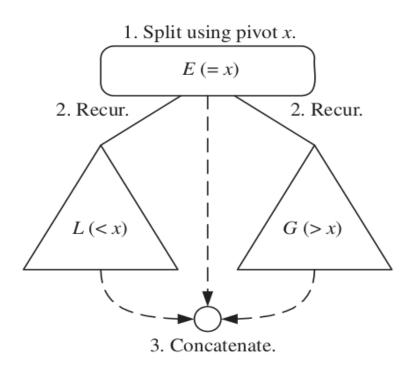
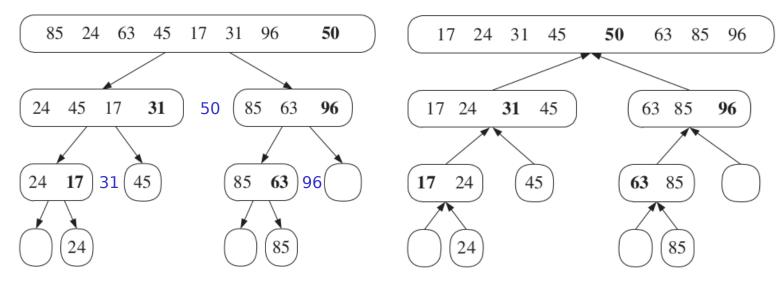
QUICK SORT - Steps:

- Divide: If S has at least two elements (nothing needs to be done if S has zero or one element), select a specific element x from S, which is called the pivot. As is common practice, choose the pivot x to be the last element in S. Remove all the elements from S and put them into three sequences:
 - L, storing the elements in S less than x
 - E, storing the elements in S equal to x
 - G, storing the elements in S greater than x.

(If the elements of S are all distinct, E holds just one element—the pivot.)

- 2. **Recur:** Recursively sort sequences L and G.
- 3. **Conquer:** Put the elements back into S in order by first inserting the elements of L, then those of E, and finally those of G.





```
{85 24 63 45 50 31 96 17}
                                                        {85 24 63 45 17 31 50 96}
L={} E={17} G={85 24 63 45 50 31 96}
                                                        L=\{85\ 24\ 63\ 45\ 31\ 50\}\ E=\{96\}\ G=\{\}
  partition (arr[], low, high)
                                                           Initial: 85 24 63 45 17 31 96 50
                                                                    low = 0, high = 7
     pivot = arr[high];
     i = (low - 1)
                                                            Pivot: arr[7] = 50
     for (j = low; j < high; j++)
                                                           i = 3
       // If current element is smaller than the pivot
                                                           i = 6
       if (arr[j] < pivot)</pre>
                                                            85 24 63 45 17 31 96 50
          i++; // increment index of smaller element
                                                            24 85 63 45 17 31 96 50
          swap arr[i] and arr[j]
                                                            24 45 63 85 17 31 96 50
                                                            24 45 17 85 63 31 96 50
     }
                                                            24 45 17 31 63 85 96 50
     swap arr[i+1] and arr[high])
                                                            24 45 17 31 50 85 96 63
     return (i + 1)
  }
                                                           24 45 17 31 | 50 | 85 96 63
   /* low --> Starting index, high --> Ending index */
   quickSort(arr[], low, high)
                                                                           {1 2 3 4 5}
     if (low < high) // i.e. if arr[] is of length >= 2
                                                                         {1 2 3 4} {5} {}
        pi = partition(arr, low, high);
                                                                     {1 2 3} {4} {}
        /* pi is partitioning index, i.e. index of pivot */
                                                                    {1'2} {3} {}
        quickSort(arr, low, pi - 1); // Before pi, i.e. L
        quickSort(arr, pi + 1, high); // After pi, i.e. G
     }
   }
                                                            Worst case time compexity:
                                                            (n-1) levels, and
   partition (arr[], low, high)
                                                            O(n) operations at each level
                                                            (counted over all partition functions)
     pivot = arr[high];
                                                            = O(n^2)
     i = low
     for (j = low; j < high; j++)
                                                            Best case time complexity:
        // If current element is smaller than the pivot
        if (arr[j] < pivot)</pre>
                                                            O(log(n)) levels, and
                                                            O(n) operations at each level
           swap arr[i] and arr[j]
                                                            (counted over all partition functions)
           i++; // increment index of smaller element
                                                            = O(n \log(n))
```

Average case time complexity:

 $O(n \log(n))$

swap arr[i] and arr[high])

return (i)