

# An Elementary Construction of Constant-Degree Expanders

Noga Alon , Oded Schwartz & Asaf Shapira

*CS6845 presentation by*  
**Amit Roy**



*Department of Computer Science & Engineering*  
*Indian Institute of Technology - Madras*

# Definition

A  $d$ -regular graph  $G(V, E)$  is  $\delta$ -edge expander if  $\forall S \subseteq V$  of at most size  $\frac{|V|}{2}$ ,

$$e(S, \bar{S}) \geq \delta d |S|$$

# Definition

We will say that a graph is  $[n, d, \delta]$ -expander if it is an  $n$ -vertex  $d$ -regular  $\delta$ -expander.

# Motivation

Why do we need constant degree expanders?

# Motivation

If we have constant -sized object then we can freely use it without having to find a nice description for it. Also we can always find constant size in constant time by brute force approach.

# Main Result

There exists a fixed  $\delta > 0$  such that for any integer  $q = 2^t$  and for any  $\frac{q^4}{40} \leq r \leq \frac{q^4}{2}$ , there is a polynomial time constructible  $[q^{4r+12}, 12, \delta]$  – *expander*.

## Theorem[ Replacement Product]

If  $E_1$  is an  $[n, D, \delta_1]$ -expander and  $E_2$  is an  $[D, d, \delta_2]$ -expander then  $E_1 \circ E_2$  is  $[nD, 2d, \frac{1}{80}\delta_1^2\delta_2]$ -expander

# Construction

- 1 Take graph  $E_1$  as  $[q^2, 3, \delta]$ -expander (By enumerating all possible graphs)
- 2 Take graph  $E_2$  as  $[q^6, q^2, \frac{1}{4}]$ -expander
- 3 Obtain  $E' = E_2 \circ E_1$  which is  $[q^8, 6, \delta']$ -expander
- 4 Take  $E_3$  as  $[q^{4r+4}, q^8, \frac{1}{4}]$ -expander
- 5 Finally obtain  $E_4 = E_3 \circ E'$  which will be  $[q^{4r+12}, 12, \delta']$ -expander



# Replacement Product

## Definition

Let  $G$  be a  $D$ -regular  $D$ -edge colorable graph on  $n$ -vertices and  $H$  be a  $d$ -regular graph on  $D$ -vertices. Replace every vertex  $v_i$  of  $G$  with a cluster of  $D$  vertices, which we denote  $C_i = \{v_1^i, v_2^i, \dots, v_D^i\}$ . For  $1 \leq i \leq n$ , we will put an edge between  $(v_p^i, v_q^i)$ , iff  $(p, q) \in E(H)$ . For every  $(p, q) \in E(G)$ , which is colored  $t$ , we will put  $d$  parallel edges between  $(v_t^p, v_t^q)$

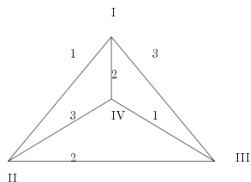
# Replacement Product

**Claim :** The resultant graph  $G \circ H$  is  $2d$ -regular on  $nD$ -vertices

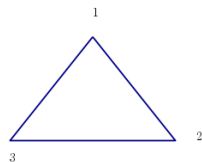
# Replacement Product

## Example

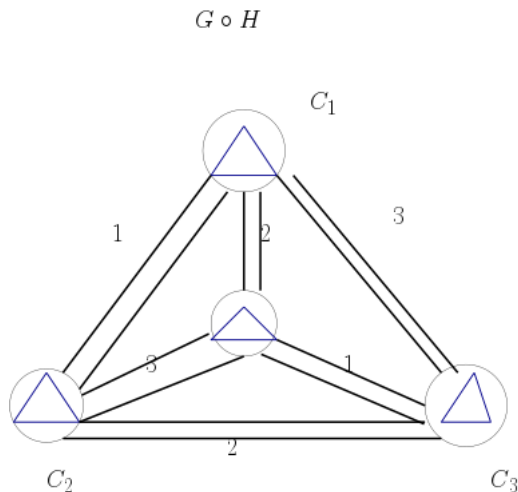
$$G[4, 3, \delta_1]$$



$$H[3, 2, \delta_2]$$



# Replacement Product



# Replacement Product

## Theorem

If  $E_1$  is an  $[n, D, \delta_1]$ -expander and  $E_2$  is an  $[D, d, \delta_2]$ -expander then  $E_1 \circ E_2$  is  $[nD, 2d, \frac{1}{80}\delta_1^2\delta_2]$ -expander

# Replacement Product

**Proof Sketch :** Take a set of vertices  $X \subseteq V$  of at most size  $\frac{nD}{2}$  and show that  $e(X, \overline{X})$  has either many edges within the same cluster  $C_i$  or between them.

# Replacement Product

**Proof :** Now we can see all the vertices as  $n$ -clusters of vertices  $C_1, C_2 \dots C_n$  each of size  $D$ .

$$\text{Goal: } e(X, \overline{X}) \geq \frac{1}{80} \delta_1^2 \delta_2 \cdot 2d \cdot |X|$$

# Replacement Product

**Definition** Let  $X \subseteq V$  of size at most  $\frac{nD}{2}$

$$X_i = C_i \cap X$$

$$I' = \{i \mid |X_i| \leq (1 - \frac{1}{4}\delta_1)D\}$$

$$I'' = \{i \mid |X_i| > (1 - \frac{1}{4}\delta_1)D\}$$

$$X' = \bigcup_{i \in I'} X_i$$

$$X'' = \bigcup_{i \in I''} X_i$$



# Replacement Product

- Consider an  $X_i$  ,  $i \in I'$
- Since  $E_2$  is an  $\delta_2$  expander , we have

$$e(X, C_i \setminus X_i) \geq \frac{1}{4} \delta_1 \delta_2 d |X_i|$$

- $e(X', \overline{X}) \geq \frac{1}{4} \delta_1 \delta_2 d |X'|$

## Observation

If  $|X'| \geq \frac{1}{10} \delta_1 |X|$  then  $e(X, \overline{X}) \geq \frac{1}{80} \delta_1^2 \delta_2 \cdot 2d \cdot |X|$  and we are done.

Suppose

$$|X'| < \frac{1}{10} \delta_1 X$$

then we will show that there are enough number of edges between the clusters.

## Proof Contd..

- We have  $|X''| \geq (1 - \frac{1}{10}\delta_1)|X|$
- $|X''| \leq |X| \leq \frac{1}{2}nD$
- $|I''| \leq \frac{2}{3}n$
- There is a set of edges  $|M| \geq \frac{1}{2}\delta_1 D |I''|$  connecting vertices from  $I''$  to  $I'$

## Proof Contd....

- The corresponding  $d|M| \geq \frac{1}{2}\delta_1 dD|I''|$  edges connect vertices from  $\bigcup_{i \in I''} C_i$  with vertices of  $\bigcup_{i \in I'} C_i$
- At most  $\frac{1}{4}\delta_1 dD|I''|$  of  $d|M|$  many edges connect  $C_i \setminus X_i$  with vertices of  $\bigcup_{i \in I'} C_i$ . Why??
- $\text{Size}(|C_i \setminus X_i|) \geq \frac{1}{4}\delta_1 D$  (*times*  $d \cdot |I''|$ )-many edges
- Hence at least  $\frac{1}{4}\delta_1 dD|I''|$  of  $d|M|$  edges connect  $\bigcup_{i \in I''} X_i$  to  $\bigcup_{i \in I'} C_i$

## Proof Contd...

- No of  $d|M|$  edges that connects vertices from  $\bigcup_{i \in I''} C_i$  with vertices of  $X'$  is clearly at most  $d|X'|$
- Since  $|X'| \leq \frac{1}{10}\delta_1|X| \leq \frac{1}{6}\delta_1 D|I''|$ , there are  $\frac{1}{6}\delta_1 dD|I''|$  edges from  $\bigcup_{i \in I''} C_i$  to  $X'$
- $\frac{1}{4}\delta_1 dD|I''| - \frac{1}{6}\delta_1 dD|I''| = \frac{1}{12}\delta_1 dD|I''|$  edges connect vertices of  $\bigcup_{i \in I''} X_i$  with vertices of  $\bigcup_{i \in I'} C_i \setminus X_i$ .

## Proof Contd..

- AS  $|I''| \geq \frac{|X''|}{D}$  and  $|X''| \geq \frac{1}{2}|X|$ , we have atleast  $\frac{1}{48}\delta_1 2d|X|$  edges as required.

# Conclusion

We have constructed a constant degree expanders by applying replacement product twice.



THANK YOU ALL

