

# FUNDAMENTALS OF COMPUTER ENGINEERING

## HOMEWORK - 4.

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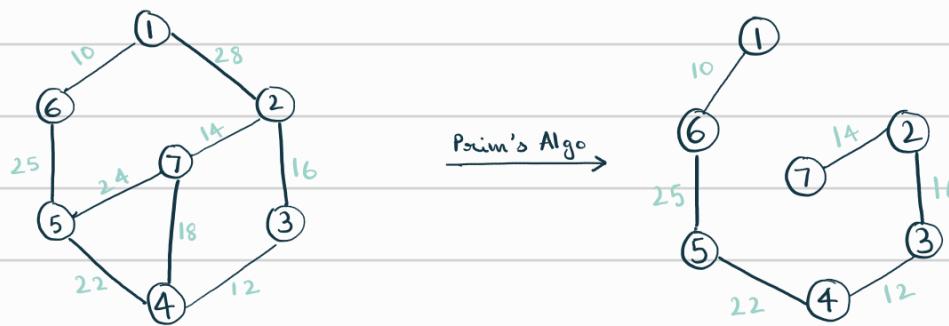
①

### PRIM'S ALGORITHM

Prim's algorithm finds the minimum spanning tree (MST) of a connected, weighted graph. It grows the MST one edge at a time, starting from an arbitrary vertex. At each step:

1. Choose the smallest edge that connects a vertex in the MST to a vertex outside it.
2. Add this edge and the corresponding vertex to the MST

e.g:



→ Prim's Algo

### Version 1: Using Adjacency Matrix and Unsorted array.

The graph is represented using an adjacency matrix and an unsorted array is used to find the minimum weight edge.

```
main.cpp
1 #include <iostream>
2 #include <vector>
3 #include <climits>
4 #include <chrono>
5 using namespace std;
6
7 // function to implement Prim's Algorithm using adjacency matrix and unsorted array
8 void primAdjMatrix(int n, vector<vector<int>>& graph) {
9     vector<int> key(n, INT_MAX); // minimum weights
10    vector<bool> inMST(n, false);
11    vector<int> parent(n, -1); // store MST edges
12    vector<vector<int>> mstMatrix(n, vector<int>(n, 0)); // store MST
13
14    key[0] = 0; // start from first vertex
15
16    for (int count = 0; count < n - 1; ++count) {
17        // finding vertex with minimum key value
18        int minKey = INT_MAX, u = -1;
19        for (int v = 0; v < n; ++v) {
20            if (!inMST[v] && key[v] < minKey) {
21                minKey = key[v];
22                u = v;
23            }
24        }
25
26        inMST[u] = true;
27
28        // updating key values of adjacent vertices
29        for (int v = 0; v < n; ++v) {
30            if (graph[u][v] && !inMST[v] && graph[u][v] < key[v]) {
31                key[v] = graph[u][v];
32                parent[v] = u;
33            }
34        }
35    }
36}
```

Find the vertex with the smallest key not in MST. Update keys of its adjacent vertices if a smaller weight is found. Then the edges and weights of the MST are printed.

```

37 // construct MST adjacency matrix
38 for (int i = 1; i < n; ++i) {
39     mstMatrix[parent[i]][i] = graph[parent[i]][i];
40     mstMatrix[i][parent[i]] = graph[parent[i]][i];
41 }
42
43 // print MST adjacency matrix
44 cout << "MST Adjacency Matrix:\n";
45 for (const auto& row : mstMatrix) {
46     for (int val : row) {
47         cout << val << " ";
48     }
49     cout << endl;
50 }
51
52 // print MST edges
53 cout << "Edges in MST:\n";
54 for (int i = 1; i < n; ++i) {
55     cout << parent[i] + 1 << " - " << i + 1 << " (Weight: " << graph[parent[i]][i] << ")\n";
56 }
57 }
58
59 int main() {
60     int n;
61     cout << "Enter the number of vertices: ";
62     cin >> n;
63
64     vector<vector<int>> graph(n, vector<int>(n));
65     cout << "Enter the adjacency matrix:\n";
66     for (int i = 0; i < n; ++i)
67         for (int j = 0; j < n; ++j)
68             cin >> graph[i][j];
69
70     auto start = chrono::high_resolution_clock::now();
71     primAdjMatrix(n, graph);
72     auto end = chrono::high_resolution_clock::now();
73
74     chrono::duration<double> elapsed = end - start;
75     cout << "Time taken: " << elapsed.count() << " seconds\n";
76     return 0;
77 }

```

```

Enter the number of vertices: 7
Enter the adjacency matrix:
0 28 0 0 0 10 0
28 0 16 0 0 0 14
0 16 0 12 0 0 0
0 0 12 0 22 0 18
0 0 0 22 0 25 24
10 0 0 0 25 0 0
0 14 0 18 24 0 0
MST Adjacency Matrix:
0 0 0 0 0 10 0
0 0 16 0 0 0 14
0 16 0 12 0 0 0
0 0 12 0 22 0 0
0 0 0 22 0 25 0
10 0 0 0 25 0 0
0 14 0 0 0 0 0
Edges in MST:
3 - 2 (Weight: 16)
4 - 3 (Weight: 12)
5 - 4 (Weight: 22)
6 - 5 (Weight: 25)
1 - 6 (Weight: 10)
2 - 7 (Weight: 14)
Time taken: 6.3711e-05 seconds

...Program finished with exit code 0
Press ENTER to exit console. []

```

Running time analysis:

Outer loop  $O(v)$  for  $V$  vertices

Inner loop  $O(v)$  for finding minimum key

Total  $O(v^2)$

Suitable for dense graphs

## Version 2 : Using Adjacency List and heap

The graph is represented using an adjacency list and a min-heap (priority queue) is used for efficient minimum weight selection.

```
main.cpp
1 #include <iostream>
2 #include <vector>
3 #include <queue>
4 #include <climits>
5 #include <chrono>
6 using namespace std;
7
8 typedef pair<int, int> pii;
9
10 // function to implement Prim's Algorithm using adjacency list and heap
11 void primAdjList(int n, vector<vector<pii>>& adj) {
12     priority_queue<pii, vector<pii>, greater<pii>> pq; // min-heap
13     vector<bool> inMST(n, false);
14     vector<int> parent(n, -1); // store MST edges
15     vector<int> key(n, INT_MAX); // minimum weights
16     vector<vector<pii>> mstAdjList(n); // store the MST adjacency list
17
18     key[0] = 0;
19     pq.push({0, 0}); // {weight, vertex}
20
21     while (!pq.empty()) {
22         int u = pq.top().second;
23         pq.pop();
24
25         if (inMST[u]) continue;
26         inMST[u] = true;
27
28         for (const auto& edge : adj[u]) {
29             int weight = edge.first;
30             int v = edge.second;
31
32             if (!inMST[v] && weight < key[v]) {
33                 key[v] = weight;
34                 parent[v] = u;
35                 pq.push({key[v], v});
36             }
37         }
38     }
39
40     // construct MST adjacency list
41     for (int i = 1; i < n; ++i) {
42         mstAdjList[parent[i]].push_back({key[i], i});
43         mstAdjList[i].push_back({key[i], parent[i]});
44     }
45
46     // print MST adjacency list
47     cout << "MST Adjacency List:\n";
48     for (int i = 0; i < n; ++i) {
49         cout << i + 1 << ": ";
50         for (const auto& edge : mstAdjList[i]) {
51             cout << "(" << edge.second + 1 << ", " << edge.first << ")";
52         }
53         cout << endl;
54     }
55
56     // print MST edges
57     cout << "Edges in MST:\n";
58     for (int i = 1; i < n; ++i) {
59         cout << parent[i] + 1 << " - " << i + 1 << " (Weight: " << key[i] << ")\n";
60     }
61
62     int main() {
63         int n, m;
64         cout << "Enter the number of vertices and edges: ";
65         cin >> n >> m;
66
67         vector<vector<pii>> adj(n);
68         cout << "Enter the edges (u v w):\n";
69         for (int i = 0; i < m; ++i) {
70             int u, v, w;
71             cin >> u >> v >> w;
72             adj[u - 1].push_back({w, v - 1});
73             adj[v - 1].push_back({w, u - 1});
74         }
75
76         auto start = chrono::high_resolution_clock::now();
77         primAdjList(n, adj);
78         auto end = chrono::high_resolution_clock::now();
79
80         chrono::duration<double> elapsed = end - start;
81         cout << "Time taken: " << elapsed.count() << " seconds\n";
82     }
}
```

Use a priority queue to efficiently find the minimum edge and then update key and parents based on the adjacency list. Finally print the edges & weights of the MST

```
Enter the number of vertices and edges: 7 9
Enter the edges (u v w):
1 2 28
2 3 16
3 4 12
4 5 22
5 6 25
6 1 10
5 7 24
4 7 18
7 2 14
MST Adjacency List:
1: (6, 10)
2: (3, 16) (7, 14)
3: (2, 16) (4, 12)
4: (3, 12) (5, 22)
5: (4, 22) (6, 25)
6: (5, 25) (1, 10)
7: (2, 14)
Edges in MST:
3 - 2 (Weight: 16)
4 - 3 (Weight: 12)
5 - 4 (Weight: 22)
6 - 5 (Weight: 25)
1 - 6 (Weight: 10)
2 - 7 (Weight: 14)
Time taken: 8.1331e-05 seconds
```

Running Time analysis.

Min-heap Operations :  $O(\log V)$

For all edges :  $O(E \log V)$

Total :  $O((V+E) \log V)$

Suitable for sparse graphs.

## ② Johnson's Algorithm

Johnson's Algorithm is typically used for finding the shortest path in a graph with both positive and negative weights. It involves a combination of Bellman-Ford and Dijkstra's algorithm.

1. Bellman-Ford Algorithm computes a potential function  $h[v]$  for all vertices, ensuring there are no negative weight cycles. This is used to reweight the graph, making all edge weights non-negative.

2. Dijkstra's Algorithm uses the reweighted graph to compute shortest paths efficiently with a priority queue.

3. Johnson's Algorithm combines the results to return shortest paths for the original graph.

main.cpp

```
1 #include <iostream>
2 #include <vector>
3 #include <limits>
4 #include <queue>
5 #include <iomanip>
6
7 using namespace std;
8
9 const int INF = numeric_limits<int>::max();
10
11 struct Edge {
12     int src, dest, weight;
13 };
14
15 // Bellman-Ford algorithm
16 vector<int> bellmanFord(int V, int src, const vector<Edge>& edges) {
17     vector<int> dist(V + 1, INF); // 1-based indexing
18     dist[src] = 0;
19
20     for (int i = 1; i <= V - 1; i++) {
21         for (const auto& edge : edges) {
22             if (dist[edge.src] != INF && dist[edge.src] + edge.weight < dist[edge.dest]) {
23                 dist[edge.dest] = dist[edge.src] + edge.weight;
24             }
25         }
26     }
27
28     for (const auto& edge : edges) {
29         if (dist[edge.src] != INF && dist[edge.src] + edge.weight < dist[edge.dest]) {
30             cout << "Negative weight cycle detected!" << endl;
31             return {};
32         }
33     }
34
35     return dist;
36 }
37
38 // Dijkstra's algorithm using a priority queue (min-heap)
39 vector<int> dijkstra(int V, int src, const vector<vector<pair<int, int>>>& adj) {
40     vector<int> dist(V + 1, INF); // 1-based indexing
41     priority_queue<pair<int, int>, vector<pair<int, int>>, greater<pair<int, int>> pq;
42
43     dist[src] = 0;
44     pq.push({0, src});
45
46     while (!pq.empty()) {
47         int u = pq.top().second;
48         int d = pq.top().first;
49         pq.pop();
50
51         if (d > dist[u]) continue;
52
53         for (const auto& [v, weight] : adj[u]) {
54             if (dist[u] + weight < dist[v]) {
55                 dist[v] = dist[u] + weight;
56                 pq.push({dist[v], v});
57             }
58         }
59     }
60
61     return dist;
62 }
```

```

64 // Johnson's algorithm
65 void johnsonsAlgorithm(int V, const vector<Edge>& edges) {
66     // Step 1: Add a new vertex and connect it to every other vertex with weight 0
67     vector<Edge> newEdges = edges;
68     for (int i = 1; i <= V; i++) {
69         newEdges.push_back({V + 1, i, 0});
70     }
71
72     vector<int> h = bellmanFord(V + 1, V + 1, newEdges);
73     if (h.empty()) return;
74
75     cout << "Reweighted graph (h values):" << endl;
76     for (int i = 1; i <= V; i++) {
77         cout << "h[" << i << "] = " << h[i] << endl;
78     }
79
80     // Step 2: Reweight edges based on h values and print them
81     vector<vector<pair<int, int>>> adj(V + 1);
82     cout << "\nReweighted edges:" << endl;
83     for (const auto& edge : edges) {
84         int newWeight = edge.weight + h[edge.src] - h[edge.dest];
85         adj[edge.src].push_back({edge.dest, newWeight});
86         cout << "Edge (" << edge.src << " -> " << edge.dest << ") : New Weight = " << newWeight << endl;
87     }
88
89     // Step 3: Run Dijkstra's algorithm for each vertex and print results
90     cout << "\nShortest paths using Dijkstra's (Johnson's reweighted):" << endl;
91     for (int i = 1; i <= V; i++) {
92         vector<int> dist = dijkstra(V, i, adj);
93         cout << "From vertex " << i << ":" << endl;
94         for (int j = 1; j <= V; j++) {
95             if (dist[j] == INF) {
96                 cout << "To vertex " << j << " = INF" << endl;
97             } else {
98                 cout << "To vertex " << j << " = " << dist[j] << endl;
99             }
100        }
101    }
102 }
103 }
```

```

105 int main() {
106     int V, E;
107     cout << "Enter number of vertices and edges: ";
108     cin >> V >> E;
109
110     vector<Edge> edges(E);
111     cout << "Enter each edge (src dest weight):" << endl;
112     for (int i = 0; i < E; i++) {
113         cin >> edges[i].src >> edges[i].dest >> edges[i].weight;
114     }
115
116     int choice;
117     cout << "\nOptions:\n1. Bellman-Ford\n2. Dijkstra's\n3. Johnson's Algorithm\nChoose: ";
118     cin >> choice;
119
120     if (choice == 1) {
121         int src;
122         cout << "Enter source vertex for Bellman-Ford: ";
123         cin >> src;
124         vector<int> dist = bellmanFord(V, src, edges);
125         if (!dist.empty()) {
126             cout << "Distances from source:" << endl;
127             for (int i = 1; i <= V; i++) {
128                 if (dist[i] == INF) {
129                     cout << "To vertex " << i << " = INF" << endl;
130                 } else {
131                     cout << "To vertex " << i << " = " << dist[i] << endl;
132                 }
133             }
134         }
135     } else if (choice == 2) {
136         int src;
137         cout << "Enter source vertex for Dijkstra's: ";
138         cin >> src;
139
140         vector<vector<pair<int, int>>> adj(V + 1); // 1-based indexing
141         for (const auto& edge : edges) {
142             adj[edge.src].push_back({edge.dest, edge.weight});
143         }
144
145         vector<int> dist = dijkstra(V, src, adj);
146         cout << "Distances from source vertex " << src << ":" << endl;
147         for (int i = 1; i <= V; i++) {
148             if (dist[i] == INF) {
149                 cout << "To vertex " << i << " = INF" << endl;
150             } else {
151                 cout << "To vertex " << i << " = " << dist[i] << endl;
152             }
153         }
154     } else if (choice == 3) {
155         johnsonsAlgorithm(V, edges);
156     }
157
158     return 0;
159 }
```

Johnson's algorithm

rewrites the graph

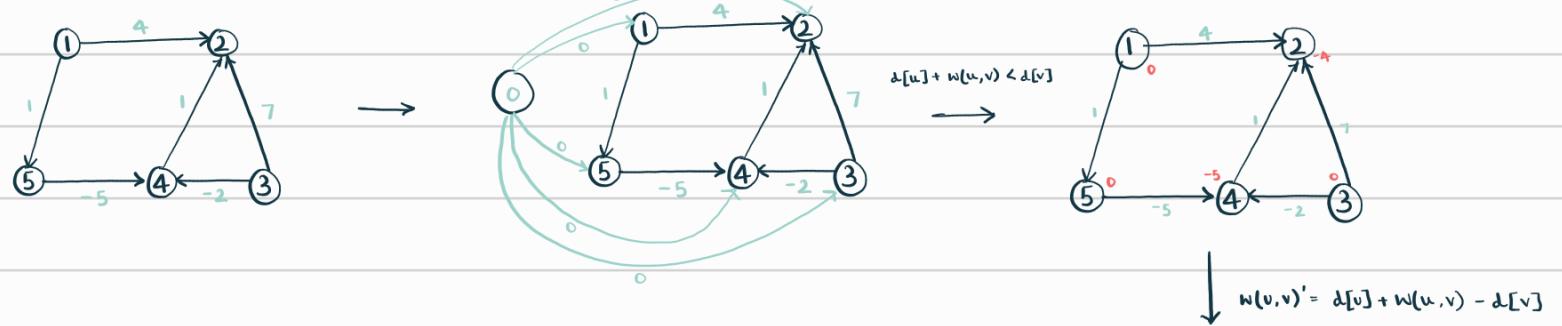
using Bellman-Ford algorithm

then runs Dijkstra's algorithm

for each vertex and prints the

shortest paths

## Example of Johnson's Algorithm



```

Enter number of vertices and edges: 5 6
Enter each edge (src dest weight):
1 2 4
1 5 1
5 4 -5
3 4 -2
3 2 7
4 2 1

Options:
1. Bellman-Ford
2. Dijkstra's
3. Johnson's Algorithm
Choose: 3
Reweighted graph (h values):
h[1] = 0
h[2] = -4
h[3] = 0
h[4] = -5
h[5] = 0

Reweighted edges:
Edge (1 -> 2) : New Weight = 8
Edge (1 -> 5) : New Weight = 1
Edge (5 -> 4) : New Weight = 0
Edge (3 -> 4) : New Weight = 3
Edge (3 -> 2) : New Weight = 11
Edge (4 -> 2) : New Weight = 0

Shortest paths using Dijkstra's (Johnson's reweighted):
From vertex 1:
To vertex 1 = 0
To vertex 2 = 1
To vertex 3 = INF
To vertex 4 = 1
To vertex 5 = 1

From vertex 2:
To vertex 1 = INF
To vertex 2 = 0
To vertex 3 = INF
To vertex 4 = INF
To vertex 5 = INF

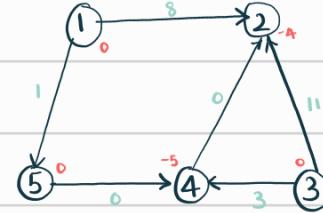
From vertex 3:
To vertex 1 = INF
To vertex 2 = 3
To vertex 3 = 0
To vertex 4 = 3
To vertex 5 = INF

From vertex 4:
To vertex 1 = INF
To vertex 2 = 0
To vertex 3 = INF
To vertex 4 = 0
To vertex 5 = INF

From vertex 5:
To vertex 1 = INF
To vertex 2 = 0
To vertex 3 = INF
To vertex 4 = 0
To vertex 5 = 0

Johnson's Algorithm execution time: 130 microseconds
  
```

Distance from 0 to 1,2,3,4,5 are 0, -4, 0, -5, 0 respectively



Running Time Analysis:

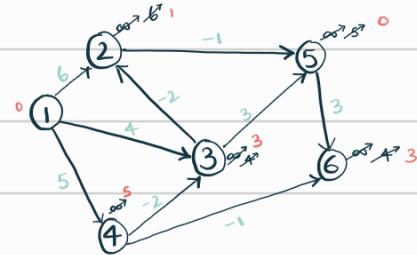
Bellman-Ford Algorithm takes  $O(V \cdot E)$

Dijkstra's Algorithm runs for each vertex taking  $O(V \cdot (V+E) \cdot \log V) \rightarrow O(V \log V)$

$\therefore$  Time Complexity of Johnson's Algorithm

$$O(V \cdot E + V \cdot (V+E) \log V) \rightarrow O(V^2 \log V + VE)$$

### Example of Bellman-Ford Algorithm



```

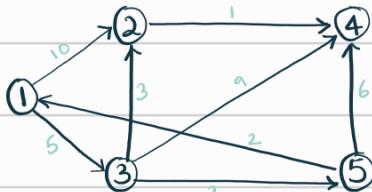
Enter number of vertices and edges: 6 9
Enter each edge (src dest weight):
1 2 6
1 3 4
1 4 5
3 2 -2
4 3 -2
2 5 -1
3 5 3
4 6 -1
5 6 3

Options:
1. Bellman-Ford
2. Dijkstra's
3. Johnson's Algorithm
Choose: 1
Enter source vertex for Bellman-Ford: 1
Distances from source vertex 1:
To vertex 1 = 0
To vertex 2 = 1
To vertex 3 = 3
To vertex 4 = 5
To vertex 5 = 0
To vertex 6 = 3
Bellman-Ford execution time: 2065474 microseconds
  
```

Bellman-Ford computes shortest distances from a single source and detects negative weight cycles.

Relaxes all edges  $V-1$  times and each relaxation step takes  $O(E)$   $\therefore$  Time Complexity:  $O(VE)$

### Example of Dijkstra's Algorithm



	1	2	3	4	5
0	$\infty$	$\infty$	$\infty$	$\infty$	
1	10	5	$\infty$	$\infty$	
2	8		14	7	
3			13		
4				9	
5					

```

Enter number of vertices and edges: 5 7
Enter each edge (src dest weight):
1 2 10
1 3 5
2 4 1
3 2 3
3 5 2
5 4 6
3 4 9

Options:
1. Bellman-Ford
2. Dijkstra's
3. Johnson's Algorithm
Choose: 2
Enter source vertex for Dijkstra's: 1
Distances from source vertex 1:
To vertex 1 = 0
To vertex 2 = 8
To vertex 3 = 5
To vertex 4 = 9
To vertex 5 = 7
Dijkstra's execution time: 987897 microseconds
  
```

Dijkstra's algorithm finds shortest path from a source vertex using an unsorted priority queue. Each vertex is processed once & each edge is relaxed once using a min-heap.  $\therefore$  Time Complexity:  $O((V+E)\log V)$