

Binary Codes

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- Source:

Chapter 1 of Z. Kohavi and N. Jha, Switching and Finite Automata Theory, 3rd Ed., Cambridge University Press, 2010

Binary codes

- To simplify the problem of communication between human and machine, several codes have been devised in which decimal digits are represented by sequences of binary digits.
 - Weighted Code
 - BCD
 - Self-complementing code
 - Non-weighted code
 - Excess-3 code
 - Cycle code
 - Gray code
 - Reflected code

Weighted codes

- Each binary digit is assigned a decimal “weight,” and, for each group of four bits, the sum of the weights of those binary digits whose value is 1 is equal to the decimal digit which they represent.
- If w_1, w_2, w_3 , and w_4 are the given weights of the binary digits and x_1, x_2, x_3, x_4 the corresponding digit values then the decimal digit
- $N = w_4 x_4 + w_3 x_3 + w_2 x_2 + w_1 x_1$ can be represented by the binary sequence $x_4 x_3 x_2 x_1$.
- The sequence of binary digits that represents a decimal digit is called a code word.

Self-complementing codes

- It is apparent that the representations of some decimal numbers in the (2, 4, 2, 1) and (6, 4, 2, -3) codes are not unique.
 - For example, in the (2, 4, 2, 1) code, decimal 7 may be represented by 1101 as well as 0111.
- Adopting the representations shown in Table causes the codes to become self-complementing.
- A code is said to be *self-complementing* if the code word of the “9’s complement of N ”, i.e., $9 - N$, can be obtained from the code word of N by interchanging all the 1’s and 0’s.
 - For example, in the (6, 4, 2, -3) code, decimal 3 is represented by 1001 while decimal 6 is represented by 0110.

Self-complementing codes

- BCD code (8, 4, 2, 1) is not self-complementing.
- A necessary condition for a weighted code to be self-complementing is that the sum of the weights must equal 9.
 - There exist only four positively weighted self-complementing codes, namely, (2, 4, 2, 1), (3, 3, 2, 1), (4, 3, 1, 1), and (5, 2, 1, 1).
 - there exist 13 self-complementing codes with positive and negative weights.

Nonweighted codes

- There are many nonweighted binary codes, such as Excess-3 Code, Cyclic Code, Gray Code etc..
- The Excess-3 code is formed by adding 0011 to each BCD code word.

Decimal digit	Excess-3				Cyclic			
0	0	0	1	1	0	0	0	0
1	0	1	0	0	0	0	0	1
2	0	1	0	1	0	0	1	1
3	0	1	1	0	0	0	1	0
4	0	1	1	1	0	1	1	0
5	1	0	0	0	1	1	1	0
6	1	0	0	1	1	0	1	0
7	1	0	1	0	1	0	0	0
8	1	0	1	1	1	1	0	0
9	1	1	0	0	0	1	0	0

Cyclic codes

- Code word having property of **successive decimal integers differ in only one digit** are referred to as cyclic codes.
- Gray code also have property of successive decimal integers differ by 1 bit, so it is also a cyclic code.
- Code useful is the simplicity of the procedure for converting from the binary number system into the Gray code

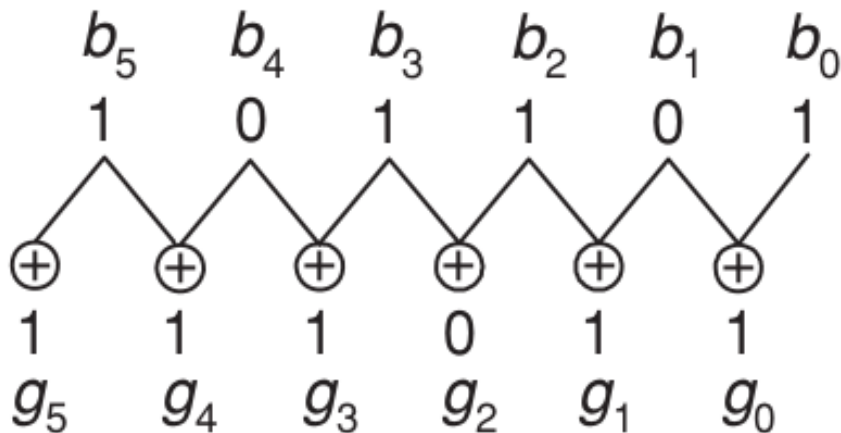
Gray Code

Table 1.5 Decimal numbers in the complete four-bit Gray code and in binary

Decimal number	Gray				Binary			
	g_3	g_2	g_1	g_0	b_3	b_2	b_1	b_0
0	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	1
2	0	0	1	1	0	0	1	0
3	0	0	1	0	0	0	1	1
4	0	1	1	0	0	1	0	0
5	0	1	1	1	0	1	0	1
6	0	1	0	1	0	1	1	0
7	0	1	0	0	0	1	1	1
8	1	1	0	0	1	0	0	0
9	1	1	0	1	1	0	0	1
10	1	1	1	1	1	0	1	0
11	1	1	1	0	1	0	1	1
12	1	0	1	0	1	1	0	0
13	1	0	1	1	1	1	0	1
14	1	0	0	1	1	1	1	0
15	1	0	0	0	1	1	1	1

Binary code to gray code conversion

- Let $g_n \cdots g_2 g_1 g_0$ denote a code word in the $(n + 1)^{\text{th}}$ -bit Gray code, and let $b_n \cdots b_2 b_1 b_0$ designate the corresponding binary number.



$$g_i = b_i \oplus b_{i+1}, \quad 0 \leq i \leq n - 1, \\ g_n = b_n,$$

Gray Code to Binary

- Start with the leftmost digit and proceed to the least significant digit, setting $bi = gi$ if the number of 1's preceding gi is even and setting $bi = gi'$ if the number of 1's preceding gi is odd
- Gray code word 1001011 represents the binary number 1110010

Reflected Code

- The term “reflected” is used to designate codes which have the property that the n -bit code can be generated by reflecting the $(n - 1)$ th-bit code.
- Three-bit Gray code can be obtained by reflecting the two-bit code about an axis at the end of the code and assigning a most significant bit of 0 above the axis and 1 below the axis.

00	0	00	0	000
01	0	01	0	001
11	0	11	0	011
10	0	10	0	010
<hr/>			1	100
	1	10	1	101
	1	11	1	111
	1	01	1	110
	1	00	1	1010
<hr/>			1	011
			1	001
			1	000