

Plan

- Directions of Fea(LPP)
- Extreme directions of Fea(LPP)
- Representation Theorem for Fea(LPP)
- Necessary and sufficient conditions for existence of optimal solutions
- Optimal solutions in atleast one corner point

- $Fea(LPP) = \{\mathbf{x} \in \mathbb{R}^n : A_{m \times n} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}.$
- **Observation 5:** Suppose if a LPP has an **unbounded feasible region**, then there exists a vector $\mathbf{d} \neq \mathbf{0}$ such that starting from any point of the feasible region if you move in the positive direction of \mathbf{d} , then you will always remain in the feasible region.

That is for any $\mathbf{x} \in Fea(LPP)$, $\mathbf{x} + \alpha \mathbf{d} \in Fea(LPP)$ for all $\alpha \geq 0$.

Then $\mathbf{d} \neq \mathbf{0}$ is called a **direction** of $S = Fea(LPP)$.

Throughout our discussion, \mathbf{d} will denote a column vector given by $\mathbf{d} = [d_1, \dots, d_n]^T$.

- **Definition:** Given a non empty convex set S , $S \subset \mathbb{R}^n$, $\mathbf{d} \neq \mathbf{0}$ is called a **direction** of S if for all $\mathbf{x} \in S$, $\mathbf{x} + \alpha \mathbf{d} \in S$ for all $\alpha \geq 0$.
- If \mathbf{d} is a **direction** of a convex set S , then for all $\gamma > 0$, $\mathbf{x} + \alpha \mathbf{d} = \mathbf{x} + \left(\frac{\alpha}{\gamma}\right) \gamma \mathbf{d} \in S$ for all $\alpha > 0$,
 $\Rightarrow \gamma \mathbf{d}$ is again a **direction** for all $\gamma > 0$.

- Two directions $\mathbf{d}_1, \mathbf{d}_2$ of S are said to be **distinct** if $\mathbf{d}_1 \neq \gamma \mathbf{d}_2$ for any $\gamma > 0$
(or equivalently $\mathbf{d}_2 \neq \beta \mathbf{d}_1$ for any $\beta > 0$).
- Example 2: (revisited)** Consider the problem,

$$\text{Min } -x + 2y$$
subject to

$$x + 2y \geq 1$$

$$-x + y \leq 1,$$

$$x \geq 0, y \geq 0.$$

Note that $\mathbf{d}_1 = [1, 1]^T, \mathbf{d}'_1 = [2, 2]^T, \dots$
are all equal as directions of $\text{Fea}(LPP)$.
Similarly $\mathbf{d}_2 = [1, 0]^T, \mathbf{d}'_2 = [2, 0]^T, \dots$
are all equal as directions of $\text{Fea}(LPP)$.
Whereas $\mathbf{d}_1 = [1, 1]^T, \mathbf{d}_2 = [1, 0]^T$ give two **distinct directions**.

- **Result:** The set of all directions of $S = \{\mathbf{x} \in \mathbb{R}^n : A_{m \times n} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$ is given by $D = \{\mathbf{d} \in \mathbb{R}^n : \mathbf{d} \neq \mathbf{0}, A_{m \times n} \mathbf{d} \leq \mathbf{0}, \mathbf{d} \geq \mathbf{0}\}$ or by $D = \{\mathbf{d} \in \mathbb{R}^n : \mathbf{d} \neq \mathbf{0}, \mathbf{a}_i^T \mathbf{d} \leq 0, \text{ for all } i = 1, 2, \dots, m, \mathbf{d} \geq \mathbf{0}\}$.

- **Remark:** Note that the set of all directions of $S = \text{Fea}(LPP)$ is a **convex set**.

So if $\mathbf{d}_1, \mathbf{d}_2$ are two directions of S , then $\mathbf{d} = \alpha \mathbf{d}_1 + \beta \mathbf{d}_2$ will again be a direction of S , for any α, β non negative (as long as both α, β are not equal to zero, or $\alpha + \beta \neq 0$).

- **Definition:** A direction \mathbf{d} of S is called an **extreme direction** of S , if it **cannot** be written as a **positive linear combination** of **two distinct directions** of S , that is, if \mathbf{d} an **extreme direction** of S and

$$\mathbf{d} = \alpha \mathbf{d}_1 + \beta \mathbf{d}_2, \text{ for } \alpha, \beta > 0 \text{ and } \mathbf{d}_1, \mathbf{d}_2 \in D$$

then $\mathbf{d}_1 = \gamma \mathbf{d}_2$ for some $\gamma > 0$.

- If D denotes the set of all directions of S (which will be the empty set if S is bounded) then $D' = \{\mathbf{d} \in \mathbb{R}^n : \mathbf{d} \geq \mathbf{0}, A\mathbf{d} \leq \mathbf{0}, \sum_i d_i = 1\}$ is a set of all **distinct directions** of S .
- Also each $\mathbf{d} \in D$ is of the form $\mathbf{d} = \alpha \mathbf{d}'$ for some $\mathbf{d}' \in D'$ where $\alpha = \sum_i d_i (> 0)$.
- D' can be written as

$$D' = \left\{ \mathbf{d} \in \mathbb{R}^n : \mathbf{d} \geq \mathbf{0}, \begin{bmatrix} 1 & 1, \dots, & 1 \\ -1 & -1, \dots, & -1 \end{bmatrix} \mathbf{d} \leq \begin{bmatrix} 0 \\ 1 \\ -1 \end{bmatrix} \right\}.$$

The set D' now looks exactly like the feasible region of an LPP.

- If D' is **non empty** then D' has at least one **extreme point** (why?).
- **Result:** \mathbf{d} is an extreme direction of S if and only if $\mathbf{d}' = \frac{\mathbf{d}}{\sum_i d_i}$ is an extreme point of D' .
- **Remark:** Hence the number of **distinct extreme directions** of S is **finite** (why?).
- Also since D' is like the set, $Fea(LPP) = S$, if $D' \neq \phi$, then D' must have atleast one **extreme point**.
- Hence if $Fea(LPP) = S$ is **unbounded** then S must have atleast one **extreme direction**.

- The **extreme directions of S** which are **extreme points** of D' (after suitable normalization) will lie on **n LI hyperplanes** defining D' .
- Since $\mathbf{d} \in \mathbb{R}^n$, $\mathbf{d} \neq \mathbf{0}$ cannot be orthogonal to n LI vectors, so \mathbf{d} cannot lie on **n LI hyperplanes** of the $(m + n)$ hyperplanes given by,
 $\{\mathbf{d} \in \mathbb{R}^n : \mathbf{a}_i^T \mathbf{d} = 0\}$ for $i = 1, 2, \dots, m$, and
 $\{\mathbf{d} \in \mathbb{R}^n : -\mathbf{e}_j^T \mathbf{d} = 0\}$ for $j = 1, 2, \dots, n$.
- If $\mathbf{d} \in D'$, is an **extreme direction** of S then it should lie on **$(n - 1)$ LI hyperplanes** of the above mentioned **$(m + n)$ hyperplanes**, and the hyperplane
 $\{\mathbf{d} \in \mathbb{R}^n : [1, 1, \dots, 1] \mathbf{d} = 1\}$
gives a collection of **n LI hyperplanes**, on which \mathbf{d} lies.

- Any $\mathbf{d} \in D$, which lies on $(n - 1)$ LI hyperplanes out of the $(m + n)$ hyperplanes given by $\{\mathbf{d} \in \mathbb{R}^n : \mathbf{a}_i^T \mathbf{d} = 0\}$ for $i = 1, 2, \dots, m$, and $\{\mathbf{d} \in \mathbb{R}^n : -\mathbf{e}_j^T \mathbf{d} = 0\}$ for $j = 1, 2, \dots, n$, is an extreme direction of S .
- Exercise:** Check that if a $\mathbf{d} \in D$ lies on $(n - 1)$ LI hyperplanes (out of the $(m + n)$ defining hyperplanes of D) given by $\{H_1, \dots, H_{n-1}\}$, then $\{H, H_1, \dots, H_{n-1}\}$ is LI where $H = \{\mathbf{d} \in \mathbb{R}^n : [1, 1, \dots, 1]\mathbf{d} = 1\}$.

- **Example 2: (revisited)** Consider the problem,

$$\text{Min } -x + 2y$$

subject to

$$x + 2y \geq 1$$

$$-x + y \leq 1,$$

$$x \geq 0, y \geq 0.$$

Note that here the set of all **directions** of S is given by

- $D = \{\mathbf{d} \in \mathbb{R}^2 : [-1, -2]\mathbf{d} \leq 0, [-1, 1]\mathbf{d} \leq 0, \mathbf{d} \geq \mathbf{0}\}.$

Also if $\mathbf{d} \in D$ is an **extreme direction** of S then it has to lie on exactly one of the hyperplanes given by

- (i) $\{\mathbf{d} \in \mathbb{R}^2 : [-1, -2]\mathbf{d} = 0\},$
- (ii) $\{\mathbf{d} \in \mathbb{R}^2 : [-1, 1]\mathbf{d} = 0\},$
- (iii) $\{\mathbf{d} \in \mathbb{R}^2 : d_1 = 0\},$
- (iv) $\{\mathbf{d} \in \mathbb{R}^2 : d_2 = 0\}.$

- Note that there exists no $\mathbf{d} \geq \mathbf{0}, \mathbf{d} \neq \mathbf{0}$ such that $[-1, -2]\mathbf{d} = 0$.
- Also if $\mathbf{d} \geq \mathbf{0}, \mathbf{d} \neq \mathbf{0}$, satisfies the condition $d_1 = 0$, then $[-1, 1]\mathbf{d} \leq 0$ cannot be satisfied, hence such a \mathbf{d} does not belong to D .
- Hence if $\mathbf{d} \in D$, is an **extreme direction** of S then it lies on either the hyperplane
- $\{\mathbf{d} \in \mathbb{R}^2 : [-1, 1]\mathbf{d} = 0\}$, or in $\{\mathbf{d} \in \mathbb{R}^2 : d_2 = 0\}$.
- Check that $\mathbf{d}' = [1, 1]^T$ and any **positive scalar multiple** of \mathbf{d}' , and $\mathbf{d}'' = [1, 0]^T$ and any **positive scalar multiple** of \mathbf{d}'' , are the only **extreme directions** of the above $S = \text{Fea}(LPP)$.

- **Theorem:**

If $S = \text{Fea}(LPP) = \{\mathbf{x} \in \mathbb{R}^n : A_{m \times n} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$ is nonempty, then S has at least one **extreme point**.

- **Remark:** Note that the above result is **not necessarily true** for all **polyhedral sets**.

For example take any **single half space**, or say a **straight line** in \mathbb{R}^n , which are **polyhedral sets**, but does not have any **extreme point**.

- The theorem works for $\text{Fea}(LPP)$ because of the **non negativity constraints**, that is because $\text{Fea}(LPP)$ is given a supply of **n LI hyperplanes**, among the **$(m+n)$ defining hyperplanes** of S .

- **Exercise:** Can you find a nonempty polyhedral set S , $S \subset \mathbb{R}^3$ which has two defining hyperplanes but does not have any extreme point.
- **Exercise:** Can you find a nonempty polyhedral set S , $S \subset \mathbb{R}^3$ which has three LI defining hyperplanes (not necessarily the nonnegativity constraints) but does not have any extreme point.
- **Definition:** Given S , a nonempty subset of \mathbb{R}^n , and $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k \in S$, $\sum_{i=1}^k \lambda_i \mathbf{x}_i$, is called a convex combination of $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k$, where $0 \leq \lambda_i \leq 1$ for all $i = 1, 2, \dots, k$, and $\sum_{i=1}^k \lambda_i = 1$.

- All possible **convex combinations** of **two distinct points** gives a **straight line segment** with those two points as boundary points.
- All possible convex combinations of **three non colinear points** gives a **triangle** with those points as corner points.
- All possible convex combinations of **four points** no three of which are colinear gives a **quadrilateral**.
- **Result:** Given $\phi \neq S \subset \mathbb{R}^n$, S is a **convex set** if and only if for all $k \in \mathbb{N}$, the **convex combination of any k elements** of S is again an element of S .

● **Theorem: (Representation Theorem)** If

$S = \{\mathbf{x} \in \mathbb{R}^n : A_{m \times n} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$ is nonempty and if $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k$ are the extreme points of S and $\mathbf{d}_1, \mathbf{d}_2, \dots, \mathbf{d}_r$ are the distinct extreme directions of S (the set of directions is empty if S is bounded) then $\mathbf{x} \in S$ if and only if

$$\mathbf{x} = \sum_{i=1}^k \lambda_i \mathbf{x}_i + \sum_{j=1}^r \mu_j \mathbf{d}_j$$

where $0 \leq \lambda_i \leq 1$ for all $i = 1, 2, \dots, k$, $\sum_i \lambda_i = 1$, and $\mu_j \geq 0$, for all $j = 1, 2, \dots, r$.

- That is, $\mathbf{x} \in S \Leftrightarrow \mathbf{x}$ can be written as a convex combination of the extreme points of S plus a non negative linear combination of the extreme directions of S .

- **Observation 6:** If $S = \text{Fea}(LPP)$ is a nonempty bounded set then any $\mathbf{x} \in S$ can be written as a convex combination of the extreme points of S .
- **Observation 7:** Since D' , the set of distinct directions of S (if it is nonempty) is a bounded set ($\mathbf{d} \geq \mathbf{0}$ and $\sum_{i=1}^n d_i = 1$),
so any $\mathbf{d} \in D'$ can be written as a convex combination of the extreme points of D' .
- So any direction $\mathbf{d} \in D$ of S can be written as a nonnegative linear combination of the extreme directions of S .

- **Observation 8:** Note that if there exists a $\mathbf{d} \in D$ such that $\mathbf{c}^T \mathbf{d} < 0$ then the LPP(*)

((*) Min $\mathbf{c}^T \mathbf{x}$, subject to $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$)

does not have an optimal solution.

Since for any given $\mathbf{x} \in S$, $\mathbf{c}^T(\mathbf{x} + \alpha \mathbf{d}) = \mathbf{c}^T \mathbf{x} + \alpha \mathbf{c}^T \mathbf{d}$ can be made smaller than any real M , by choosing $\alpha > 0$ sufficiently large.

- **Exercise:** If $\mathbf{c}^T \mathbf{d}_j \geq 0$ for all extreme directions \mathbf{d}_j of the nonempty and unbounded feasible region S of a LPP, then does it imply that $\mathbf{c}^T \mathbf{d} \geq 0$ for all directions $\mathbf{d} \in D$, of the feasible region S ?

Ans is yes.

- **Observation 9:** From the representation theorem of S we can see that if $S \neq \phi$ and $\mathbf{c}^T \mathbf{d}_j \geq 0$ for all $j = 1, 2, \dots, r$, then LPP(*) has an optimal solution, and atleast one optimal solution is attained at an extreme point of S .

- **Observation 10:** From the representation theorem of S we can also see that if $S = \text{Fea}(LPP)$ is nonempty and bounded then the $LPP(*)$ has an optimal solution and the optimal value is attained in at least one extreme point.

From the above observations we can conclude the following:

- **Conclusion 1:** If $S = \text{Fea}(LPP) \neq \phi$, then the $LPP (*)$ has an optimal solution if and only if one of the following is true:
- (i) $S = \text{Fea}(LPP)$ is bounded (also seen before by using extreme value theorem)
- (ii) $S = \text{Fea}(LPP)$ is unbounded and $\mathbf{c}^T \mathbf{d}_j \geq 0$ for all extreme directions \mathbf{d}_j of the feasible region S .

- **Conclusion 2:** If LPP (*) has an **optimal solution** then there exists an **extreme point** of the feasible region S , which is an **optimal solution**.
- **Exercise:** Give an example of a **nonlinear** function $f : S \rightarrow \mathbb{R}$, where $S \subset \mathbb{R}$ is a **closed** and **bounded polyhedral** subset of \mathbb{R} , (what are these sets?) such that f has a **minimum value** in S but the minimum value is not attained at any **extreme point** of S .
- **Conclusion 3:** If $S = \text{Fea}(LPP)$ is nonempty, and there exists an $M \in \mathbb{R}$ such that for all $\mathbf{x} \in S$, $\mathbf{c}^T \mathbf{x} \geq M$, then the LPP (*) has an optimal solution.

- To understand the significance of the previous result solve the following problems.
- **Exercise:** Give an example of a **linear** function $f : S \rightarrow \mathbb{R}$, where $S \subset \mathbb{R}$ is not a polyhedral subset of \mathbb{R} , such that $f(x) \geq 1$ but f **does not** have a minimum value in S .
- **Exercise:** Give an example of a **nonlinear** function $f : S \rightarrow \mathbb{R}$, where $S \subset \mathbb{R}$ is a **polyhedral** subset of \mathbb{R} , such that $f(x) \geq 1$ but f **does not** have a minimum value in S .

- We can come to similar conclusions if we consider a linear programming problem, LPP(**) as
 (**) $\text{Max } \mathbf{c}^T \mathbf{x}$
 subject to $\mathbf{Ax} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$.
- **Conclusion 1a:** If $S = \text{Fea}(LPP) \neq \phi$, then the LPP (**) has an optimal solution if and only if one of the following is true:
 - (i) $S = \text{Fea}(LPP)$ is bounded
 - (ii) $S = \text{Fea}(LPP)$ is unbounded and $\mathbf{c}^T \mathbf{d}_j \leq 0$ for all extreme directions \mathbf{d}_j of the feasible region S .
- **Conclusion 2a:** If a LPP (**) has an optimal solution then there exists an extreme point of the feasible region S , which is an optimal solution.
- **Conclusion 3a:** If $S = \text{Fea}(LPP)$ is nonempty, and there exists an $M \in \mathbb{R}$ such that for all $\mathbf{x} \in S, \mathbf{c}^T \mathbf{x} \leq M$, then the LPP (**) has an optimal solution.