



Formal Model Driven Engineering (FM & MDE)

Habilitation à Diriger des Recherches



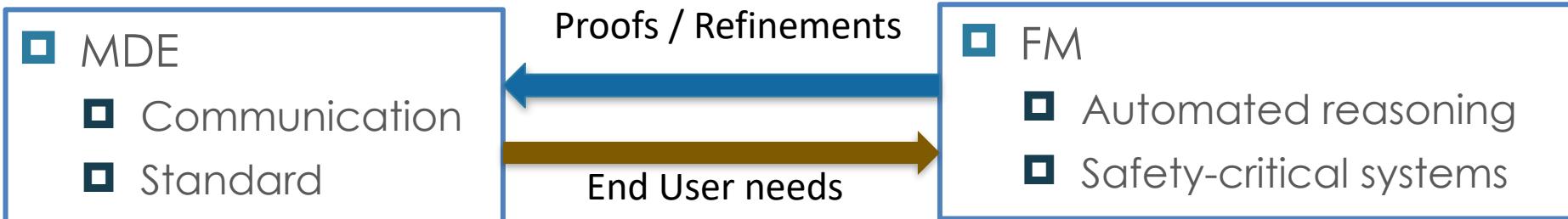
Akram Idani

26 Mai 2023



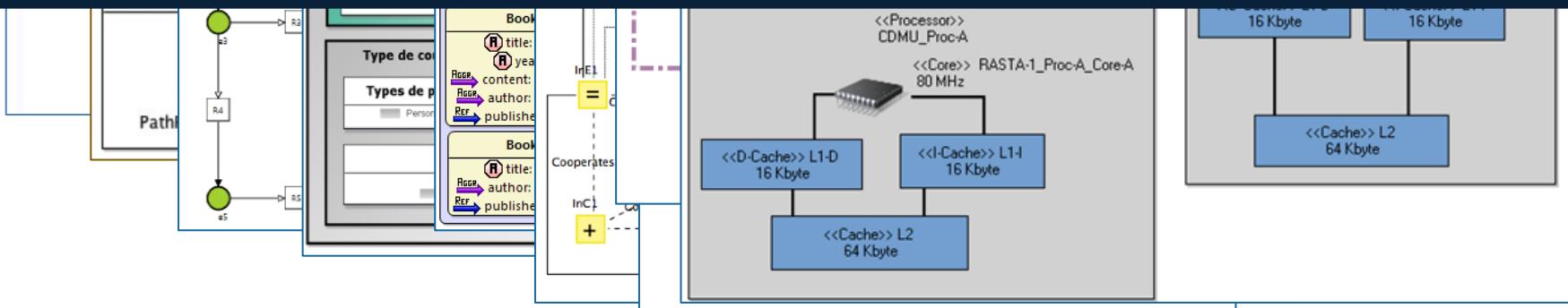
“It is easier to perceive error than to find truth, for the former lies on the surface and is easily seen, while the latter lies in the depth, where few are willing to search for it.” – Johann Wolfgang von Goethe

Overview



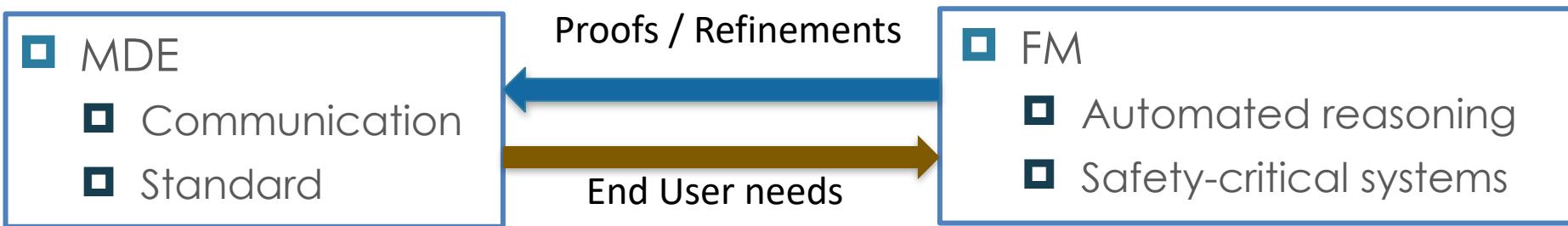
« The picture changes dramatically for **safety-critical**, high-assurance software. Here, validation by testing reaches its limits and needs to be **complemented or even replaced** by the use of formal methods such as model checking, static analysis, and program proof. »

Leroy, X.: Formal verification of a realistic compiler. Communications of the ACM (2009)



“It is easier to perceive error than to find truth, for the former lies on the surface and is easily seen, while the latter lies in the depth, where few are willing to search for it.” – Johann Wolfgang von Goethe

Overview



« [...] the learning curve of formal methods is steep, whereas the learning curve for drawing diagrams on the black board is very low. »

Andova, S. et al.,: MDE basics with a DSL focus. In: Formal Methods for Model-Driven Engineering. LNCS, vol. 7320 (2012).

The screenshot shows a software interface for formal method development. On the left, there is a code editor with the following text:

```

n1 = name1 ∧
¬ (Ǝ n2: {name1} ∪ {name2} • (¬ name1 = n2))
⇒
(Ǝ NAMES: F≈1≈ NAME •
(∀ n1_0: NAMES • (Ǝ n2_0: NAMES • ¬ n1_0 = n2_0)))
    
```

To the right of the code editor is a vertical toolbar with the following options:

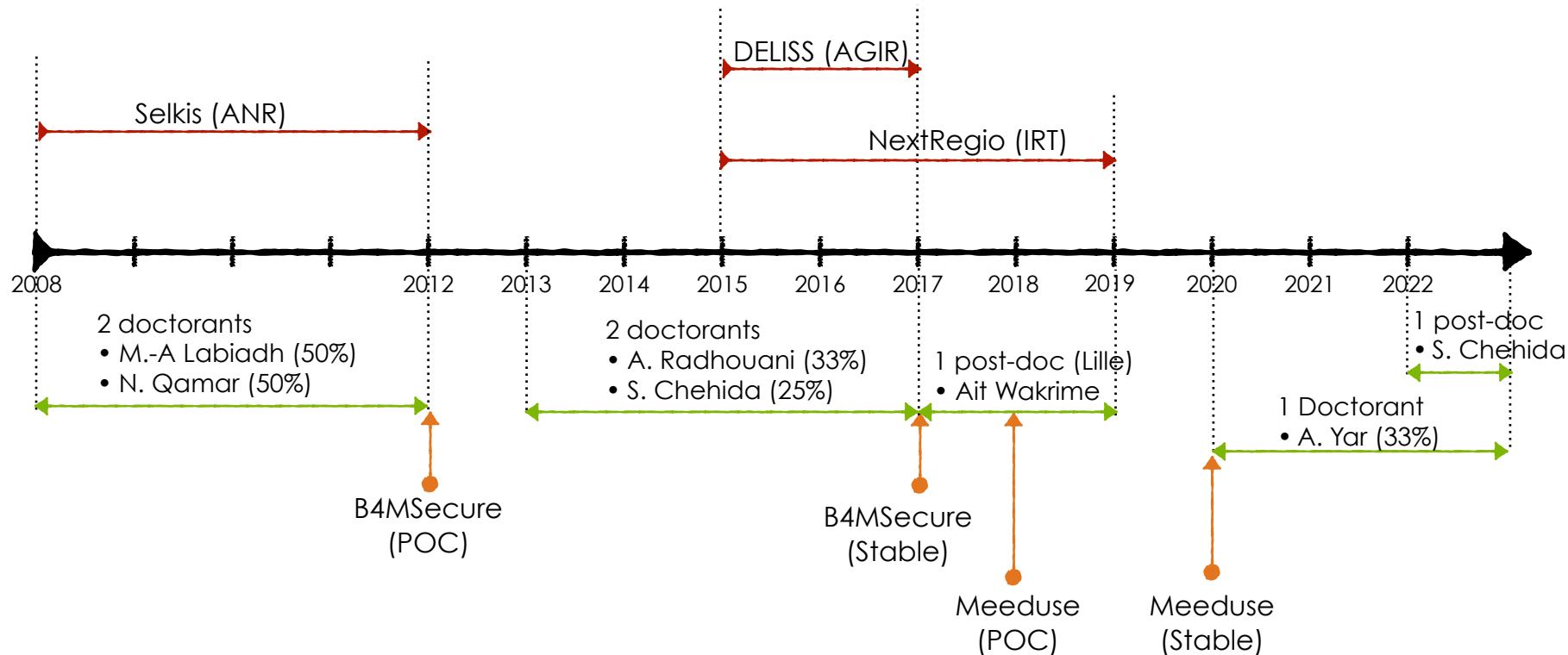
- Prove by Reduce
- Reduce
- Rewrite
- Simplify
- Invoke
- Trivial Rewrite

The main workspace shows a proof state with a yellow highlighted area containing the text "v(3) Ok". Below the workspace, there are two small numbers: 7, 8 and 4, 8, each with a crossed-out symbol above them.

Outline

- Introduction
- **Model Driven Security**
 - UML / SecureUML
 - Insider attacks
- **Executable DSLs**
 - Motivations and results
 - Applications
- Conclusion

Introduction



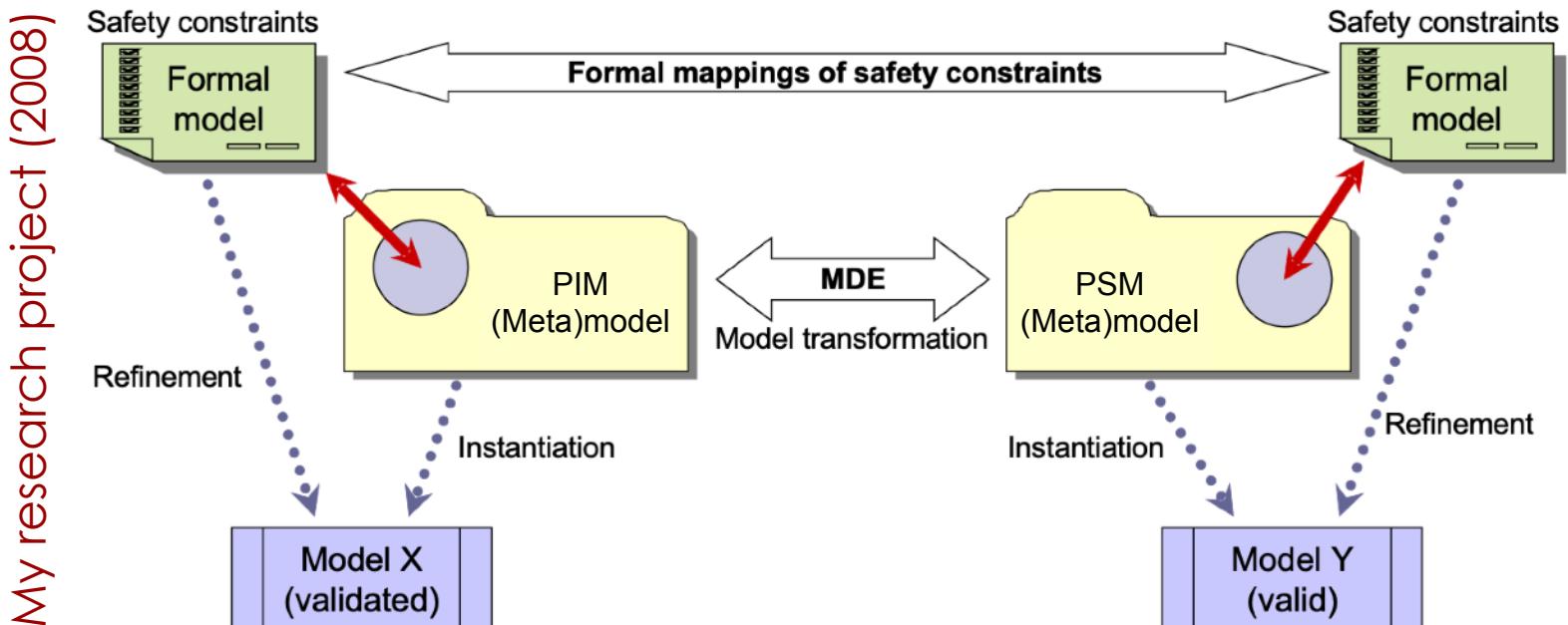
- **Selkis** (projet ANR, 2008/2012) : A development method of secure health care networks information systems : from requirements engineering to implementation
- **DELISS** (Alpes Grenoble Innovation Recherche, 2015/2017) : Déploiement validé de politiques de Sécurité en systèmes d'Information.
- **NextRegio** (IRT Railenium, 2015/2019) : solutions d'exploitation pour les lignes de desserte fine du territoire.

Introduction

Robert France, Bernhard Rumpe., Model-driven Development of Complex Software: A Research Roadmap. *International Conference on Software Engineering 2007.*

« Members of the FST and the MDE communities need to collaborate. »

« Modeling languages must have formally defined semantics if they are to be used to create analyzable models. »



Introduction

- Two schools
 - The Extensible General-Purpose Modeling Language School
 - e.g. SysML, OntoML, UML-RT
 - Provide a language with extension mechanisms
 - **B4MSecure: Model-Driven Security**
 - The Domain Specific Modeling Language School
 - e.g. Capella, OWL, Lustre
 - Tool support for engineering modeling languages
 - **Meeduse: grammars and Meta(-meta)-modeling**

Robert France, Bernhard Rumpe., Model-driven Development of Complex Software: A Research Roadmap. *International Conference on Software Engineering 2007.*

« [...] can both play vital roles in an MDE environment. We envisage that research in both schools will provide valuable insights and research results that will lead to a convergence of ideas. »

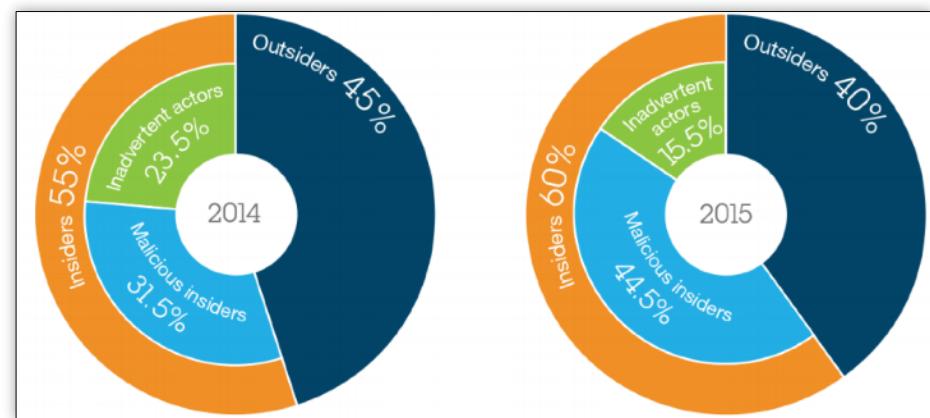
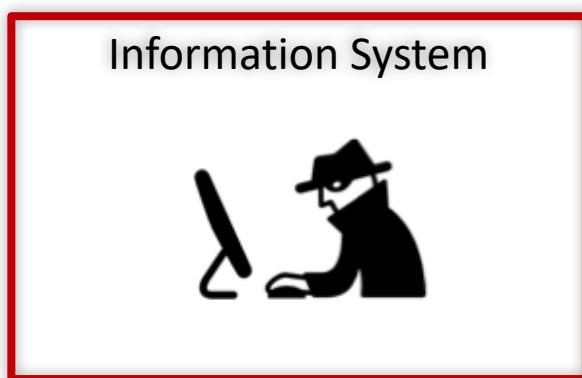
Outline

- Introduction
- **Model Driven Security**
 - UML / SecureUML
 - Insider attacks
- Executable DSLs
 - Motivations and results
 - Applications
- Conclusion

Access control

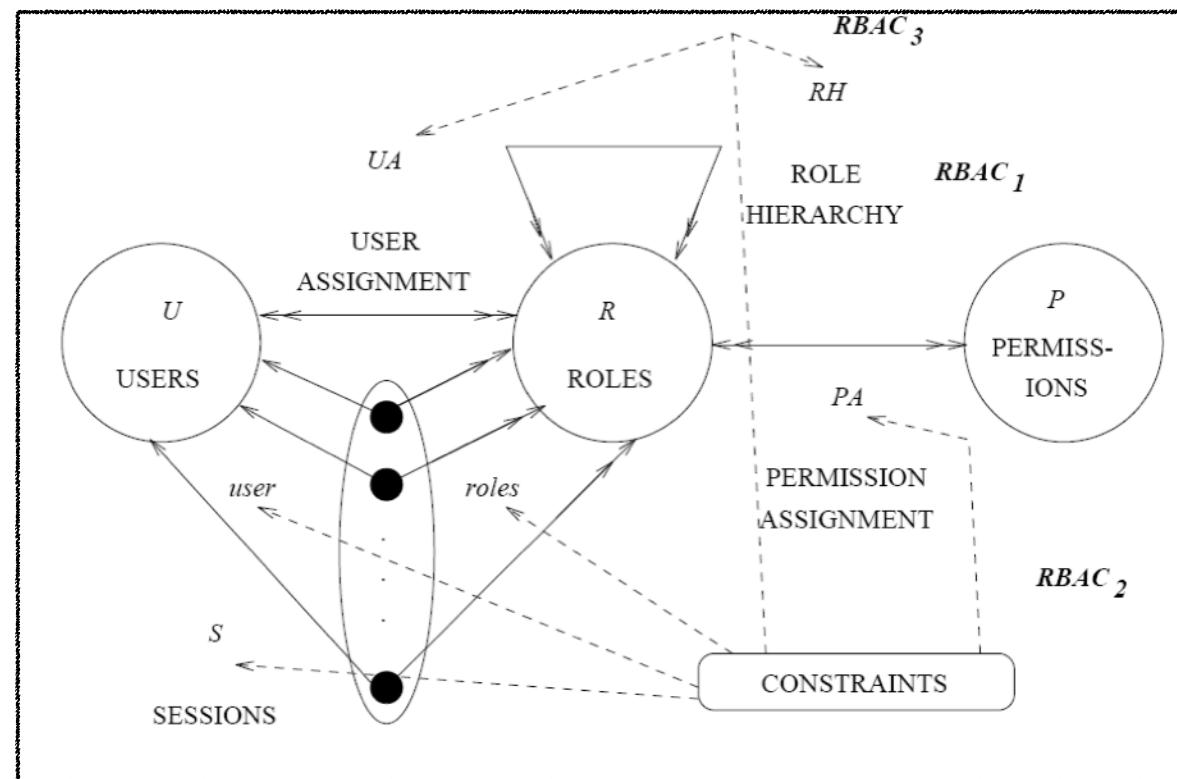


- Access control is horizontal :
 - Ad-hoc integration is error prone and costly
- High-level of abstraction :
 - modeling structure and behavior
- Correctness :
 - Testing / model-checking / Proofs
- Prevention against attacks



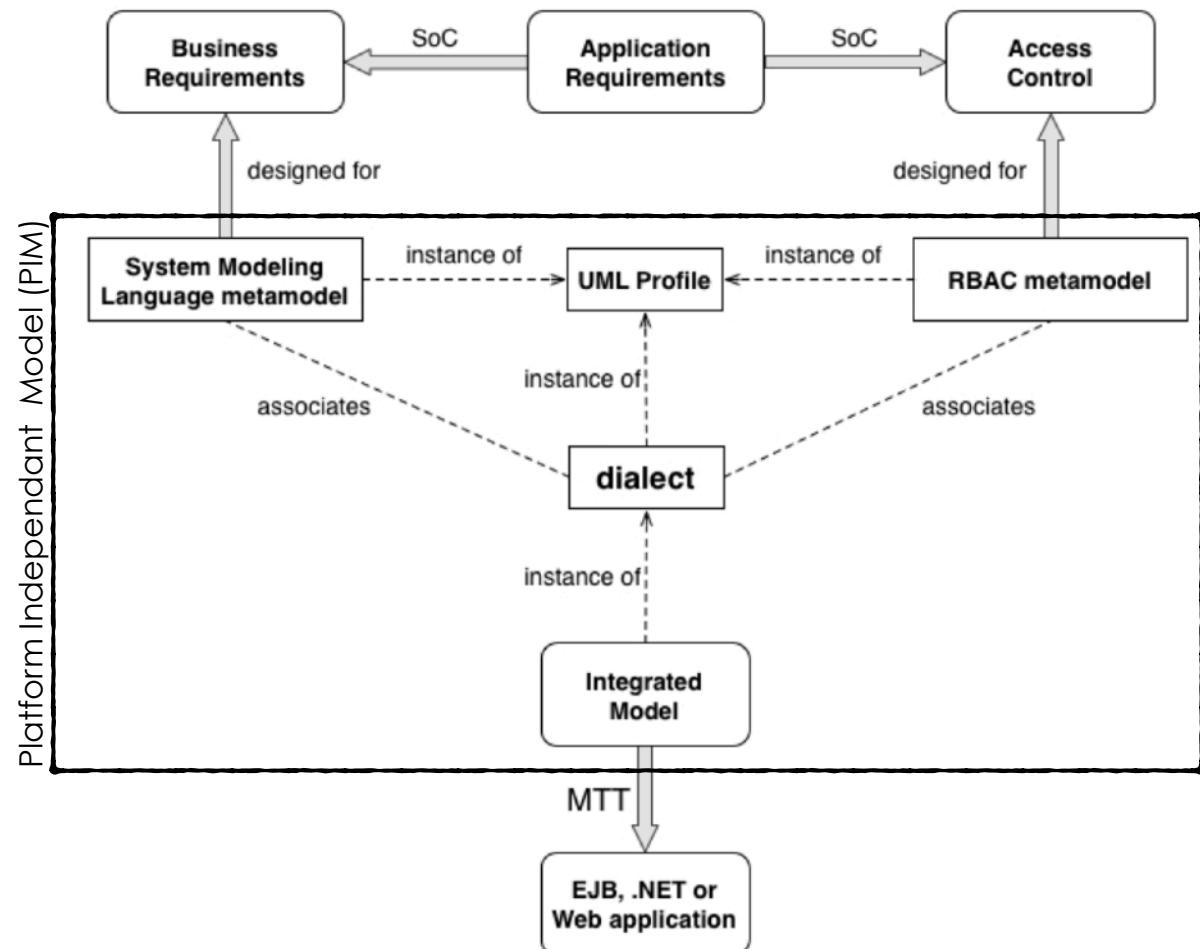
Role-Based Access Control

- RBAC : 1992 by D. Ferraiolo and R. Kuhn ↗ 2000
- Constraints
 - Static Separation of duties
 - Dynamic Separation of duties
 - Contextual
 - **Data-centric too!**
- ABAC : 2011



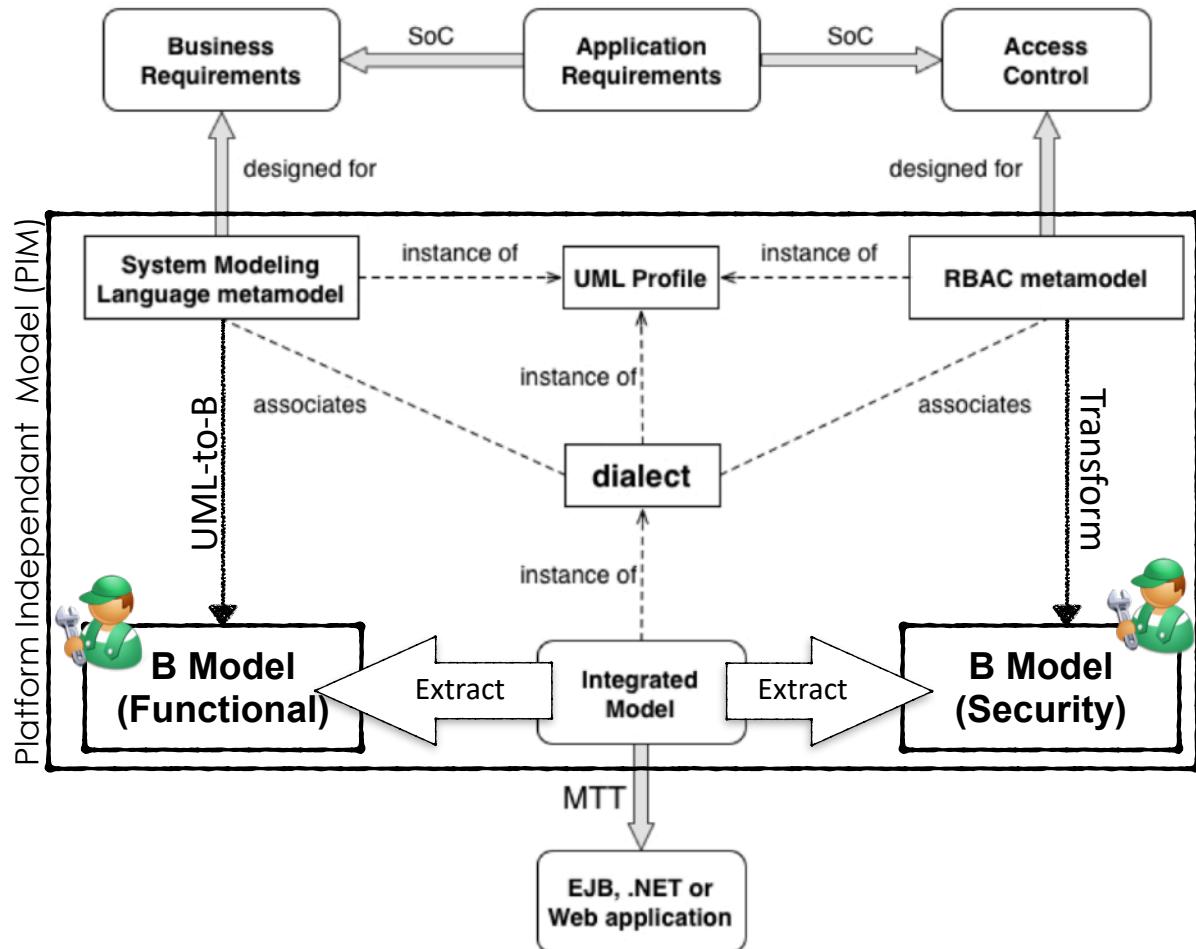
MDE for Security

- Separation of concerns
- e.g. SecureUML



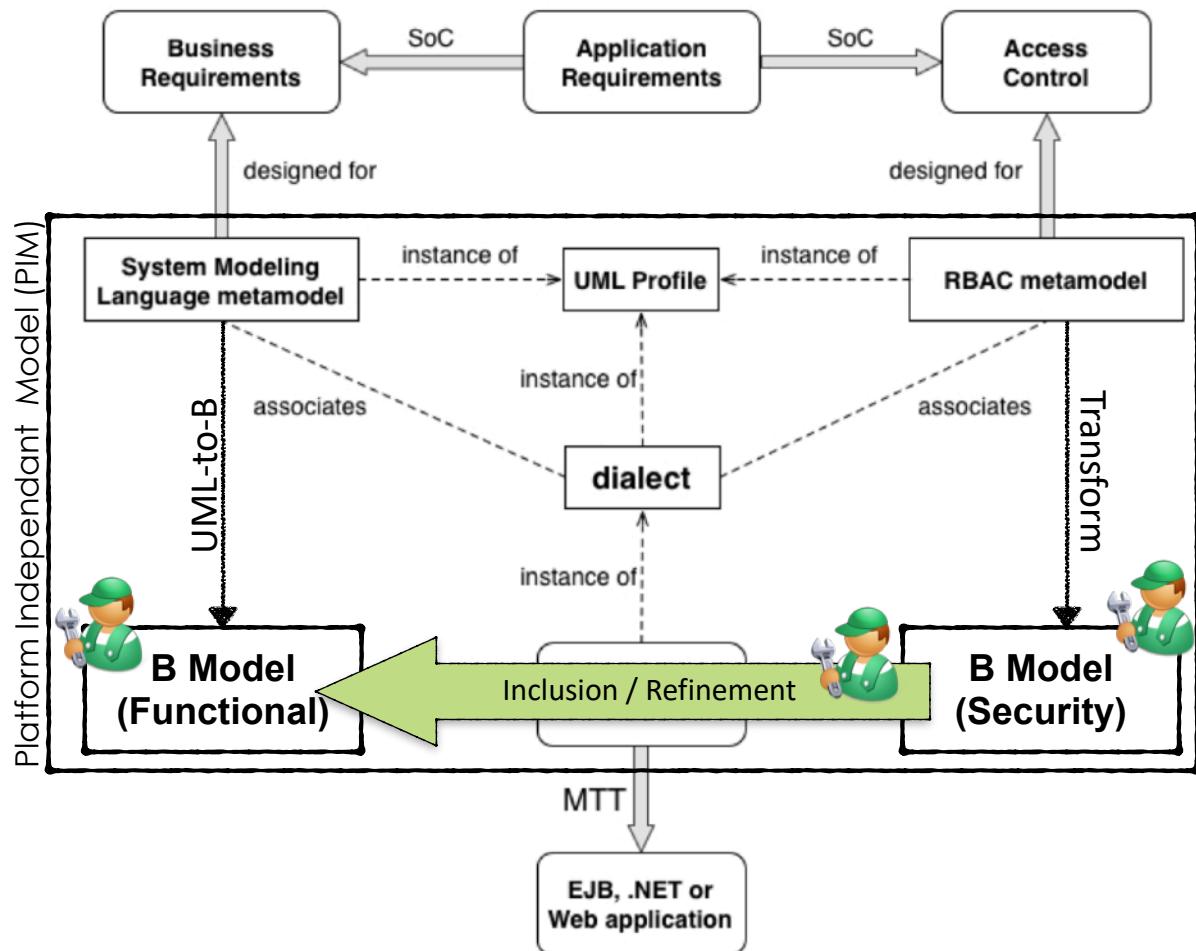
MDE for Security

- Separation of concerns
- e.g. SecureUML
- Contribution:
 - Formal MDS
 - Proof
 - Test/Animation
 - Model-checking

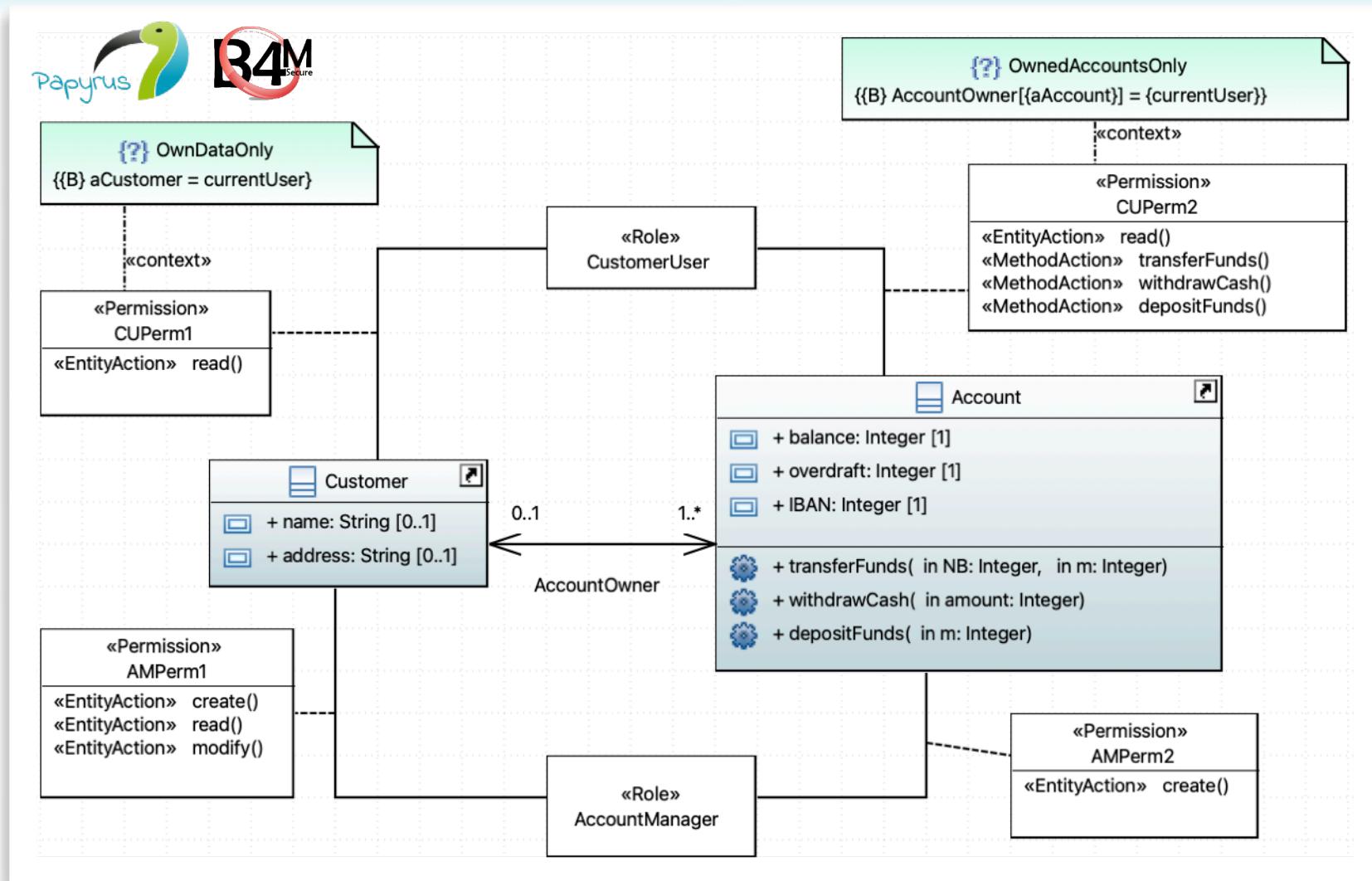


MDE for Security

- Separation of concerns
- e.g. SecureUML
- Contribution:
 - Formal MDS
 - Proof
 - Test/Animation
 - Model-checking



A Simple Example

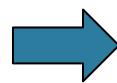


Verification & Validation (1/2)

- Static V&V
 - USE [Kuhlmann et al. 2013] , SecureMOVA [Basin et al. 2009], Alloy-Analyser [Zao et al. 2003, Ahn et Hu 2007], etc.
 - Evaluation is done in a given state
 - Given an atomic action, which roles can perform this action?
 - Given a role and an atomic action, under which circumstances can a user in this role perform this action?
 - Do two permissions overlap?



Is the manager able to transfer funds from a customer account?



A static query: **NO**

Verification & Validation (2/2)

- Dynamic V&V
 - **Jaza/animation [PhD of N. Qamar (50%)]: static queries + animation**
 - USE/OCLE [Yu et al. 2009], generation of execution schemas.
 - RW [Zhang et al. 2008], model-checking
 - SMP (Logic for State Modifying Policies) [Becker et Nana 2010]
 - Theorem proving [A. Mammar et al., FAC'15]

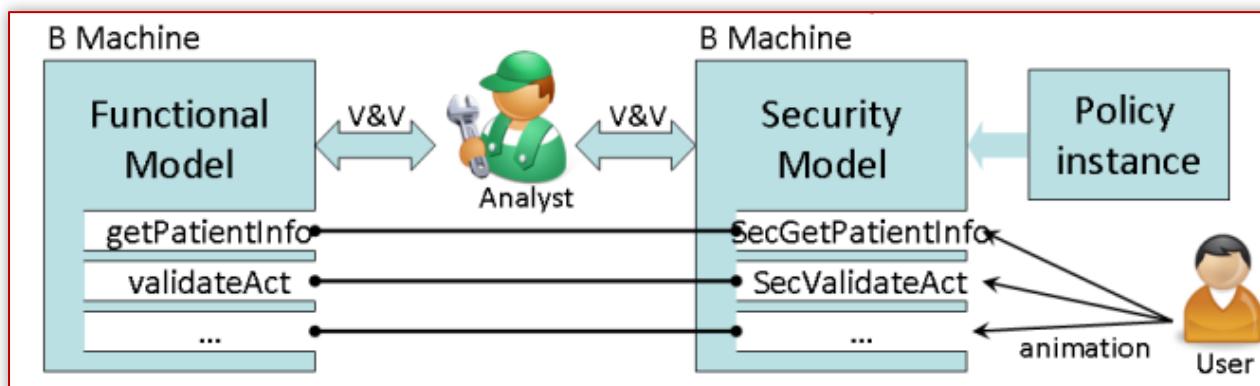
Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?



Research directions

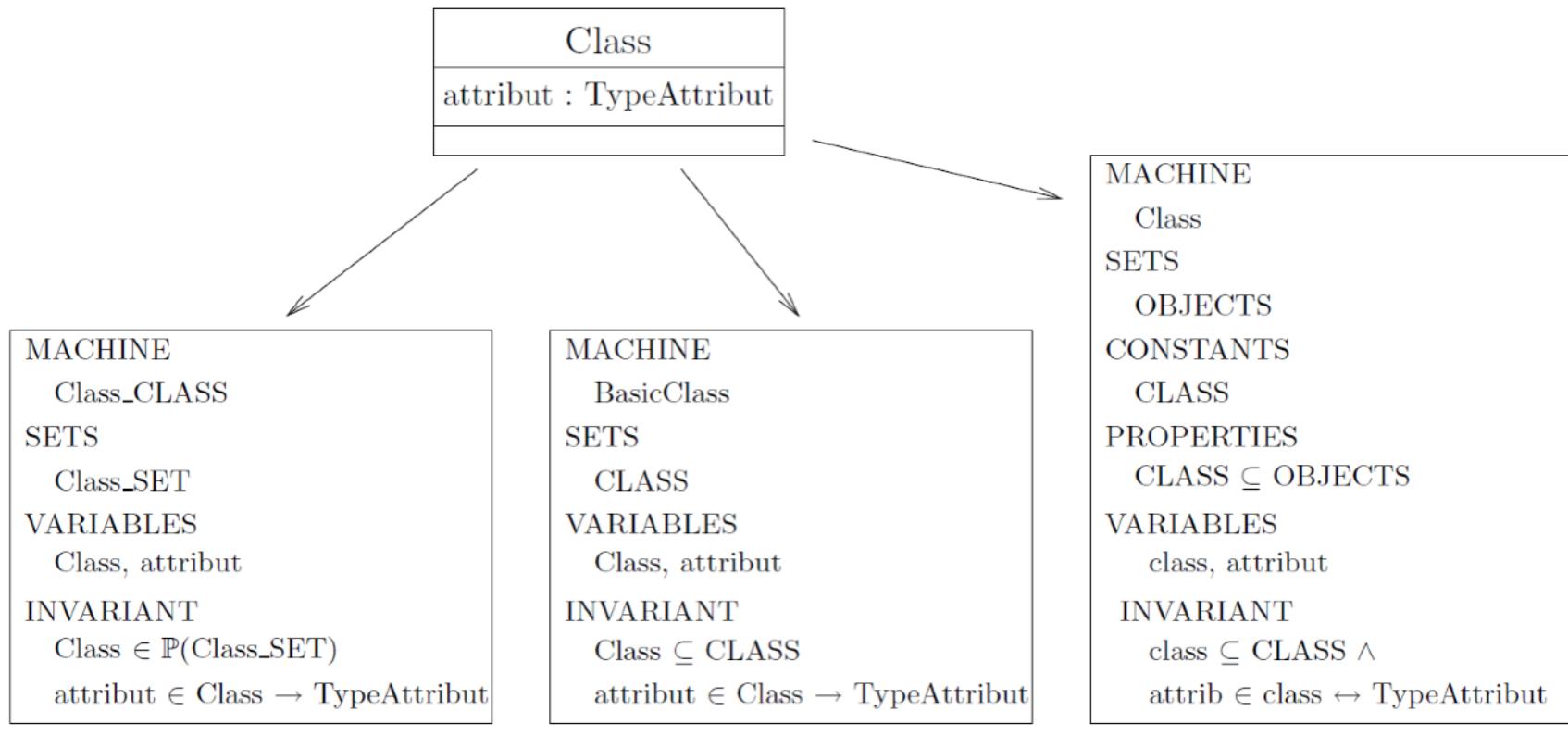
- Modeling/Testing
 - **Configurable transformations**
 - Abstraction levels: M1 ; M2 and M1/M2
 - PhD M.-A. Labiadh (50%) + 1 M2 + 1 PFE
- Attacks
 - Forward search
 - **Guided model-checking** (2 M2)
 - Ant-colony optimisation (2 IRL-ENSIMAG)
 - Backward search
 - **Proof** and constraint solving
 - PhD A. Radhouani (33%) + 1 M2
- Business processes
 - Task-based access control: PhD S. Chehida (25%)
 - BPMN: 2 M2, collaboration with SIGMA/LIG

Modeling/Testing



Modeling/Testing (1/3)

- Various tools and approaches, but
 - Most of them are unavailable
 - Should be revisited/updated



(Snook *et al.*, 2004)

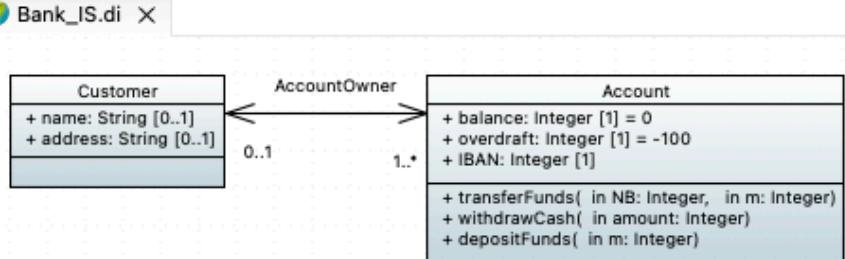
(Laleau, 2002)

(Meyer, 2001)

Modeling/Testing (2/3)

	R. Laleau A. Mammar	E. Meyer H. Ledang	C. Snook M. Butler
Classes	X	X	X
Classes (fixed instances)	-	-	X
Attributes	X	X	X
(Single/Multi)-valued attributes	X	-	-
Inheritance	X	X	X
Multiplicities	X	X	X
Navigation	-	X	X
Roles	X	-	X
Association constraints	X	X	-
Fixed/Non-fixed associations	X	-	-
Association classes	X	X	-
Association classes with inheritance or other relationships	X	-	
Parameterized classes	-	-	X

Modeling/Testing (3/3)



```

classDiagram
    class Customer {
        +name: String [0..1]
        +address: String [0..1]
    }
    class Account {
        +balance: Integer [1] = 0
        +overdraft: Integer [1] = -100
        +IBAN: Integer [1]
        +transferFunds( in NB: Integer, in m: Integer)
        +withdrawCash( in amount: Integer)
        +depositFunds( in m: Integer)
    }
    Customer "0..1" --> "1..*" Account : AccountOwner
  
```

Functional.mch

- Customer_SetName(aCustomer,aName)
 - Customer_SetName(aCustomer=Paul,aName="STRING1")
 - Customer_SetName(aCustomer=Paul,aName="STRING2")
 - Customer_SetName(aCustomer=Martin,aName="STRING1")
 - Customer_SetName(aCustomer=Martin,aName="STRING2")
- result <-- Account_GetOverdraft(aAccount)
 - 100 <-- Account_GetOverdraft(aAccount=cpt1)
 - 100 <-- Account_GetOverdraft(aAccount=cpt2)
- Customer_Free(aCustomer)
 - Customer_Free(aCustomer=Paul)
 - Customer_Free(aCustomer=Martin)
- Account_transferFunds(Instance,NB,m)
 - Account_transferFunds(Instance=cpt1,NB=222,m=100)

Invariant=OK

Completed=NO

Variables	Expressions
variable	value
> Account	{ cpt2 , cpt1 }
> Account_balance	{ (cpt1 -> 300) , (cpt2 -> -100) }
> Customer	{ Paul , Martin }
> Customer_address	()
> Account_overdraft	{ (cpt1 -> -100) , (cpt2 -> -100) }
> Customer_name	()
> Account_IBAN	{ (cpt1 -> 111) , (cpt2 -> 222) }
> AccountOwner	{ (cpt1 -> Paul) , (cpt2 -> Martin) }

Execution View

Search:

Functional.mch

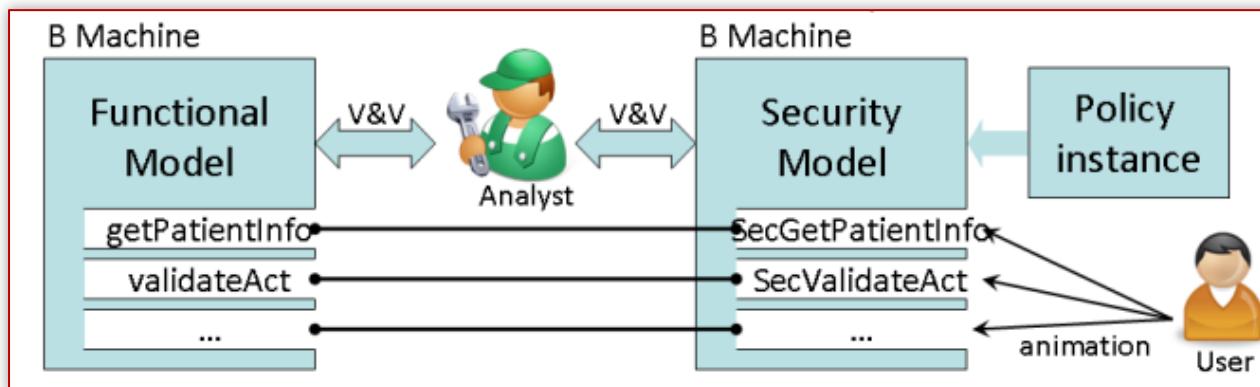
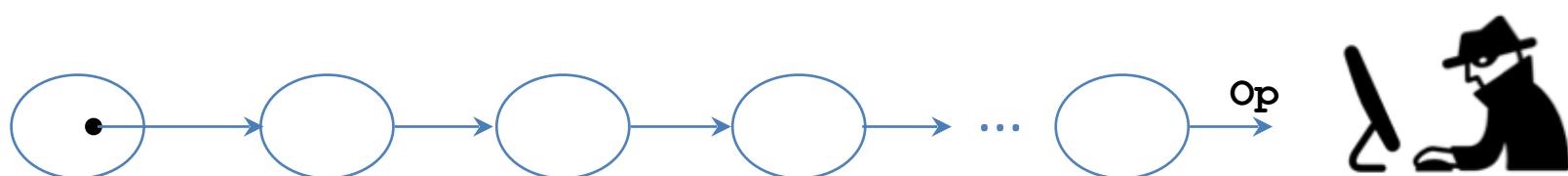
- Customer_SetName(aCustomer,aName)
 - Customer_SetName(aCustomer=Paul,aName="STRING1")
 - Customer_SetName(aCustomer=Paul,aName="STRING2")
 - Customer_SetName(aCustomer=Martin,aName="STRING1")
 - Customer_SetName(aCustomer=Martin,aName="STRING2")
- result <-- Account_GetOverdraft(aAccount)
 - 100 <-- Account_GetOverdraft(aAccount=cpt1)
 - 100 <-- Account_GetOverdraft(aAccount=cpt2)
- Customer_Free(aCustomer)
 - Customer_Free(aCustomer=Paul)
 - Customer_Free(aCustomer=Martin)
- Account_transferFunds(Instance,NB,m)
 - Account_transferFunds(Instance=cpt1,NB=222,m=100)

Execution History View

Functional.mch

- Account_transferFunds(Instance=cpt2,NB=111,m=300)
- Account_depositFunds(Instance=cpt2,m=200)
- Customer_NEW(aCustomer=Martin,theAccount={ cpt2 })
- Account_NEW(aAccount=cpt2,alBAN=222)
- Customer_NEW(aCustomer=Paul,theAccount={ cpt1 })
- Account_NEW(aAccount=cpt1,alBAN=111)

Attacks



Insider attacks: forward search (1/3)

- Guided model checking (CSP | | B)

Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?



```
MAIN = UI
```

```
UI = (Connect?user!{AccountManager} -> setCurrentUser(user) -> MANAGER_FUNC  
      [] Connect?user!{CustomerUser} -> setCurrentUser(user) -> CLIENT_FUNC)  
      ; disconnectUser -> UI
```

```
MANAGER_FUNC =
```

```
    CREATE_ACCOUNT [] CREATE_CUSTOMER [] UPDATE_CUSTOMER [] SKIP
```

```
CREATE_ACCOUNT =
```

```
    secure_Account_NEW -> (CREATE_ACCOUNT [] MANAGER_FUNC)
```

```
CREATE_CUSTOMER =
```

```
    secure_Customer_NEW?customer -> secure_Customer__SetName!customer  
    -> (ADD_CUSTOMER_ACCOUNT(customer) [] CREATE_CUSTOMER [] MANAGER_FUNC)
```

Insider attacks: forward search (2/3)

- Guided model checking (CSP | | B)

Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?



```
MAIN = UI [|{| Connect, secure_Account_transferFunds |}|] ATTACK
```

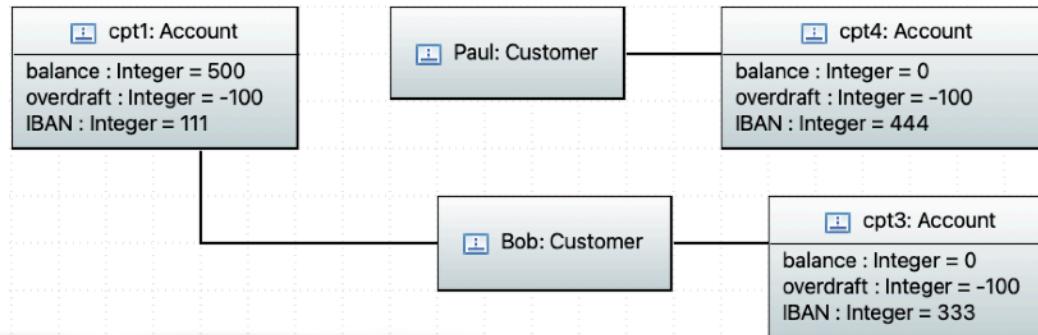
```
ATTACK = ATTACKER ||| secure_Account_transferFunds!cpt1 -> goal -> SKIP
```

```
ATTACKER = []role:Set(ROLES) @ Connect!Bob!role -> ATTACKER
```

Insider attacks: forward search (3/3)

- Guided model checking (CSP | | B)

Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?



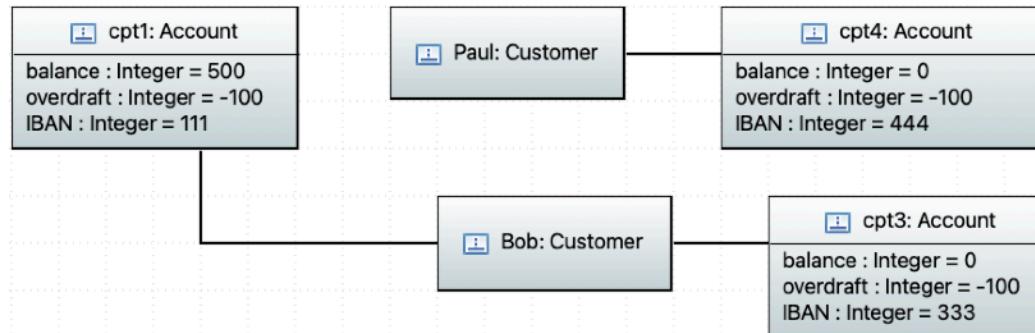
```

Connect(Bob, {AccountManager}) ;
setCurrentUser(Bob) ;
secure_Account_NEW(cpt3, 333) ;
secure_Customer_NEW(Bob,{cpt3}) ;
secure_Customer_SetName(Bob, "...") ;
  
```

Insider attacks: forward search (3/3)

- Guided model checking (CSP | | B)

Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?

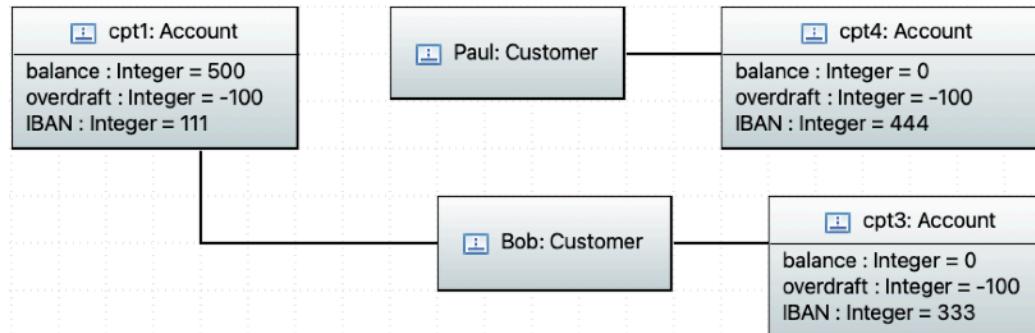


```
Connect(Bob, {AccountManager}) ;
setCurrentUser(Bob) ;
secure_Account_NEW(cpt2, 333) ;
secure_Account_NEW(cpt4, 444) ;
secure_Customer_AddAccount(Paul,{cpt4}) ;
secure_Customer_RemoveAccount(Paul,{cpt1}) ;
secure_Customer_AddAccount(Bob,{cpt1}) ;
```

Insider attacks: forward search (3/3)

- Guided model checking (CSP | | B)

Is there a sequence of operations that can be executed by a manager in order to transfer funds from a customer's account?



```

Connect(Bob, {AccountManager}) ;
setCurrentUser(Bob) ;
secure_Account_NEW(cpt3, 333) ;
secure_Customer_NEW(Bob, {cpt3}) ;
secure_Customer_SetName(Bob, ...) ;
secure_Customer_RemoveAccount(Paul, {cpt4}) ;
secure_Customer_AddAccount(Paul, {cpt4}) ;
disConnect(Bob) ;
Connect(Bob, {CustomerUser}) ;
secure_Account_transferFunds(cpt1, 333, 500) ;
  
```

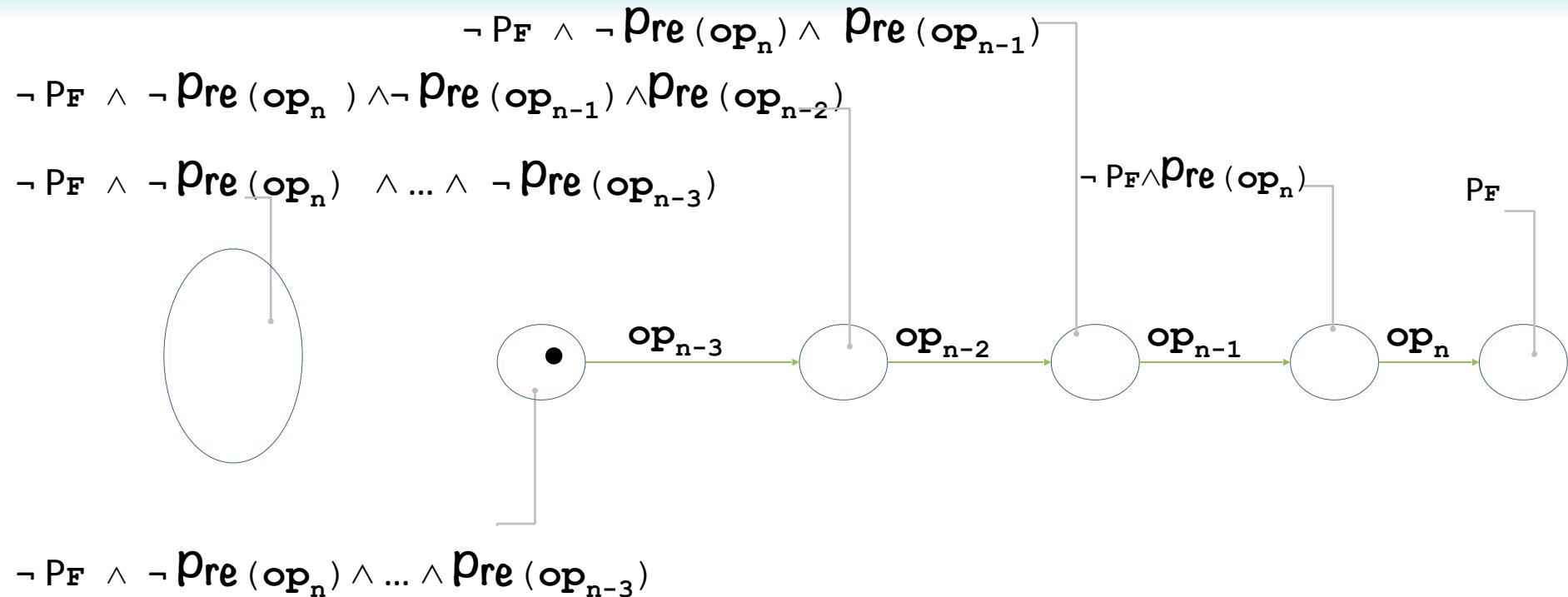
Insider attacks: backward search

States = Predicates

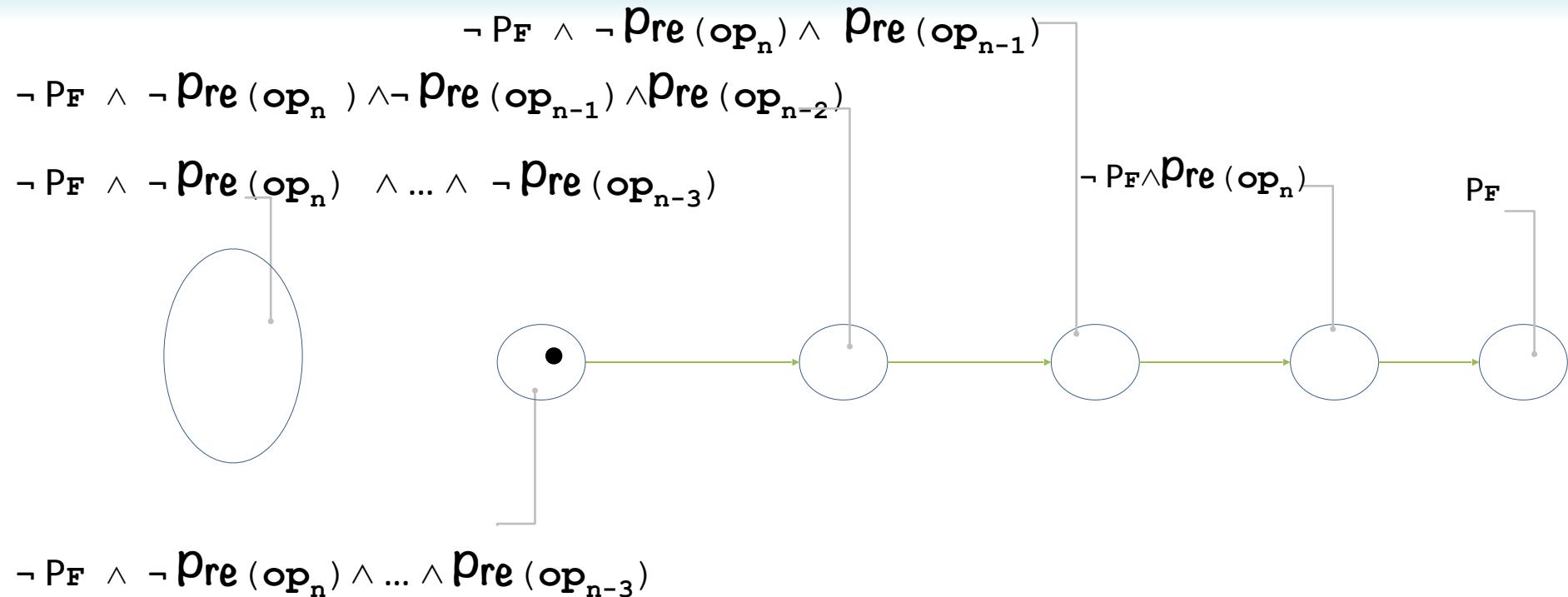


- (1) always enabled: $\forall x.E \Rightarrow \text{Pre}(op)$
- (2) never enabled: $\forall x.E \Rightarrow \neg \text{Pre}(op)$
- (3) possibly enabled ($\neg (1) \wedge \neg (2)$): $\exists x.E \wedge \text{Pre}(op)$
- (4) always reached: $\forall x.E \wedge \text{Pre}(op) \Rightarrow [\text{Action}(op)]F$
- (5) never reachable: $\forall x.E \wedge \text{Pre}(op) \Rightarrow [\text{Action}(op)]\neg F$
- (6) possibly reached ($\neg (4) \wedge \neg (5)$): $\exists x.E \wedge \text{Pre}(op) \wedge \neg [\text{Action}(op)]\neg F$

Insider attacks: backward search



Insider attacks: backward search



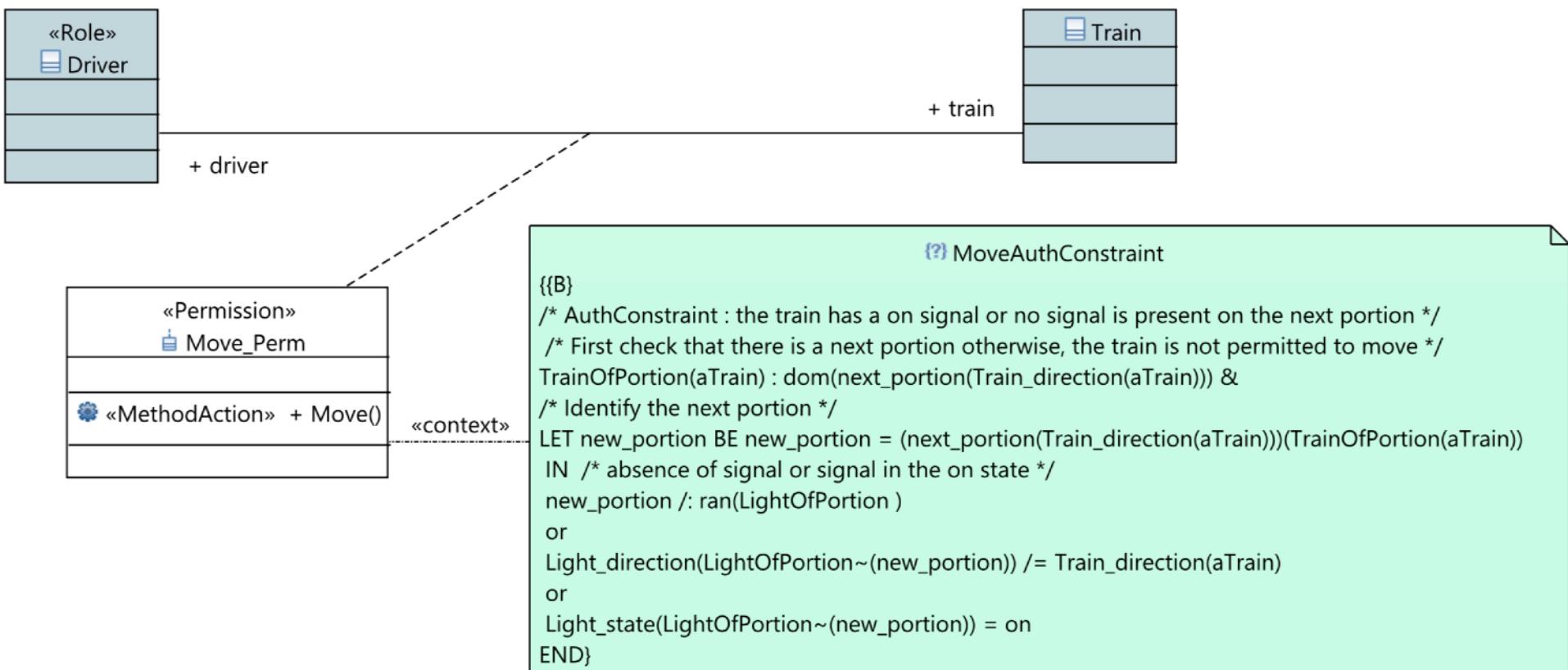
$$Q_{\text{symb}} = \langle \text{init}, \text{ op}_{n-3}, \text{ op}_{n-2}, \text{ op}_{n-1}, \text{ op}_n \rangle$$

Outline

- Introduction
- **Model Driven Security**
 - UML / SecureUML
 - Insider attacks
- **Executable DSLs**
 - Motivations and results
 - Applications
- Conclusion

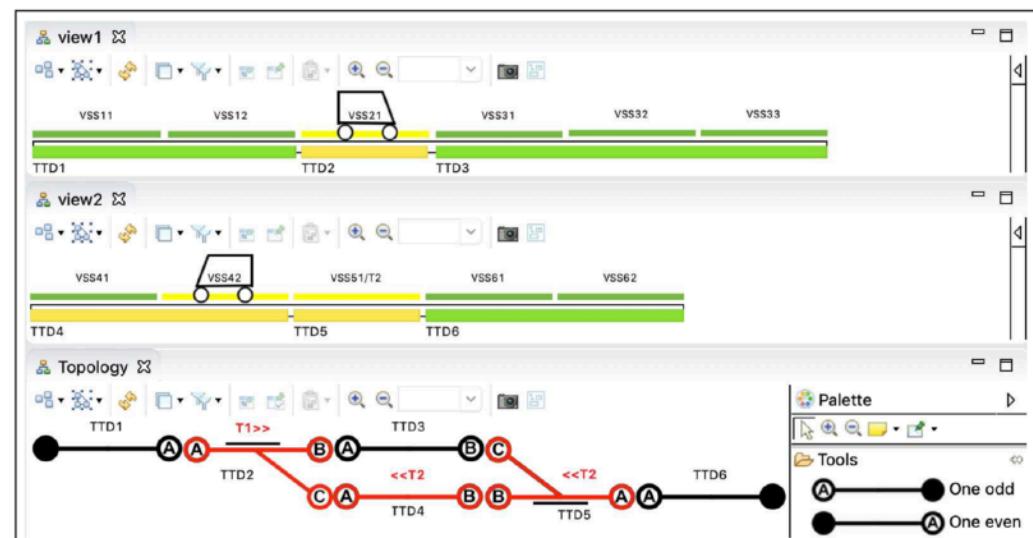
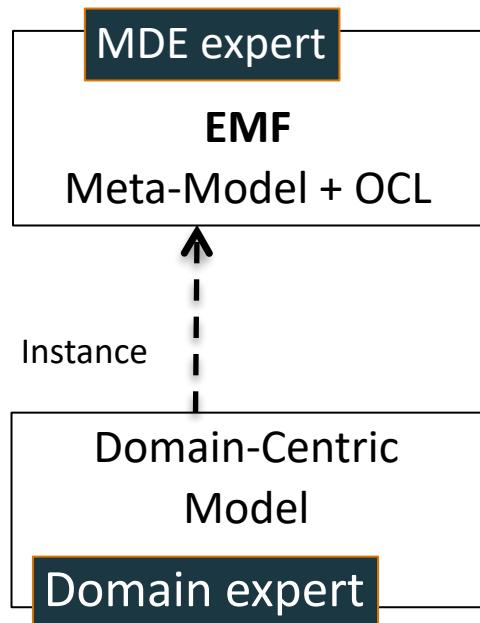
UML vs DSLs

- NextRegio (IRT Railenium, SNCF réseau)
- Undesirable scenarios + responsibilities



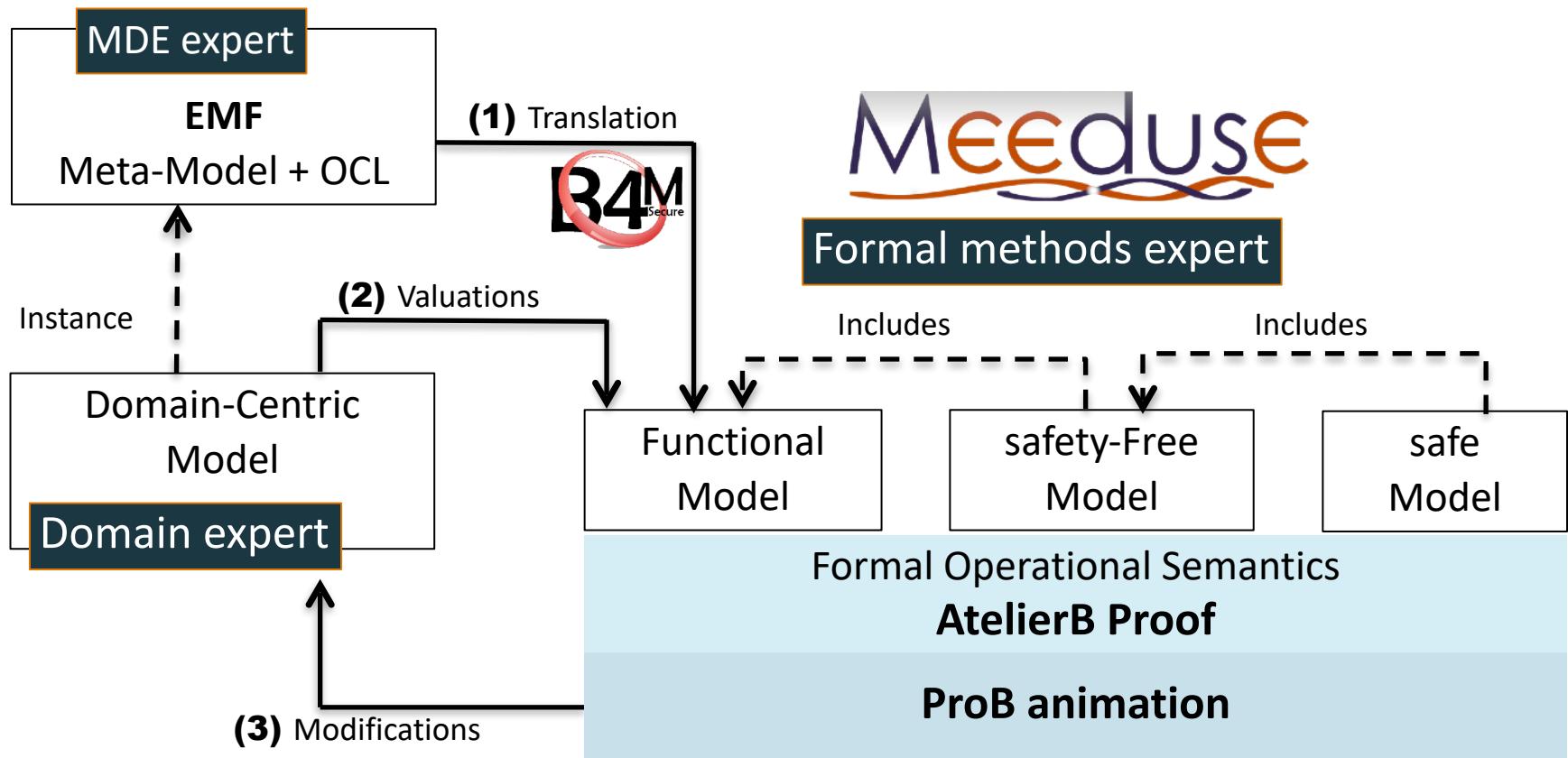
UML vs DSLs

- NextRegio (IRT Railenium, SNCF réseau)
- Undesirable scenarios + responsibilities



UML vs DSLs

- NextRegio (IRT Railenium, SNCF réseau)
- Undesirable scenarios + responsibilities



Positioning

- DSL → FM
 - Survey 1 (2006/2012):
 - T. Kosar, S. Bohra, and M. Mernik, “Domain-specific languages: A systematic mapping study,” *Information and Software Technology*
 - Survey 2 (2012/2019):
 - A. Iung, J. Carbonell, L. Marchezan, E. M. Rodrigues, M. Bernardino, F. P. Basso, and B. Medeiros, “Systematic mapping study on domain-specific language development tools,” *Empirical Software Engineering*.

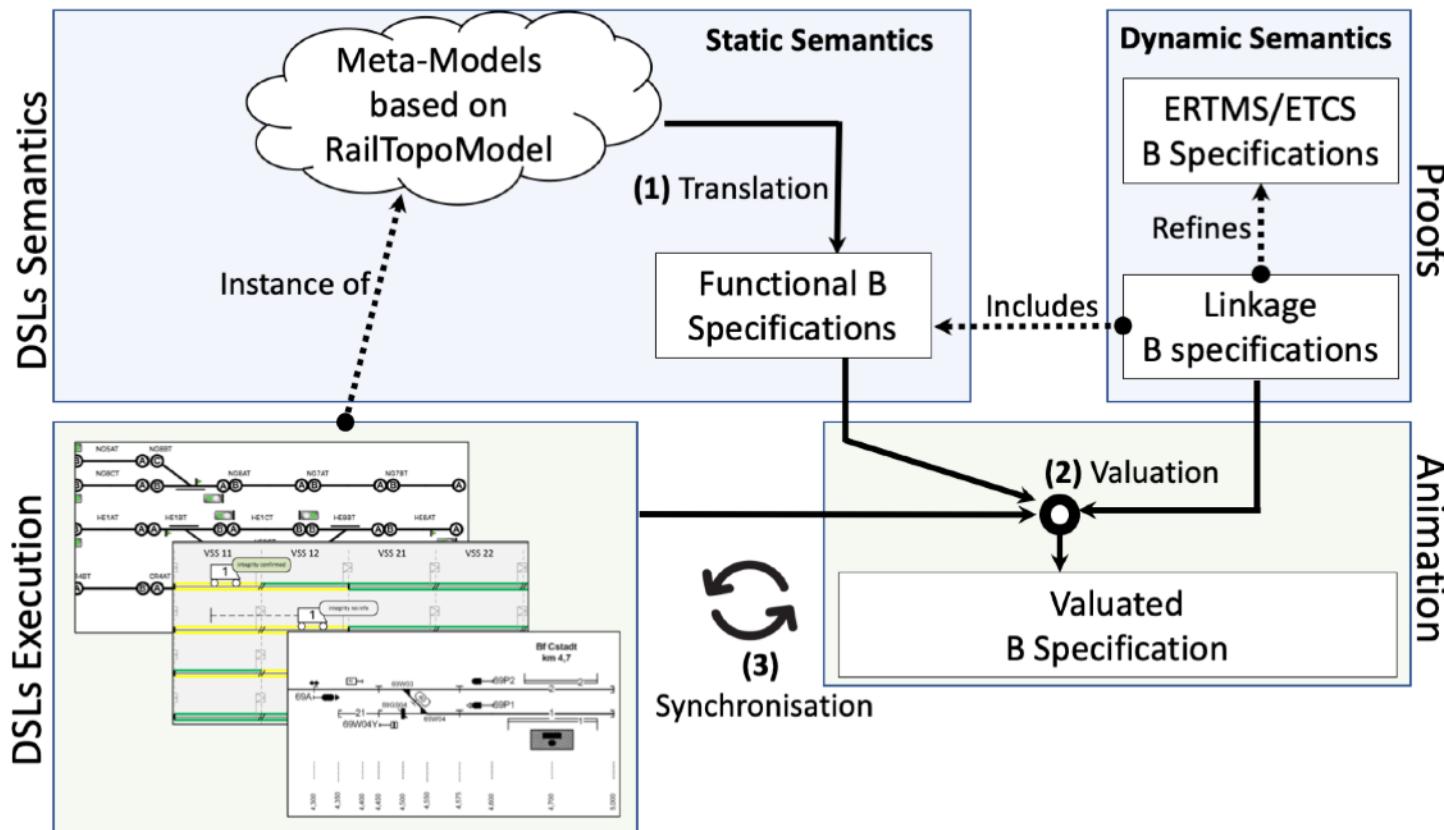
Refers to testing as “the” verification feature of Language Work-Benches

Personal opinion

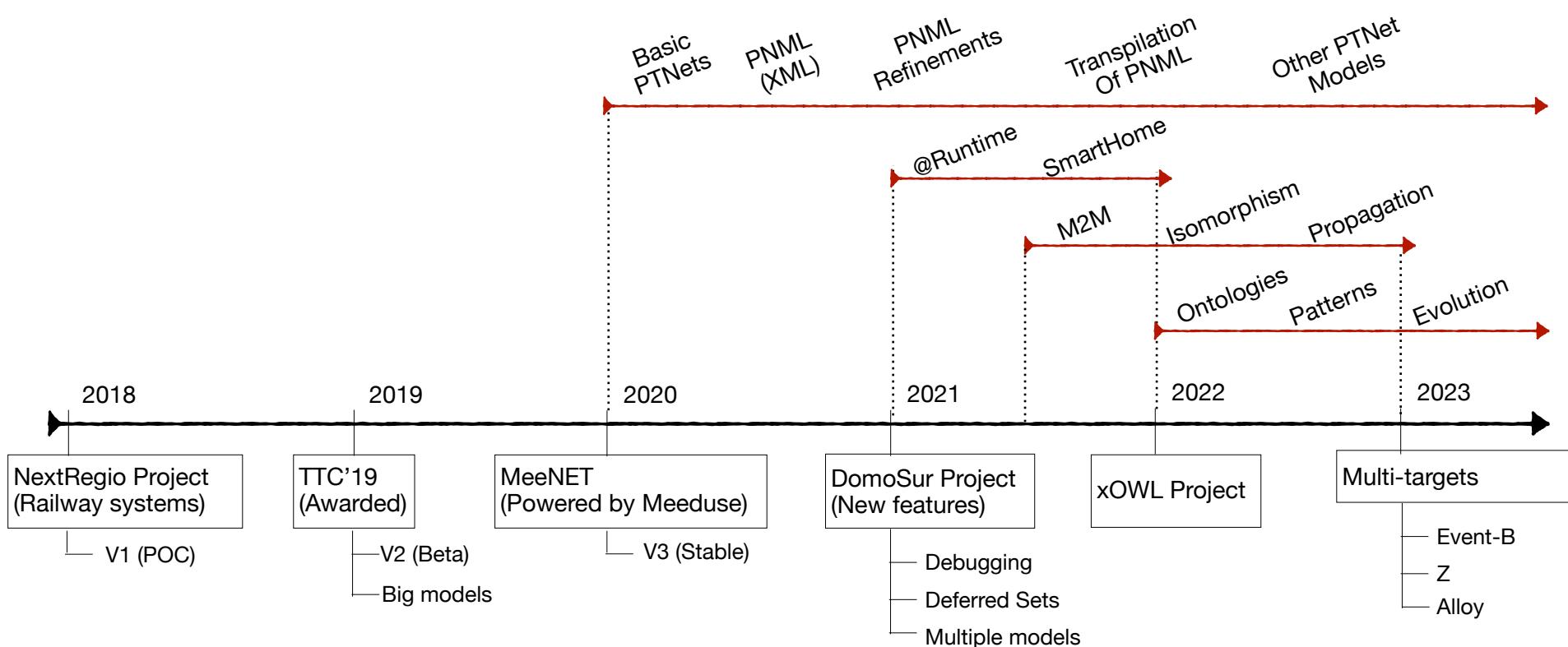
- DSL → FM
 - B. R. Bryant, J. Gray, M. Mernik, P. J. Clarke, R. B. France, and G. Karsai, "Challenges and directions in formalizing the semantics of modeling languages," Comput. Sci. Inf. Syst (2011).
 - Translational approaches:
 - **Advantages:**
 1. provide a well-defined and understood semantics, and
 2. the DSL can convey existing tools of the language into which it is translated.
 - **Limitations:**
 3. it is very challenging to correctly map the constructs of the DSL into the constructs of the target language, and
 4. the mapping of the verification results back into the DSL appears as a major issue of existing approaches.

Applications

- Railway systems - PhD A. Yar (33%)



Applications



Outline

- Introduction
- **Model Driven Security**
 - UML / SecureUML
 - Insider attacks
- **Executable DSLs**
 - Motivations and results
 - Applications
- Conclusion & Perspectives

Conclusion & perspectives

- FM & MDE
 - MDS: Healthcare IS (Security)
 - xDSLs: Railway systems (Safety)
- B4MSecure / Meeduse
- FM and MDS: is (still) an active topic
 - Conformance PIM/PSM/Application
 - Monitor/Controller synthesis
 - Align (Data, Security, business)-Models
- FM and xDSLs: not a new topic but requires further works
 - Material already exists (iFM, MoDeVVa, UML&FM, VOLT, etc)
 - Investigate other languages (Maude, ASM, Z, Alloy, etc)
 - Provide a unifying framework

Facts

- B4MSecure: Teaching (master degree, MoSIG 2)
 - Topic: Software/Hardware: quality engineering, models of computation
 - Lecture: Information Security
- Meeduse: Recent papers
 - A. Idani., The B Method meets MDE.
Research Challenges in Information Science (RCIS-2022)
 - A. Idani., Formal Model Driven Executable DSLs: Application to Petri-nets.
NASA Journal on Innovations in Systems and Soft. Engineering (ISSE-2022).
 - A. Idani et al., Alliance of model-driven engineering with a proof-based formal approach.
NASA Journal on Innovations in Systems and Soft. Engineering (ISSE-2020).
 - A. Idani., Meeduse: A Tool to Build and Run Proved DSLs.
Integrated Formal Methods (iFM-2020).
 - A. Idani et al., Incremental Development of a Safety Critical System Combining formal Methods and DSMLs - Application to a Railway System.
Formal Methods for Industrial Critical Systems (FMICS-2019).
 - A. Idani et al., Towards a Tool-Based Domain Specific Approach for Railway Systems Modeling and Validation.
Reliability, Safety and Security of Railway Systems (RSSRail-2019).

