

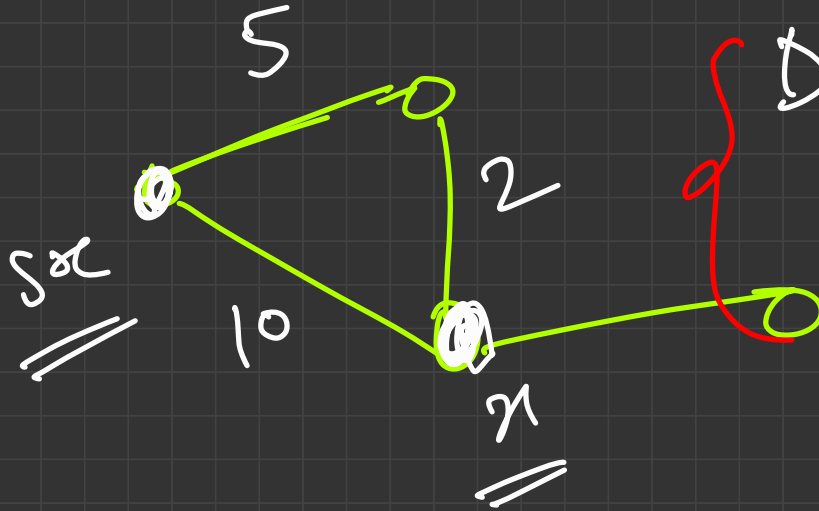
Shortest path algorithms

Graphs 2

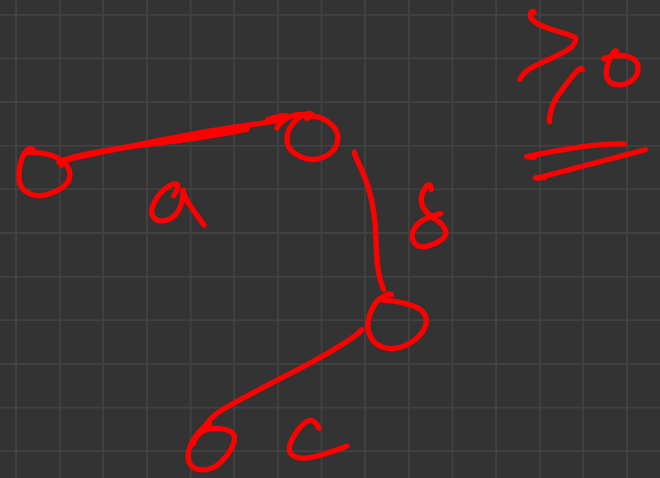
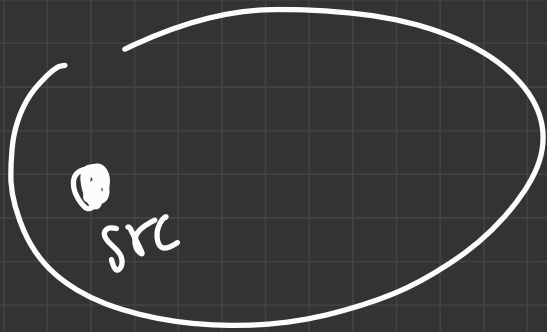
✓
Bipartite Graphs, Dijkstra, Bellman Ford

↓
7,0

-Priyansh Agarwal



Dijkstra = Bellman Ford



Bi-partite Graphs

✓ Definition

✗ Algorithm

✓ Odd Length Cycle

✓ Tree Property

✓ Problem: [Link](#)

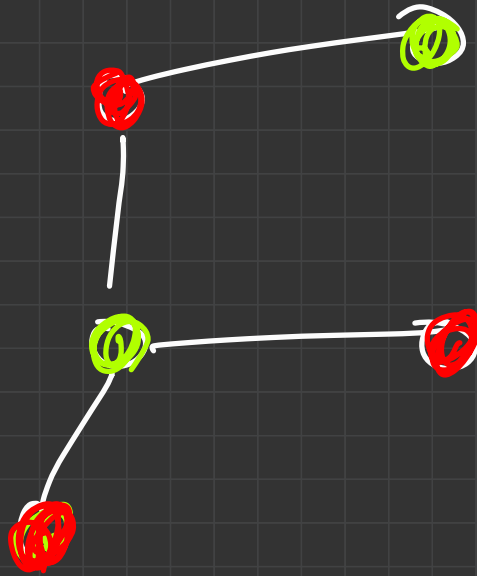
→ not bi-partitely colourable

H.W

Bipartite Graphs :

nodes

Every 2 adj.
have different
color



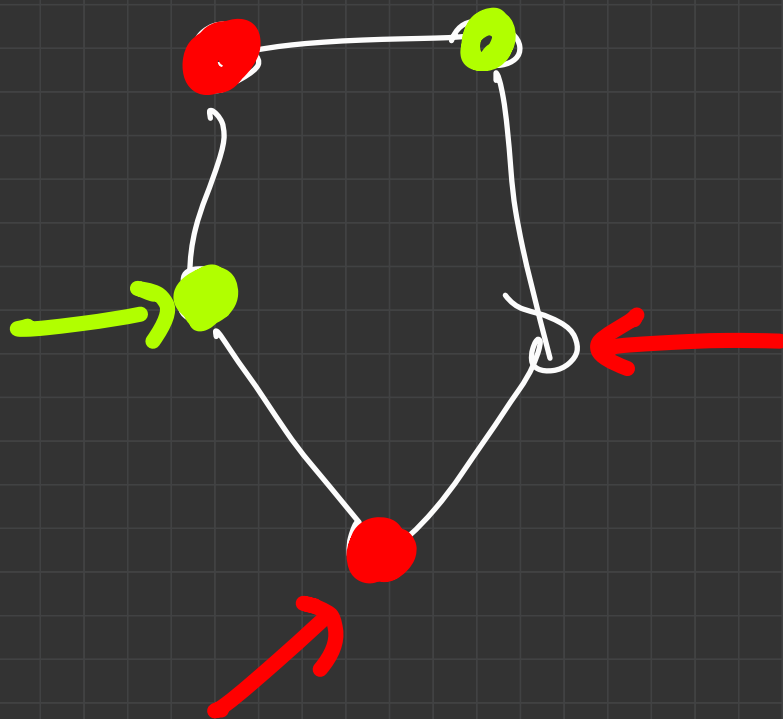
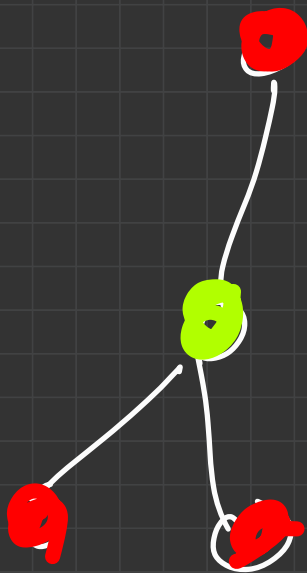
② Color

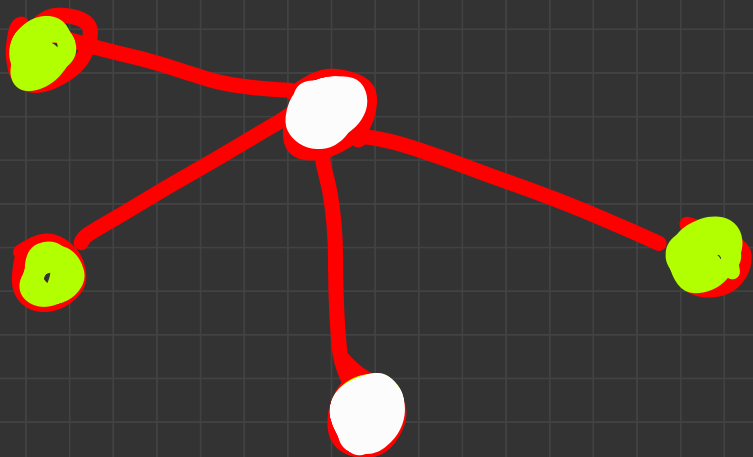


connected

Given a \wedge graph
it can only have
2 possible colorings
(bi-partite)

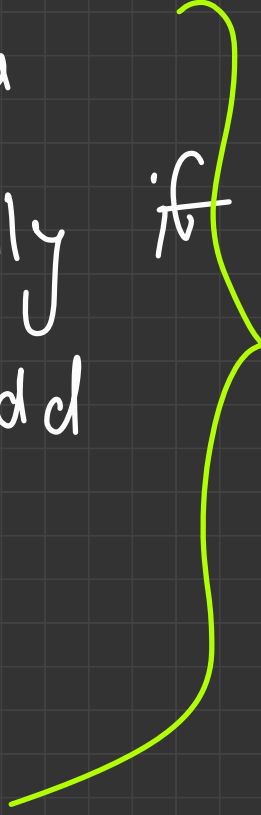
2

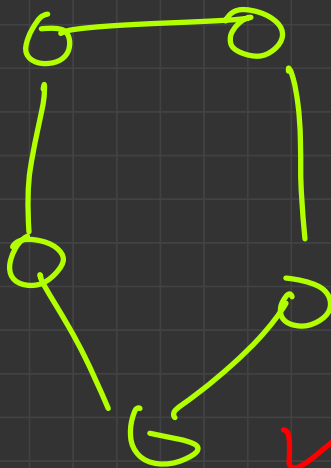
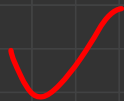
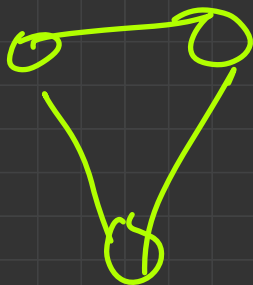




①

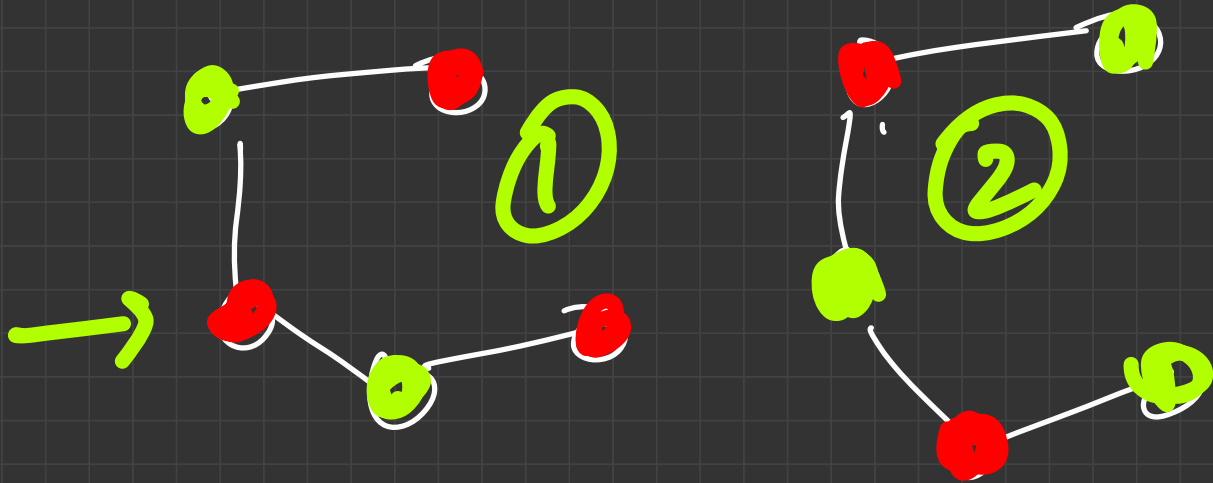
A graph can be colored in a
bi-partite manner if and only if
it does not contain an odd
length cycle



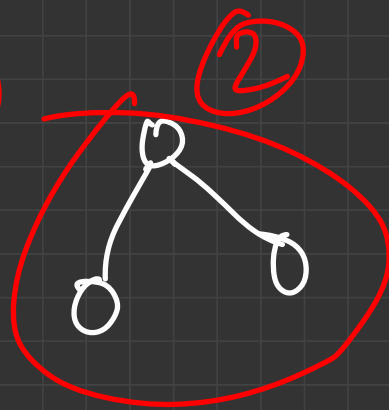
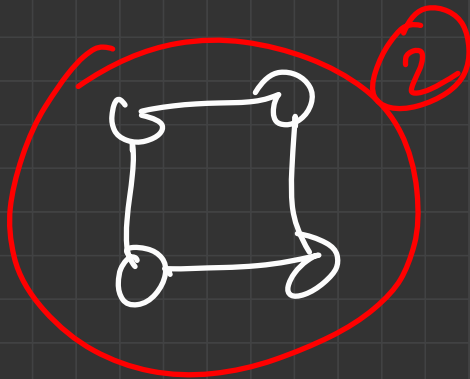
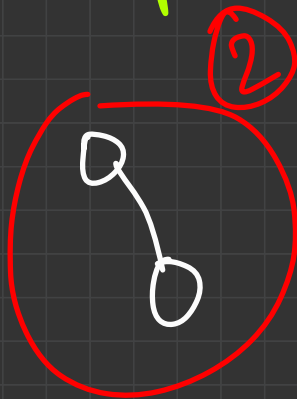


A connected graph can have either

0 bipartite coloring or $\underline{\underline{2}}$

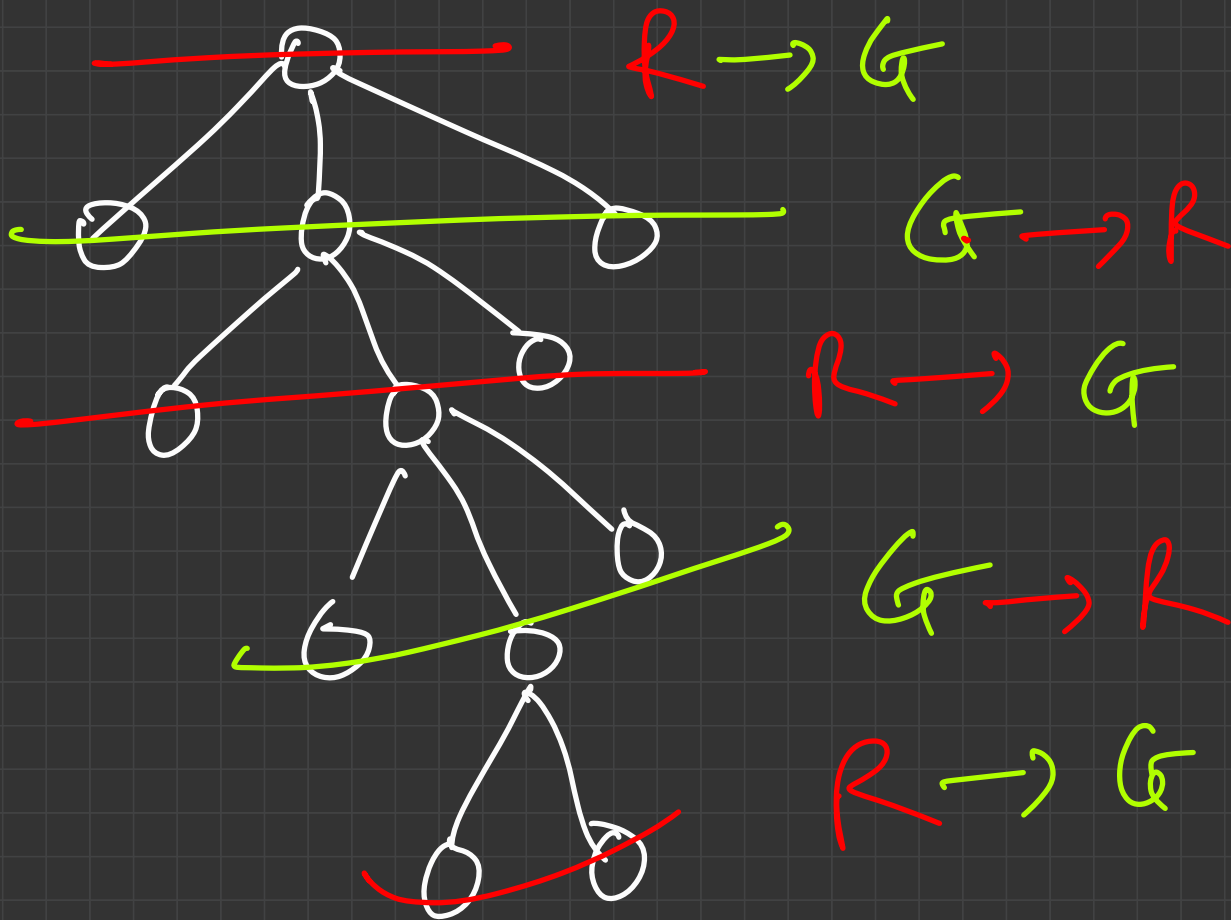


If a graph contains n
connected components how many
dis-joint colorings can you have

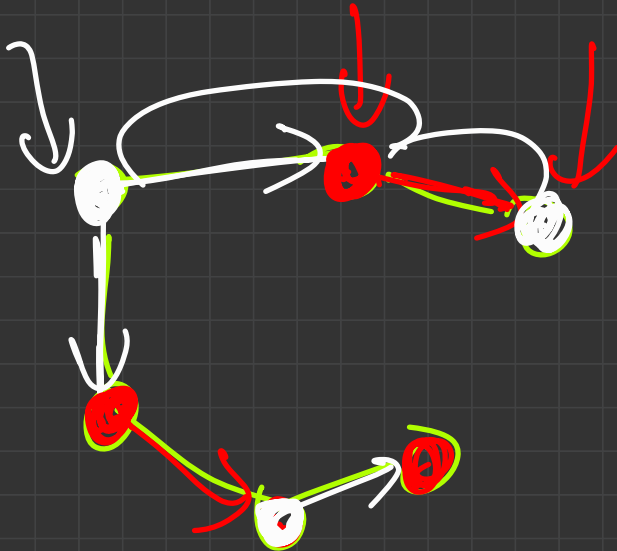


$$= 2^n$$

(2)



Choose a random node and color
it with first color



```
vector<int> colors(n, 0);
```

```
bool dfs(int curr, edges, colors)    W = 1  
                                     R = 2
```

```
{ for (int neighbour: edges[curr])
```

```
    if (colors(neighbour) == 0)
```

```
        colors(neighbour) = 1 if curr color = 2  
                             2 if curr color = 1
```

```
    curr = curr && dfs(neighbour, edges, colors)
```

```
    else if (colors(neighbour) == colors[curr])
```

return false;

return ans;

}

$O(n)$

bool ans = true

for(int i=0; i<n; i++) {
if (color[i] == 0)

color[i] = 1

ans = ans && dfs(i, edges, color) }

Dijkstra (most important)

- Single Source Shortest Path Algorithm - Idea + Visualization

- Non-negative edge weights

- Proof/Intuition:

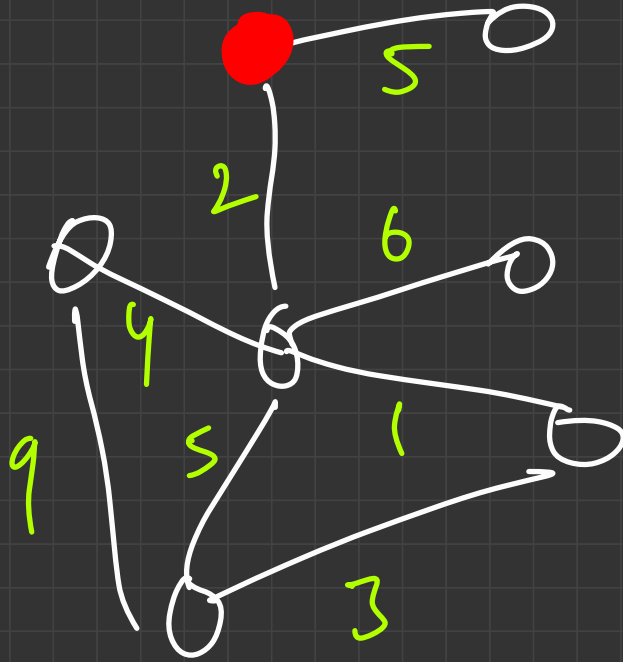
- On every iteration the marked vertex is the one that can never have a better distance later on.

- Code

- Retrieving the shortest path?

- Problems

- Google Interview Problem

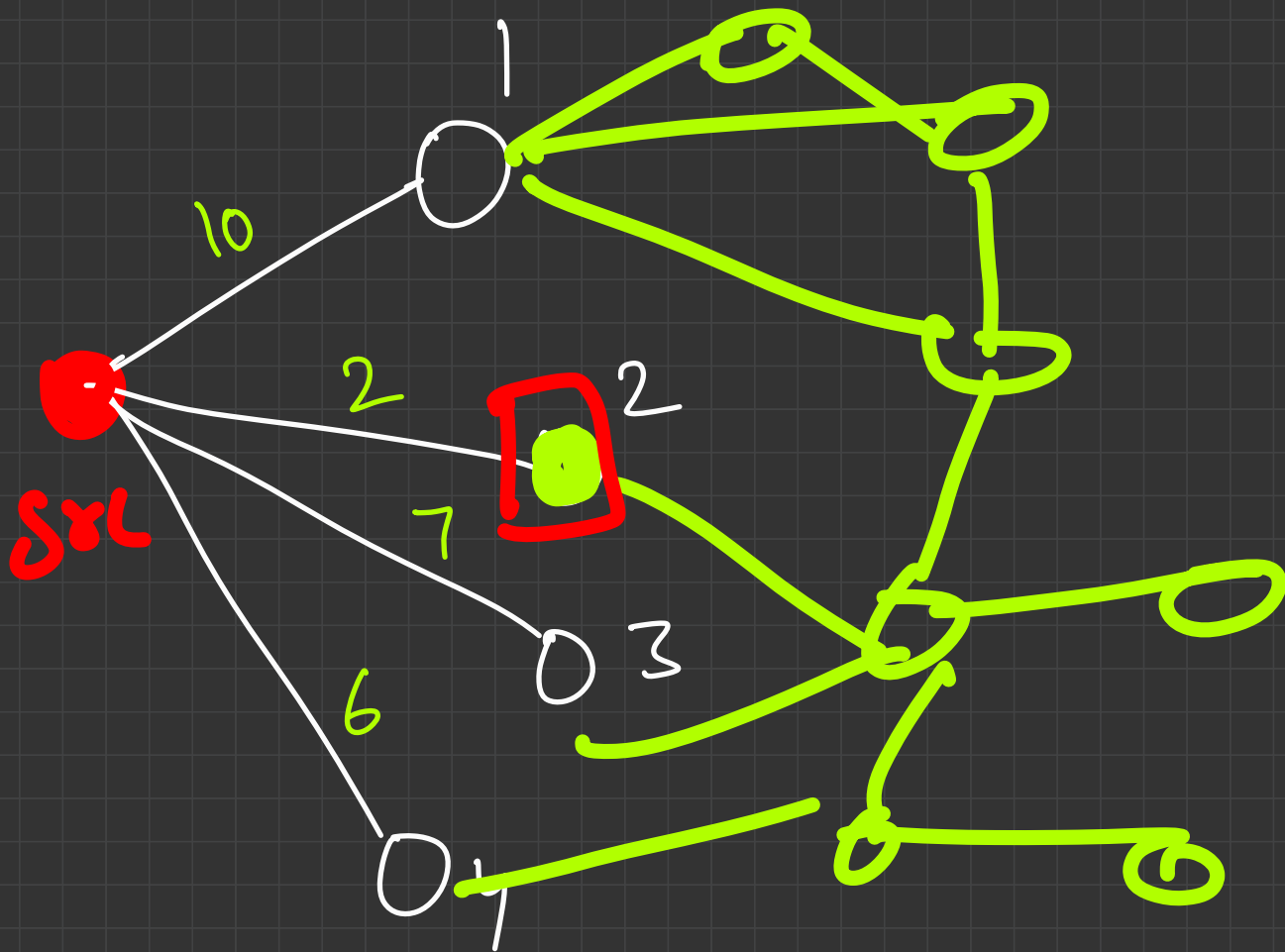


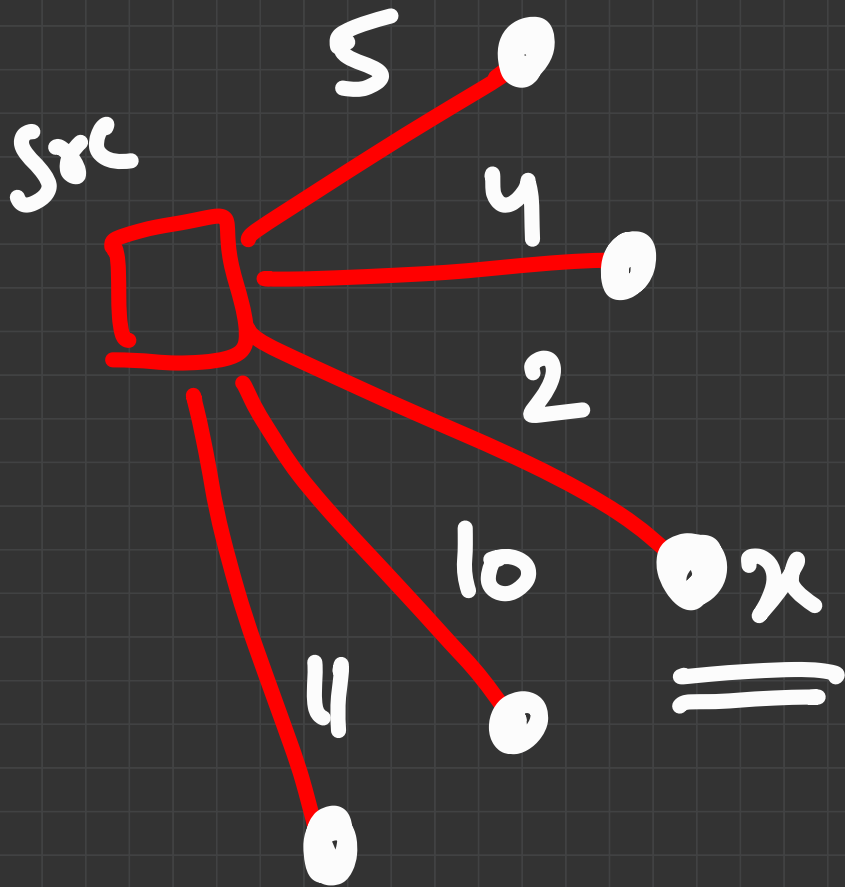
all edge weights

≥ 0

non-negative

src





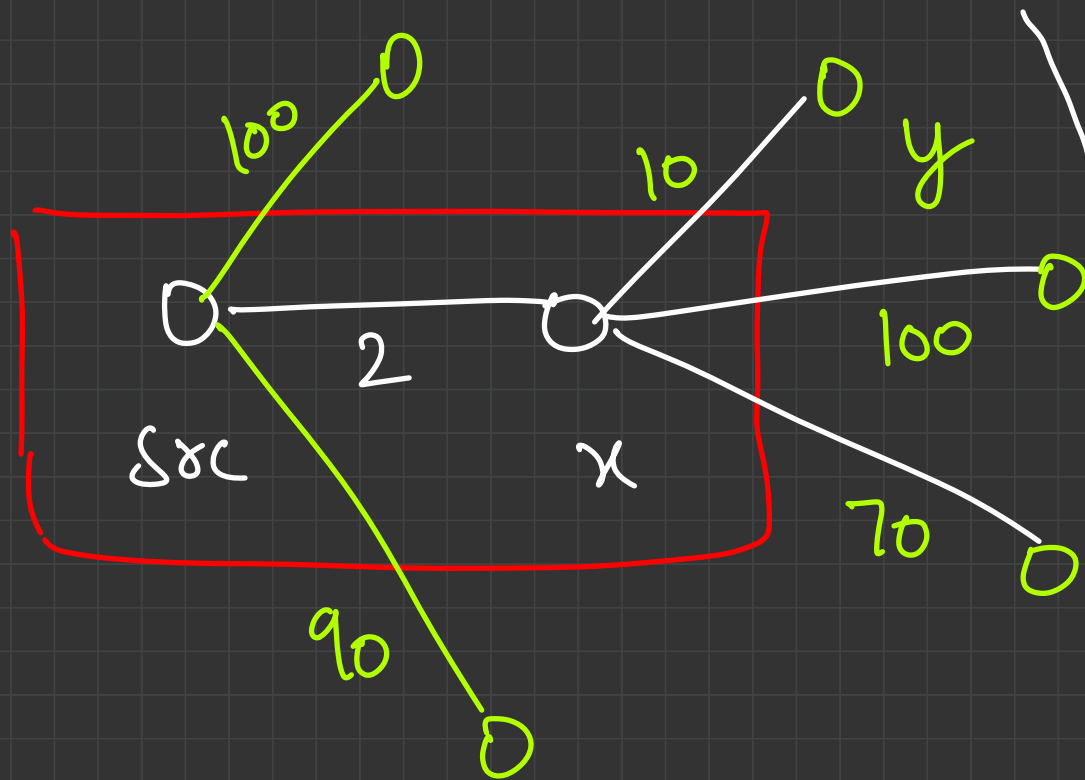
$src = 0$



2



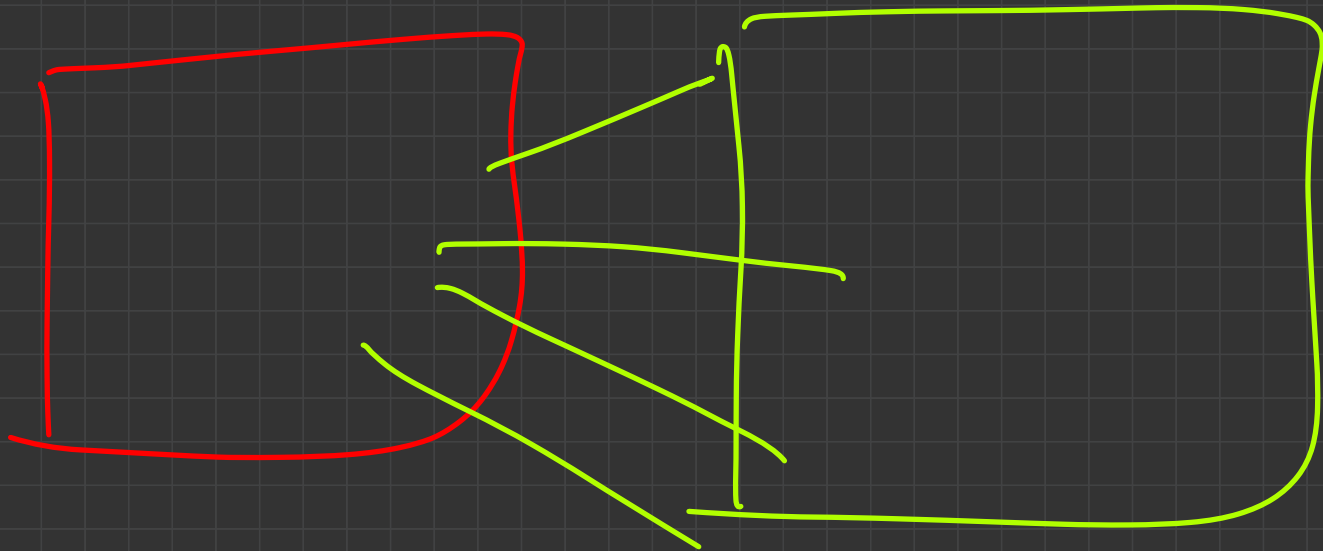
$x = 2$



① Coming out of x

② Coming out of src

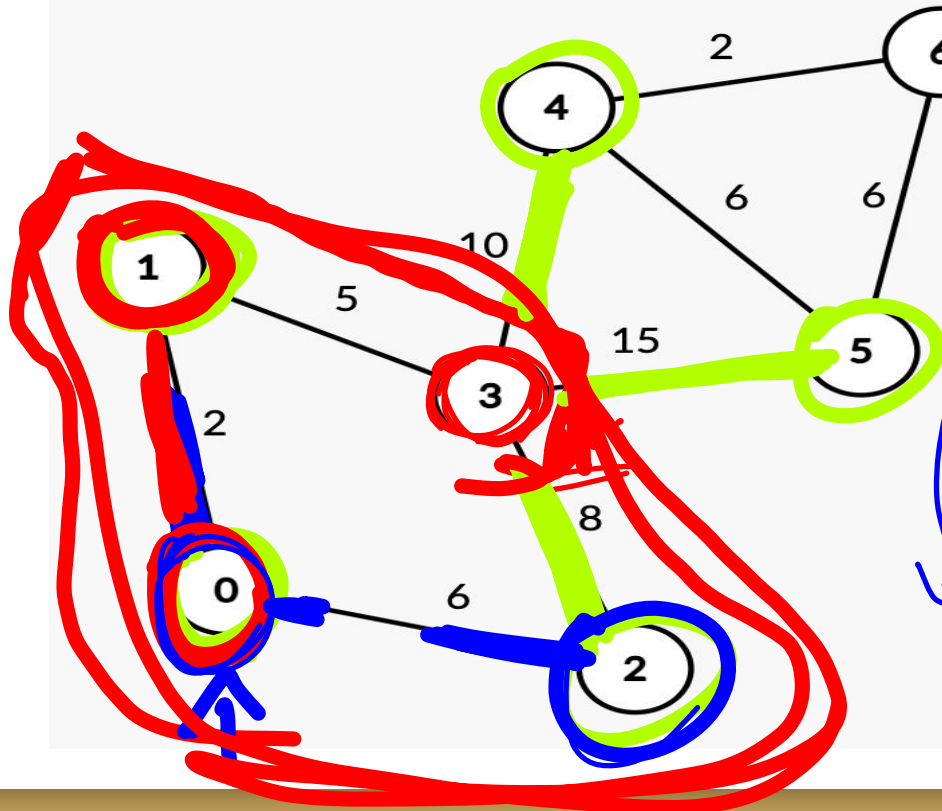
12



Dijkstra Visualization

dist 7 5 8 0 10 15 inf ..

0 1 2 3 4 5 6 ..
~~✓~~ ✓ ✓ ✓ ✓ ✓



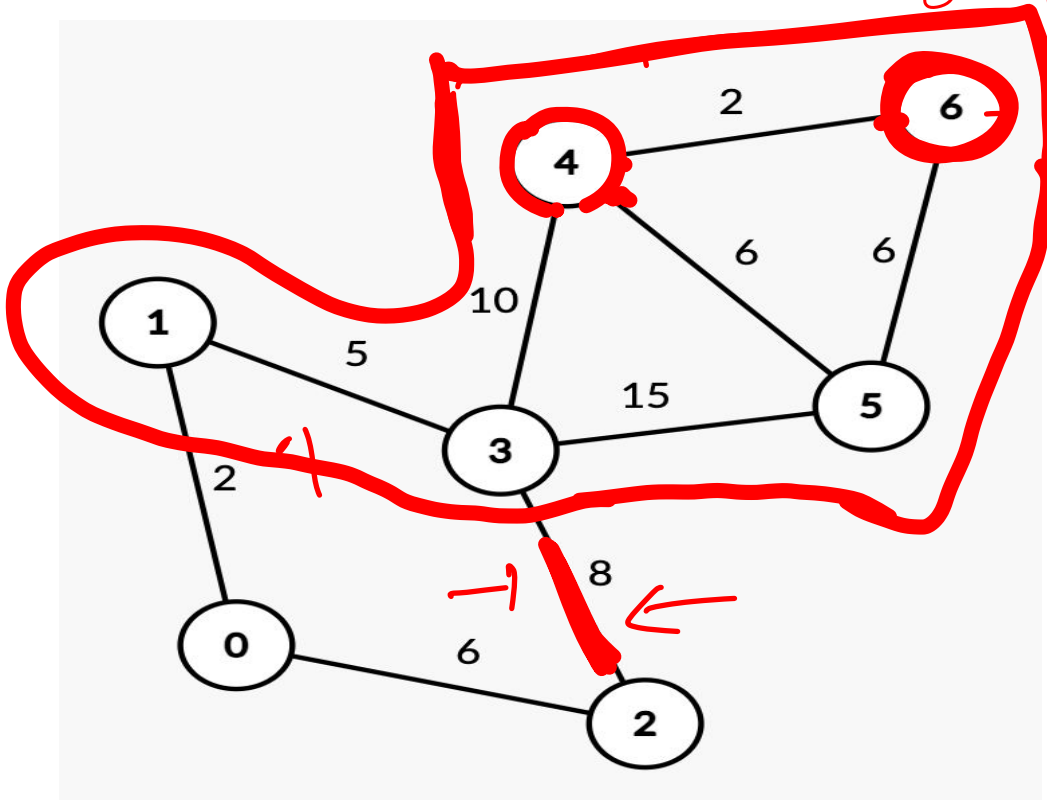
edge weight + dist(u)
 $= 2 + 5 = 7$

$6 + 7 = 13$

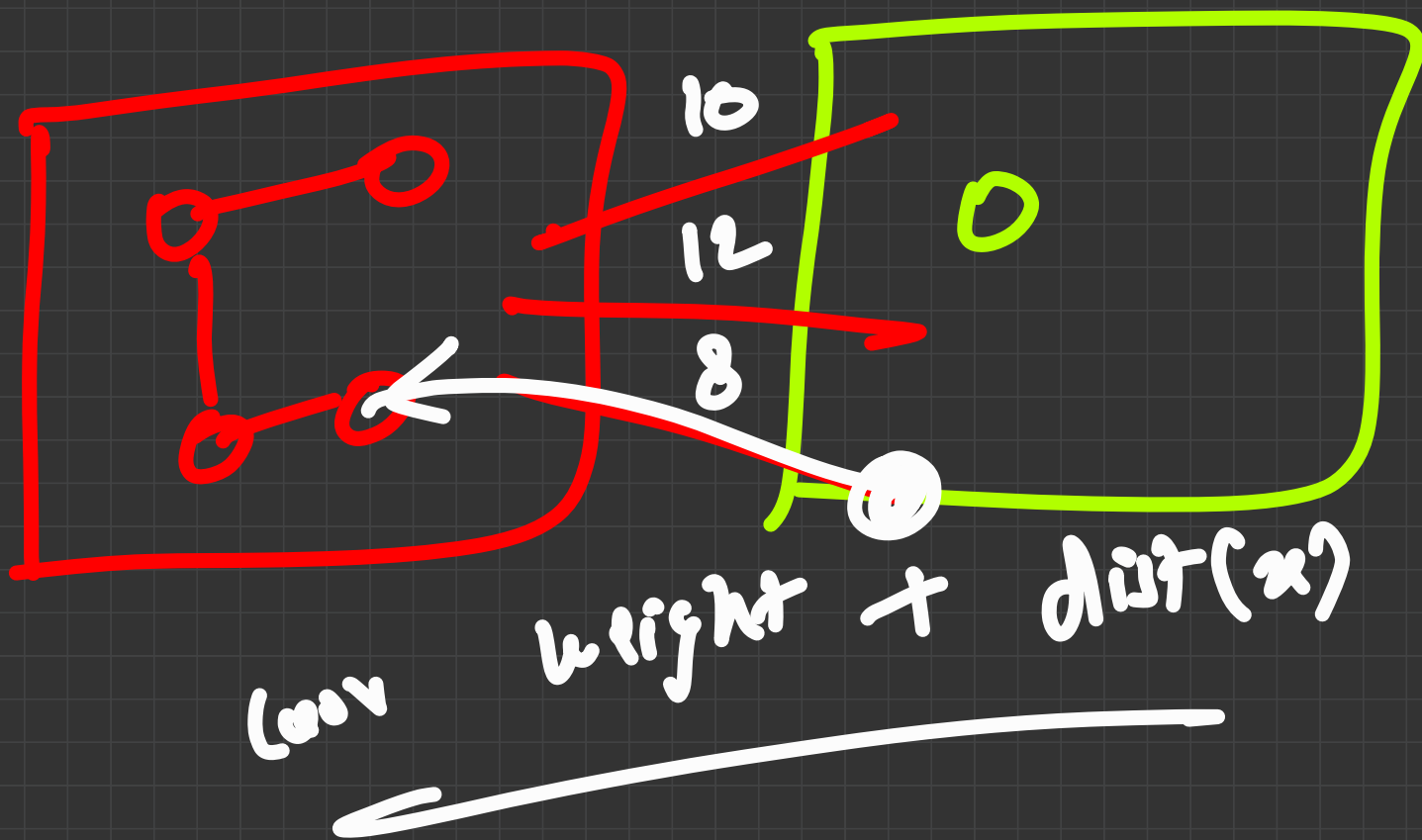
Dijkstra Visualization

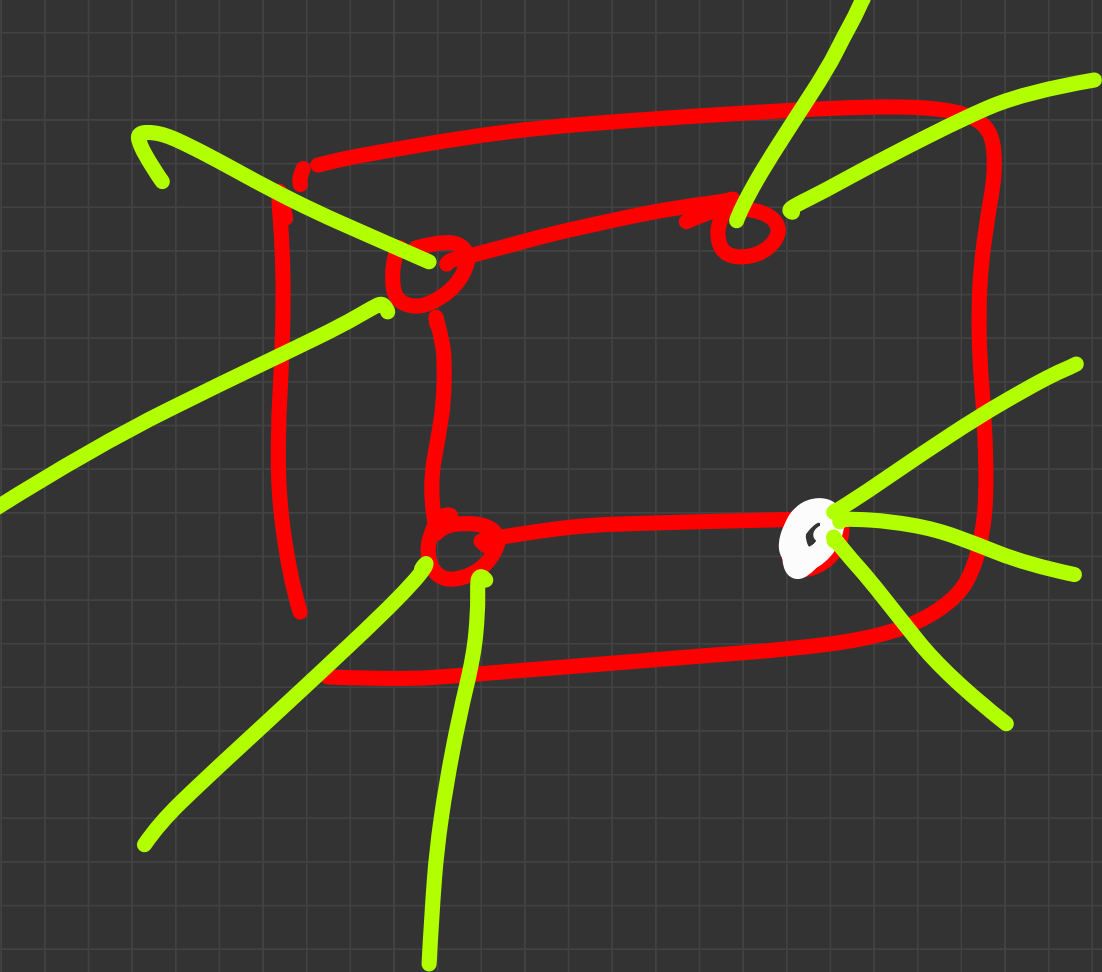
dist

10	15	18	10	0	6	2
0	1	2	3	4	5	6



$$\begin{aligned} 5 + 10 \\ \hline 8 + 10 \end{aligned}$$



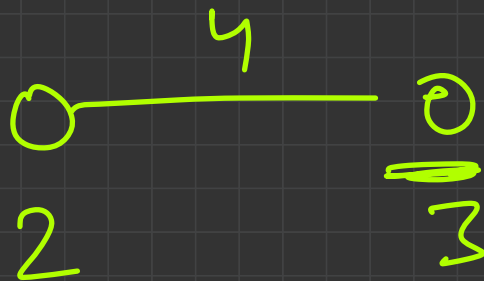


$(14, 3), (15, 3), (20, 4)$

Set

$(10, 2)$

$vis(3) = \text{true}$



① Given a weighted undirected connected Graph
find out the shortest path from source

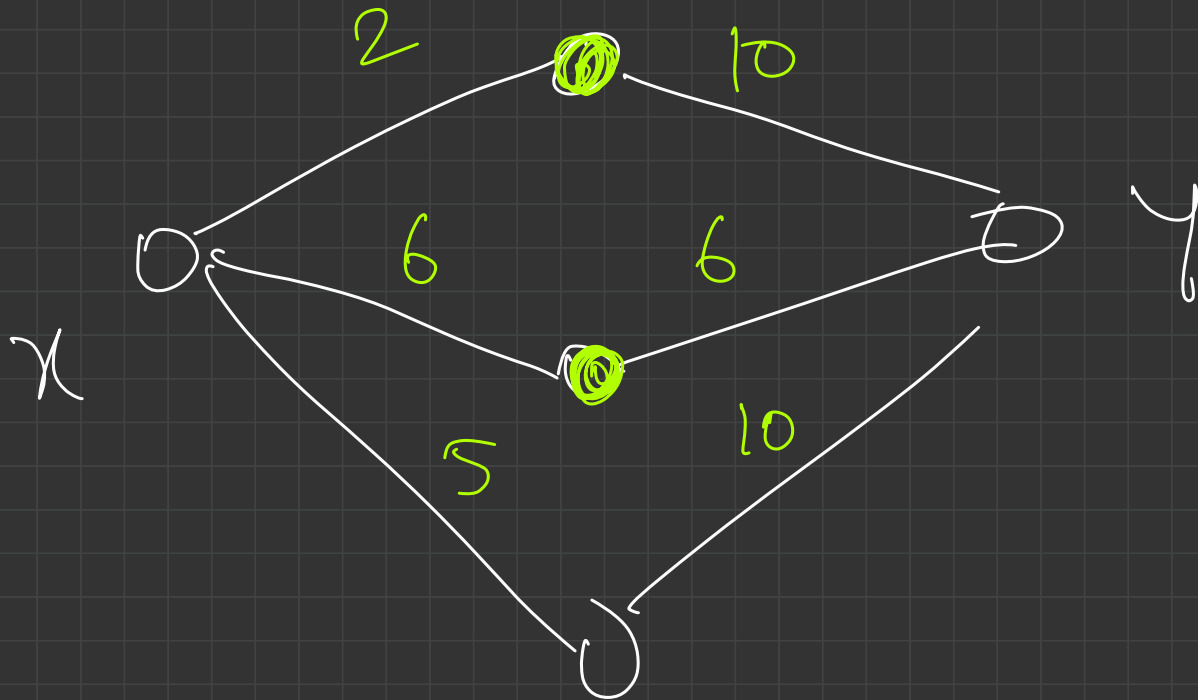
X to Y

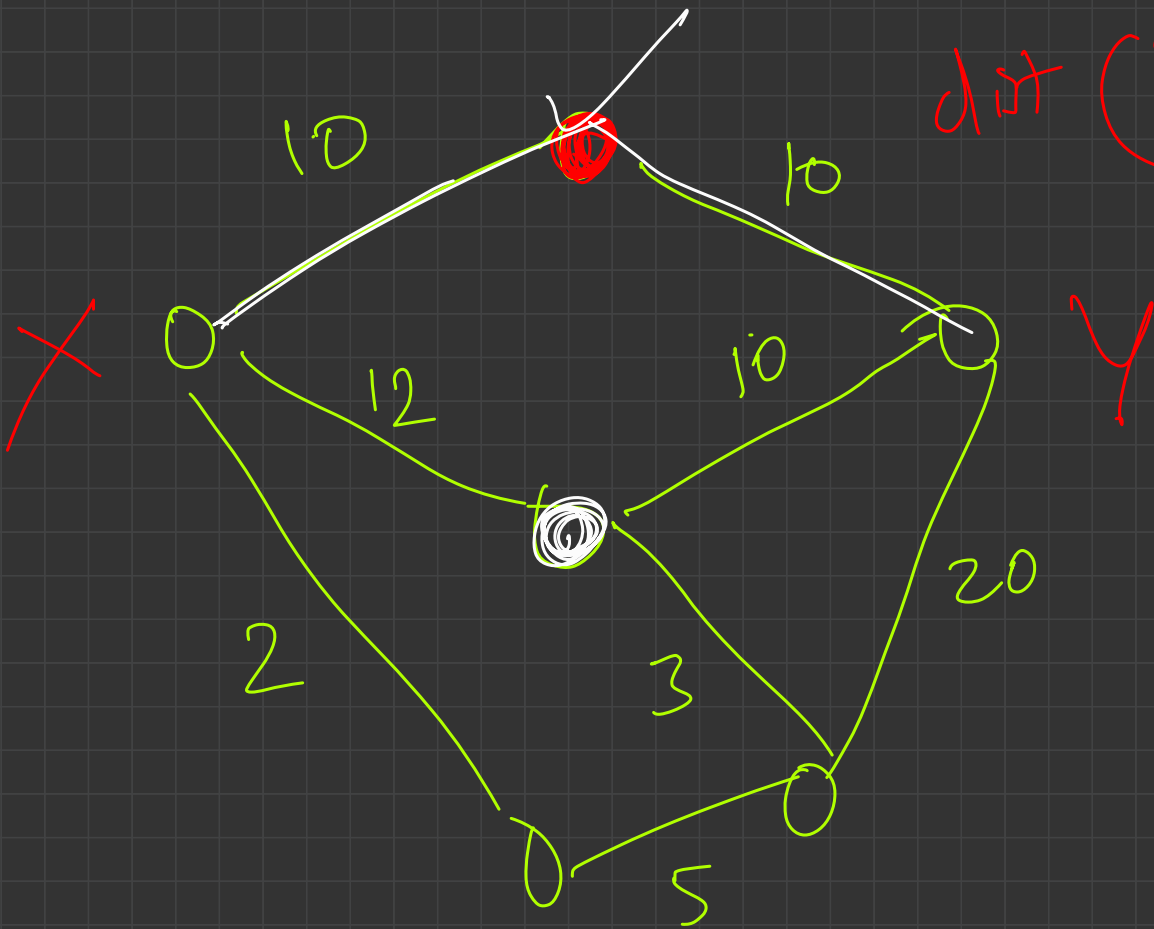


Dijkstra

$n \leq 10^5$
 $m \leq 10^5$

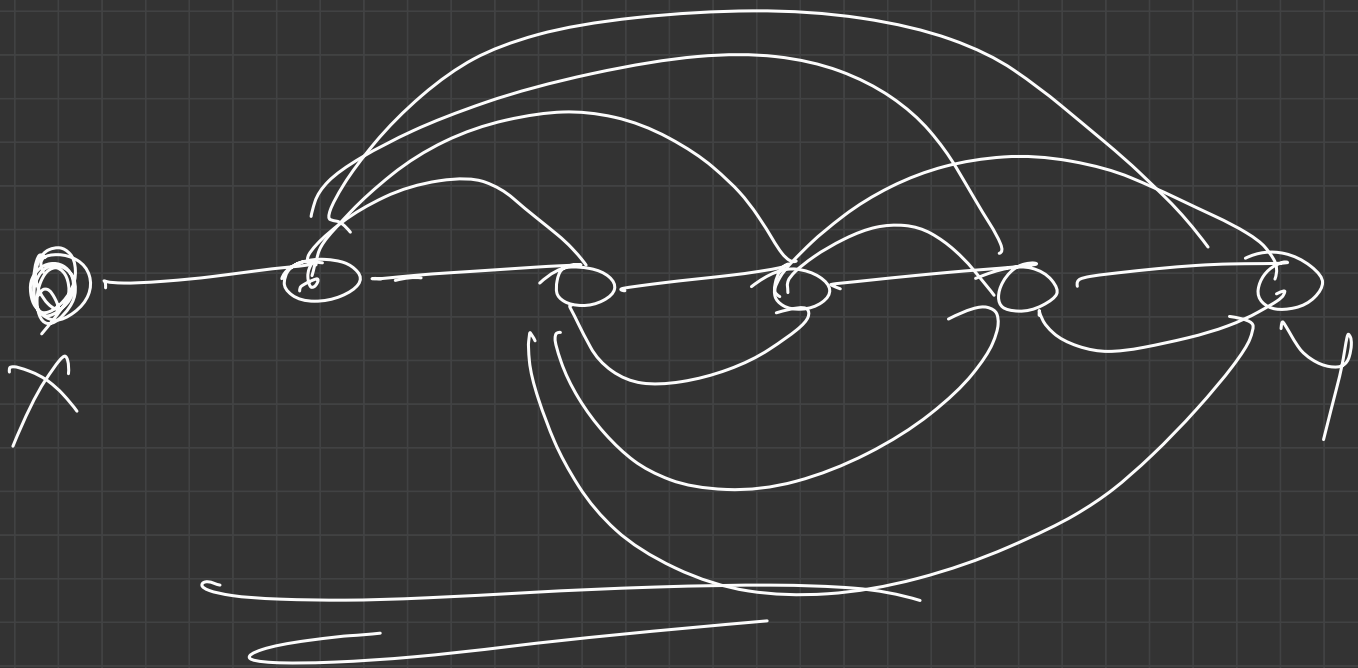
1.1) Consider all shortest paths from
X to Y and find out all the
nodes that can be on one such path



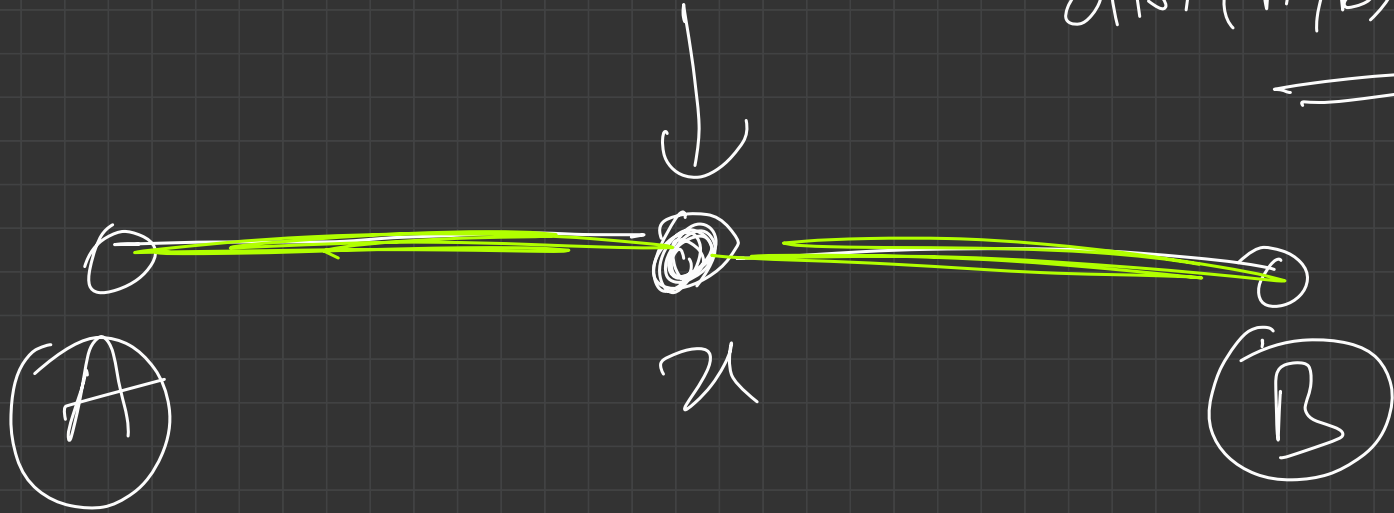


$$\text{dist}(X, Y) = 20$$

$$\underline{\underline{O(n)}}$$



$$\text{dist}(A, B) = 15$$



$$\text{dist}(A, n) + \text{dist}(n, B) = 15$$
$$\underline{\underline{= \text{dist}(A, B)}}$$

① 1 dijkstra from A



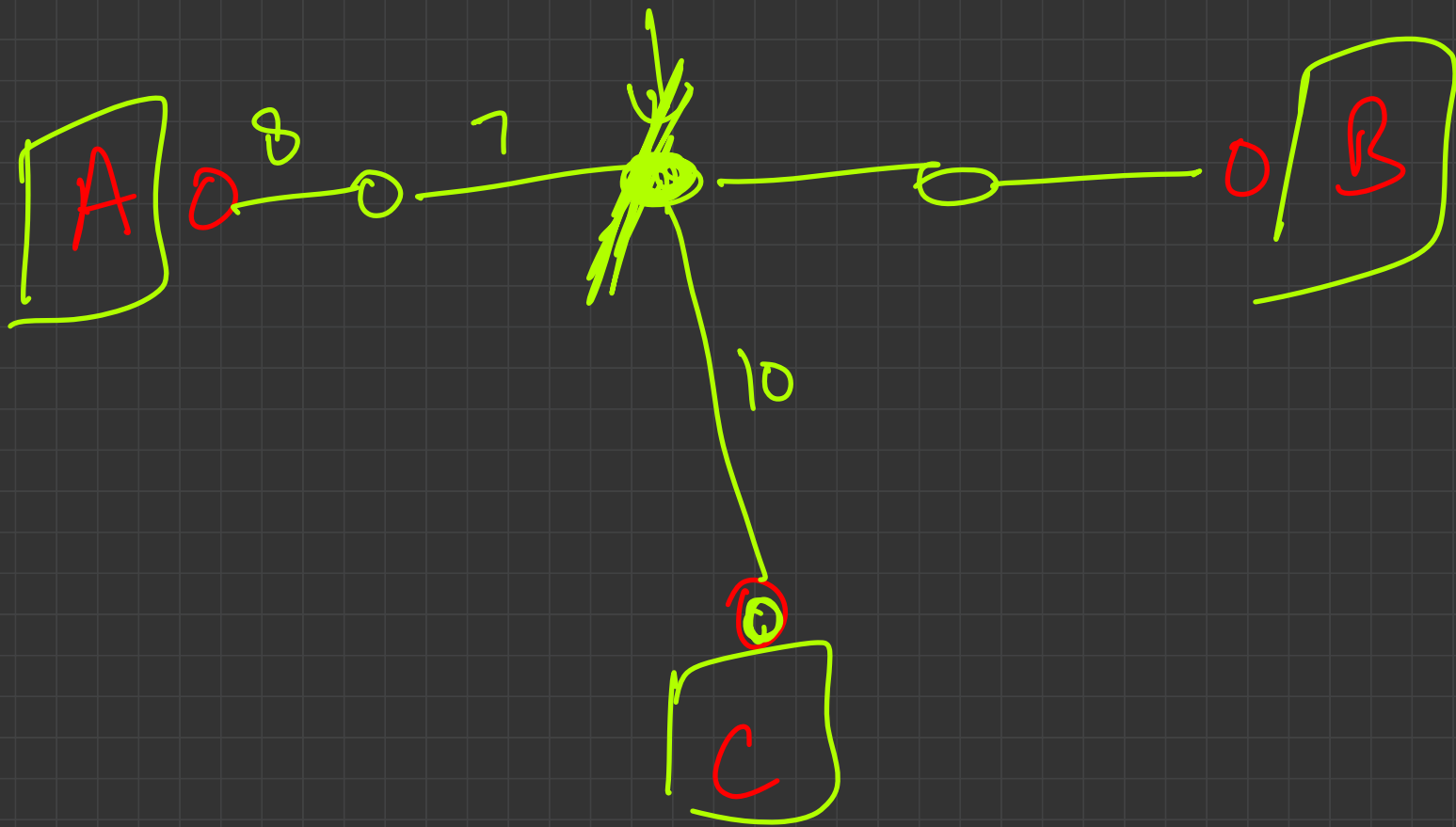
$$\text{dist } A = \begin{bmatrix} \\ x \end{bmatrix}$$

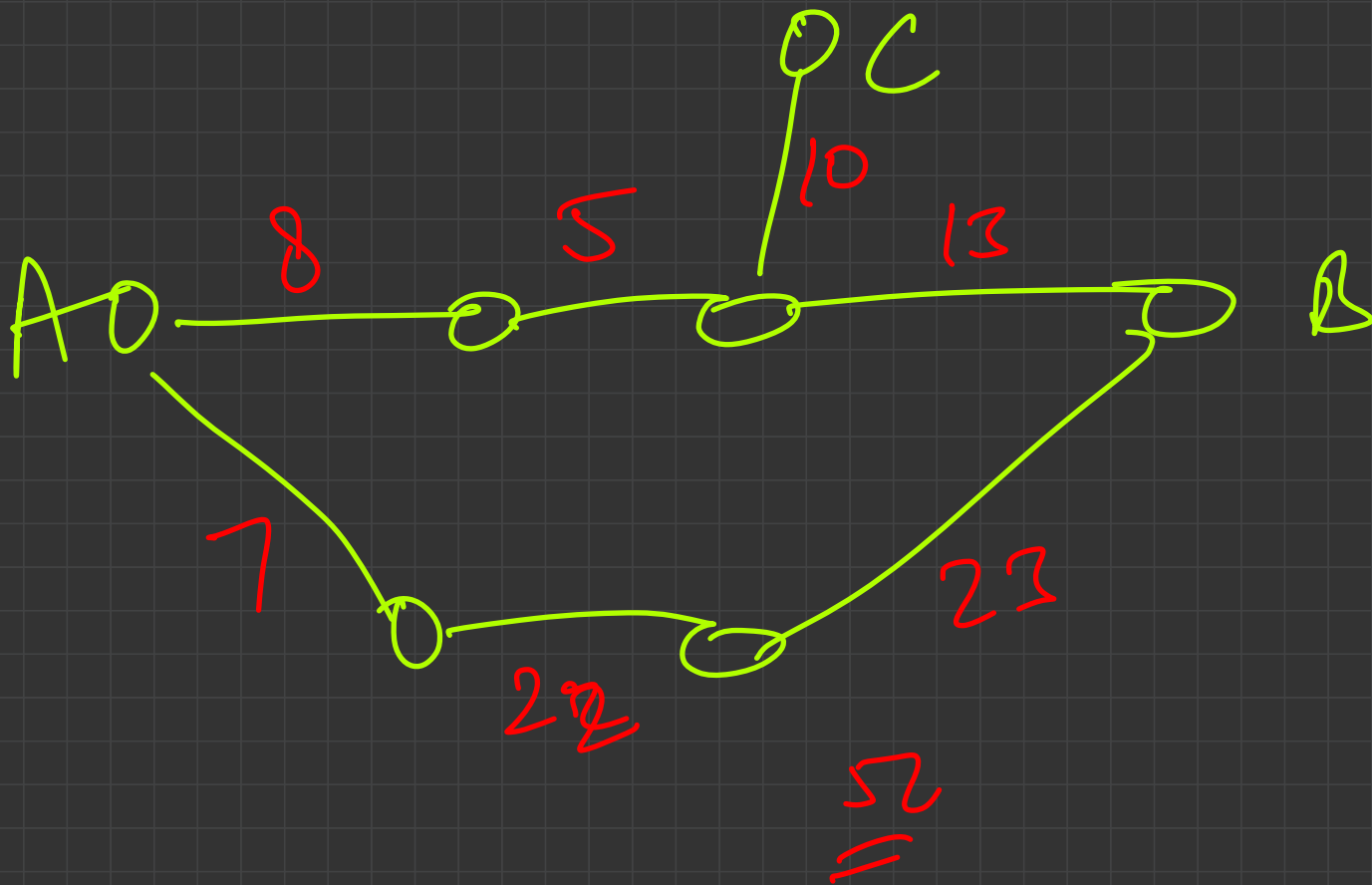
②

~~2~~ dijkstra from B

$$\text{dist } B = \begin{bmatrix} \end{bmatrix}$$

$$\text{dist } A(x) + \text{dist } B(x) = \text{dist}(A, B)$$





① Run dijkstra from A

② Run dijkstra from C

if $\text{dist}(u) < \text{dist A/u}$

remove u from the graph

③ Keep final dijkstra from A

Bellman Ford

- Single Source Shortest Path Algorithm - Idea + Intuition
- Negative edge weights (in directed) without negative cycles
- Proof/Intuition:
 - Principle of Mathematical Induction
- Code
- Breaking Early
- Retrieving the shortest path?
- Finding a negative cycle

Bellman Ford Visualization

