System Calls for Process Management

### Process Creation – The fork() System Call

- To Create a child process
- The process that invoked the system call Parent Process
- Newly Created Process Child Process

$$Pid = fork()$$

# The fork() System Call

- Will not take any arguments
- Return value
  - Negative Error in process Creation
  - Positive Created successfully (to parent it is PID of child)
  - Zero Child created (In child it is zero )

# Kernel Operation – on fork()

- It allocates a slot in the process table for new process
- Assigns a unique ID for child process
- Makes a logical copy of the context of the parent process
- Returns ID number of newly created process to parent and 0 to the child.

#### Example

```
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main()
  /* fork a process */
  fork();
  /* the child and parent will execute every line of code after the fork (each separately)*/
  printf("Hello world!\n");
  return 0;
```

```
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main()
  /* fork a process */
  int pid;
  pid= fork();
  /* the child and parent will execute every line of code after the fork (each separately)*/
  if (pid==0)
     printf("Hi I am Child - pid = %d \n", getpid());
     printf("Hi My parent id is = %d\n ", getppid());
  else if (pid>0)
   printf("Parent Process - pid = %d\n", getpid());
   printf("Hello world!\n");
```

### Exec() System Call

- The exec system call is used to execute a file which is residing in an active process.
- When exec is called the previous executable file is replaced and new file is executed.
- Replaces the old file or program from the process with a new file or program. The entire content of the process is replaced with a new program.

#### Kernel Operation – on exec()

- Current process image is overwritten with a new process image.
- New process image is the one you passed as exec argument
- The currently running process is ended
- New process image has same process ID, same environment, and same file descriptor (because process is not replaced process image is replaced)
- The CPU stat and virtual memory is affected. Virtual memory mapping of the current process image is replaced by virtual memory of new process image.

# Exec() System Call

- Takes the filename & path as the argument
- Return value +ive integer successful
- -ive Error

# Exit()- System Call

- Terminates the process normally
- Status is returned to parent program
  - 0 Successful
  - Non Zero Error

# Exit()- System Call

- exit() performs following operations.
  - \* Flushes unwritten buffered data.
  - \* Closes all open files.
  - \* Removes temporary files.
  - \* Returns an integer exit status to the operating system.