - Key distribution
- Works by taking two large prime numbers, multiplying them together and combining them with a pair of exponents
 - Ex: $15 = 3 * 5$
- Security depends on the difficulty of finding those prime factor

Lh_f Fincta&

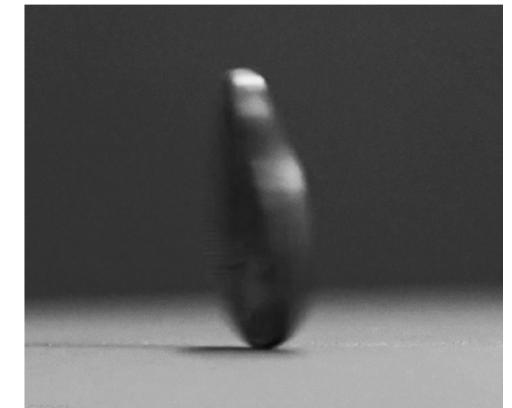
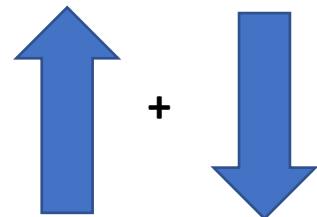
Traditional Breaking of RSA

- How – the best methods guess the factors, check, and try again!
- Super slow
 - Even with the fast ways to make good guesses!
- RSA-768 takes over 2,000 years of computer processing time break
 - 1 month in a proper computing cluster

The Quantum Computing Way

- Easy due to
 - Quantum super positioning
 - Quantum interference
- Uses Shor's Algorithm with some tricks!

Superposition



Shor's Algorithm

- Start with a guess of what might share a factor but probably doesn't
- Algorithm turns it into a better guess that might share a factor!
- Nothing quantum about this on the surface
- Can do it on a classical computer
 - However Turning a bad guess to a good one takes a long time on normal computers!
- Super fast on quantum computers



How it works?

To find a factor of N we don't have to guess a factor of N!

1. Starts with a big number, N, that we need to find the factors of to break into encrypted data
2. We make a guess, g, that is some number less than N
 - Doesn't need to be a pure factor of N but a number that shares some factors
 - Ex: 4 isn't a factor of 6 but shares some factors with it. $4 = 2 \times 2$, $6 = 2 \times 3$
3. We can guess numbers that share a factor - because of Euclid's Algorithm
 - Hard though because it is extremely unlikely a guess with share a factor

We have a trick!

- Let's convert a crappy guess into a pair of great guesses
 - $g \rightarrow g^{\frac{P}{2}} \mp 1$
- For any pair of whole numbers that don't share a factor if you multiply one by it's self enough times you'll eventually get:
 - $A, B \rightarrow A * A * A \dots * A = \text{something} * B + 1$
 - $A^P = m * B + 1$
- Ex:
 - 7, 15
 - $7^4 = 160 * 15 + 1$

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 - $g \rightarrow g^{\frac{P}{2}} \mp 1$
- For any pair of whole numbers that don't share a factor if you multiply one by it's self enough times you'll eventually get:
 - N, g
 - $g, N \rightarrow g * g * g \dots * g = \text{something} * N + 1$
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- Ex:
 - 7, 15
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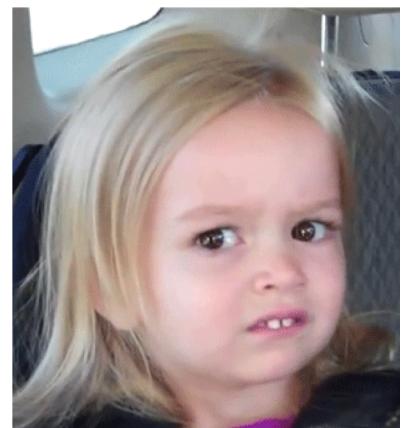
- Let's rearrange the equation
 - N, g
 - $g, N \rightarrow g * g * g \dots * g = \text{something} * N + 1$
 - $g^P = m * N + 1$
 - $g^P - 1 = m * N$
 - $(g^{\frac{P}{2}} + 1) * (g^{\frac{P}{2}} - 1) = m * N$
- Simplified:
 1. Take crappy guess
 2. Multiply itself P/2 times
 3. Then add or subtract 1

We have a trick!

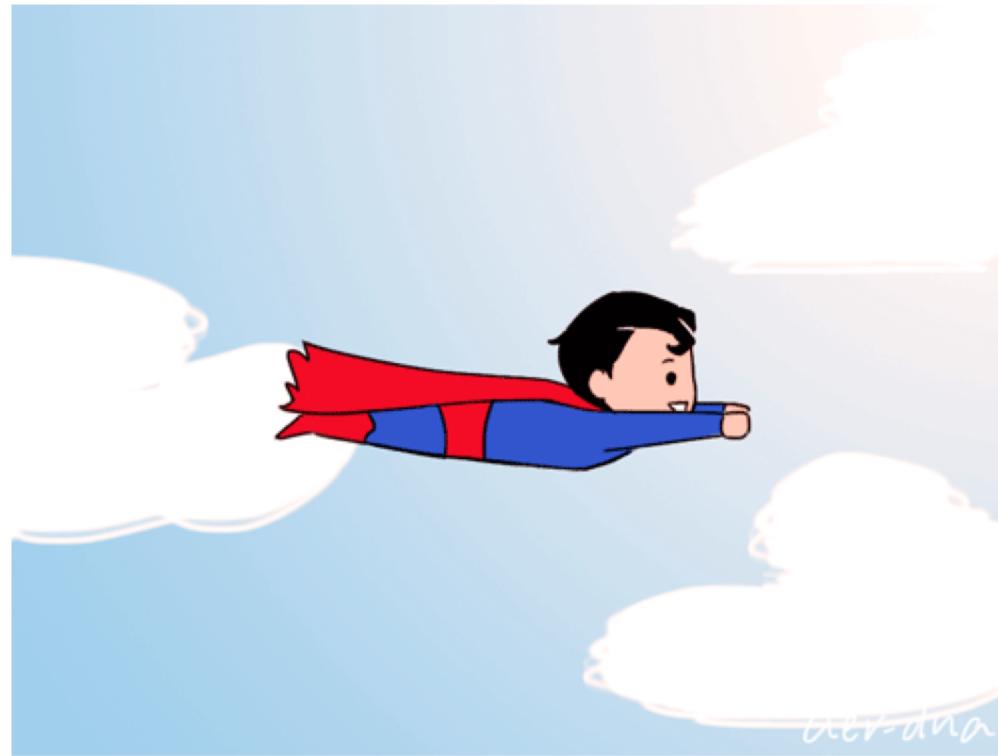
- Our equation
 - N, g
 - $\left(g^{\frac{P}{2}} + 1\right) * \left(g^{\frac{P}{2}} - 1\right) = m * N$
- Note factors on either side may be a multiple of the factor
- We can use Euclid's Algorithm to solve that
- Then the encryption is broken!

Why is this hard for classical computers?

- Problems:
 - P might be odd
 - $(g^{\frac{P}{2}} \pm 1)$ might not be a multiple of N
- However – these two problems only exist 37.5% of the time
- But....finding p takes a long time for a classical computer to guess
 - Longer than brute forcing the actual prime



Let's do it the Quantum Way!



Let's do it the Quantum Way!

- Our equation
 - Ng
 - $g^P = m * N + 1$
 - $\left(g^{\frac{P}{2}} + 1\right) * \left(g^{\frac{P}{2}} - 1\right) = m * N$
- Setup a quantum superposition to represent all possible P's for our g
 - We need to keep track of all P's and all g^P 's
 - $|P, g^P\rangle$
- Then superposition of how much bigger than $m*N$ it is.
 - $|P, +r\rangle$
- Note – superpositions allow us to simultaneously represent all possibilities

Let's do it the Quantum Way!

- Problem – We can't just measure to get the outcome because superpositions are a probability. We would be a random element of the superposition steps.
- We have to get all the bad answers to cancel out while in the quantum superposition to leave the right answer.
 - $|P, +1 >$

Some relationships to help...

- Our guess, g , to some random power, x , is most likely some multiple of N plus some other number
 - $g^x = m * N + r$
- But something cool....
 - $g^{x+P} = m_1 * N + r$
 - $g^{x+2P} = m_2 * N + r$
- So....power P has a repeating property where the “r” stays the same
- We can take advantage of a superposition of all possible powers to take advantage of this!
- Those numbers in the superposition will repeat with a period of P .
 - Frequency $\rightarrow f = \frac{1}{P}$
 - If we find the frequency we break the encryption!

Quantum Fourier Transform!

- Best thing to find the frequency of things is a Fourier Transform
- We use the quantum version of the formula to make all the frequencies that don't exist to interfere and leave the single quantum state
 - $|\frac{1}{P}\rangle$
- How does Quantum Fourier Transform work?
 - A single number inputted will represent a superpositions of numbers with weights - that looks like a sinewave with the frequency of the inputted number
 - When you put in a superposition of numbers you get out a superposition of superpositions where the sinewaves subtract and cancel out.
 - This then leaves you with the single quantum state $|\frac{1}{P}\rangle$
- Now finally...as long as P is even
 - We take out guess, g, to compute $(g^{\frac{P}{2}} \pm 1)$
 - As long as $(g^{\frac{P}{2}} \pm 1)$ is not a multiple of N then that shares some factor with N
 - Then using Euclid's algorithm we can quickly find N $\rightarrow N = a * b$
 - We now have it decrypted!!

The reality of the risk

- We only have computers with enough qubits for unknown primes in the thousands (less than 50bits).
- Typical RSA encryption has around 2048 to 4096 bits of encryption
....Thus we don't expect this attack to be viable for ~10 years.



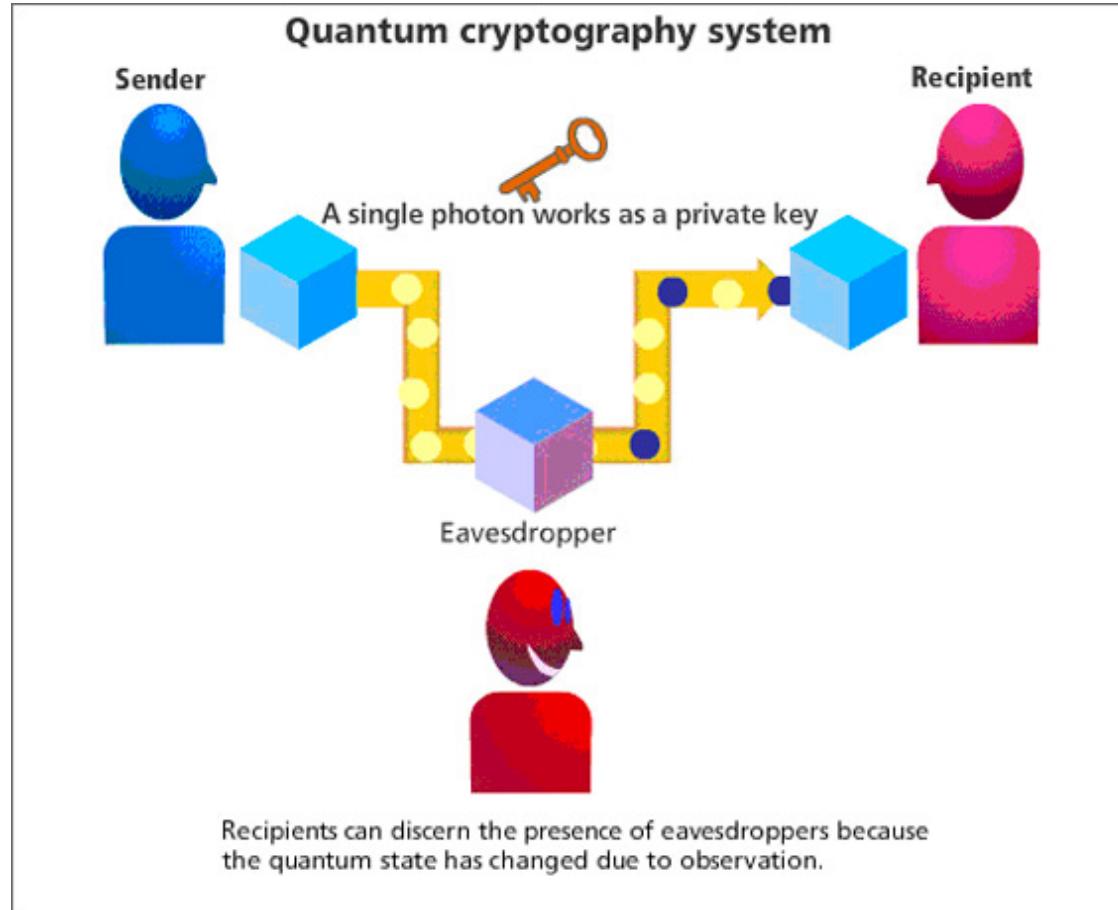
Want to try this on your own?

- Signup for computing time on IBM Q
 - <https://quantum-computing.ibm.com>
- Run the implementation of shor's algorithm for quantum computers
 - <https://github.com/ttlion/ShorAlgQiskit> – python
- Remember it only works for small numbers (ex: 15) because the amount of qubits that quantum computer has but it's still cool!



Quantum Computing Can Be
Used to Improve Cryptographic
Systems!

Using Quantum Mechanics to Improve Cryptography

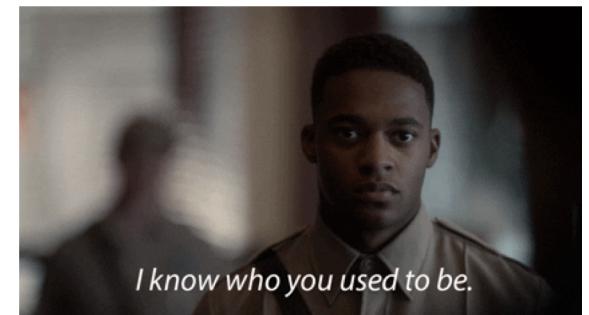


Post-Quantum Cryptography

Quantum Resistant Algorithms

Post-Quantum Cryptography

- We have some traditional cryptography methods still
 - Symmetric key quantum resistance
 - Hash-based cryptography
- And...we also have new more symmetric like algorithms
 - Lattice-based cryptography
 - Multivariate cryptography
 - Supersingular elliptic curve isogeny cryptography



I know who you used to be.

Post-Quantum Cryptography

Symmetric key quantum resistance

- Provided one uses sufficiently large key sizes, the symmetric key cryptographic systems like AES are already resistant to attack by a quantum computer!
- As a general rule, for 128 bits of security in a symmetric-key-based system, one can safely use key sizes of 256 bits.
- Grover's algorithm allows brute-forcing a key on quantum computers but....
 - The time to brute force a n-bit key is speed up only to $O(2^{\frac{n}{2}})$ time instead of $O(2^n)$ with traditional computers
 - A speed up but not as significant as the asymmetric attacks!

Reminder – Pros vs Cons of Symmetric Crypto

- Pros
 - Fast and simple computations
 - Effective confidentiality
 - Strong algorithms that are normally only cracked through “brute force”
 - Good for bulk data
- Cons
 - Difficult to securely store and share the secret key
 - Often lacks authentication and data integrity
 - Data retention cases require key retention which can result in many secret keys

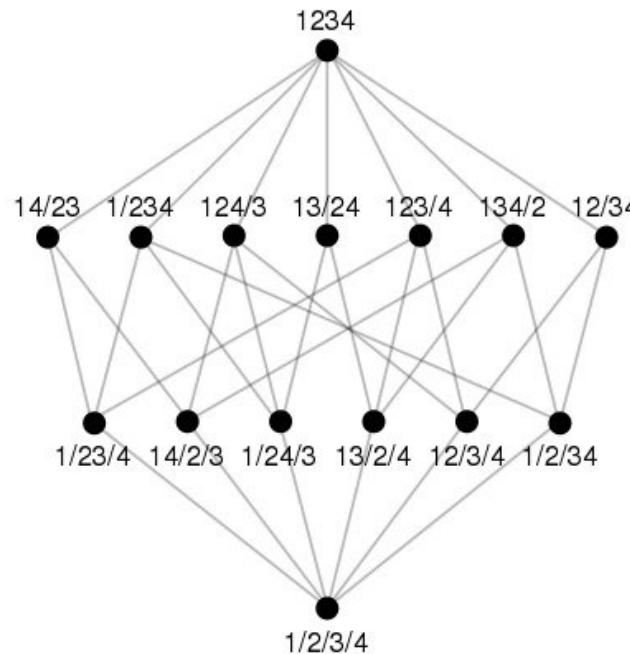
Post-Quantum Cryptography

Hash-based cryptography

- Single direction functions like traditional hash encryption
- Limited to digital signatures

Post-Quantum Cryptography

- Lattice-based cryptography
 - A partially ordered set of numbers in which every two elements have a unique upper bound and greatest lower bound.



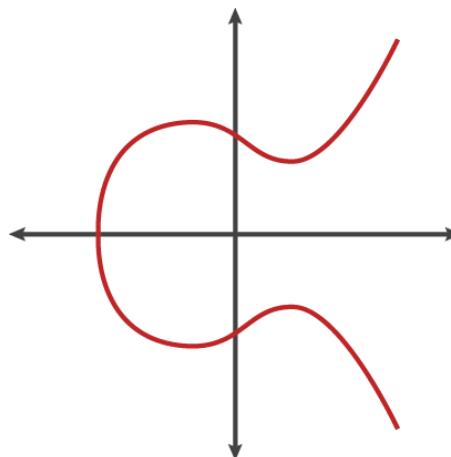
Post-Quantum Cryptography

- Multivariate cryptography
 - Asymmetric cryptography using multivariate polynomials (function such as $x^2 - 4x + 7$) over a finite field

Post-Quantum Cryptography

Supersingular elliptic curve isogeny cryptography

- Relies on the properties of supersingular elliptic curves and supersingular isogeny graphs to create a Diffie-Hellman replacement with forward secrecy.
- Supersingular elliptic curves form a certain class of elliptic curves over a field of characteristic $p > 0$ with unusually large endomorphism rings.
- An elliptic curve is a plane algebraic curve defined by an equation of the form



Post-Quantum Cryptography

- NIST is currently reviewing the new algorithms for standardization
 - <https://csrc.nist.gov/Projects/Post-Quantum-Cryptography>
 - Expected draft standards available in 2022-2024
- TLS 1.3 already supports quantum resistant algorithms that are draft standards by NIST
- Cloudflare's edge implementation of TLS 1.3 supports:
 - One based on lattices (NTRU-HRSS: big key, low CPU)
 - One based on isogenies (SIDH: small key, high CPU)
- Evolving into production use but...
 - A lot of work is being done still in this area
 - Some flaws may not have been discovered yet

The Risk - Cryptographic Recommendations

Algorithm Type	Data Protection for < 10 years	Data Protection for 10+ years
Symmetric Use cases	AES-256	AES-256
Asymmetric Use cases	RSA-4096	SIDH or NTRU-HRSS algorithms

Remember! Cryptography is a deterrent. It is not a guarantee that someone can't get it. Typically cryptography makes it not worth someone's time to try to break your encryption.

Think locks and bars on the windows/doors of a house.

In summary...

We are not doomed 😊



Questions?

Anthony.Kosednar@gmail.com

Some Useful Resources

- <https://github.com/ttlion/ShorAlgQiskit>
- <https://www.gcppodcast.com/post/episode-123-post-quantum-cryptography-with-nick-sullivan-and-adam-langley/>
- <https://blog.cloudflare.com/disruptive-cryptography/>
- <http://www.ams.org/journals/notices/201605/rnoti-p508.pdf>
- <https://csrc.nist.gov/Projects/Post-Quantum-Cryptography/Workshops-and-Timeline>
- <http://hdc.amongbytes.com/post/201810-baby-steps-to-pq-https-server/>
- <http://hdc.amongbytes.com/post/20190130-pqc-round2/>
- <https://www.youtube.com/watch?v=q7v5NtV8v6I&feature=youtu.be>
- <https://www.youtube.com/watch?v=SCdaYAFx7hw>
- https://twitter.com/_henrycase
- <https://twitter.com/grittygrease/status/1086339192416743424>
- <https://twitter.com/grittygrease?lang=en>
- <https://blog.cloudflare.com/tls-1-3-explained-by-the-cloudflare-crypto-team-at-33c3/>
- <https://csrc.nist.gov/Projects/Post-Quantum-Cryptography/Round-1-Submissions>
- <https://blog.cloudflare.com/sidh-go/>
- <http://www.ams.org/journals/notices/201605/rnoti-p508.pdf>