# The cohesive principle and the Bolzano-Weierstraß principle

Alexander P. Kreuzer

Technische Universität Darmstadt, Germany

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# The Bolzano-Weierstraß principle

A real number is a sequence of rational numbers with Cauchy-rate  $2^{-n}$ .

#### **Definition**

(BW):

Every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  has a cluster point.

Equivalently, every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  contains a Cauchy-subsequence  $(y_n)$  with Cauchy-rate  $2^{-n}$ , i.e. with

$$\forall n \, \forall i, j \ge n \, \left( |y_i - y_j| < 2^{-n} \right).$$

#### Definition

 $(BW_{weak})$ :

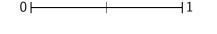
Every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  contains a Cauchy-subsequence  $(y_n)_n$ , i.e.

$$\forall n \,\exists k \,\forall i, j \geq k \, \left( |y_i - y_j| < 2^{-n} \right).$$

# Computing BW

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ .

Goal: Construct a subsequence  $(y_n)$  with  $\forall n \forall i, j \geq n \ (|y_i - y_j| < 2^{-n})$ .



 $0 \mid \frac{1}{2} \quad \frac{1}{2} \mid \frac{1}{2} \mid$ 

bi-partition argument

•  $y_n := \begin{cases} \text{next element of } (x_n) \\ \text{in the } n\text{-th partition} \end{cases}$ 

$$0 \longmapsto \frac{1}{4} \frac{1}{4} \longmapsto \frac{1}{2} \quad \frac{1}{2} \longmapsto \frac{3}{4} \frac{3}{4} \longmapsto -1$$

$$\frac{1}{2}$$
  $\frac{1}{2}$   $+$   $+$   $\frac{3}{4}$   $\frac{3}{4}$   $+$ 

; ; ;

- The partitions form a  $\Pi_2^0$ -0/1-tree.
- This is a  $\Pi_1^0$ -0/1-tree in 0'.
- WKL relativized to 0' yields an infinite branch and therefore computes the sequence of partitions.

## Theorem (Kohlenbach '98, Kohlenbach, Safarik '10, K. '10)

- For each computable sequence  $(x_n)$  there is a 0'-computable 0/1-tree T, such that an infinite branch of T computes a cluster point, and vice versa.
- Over RCA<sub>0</sub> the principles BW and WKL for  $\Sigma_1^0$ -trees are instance-wise equivalent.

By the low basis theorem:

## Corollary

BW has for computable instances a solution low relative to 0', i.e. the first Turing jump of a solution is computable in 0''.

# Computing BW<sub>weak</sub>

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ .

Goal: Construct a subsequence  $(y_n)$  with  $\forall n \exists k \ \forall i, j \geq k \ (|y_i - y_j| < 2^{-n})$  and compute the Turing jump of  $(y_n)$ .

It is clear that

$$\Phi_e^{(y_n)_n} \downarrow \quad \text{iff} \quad \exists k \, \Phi_e^{(y_n)_{n < k}} \downarrow.$$

Suppose that  $(y_n)_{n < m}$  is an initial segment that has already been computed. Deciding, whether there is an extension  $(y_n)_{n < l}$ , such that

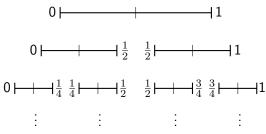
$$\Phi_e^{(y_n)_{n< l}}\downarrow$$

can be done in 0'.

# Computing BW<sub>weak</sub>

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ .

Goal: Construct a subsequence  $(y_n)$  with  $\forall n \exists k \ \forall i, j \geq k \ (|y_i - y_j| < 2^{-n})$  and compute the Turing jump of  $(y_n)$ .



- bi-partition argument
- now add at each step not only one element but finitely many elements of the chosen interval to (y<sub>n</sub>).
- Let  $(y_n)_{n < m}$  be the initial segment of  $(y_n)$  computed up to the k-th step.
  - At the k-th step extend this to  $(y_n)_{n< l}$  by elements in the k-th chosen interval, such that  $\Phi_k^{(y_n)_{n< l}}\downarrow$ , if possible.
- Then extend this by another element of the interval.

# Theorem (K.)

For each bounded, computable sequence  $(x_n)$  there is a Cauchy-subsequence  $(y_n)$ , such that  $(y_n)$  and  $(y_n)'$  are computable in a Turing degree that contains infinite branches of 0'-computable 0/1-trees.

## Corollary

BW<sub>weak</sub> has  $low_2$  solutions, i.e.  $(y_n)''$  is computable in 0''.

#### Proof.

$$(y_n)' \leq_T 0' + \mathsf{WKL} \implies (y_n)'' \leq_T 0''$$

 $\mathsf{BW}_\mathsf{weak}$  does not compute 0' and is therefore strictly weaker than  $\mathsf{BW}$ .

# The cohesive principle

Write  $X \subseteq^* Y$  if  $X \setminus Y$  is finite.

#### Definition

• A set X is *cohesive* for a sequence of set  $(R_n)_n \subseteq 2^{\mathbb{N}}$  if

$$X \subseteq^* R_n \vee X \subseteq^* \overline{R_n}$$
 for each  $n$ .

• The cohesive principle (COH) states that for each  $(R_n)_n$  there is an infinite cohesive set X.

## Theorem (K.)

- For each sequence  $(x_n)_n \subseteq \mathbb{R}$  there exists  $(R_n)_n \subseteq 2^{\mathbb{N}}$ , such that from an infinite cohesive set for  $(R_n)$  one can compute a Cauchy-subsequence of  $(x_n)$  and vice versa.
- RCA<sub>0</sub>  $\vdash$  BW<sub>weak</sub>  $\leftrightarrow$  COH  $\land$   $B\Sigma_2^0$ Moreover, this equivalence also holds instance-wise.

#### Theorem

- COH and hence also BW<sub>weak</sub> do not compute solutions to WKL in general. (Cholak, Jockusch, Slaman '01)
- There are instance of these principle which have no low solutions. (Jockusch, Stephan '93)

Proof of the  $low_2$ -ness of BW<sub>weak</sub> is a streamlined version of the  $low_2$ -ness of COH (Jockusch, Stephan '93).

# Theorem (Chong, Slaman, Yang '10)

 $\mathsf{RCA}_0 + \mathsf{COH} + B\Sigma_2^0$  and thus  $\mathsf{RCA}_0 + \mathsf{BW}_\mathsf{weak}$  are  $\Pi_1^1$ -conservative over  $\mathsf{RCA}_0 + B\Sigma_2^0$ .

## Theorem (K., Kohlenbach '10)

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If \mathsf{WKL}_0 + \mathsf{BW}_\mathsf{weak} \vdash \forall f \, \exists y \, \phi(f,y) for quantifier free \phi, then one can extract from a given proof a primitive recursive function(al) t such that \forall f \, \phi(f,t(f)).
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"Proof mining"

# Bolzano-Weierstraß in the weak topology

We consider the Hilbert space  $\ell_2=(\mathbb{R}^\mathbb{N},\langle\cdot,\cdot\rangle)$ . An element of  $\ell_2$  is given by a Cauchy-sequence  $(w_n)_n$  of finite dimensional and rational approximations, i.e.  $w_n\in\mathbb{Q}^{<\mathbb{N}}$ , with Cauchy-rate  $2^{-n}$  with respect to  $\|\cdot\|$ .

#### Definition

(weak-BW): Every  $\|\cdot\|$ -bounded sequence  $(x_n) \subseteq \ell_2$  has a weak cluster point x, i.e.  $\forall y \in \ell_2 \lim_{n \to \infty} \langle y, x_n \rangle = \langle y, x \rangle$ .

# Theorem (K.)

- For each bounded sequence  $(x_n) \subseteq \ell_2$  there is a weak cluster point x computable in 0''.
- There is a bounded and computable sequence  $(x_n) \subseteq \ell_2$ , such that each weak cluster point of it computes 0''.
- Over RCA<sub>0</sub> the principles Π<sub>2</sub>-CA and weak-BW are instance-wise equivalent.

# Summary

- BW is equivalent to WKL for 0'-computable trees.
- BW<sub>weak</sub> is equivalent to COH.
  - Hence, it does not imply 0'.
  - It admits extraction of primitive recursive terms.
- weak-BW is equivalent to 0''.

### References



Alexander P. Kreuzer

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Alexander P. Kreuzer,

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