

CSC1135 Secure Programming Lab 02

Darragh O'Brien B.Sc. Ph.D. FHEA
School of Computing
Dublin City University

Overview

In this lab we will:

- Introduce gdb
- Explore process layout in memory
- Examine stack layout in memory
- Disassemble some C and annotate the resultant assembly
- Reverse engineer some C from assembly
- Rewrite a return address to control program execution

A. Process layout in memory

Verify you are using gcc version 3.4:

```
$ gcc -v
```

Download a copy of `program.c`. Build an executable as follows:

```
$ gcc -g -o program program.c
```

Sketch an outline of the program's process address space. It should stretch from 0x00000000 at the bottom to 0xFFFFFFFF at the top. Use gdb to locate (as best you can since it may not be possible to be precise in all cases) the following and mark them in your diagram:

- .text section start and end
- heap base address (approximately)
- ptr (where ptr points to)
- &ptr (where ptr lives)
- argc (where argc lives)
- argv (where argv lives)
- argv[0] (where argv points to)
- program entry point
- .rodata section start and end
- I am a C program (where the string lives)
- big_array (where big_array lives)
- .data section start and end
- stack start address (approximately)
- buffer (where buffer lives)
- malloc (where malloc is defined)
- strcpy (where strcpy is defined)
- main (where main is defined)

- `sumup` (where `sumup` is defined)
- `.bss` section start and end
- `swap` (where `swap` is defined)
- `local` (where `local` lives)

Here is an example [gdb session](#) which you might find useful in getting started. For extra help, here are some [gdb notes](#), a [gdb tutorial](#) and a [gdb quick reference](#). Or you could try the help available in `gdb` itself.

B. Stack layout in memory

Set a breakpoint on the `swap` function in your program. Draw a diagram of the stack at the time the breakpoint is encountered. Your diagram should include arguments passed by the caller, the return address, the caller's frame pointer, the current frame pointer, local variables and any saved registers.

C. Disassembling code

Disassemble the `swap` function and annotate each line with a description of what it does.

D. Reverse engineering C from assembly

Translate each of the two assembly routines [here](#) into equivalent C functions. Each function returns an integer. Once you've made an attempt you can compare your answers to [mine](#).

E. Rewriting a return address

Rather than having control return to the `main` function, add C code to the `swap` function such that the `foo` function is executed when `swap` returns.

Hint: Declare a new pointer variable in `swap` and set it pointing to some other local variable; add an offset to it in order to have it point at the return address on the stack. Overwrite the return address with the address of the `foo` function. Buffer overflow attacks work in a similar way. Once you've made an attempt you can compare your approach to [mine](#).