

Relational Model

Definition



The **Relational Model** is the foundation of modern databases and was introduced by **E.F. Codd** in 1970.



It organizes data into **relations**, which are represented as **tables**. Each table consists of **rows (tuples)** and **columns (attributes)**.

Key Features

Simplicity

- Data is stored in a tabular format, making it easy to understand

Mathematical Foundation

- Based on set theory

Data Independence

- Logical and physical data are separated, ensuring flexibility

Key Concepts

Domain

Attribute

Tuple

Relation

Domain

A **domain** is the set of permissible values for an attribute

Think of it as the **data type** or the range of acceptable values

Examples

- A StudentID might have a domain of integers from 1 to 9999
- A DateOfBirth might have a domain of valid dates

Attribute

An **attribute** is a **column** in a table and represents a specific property of the entity

Examples

name, salary, joining_date can be termed as attributes of an **employee**

author, genre, price can be termed as attributes of a **book**

brand, RAM, processor can be termed as attributes of a **PC**

Tuple

A **tuple** is a single **row** in a table, representing a single record

Examples

A tuple in a student table can be (1, 'Alice', 17, 'CSE', 'AEC')

A tuple in a book table can be (174, 'Wings Of Fire', 'Autobiography', 1500)

A tuple in a item table can be (12, 'Laptop', 55000)

Relation

A **relation** is a **table** consisting of tuples (rows) and attributes (columns)

It's characterized by

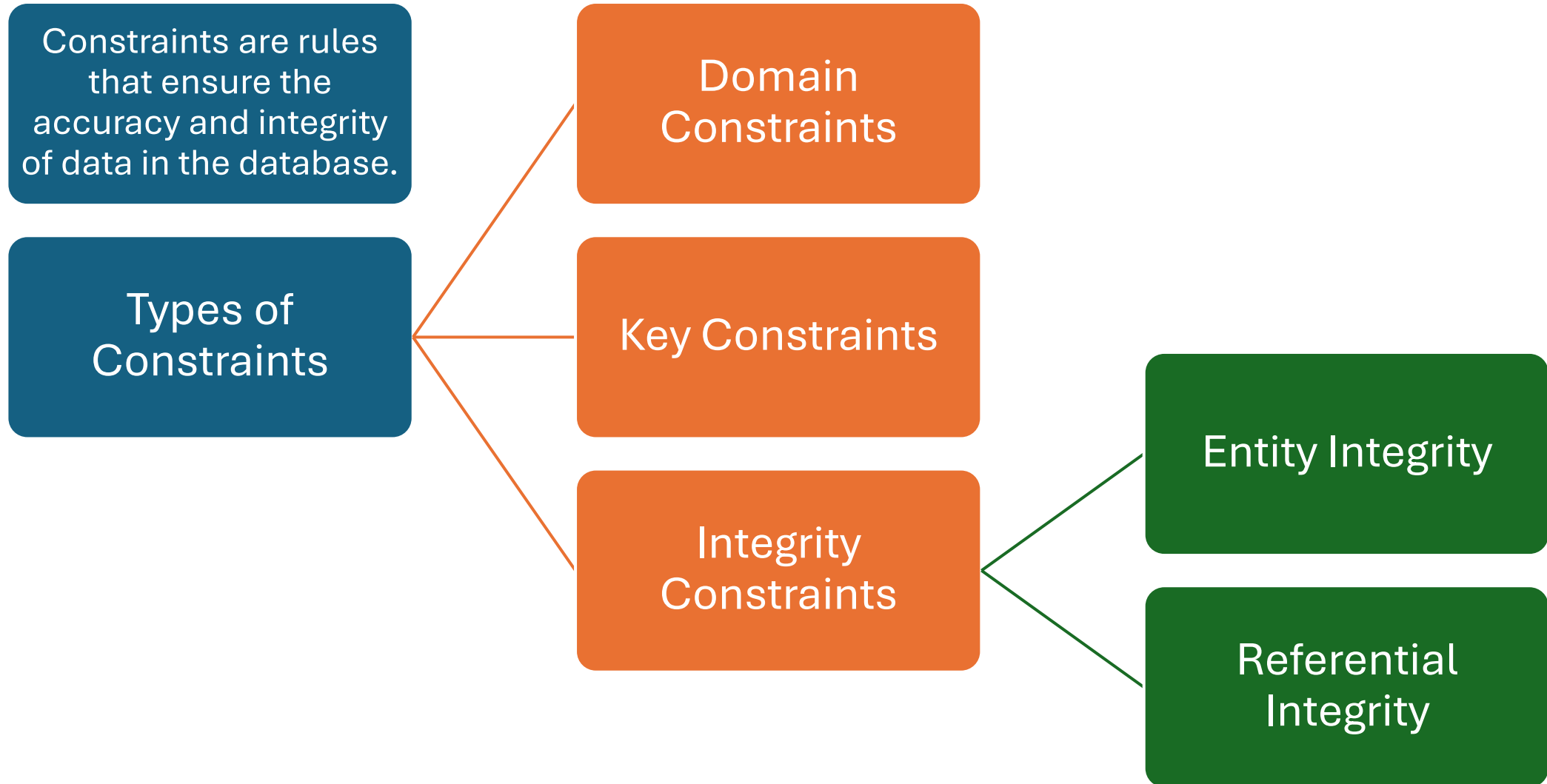
A unique name

Each tuple being distinct

Attribute values conforming to their domain

Constraints in Relational Model

Constraints in Relational Model



Domain Constraints

Ensure that attribute values fall within the specified domain

Examples

- If the **age** attribute has a domain of integers from 18 to 60, any value outside this range is invalid (Can be enforced with **CHECK()**)
- Ensuring **joining_date** attribute of an employee is greater than or equals to the **date_of_establishment** of organization
- Ensuring **grade** attribute of a student falls under a valid grade of **1 to 10**
- Ensuring **rating** attribute of a movie will never have a value less than 0 and greater than 10

Key Constraints

Ensure that each tuple in a relation is uniquely identifiable

Candidate Key: A set of attributes that are capable uniquely identifying a tuple

Primary Key: One chosen candidate key

Examples

- A set of **candidate keys** to identify an employee can be
 - {emp_id, aadhaar_number, PAN, email (If set to UNIQUE)}
- One of these set of candidate keys can be chosen as a **Primary Key** upon developers choice

Integrity Constraints

Entity Integrity

- Primary Key attributes cannot have null values
- Ensuring every tuple is uniquely identifiable

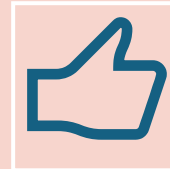
Referential Integrity

- Maintains consistency between two related tables
- Often achieved through **foreign key** that refers to a **primary key** or any other **unique** valued attribute from another table
- Ensuring only existing user from **users** table can make **orders**
- Ensuring only existing student from **students** table able to **enroll** into an existing course from **courses** table

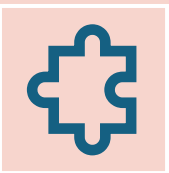
Real World Analogy



Domain: Defines acceptable inputs, e.g., roll numbers should be integers, and names should be strings, age cannot be more than 100 and so on



Attributes: Represent student details like roll number, name, and branch



Tuples: Each student is a unique row



Constraints: Ensure data is consistent. For instance, a course enrollment record must refer to a valid student and course

Relational Algebra

Relational Algebra (Procedural Query Language)

Definition: Relational Algebra is a **formal query language** that is used in relational databases

Relational Algebra uses a set of operations to manipulate relations (tables) to produce desired results

Basic Symbols (Operators) in Relation Algebraic Expressions

Operator Name	Symbol Used	Description
Selection	σ	Filters rows based on a condition
Projection	π	Selects specific columns from a table
Union	\cup	Combines tuples from two relations (tables) without duplicates
Intersection	\cap	Retrieves common tuples from two relations
Difference	$-$	Retrieves tuples in one relation but not in the other
Cartesian Product	\times	Combines tuples from two relations (tables) in every possible way
Join	\bowtie	Combines tuples from two relations based on a condition
Rename	ρ	Renames a relation or it's attributes
Grouping / Aggregation	γ	Used to represent groups and aggregations
Ordering	τ	Used to perform ordering of projected data
AND, OR, NOT	\wedge, \vee, \neg	Combines conditions in Selection or Join

LaTeX Commands for the Symbols used in Relation Algebra

Operator Name	Symbol Used	Latex Command
Selection	σ	<code>\sigma_{condition}</code>
Projection	π	<code>\pi_{attributes}</code>
Union	\cup	<code>\cup</code>
Intersection	\cap	<code>\cap</code>
Difference	$-$	<code>-</code>
Cartesian Product	\times	<code>\times</code>
Join	\bowtie	<code>\bowtie</code>
Rename	ρ	<code>\rho_{new_name}</code>
Grouping / Aggregation	γ	<code>\gamma_{group_by, agg}</code>
Ordering	τ	<code>\tau_{attributes}</code>
AND, OR, NOT	\wedge, \vee, \neg	<code>\land, \lor, \neg</code>