

CS221: Digital Design

FSM- Correctness and Completeness

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Outline

- Finite State Machine
- FSM Completeness
- FSM Correctness

FSM Completeness

- Finite State Machine Transition Function

$$\delta : S \times I \rightarrow S$$

- δ : For all the states of FSM and For all the type of input
- FSM Completeness
 - If for all the states, all the input combination specified in transition function δ

FSM Correctness

- Finite State Machine Transition Function

$$\delta : S \times I \rightarrow S$$

- FSM is correct if For all the state of FSM
 - **AND operation of every pair of out going edges of a state =0**
 - *Meaning: for some condition FSM should not transition to more than one state*
 - **OR operation all out going edges of a state =1**
 - *Meaning: there should be at least one transition for every condition*

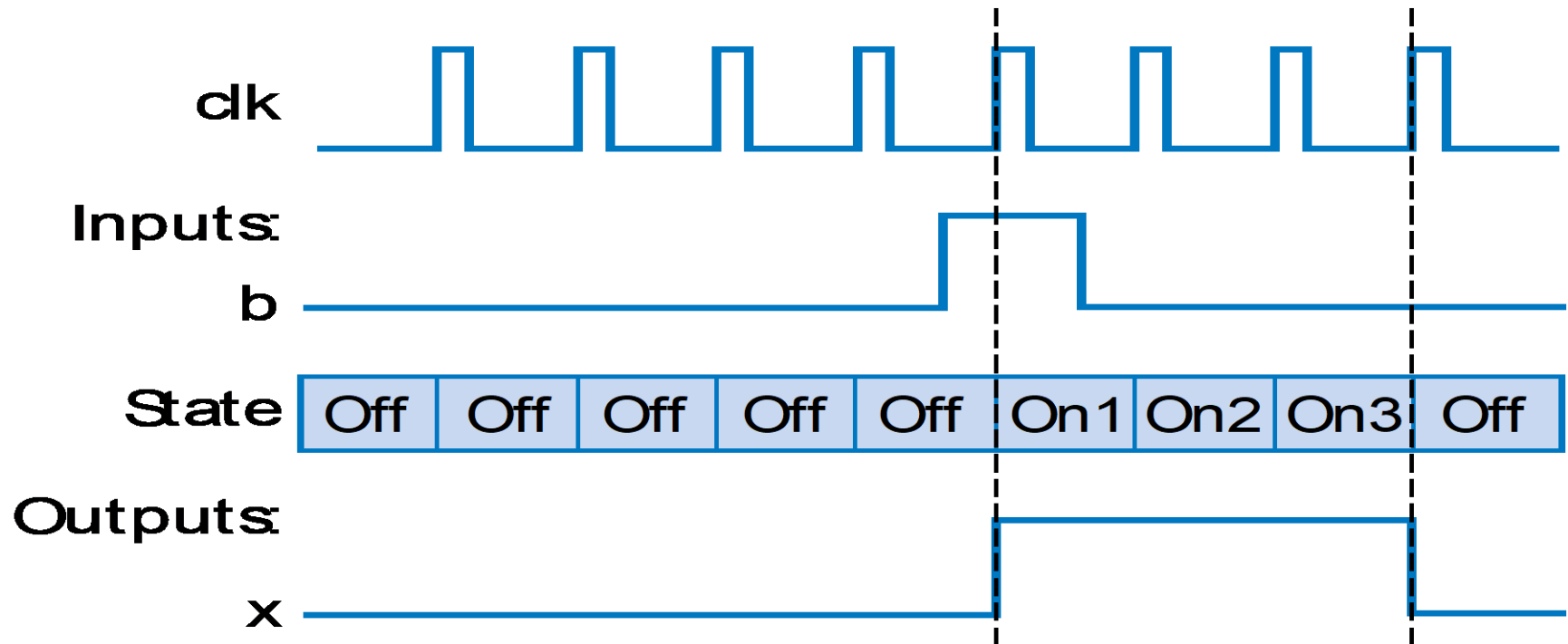
FSM Completeness

Example Three-Cycles High Laser Timer

Extend FSM to Three-Cycles High Laser Timer

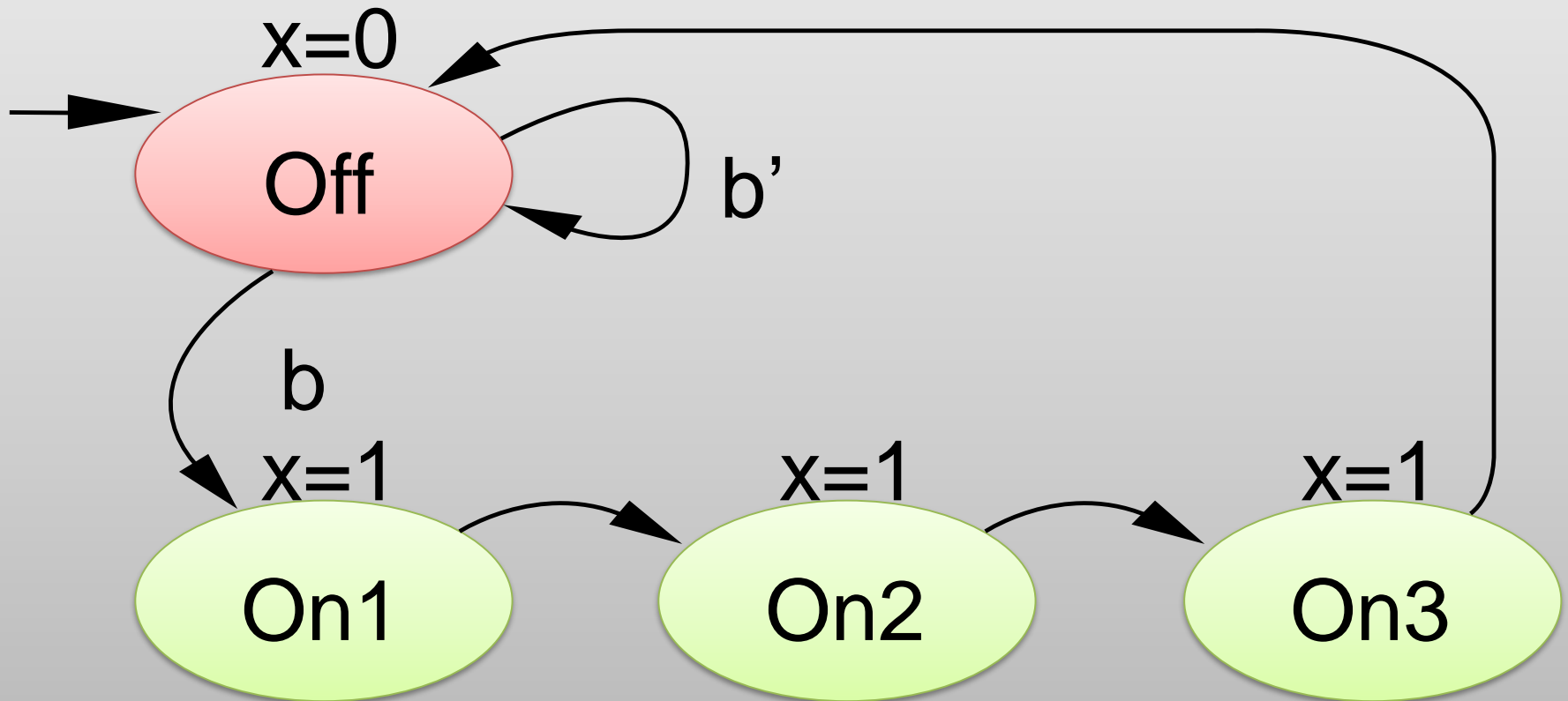
- Four states: Wait in “Off” state while b is 0 (b')
- When $b=1$ (& rising clock edge), transition to On1
 - Sets $X=1$
 - On next two clock edges, transition to On2, then On3, which also set $x=1$
- So $x=1$ for three cycles after button pressed

Extend FSM to Three-Cycles High Laser Timer



FSM of Three-Cycles High Laser Timer

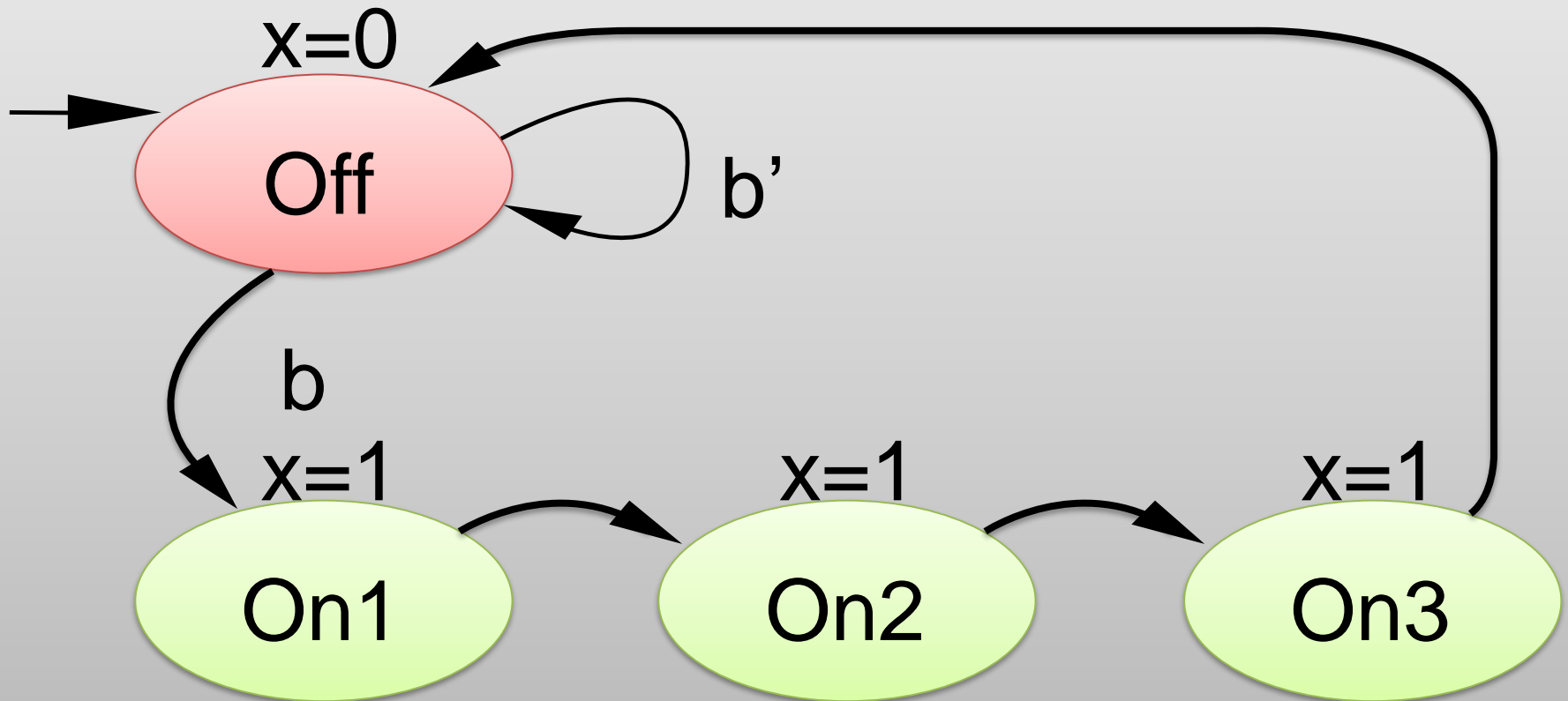
Inputs: b ; Outputs: x



Note: Transition with no associated condition thus transistions to next state on next clock cycle

FSM Completeness

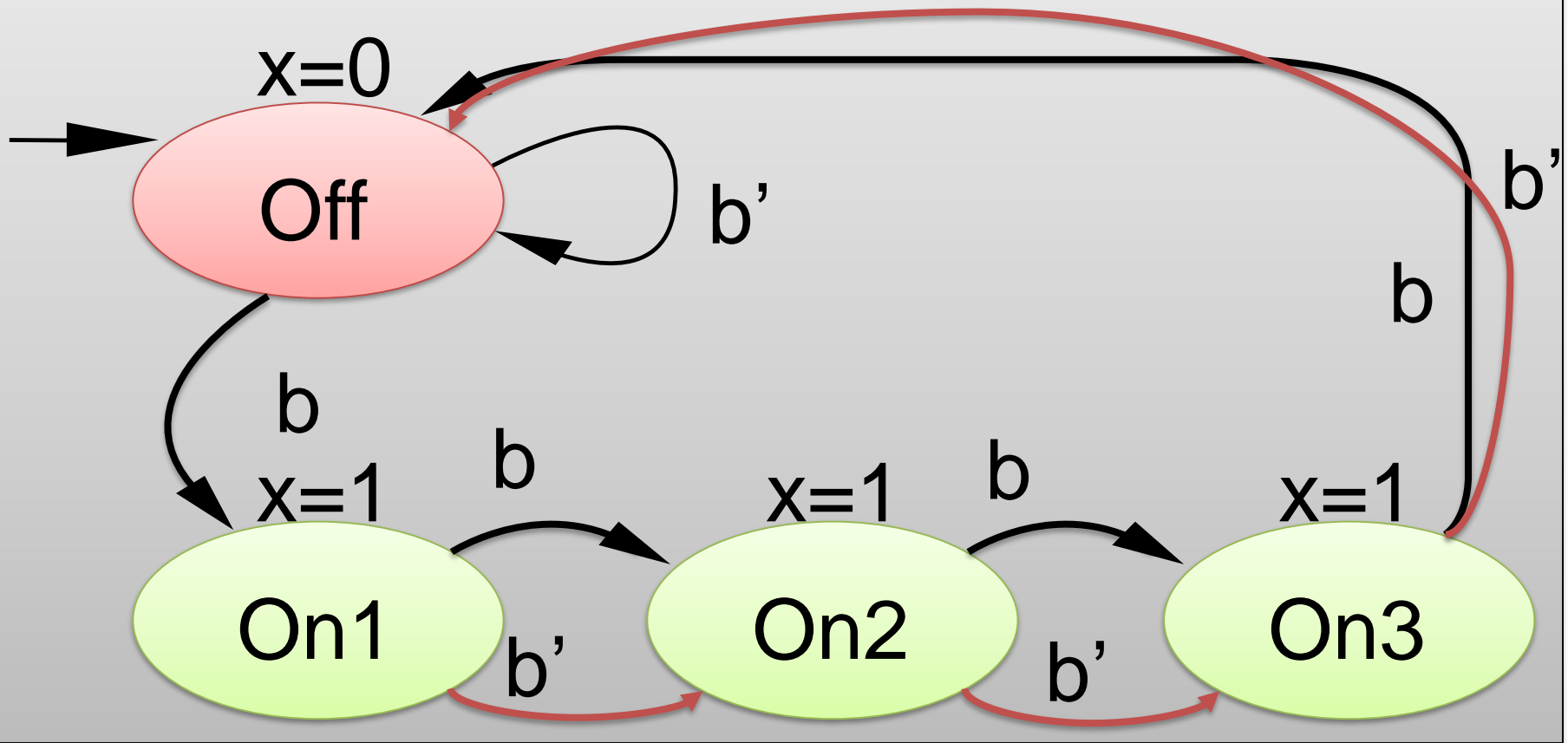
Inputs: b ; Outputs: x



Value of $b=1$: 0111..repeat, **Is this FSM complete?**

FSM Completeness

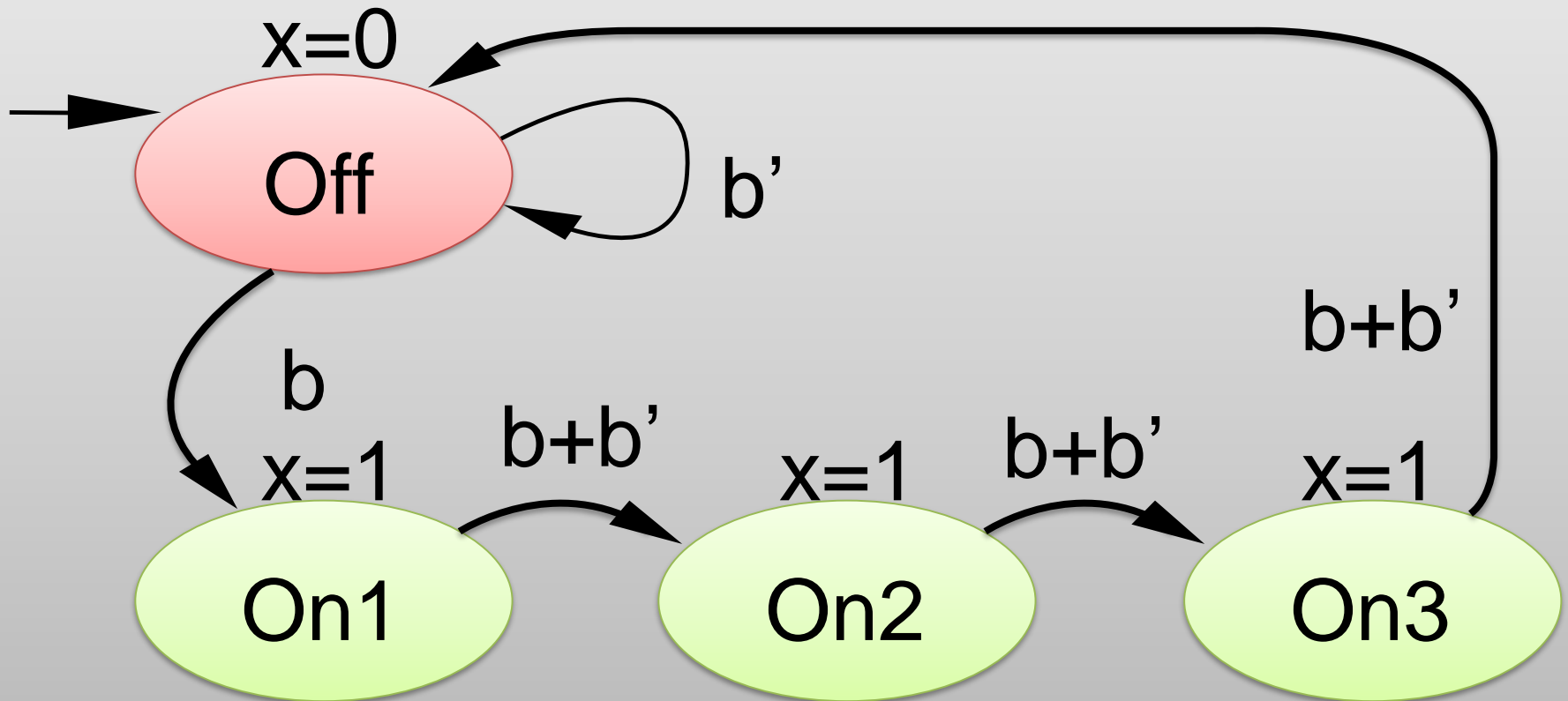
Inputs: b ; Outputs: x



Value of $b=1$: 0111..repeat, **Is this FSM complete?**

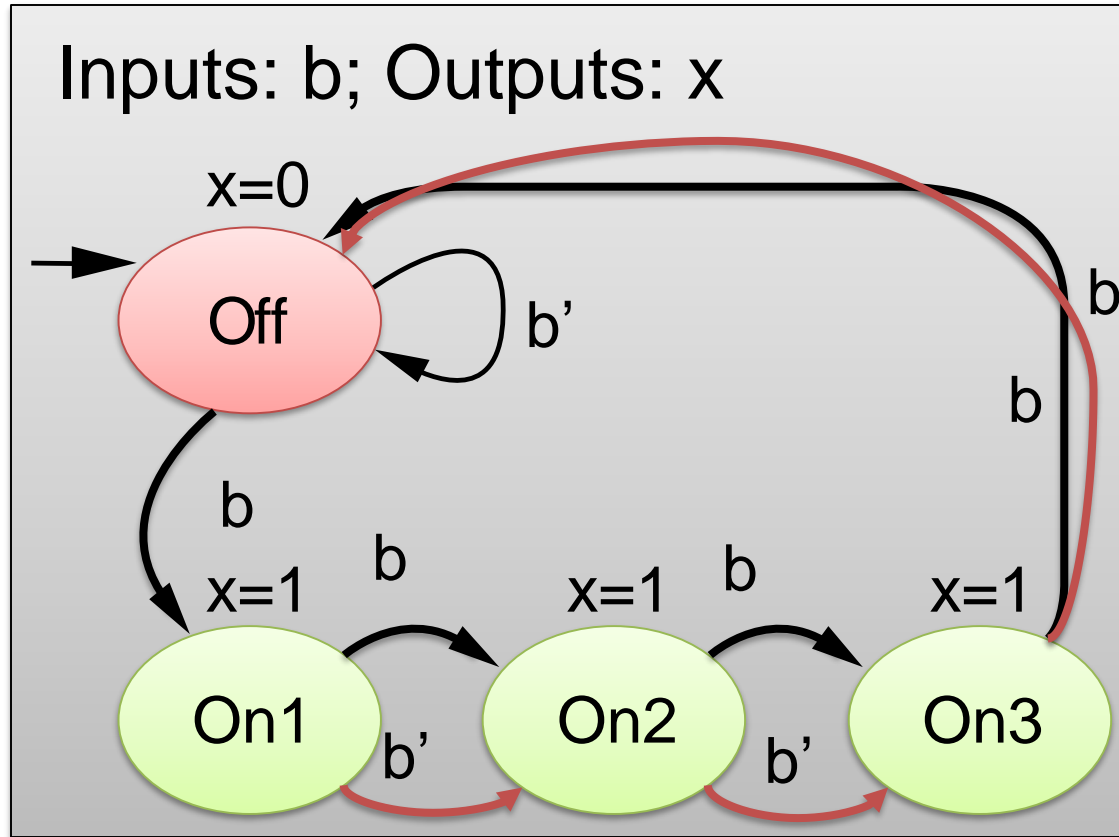
FSM Completeness

Inputs: b ; Outputs: x



Value of $b=1$: 0111..repeat, **Is this FSM complete?**

FSM Correctness



AND operation
every pair of out
going edges of a
state = 0

$$b.b' = 0$$

OR operation all out
going edges of a
state = 1

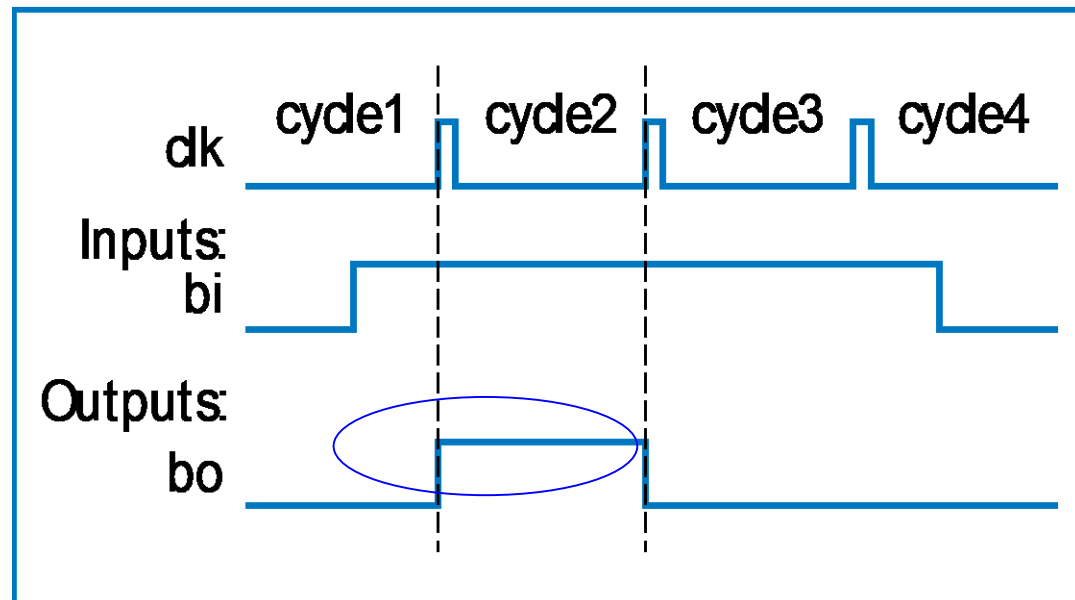
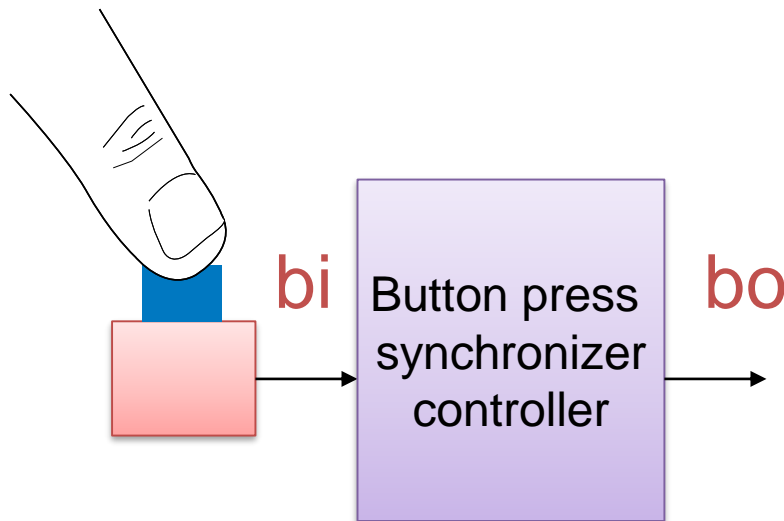
$$b + b' = 1$$

Is this FSM is correctly specified: based on rule

FSM Example 9: Button Press Synchronizer

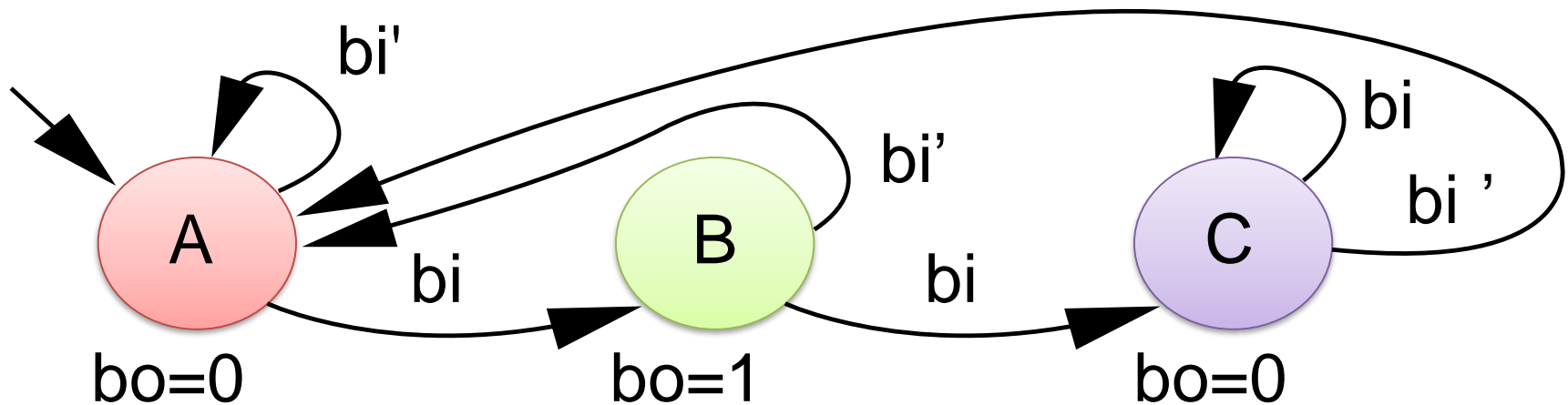
Example 9 : Button Press Synchronizer

- English Language Specification
- All most all the keyboards use this method
- We want simple sequential circuit
 - Converts button press to single cycle duration
 - Regardless of length of time that button actually pressed



FSM Example 9 : Button Press Synchronizer

FSM inputs: bi ; FSM outputs: bo



I am Off
When $B=0$

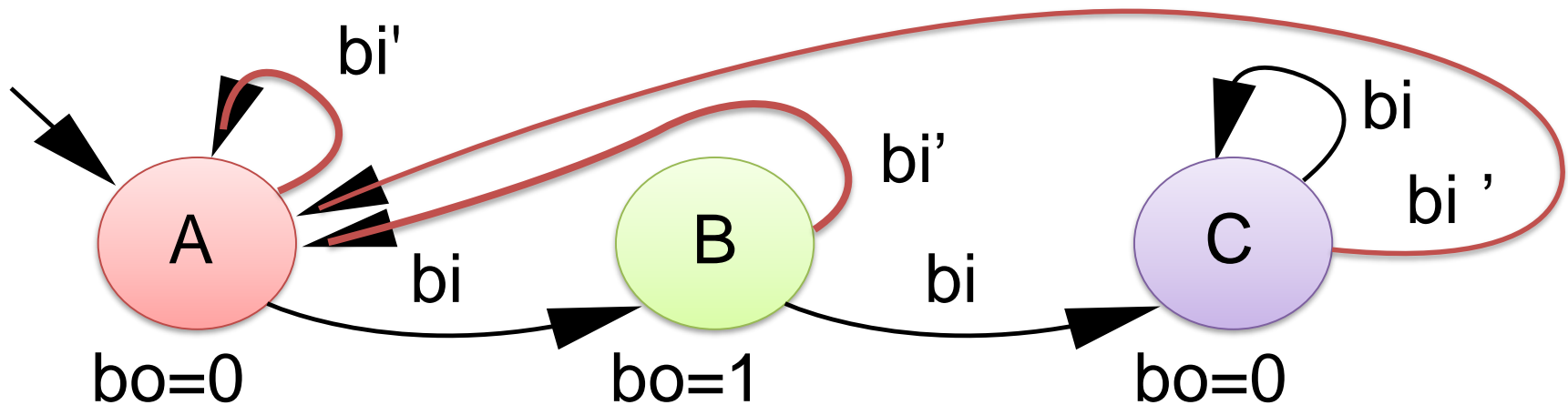
I am On in First
CLK and
When $B=1$

I am Off
Even if $B=1$

Step 1: Design FSM

Button Press Synchronizer: Completeness & correctness

FSM inputs: bi ; FSM outputs: bo



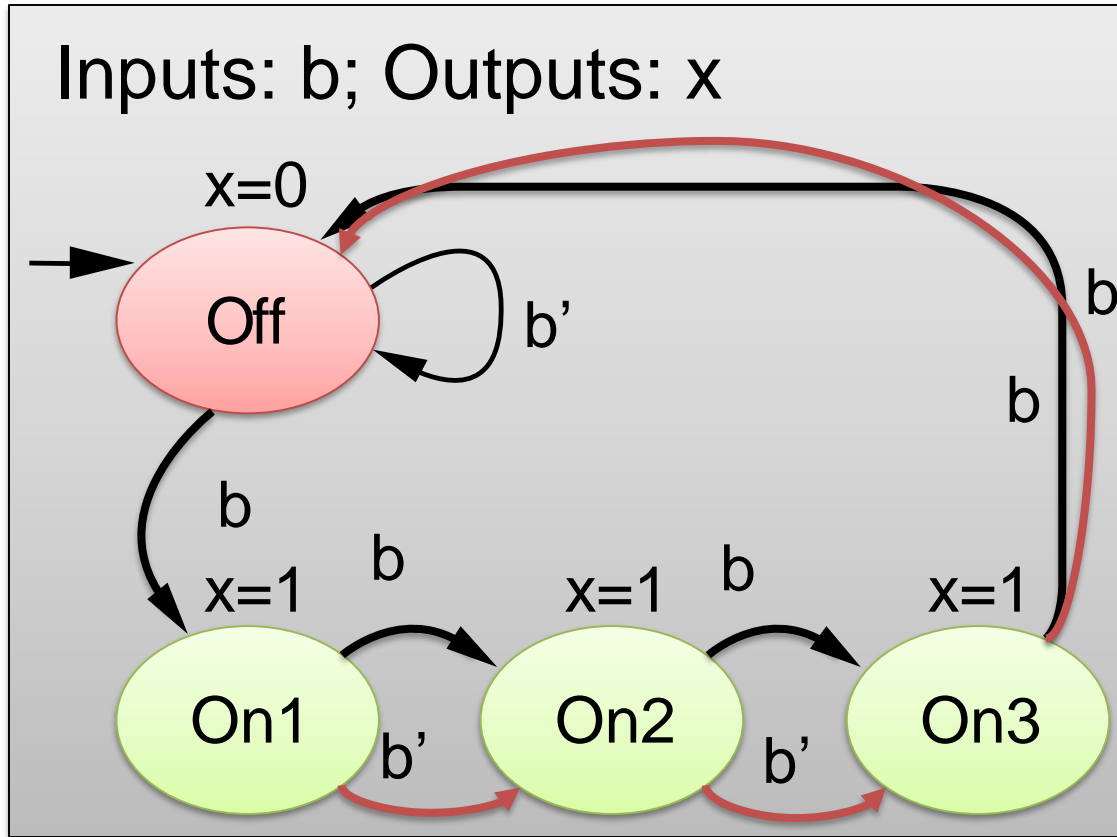
For all the state exactly two out going transitions, one for b and other for b'

$$b \cdot b' = 0$$

$$b + b' = 1$$

**FSM With Transition specified in Boolean
expression or Compressed form**

FSM Correctness



AND operation
every pair of out
going edges of a
state = 0

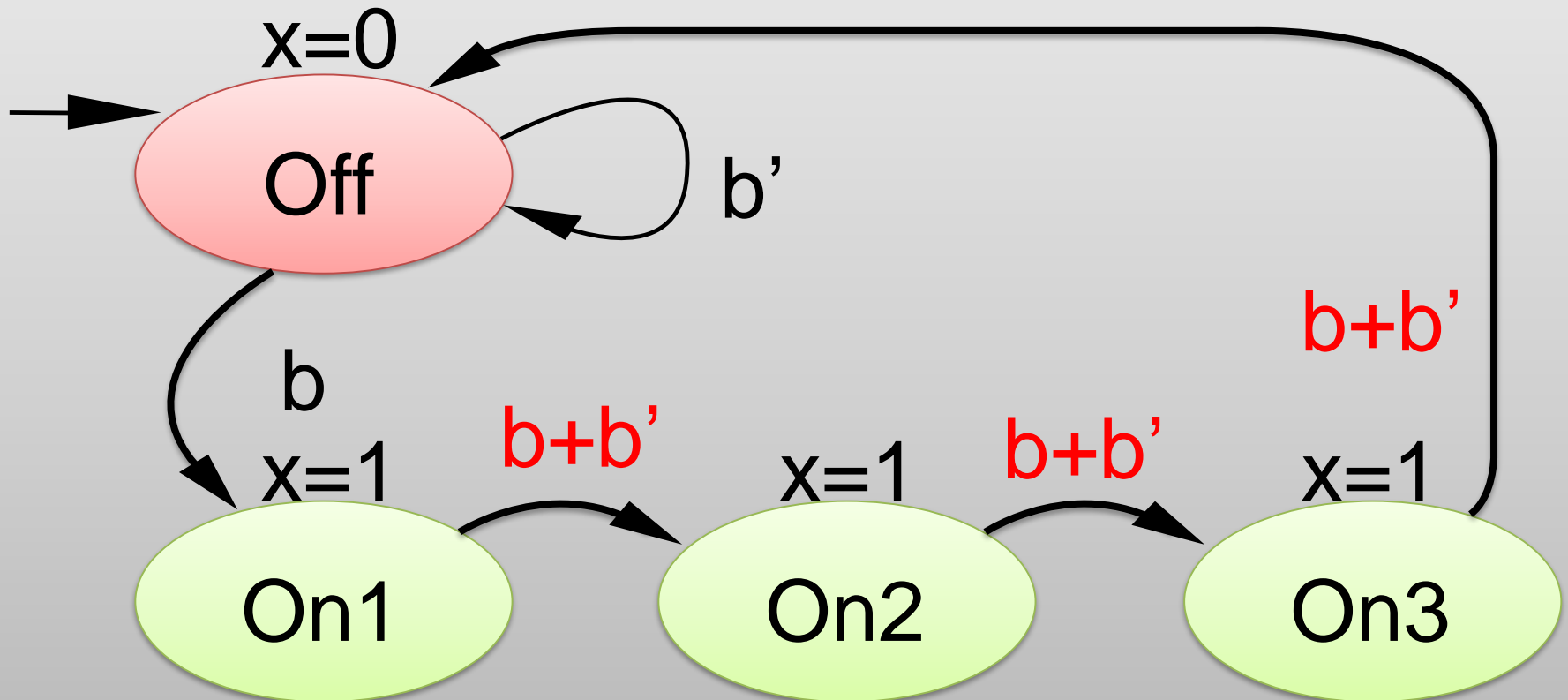
$$b.b' = 0$$

OR operation all out
going edges of a
state = 1

$$b + b' = 1$$

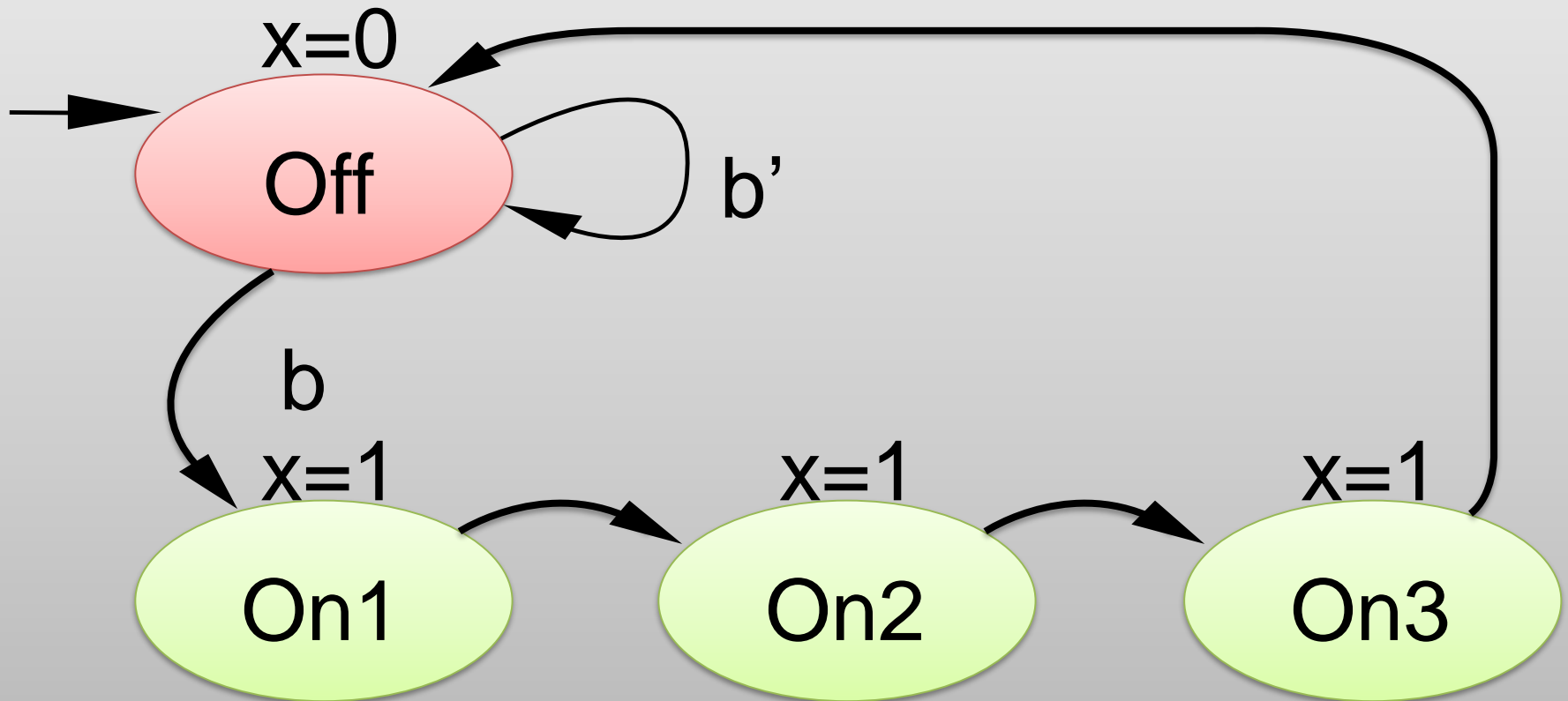
FSM Completeness

Inputs: b ; Outputs: x



FSM Completeness

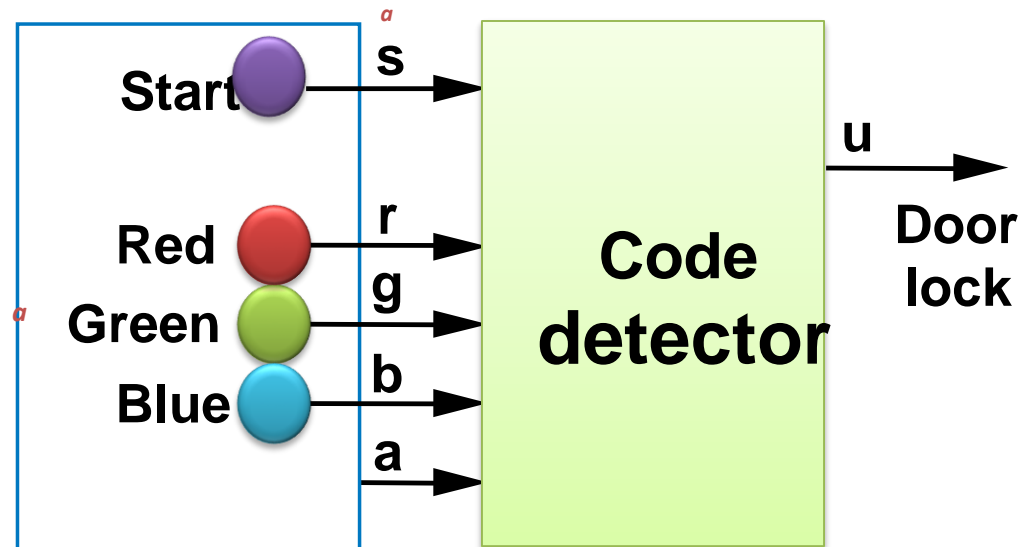
Inputs: b ; Outputs: x



FSM Example 11: Code Detector

FSM Example : Code Detector

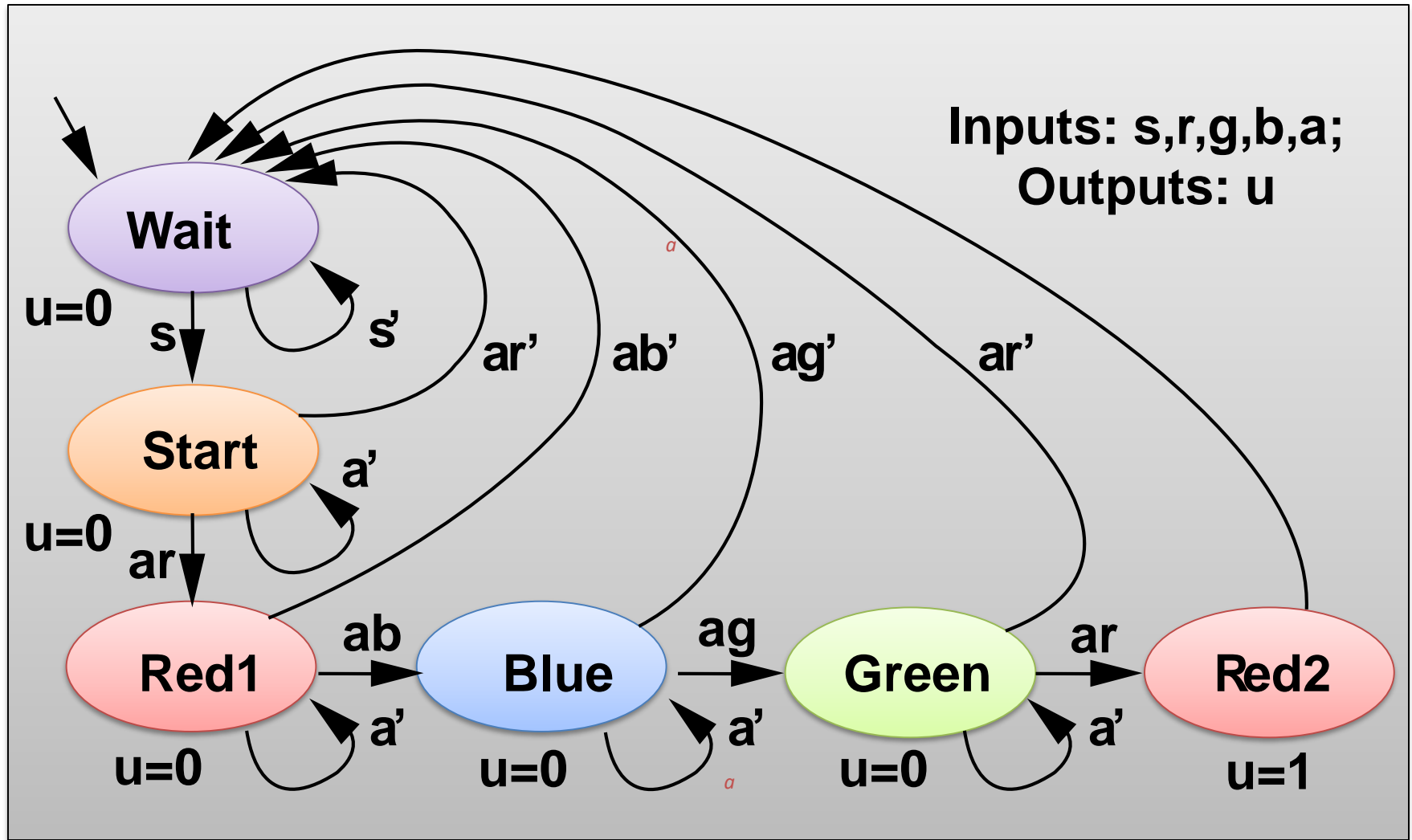
- Unlock door ($u=1$) only when buttons pressed in sequence:
 - **start, then red, blue, green, red**
- Input from each button: s, r, g, b
 - Also, output a indicates that some colored button pressed



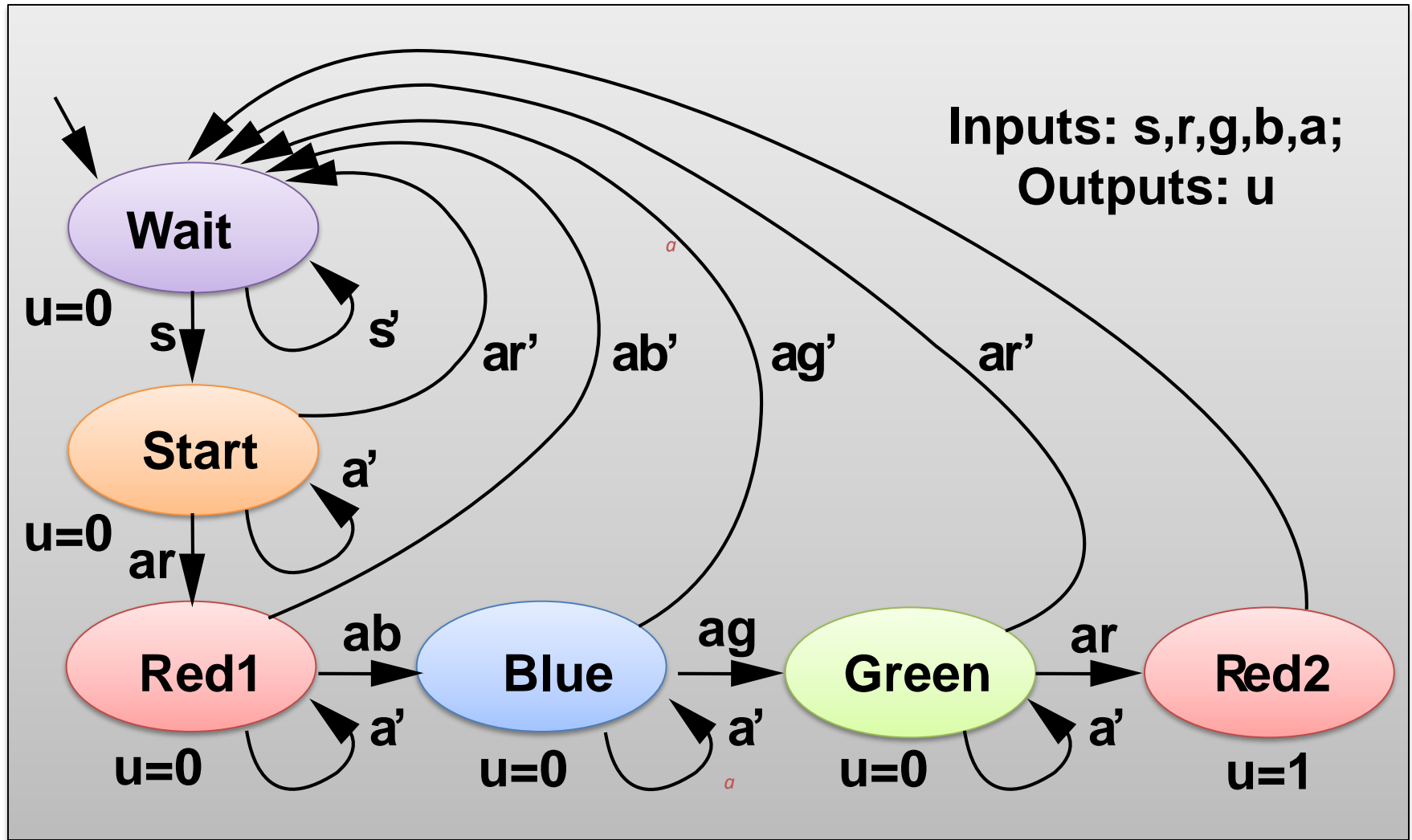
FSM Example 11: Code Detector

- Wait for start ($s=1$) in “Wait”,
- **Once started (“Start”)**
 - If see red, go to “Red1”
 - Then, if see blue, go to “Blue”, Then, if see green, go to “Green”, Then, if see red, go to “Red2”
 - In that state, open the door ($u=1$)
 - Wrong button at any step, return to “Wait”

FSM Example 11 : Code Detector



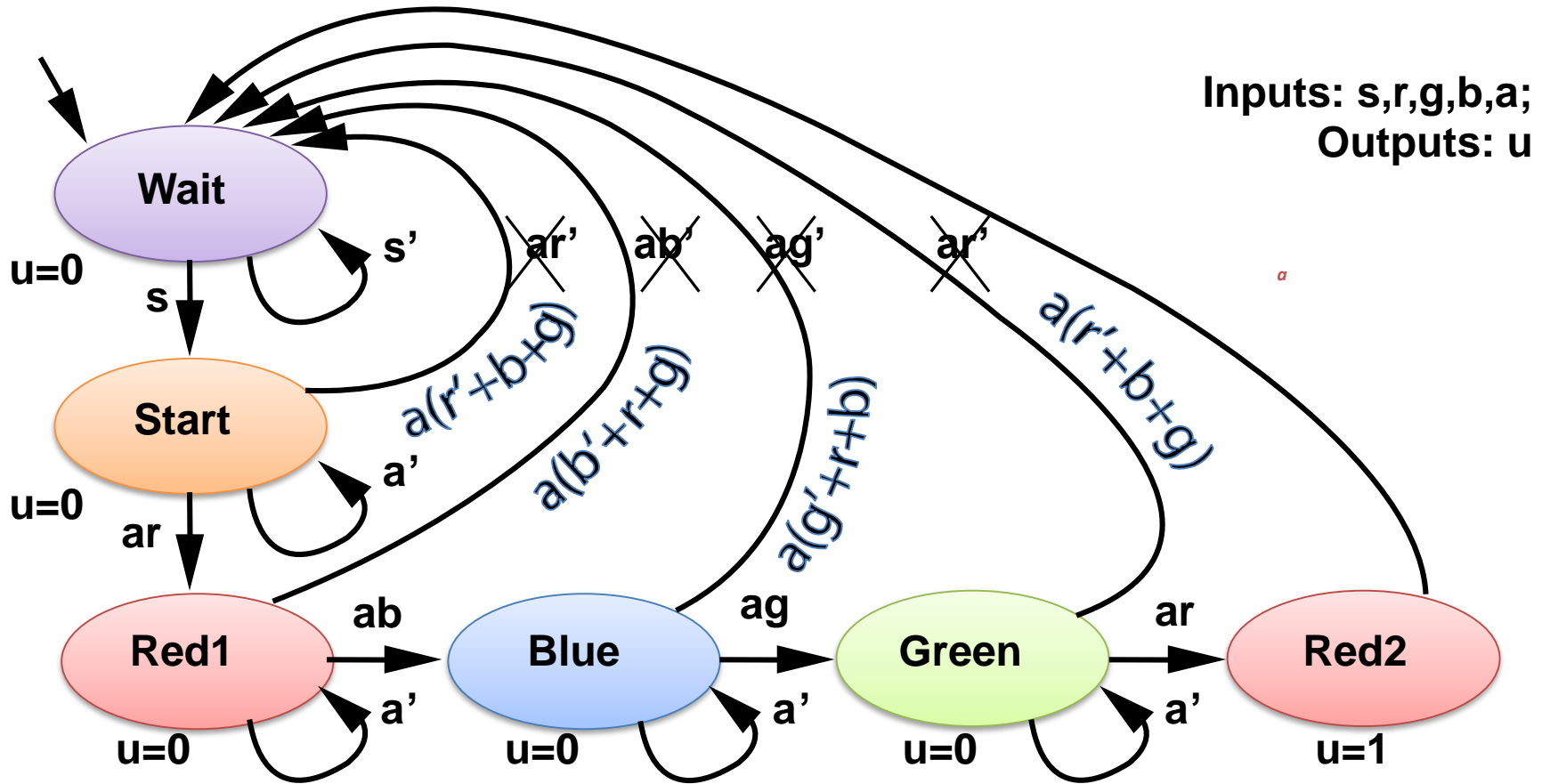
FSM Example 11 : Code Detector



Q: Can you trick this FSM to open the door, without knowing the code?

A: Yes, hold all buttons simultaneously

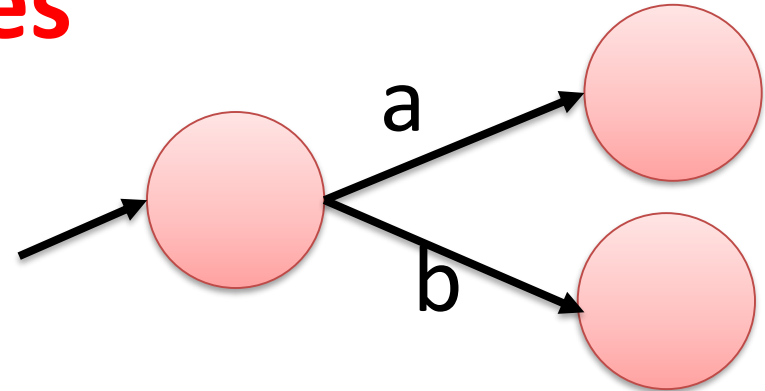
Improve FSM for Code Detector



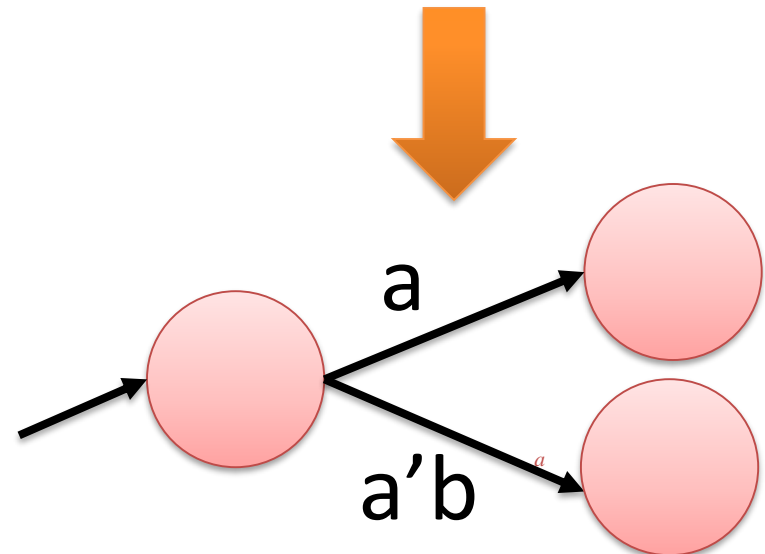
- **New transition conditions** detect if wrong button pressed, returns to “Wait”
- **FSM provides formal, concrete means to accurately define desired behavior**

Common Pitfalls Regarding Transition Properties

- *Only* one condition should be true
 - For all transitions leaving a state
 - Else, which one?

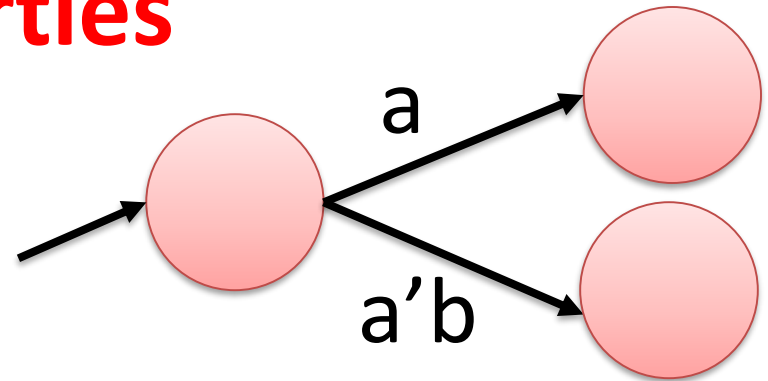


$ab=11$
Next state?

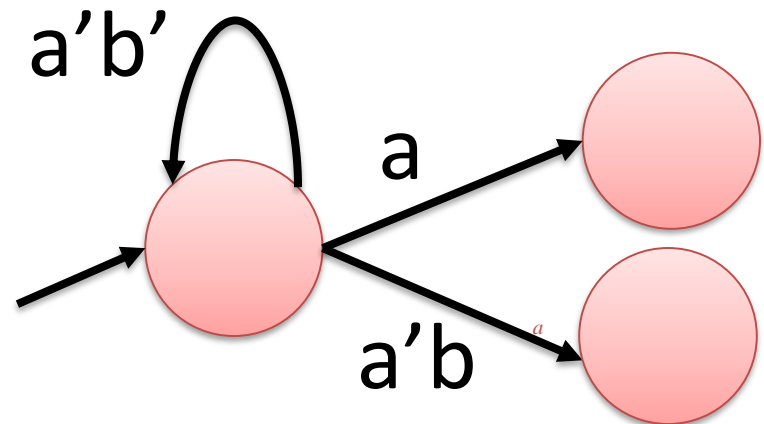


Common Pitfalls Regarding Transition Properties

- *One* condition must be true
 - For all transitions leaving a state
 - Else, where go?

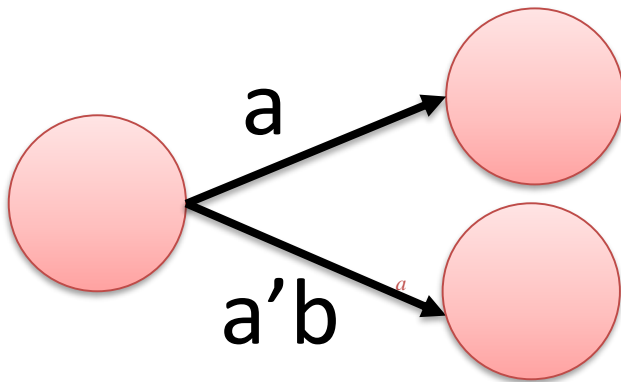


What if $ab=00$?



Verifying Correct Transition Properties

- Can verify using Boolean algebra
 - Only one condition true: AND of each condition pair (for transitions leaving a state) should equal 0
 - \rightarrow proves pair can never simultaneously be true
 - Example

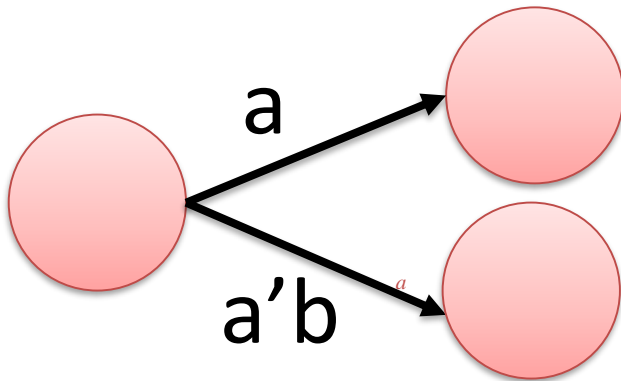


Answer:

$$\begin{aligned} & a * a'b \\ &= (a * a') * b \\ &= 0 * b \\ &= 0 \\ &\text{OK!} \end{aligned}$$

Verifying Correct Transition Properties

- Can verify using Boolean algebra
 - One condition true: OR of all conditions of transitions leaving a state) should equal 1
 - \rightarrow proves at least one condition must be true
 - Example



$$\begin{aligned} & a + a'b \\ &= a*(1+b) + a'b \\ &= a + ab + a'b \\ &= a + (a+a')b \\ &= a + b \end{aligned}$$

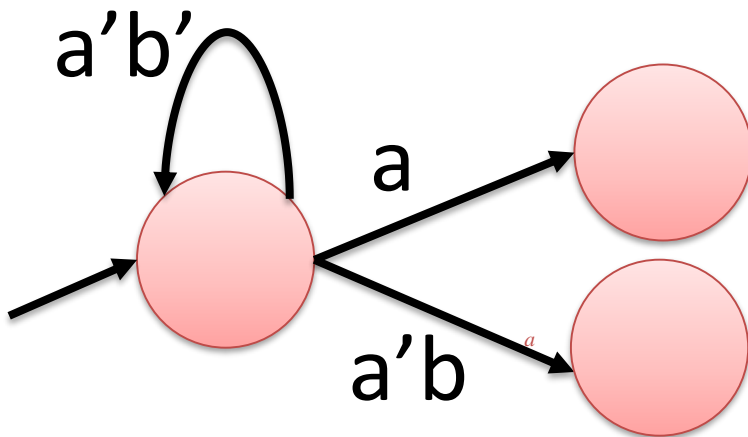
Fails! Might not be 1 (i.e., $a=0, b=0$)

Verifying Correct Transition Properties

- Can verify using Boolean algebra
 - Only one condition true : among all pairs of transition from a state
 - One condition true : All the transitions from a state ^a

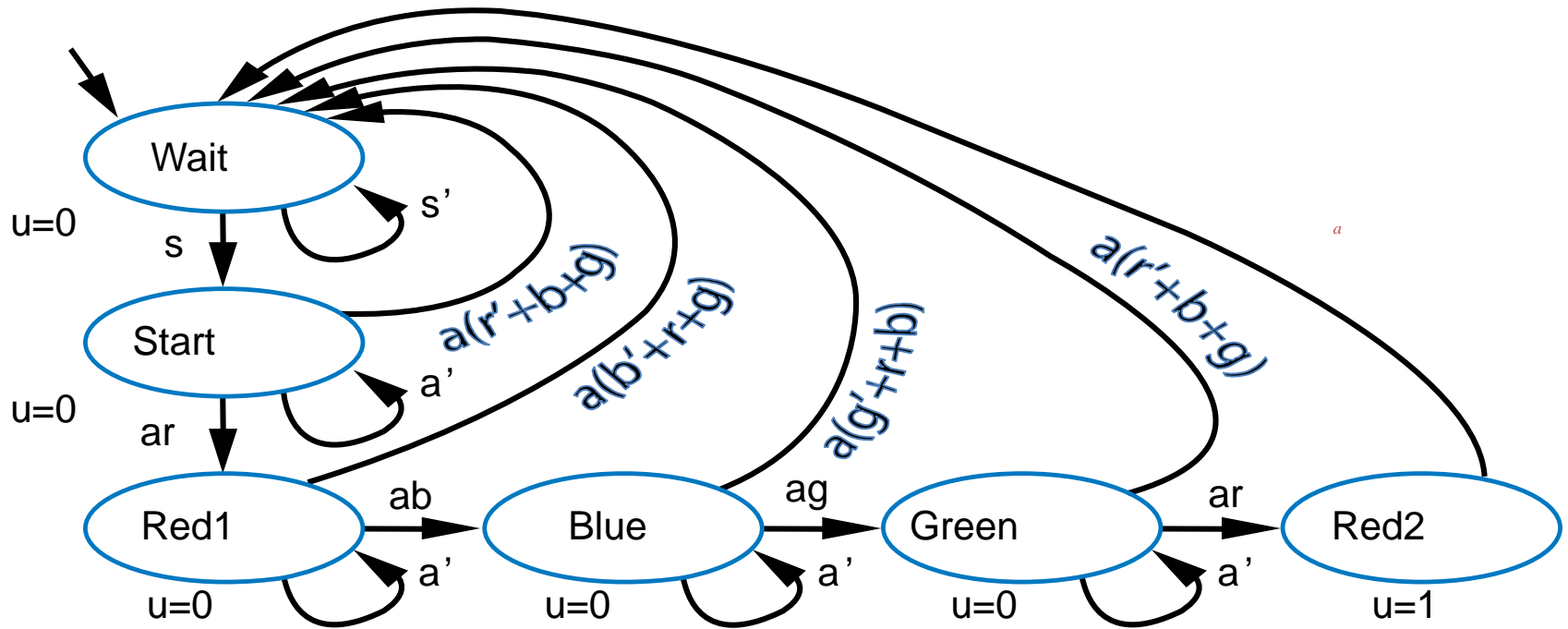
Q: For shown transitions, prove whether:

- * Only one condition true (AND of each pair is always 0)
- * One condition true (OR of all transitions is always 1)



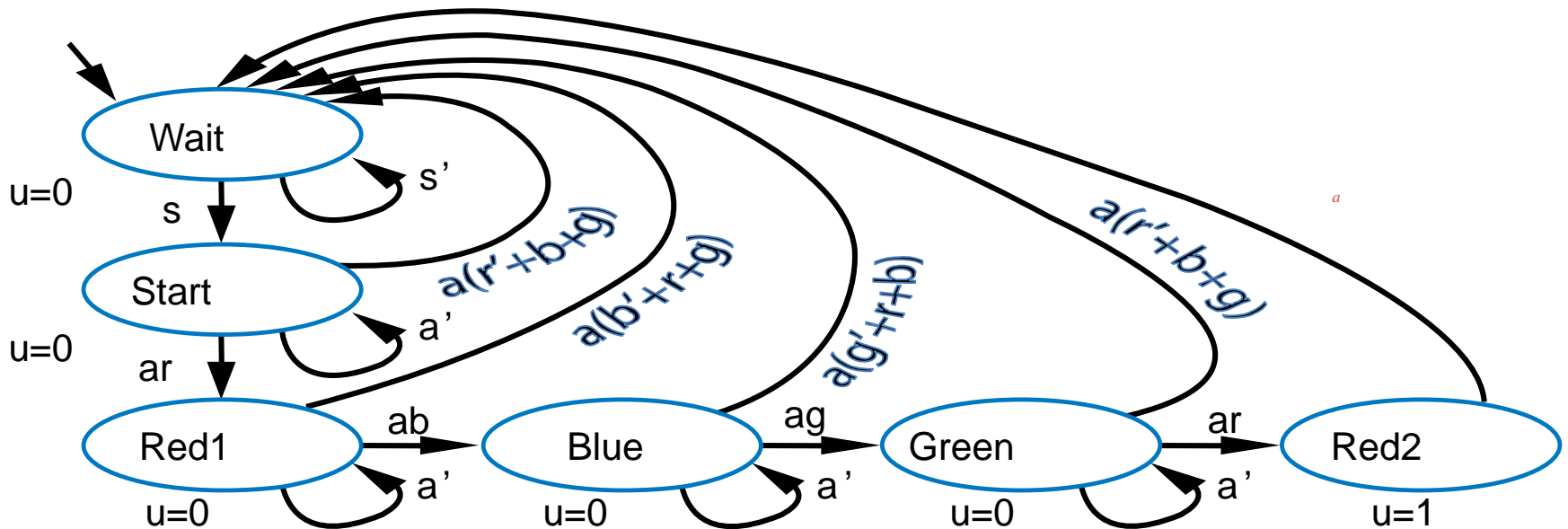
$$a + a'b + a'b' = 1$$

Evidence that Pitfall is Common



- Recall code detector FSM
 - We “fixed” a problem with the transition conditions
 - Do the transitions obey the two required transition properties?

Evidence that Pitfall is Common



Consider transitions of state *Start*, and “only one true” property

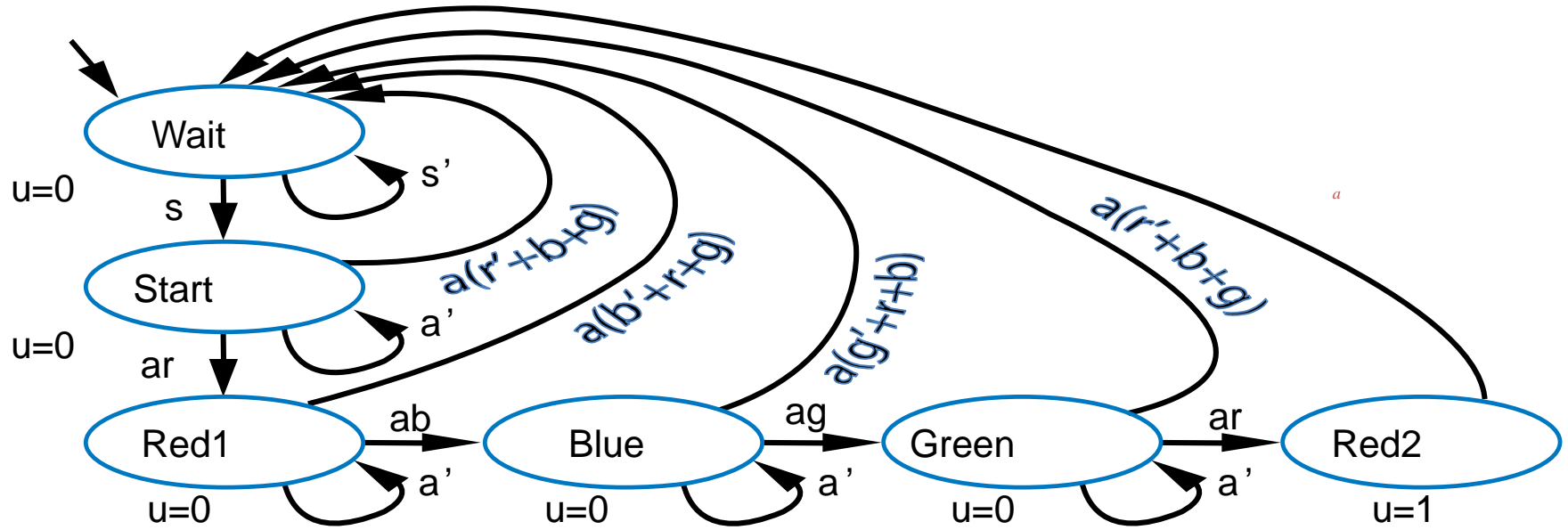
$$\begin{aligned} ar * a' \\ = (a * a')r \\ = 0 \end{aligned}$$

$$\begin{aligned} a' * a(r'+b+g) \\ = 0 * r \\ = 0 \end{aligned}$$

$$\begin{aligned} ar * a(r'+b+g) \\ = (a' * a) * (r' + b + g) &= 0 * (r' + b + g) \\ = (a * a) * r * (r' + b + g) &= * r * (r' + b + g) \\ = arr' + arb + arg \\ = 0 + arb + arg \\ = arb + arg \\ = ar(b+g) \quad // \text{ not ZERO} \end{aligned}$$

Fails! Means that two of Start's transitions could be true

Evidence that Pitfall is Common



Consider transitions of state *Start*, and “only one true” property

Intuitively: press red and blue buttons at same time:
conditions ar , and $a(r'+b+g)$ will both be true.
Which one should be taken?

Q: How to solve?

A: ar should be $arb'g'$
(likewise for ab , ag , ar)

$$arb'g' * a(r'+b+g) = 0$$

Evidence that Pitfall is Common

