# Ramsey's theorem for pairs and program extraction

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- Ramsey's Theorem for pairs  $(RT_2^2)$ Every 2-coloring of pairs of  $\mathbb N$  contains an infinite homogeneous set.
- Each sequence of real numbers contains an infinite monotone subsequence (ADS)
- Each bounded sequence of real numbers contains a slowly converging subsequence (BW<sub>weak</sub>).

$$(x_n)$$
 converges slowly:  $\forall k \, \exists n \, \forall m \geq n \, |x_n - x_m| < 2^{-k}$   
 $(x_n)$  converges fast:  $\forall k \quad \forall m \geq k \, |x_k - x_m| < 2^{-k}$ 

## Theorem (K. '10)

$$\mathsf{RCA}_0 \vdash \mathsf{BW}_{\mathsf{weak}} \leftrightarrow \mathsf{COH} + B\Sigma^0_2$$

- $\mathsf{RT}^2_2 \to \mathsf{ADS} \to \mathsf{BW}_{\mathsf{weak}}$
- RT<sub>2</sub><sup>2</sup> ← ADS ← BW<sub>weak</sub> (Hirschfeldt, Shore '07)
- BW<sub>weak</sub> proves  $I\Sigma_1^0$  and hence primitive recursion.
- Solutions to computable instances of these principles are in general not computable in 0'.
- Computable instances of these principle have  $low_2$  solutions, i.e. solutions X, such that X'' is computable in 0''. (Cholak, Jockusch, Slaman '01, Hirschfeldt, Shore '07)

## Program extraction

## Theorem (K., Kohlenbach '10)

lf

$$\mathsf{WKL}_0 + I\Sigma_2^0 + \mathsf{RT}_2^2 \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a term t of Ackermann type, such that

$$\forall x A(x, t(x)).$$

## Theorem (K., Kohlenbach '10, K. '10)

lf

$$\mathsf{WKL}_0 + B\Sigma_2^0 + \mathsf{ADS} + \mathsf{BW}_{\mathsf{weak}} \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a primitive recursive term t, such that

$$\mathsf{PRA} \vdash \forall x \, A(x, t(x)).$$

#### Motivation

- Proof mining
- Hilbert's program
- Reverse mathematics:
  New proofs for the facts that
  - the functions provable recursive by  $RT_2^2$  are already provably by  $I\Sigma_2^0$ , cf. Cholak, Jockusch, Slaman '01,
  - ADS and the chain antichain principle does *not* imply  $I\Sigma_2^0$ , cf. Chong, Slaman, Yang '10.

#### Theorem

lf

$$\mathsf{RCA}_0 + \mathsf{BW}_{\mathsf{weak}} \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a primitive recursive term t, such that

$$\mathsf{PRA} \vdash \forall x \, A(x, t(x)).$$

$$RCA_0$$

$$\mathsf{RCA}_0 :\equiv \mathsf{Basic} \; \mathsf{Arithmetic} + \Delta^0_1\text{-}\mathsf{CA} + I\Sigma^0_1$$
 If  $\mathsf{RCA}_0 + \mathsf{WKL} \vdash \forall x \, \exists y \, A(x,y)$ , then exists primitive recursive term  $t$ , such that  $\mathsf{PRA} \vdash A(x,t(x))$ .

 $RCA_0^*$ 

$$\mathsf{RCA}_0^* :\equiv \mathsf{Basic} \; \mathsf{Arithmetic} \; \mathsf{plus} \; 2^x + \Delta^0_1 \mathsf{-CA} + I\Sigma^0_0$$

$$\Pi^0_1\text{-CA}(\phi) :\equiv \exists X \, \forall n \, (n \in X \leftrightarrow \forall y \, \phi(n, y))$$

## Theorem (Kohlenbach '98)

If for closed  $\phi$ 

$$\mathsf{RCA}_0^* + \mathsf{WKL} + \mathsf{\Pi}_1^0 \mathsf{-CA}(\phi) \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a primitive recursive term t, such that

$$\mathsf{PRA} \vdash A(x, t(x)).$$

With this it is sufficient to show

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$$\mathsf{RCA}_0^* + \mathsf{BW}_{\mathsf{weak}} \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a  $\phi$ , such that

$$\mathsf{RCA}_0^* + \mathsf{WKL} + \Pi_1^0 \mathsf{-CA}(\phi) \vdash \forall x \,\exists y \, A(x,y).$$

There exists a  $\phi$ , such that

$$\mbox{RCA}_0^* + \mbox{WKL} + \Pi_1^0 - \mbox{CA}(\phi[X]) \vdash \\ \exists Y \quad Y \mbox{ codes a slowly converging subsequence of } X$$

This yields that RCA $_0^*$  + WKL +  $\Pi_1^0$ -CA $(\phi)$  proves iterations of instances of BW<sub>weak</sub>.

Together with a proof- / term-normalization the result follows.

With this it is sufficient to show

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$$\mathsf{RCA}_0^* + \mathsf{BW}_{\mathsf{weak}} \vdash \forall x \,\exists y \, A(x,y),$$

then there exists a  $\phi$ , such that

$$\mathsf{RCA}_0^* + \mathsf{WKL} + \mathsf{\Pi}_1^0 \mathsf{-CA}(\phi) \vdash \forall x \,\exists y \, A(x,y).$$

For each  $\psi$  there exists a  $\phi$ , such that

$$\begin{split} \mathsf{RCA}_0^* + \mathsf{WKL} + \Pi_1^0\text{-}\mathsf{CA}(\phi[X]) \vdash \\ &\exists Y \ \left(Y \text{ codes a slowly converging subsequence of } X \\ & \wedge \Pi_1^0\text{-}\mathsf{CA}(\psi[X,Y])\right). \end{split}$$

This yields that RCA $_0^*$  + WKL +  $\Pi_1^0$ -CA $(\phi)$  proves iterations of instances of BW $_{\rm weak}$ .

Together with a proof- / term-normalization the result follows.

#### Proofwise low

#### Theorem (K', Kohlenbach '10)

Let P be a principle of the form

$$\forall X \exists Y P(X,Y)$$

where P is  $\Pi^0_3$  and proofwise low over  $\mathsf{RCA}^*_0 + \mathsf{WKL}$ , i.e.

$$\begin{split} \mathsf{RCA}_0^* + \mathsf{WKL} + \Pi^0_1\text{-}\mathsf{CA}(\phi[X]) \vdash \exists Y \ P(X,Y) \\ & \wedge \Pi^0_1\text{-}\mathsf{CA}(\psi[X,Y]) \big). \end{split}$$

If

$$\mathsf{WKL}_0 + B\Sigma_2^0 + \mathsf{P} \vdash \forall x \,\exists y \, A(x,y)$$

then one can extract a primitive recursive term t, such that

$$\mathsf{PRA} \vdash \forall x \, A(x, t(x)).$$

# Proofwise low (for $I\Sigma_2^0$ )

## Theorem (K', Kohlenbach '10)

Let P be a principle of the form

$$\forall X \exists Y P(X,Y)$$

where P is  $\Pi_3^0$  and proofwise low over  $\mathsf{WKL}_0$ , i.e.

$$\begin{aligned} \mathsf{WKL}_0 + \Pi_1^0\text{-CA}(\phi[X]) \vdash \exists Y \ P(X,Y) \\ & \wedge \Pi_1^0\text{-CA}(\psi[X,Y]) \big). \end{aligned}$$

If

$$\mathsf{WKL}_0 + \underline{I\Sigma_2^0} + \mathsf{P} \vdash \forall x \,\exists y \, A(x,y)$$

then one can extract a term t of Ackermann type, such that

$$\forall x A(x, t(x)).$$

# Summary of the strength of the Bolzano-Weierstraß principle

- The Bolzano-Weierstraß principle implies ACA<sub>0</sub> and hence PA.
  In general no primitive recursive terms. (Friedman '76)
- Sequences of instances of the Bolzano-Weierstraß principle over RCA<sub>0</sub>\* + WKL allow extraction of primitive recursive terms. (Kohlenbach's elimination of Skolemfunctions for monotone formulas '98)
- We showed that general uses of the weak Bolzano-Weierstraß principle and ADS even over  $RCA_0 + WKL$  allow extraction of primitive recursive terms.

#### References I

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